

IIT School of Applied Technology

ILLINOIS INSTITUTE OF TECHNOLOGY

information technology & management

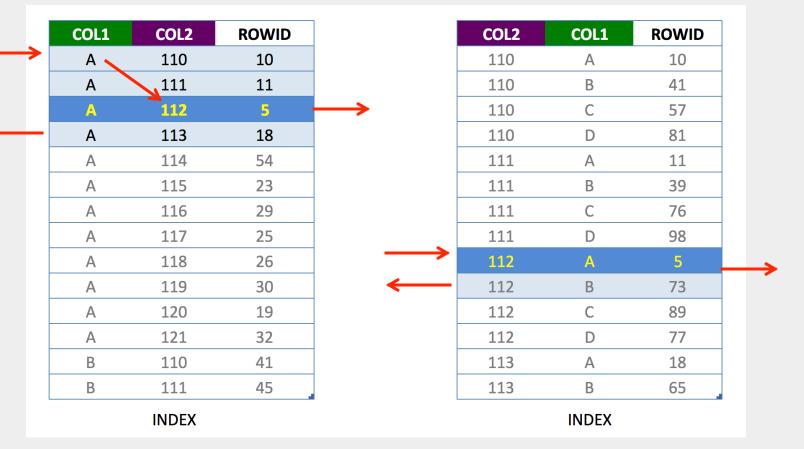
526 Data Warehousing

April 27, 2016

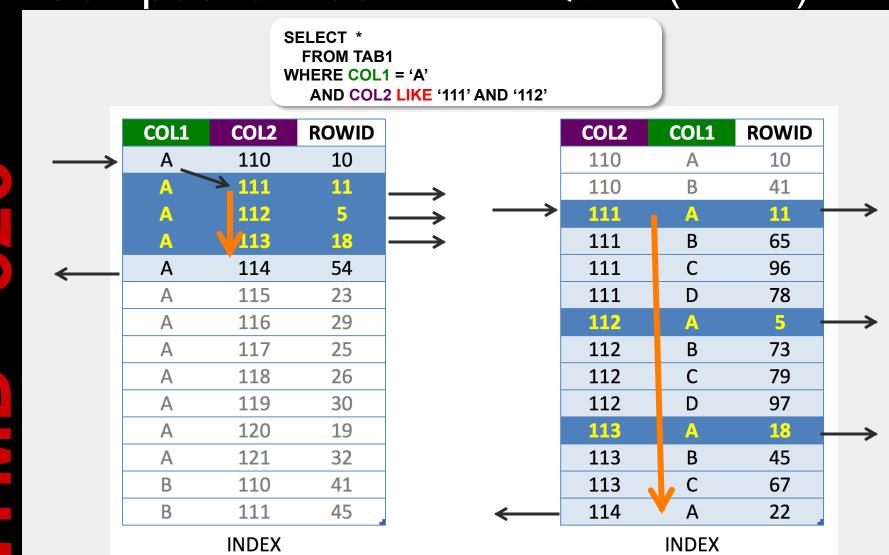
Week 14 Presentation

Composite Index with EQUAL

SELECT *
FROM TAB1
WHERE COL1 = 'A' AND COL2 = '112'



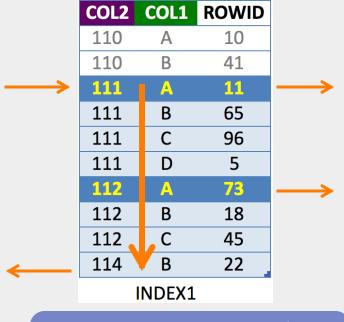
Composite Index with EQUAL (cont'd)



Composite Index: BETWEEN vs. IN

SELECT * FROM TAB1
WHERE COL1 = 'A'
AND COL2 between '111' AND '112'

SELECT * FROM TAB1 WHERE COL1 = 'A' AND COL2 in ('111', '112')



110 10 110 41 111 11 111 B 65 111 96 111 112 73 112 В 18 112 45 22 114 INDFX1

COL2 COL1 ROWID

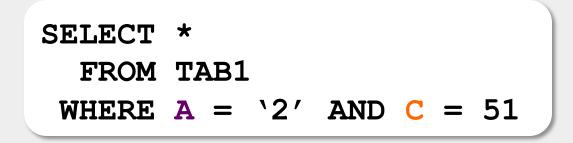
TABLE ACCESS BY ROWID TAB1

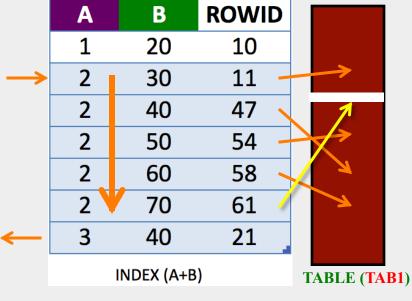
INDEX RANGE SCAN INDEX1

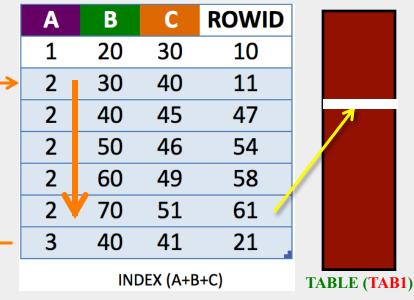
CONCATENATION

TABLE ACCESS BY ROWID TAB1
INDEX RANGE SCAN INDEX1
TABLE ACCESS BY ROWID TAB1
INDEX RANGE SCAN INDEX1

Composite Index: Adding a Column







Composite Index: Adding a Column (cont'd)

```
SELECT *
FROM Sales
WHERE zip = '60201'
AND dt = '20141005'
AND st = 'DELIVERED'
```



zip + dt

Index02

st

Index01

zip + dt + col

Index02

st

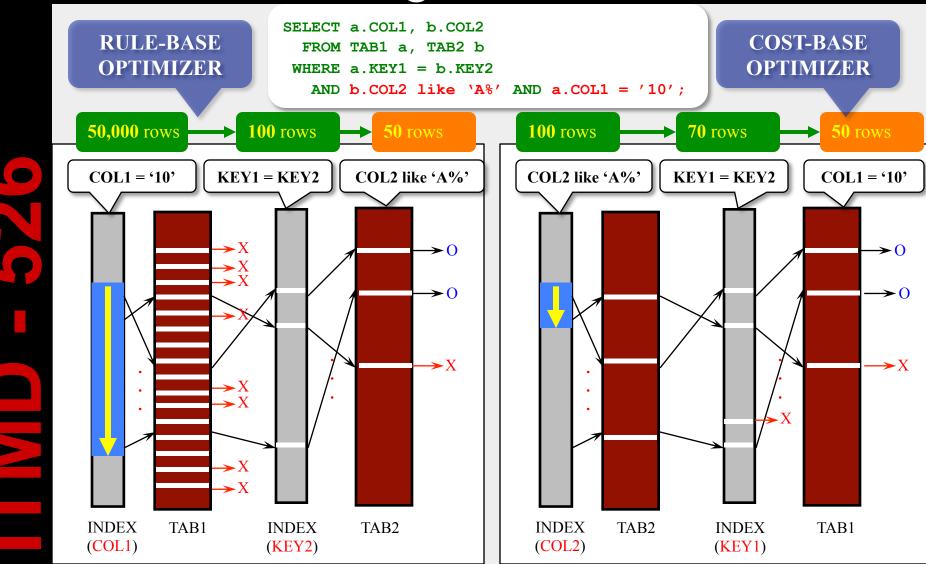


Optimizer's Pick

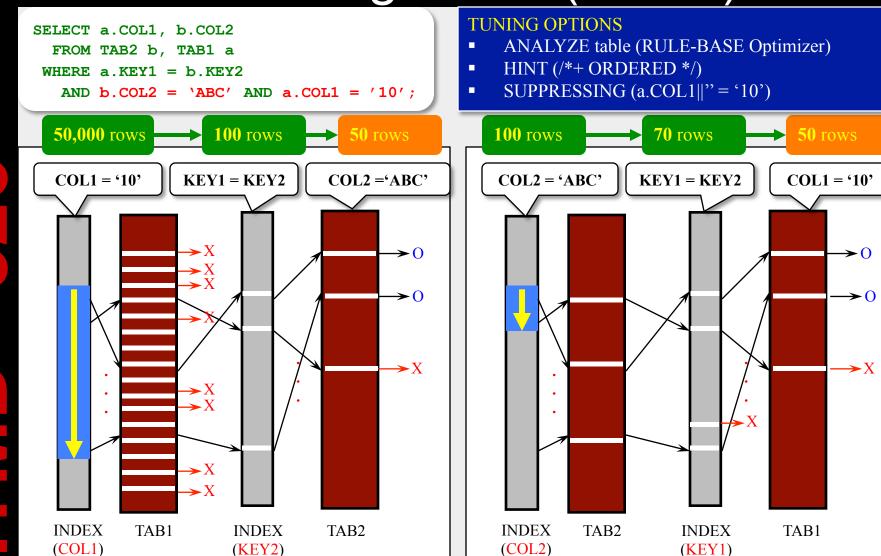
Bad



SQL Optimization for DW JOIN — Driving Table



JOIN - Driving Table (cont'd)



JOIN: Driving Table Example

```
SELECT INVNO, INVDATE, CUSTNAME
FROM CUSTOMER Y, INVOICE X
WHERE X.CUSTNO = Y.CUSTNO
AND X.INVDATE = '20151013'
AND Y.NATION = 'USA'
```

1 rows, 0.14 sec

NESTED LOOPS

TABLE ACCESS BY ROWID CUSTOMER
INDEX RANGE SCAN IX_NATION
TABLE ACCESS BY ROWID INVOICE
AND-EQUAL

INDEX RANGE SCAN IX_CUSTNO
INDEX RANGE SCAN IX_INVDATE

```
SELECT INVNO, INVDATE, CUSTNAME
```

FROM INVOICE X, CUSTOMER Y

WHERE X.CUSTNO = Y.CUSTNO

AND X.INVDATE = '20151013'

AND Y.NATION = 'USA'

1 rows, 0.03 sec

NESTED LOOPS

TABLE ACCESS BY ROWID INVOICE
INDEX RANGE SCAN IX_INVDATE
TABLE ACCESS BY ROWID CUSTOMER
INDEX UNIQUE SCAN IX_CUSTNO

JOIN: Use of Index

SELECT INVNO, INVDATE, CUSTNAME
FROM CUSTOMER Y, INVOICE X
WHERE X.CUSTNO = Y.CUSTNO
AND X.INVDATE = '20151013'
AND Y.NATION = 'USA'

1 rows, 0.14 sec

NESTED LOOPS

TABLE ACCESS BY ROWID CUSTOMER
INDEX RANGE SCAN IX_NATION
TABLE ACCESS BY ROWID INVOICE
AND-EQUAL
INDEX RANGE SCAN IX_CUSTNO
INDEX RANGE SCAN IX_INVDATE

SELECT INVNO, INVDATE, CUSTNAME FROM CUSTOMER Y, INVOICE X

WHERE RTRIM(X.CUSTNO) = Y.CUSTNO

AND X.INVDATE = '20151013'

AND Y.NATION = 'USA'

1 rows, 0.14 sec

NESTED LOOPS

TABLE ACCESS BY ROWID INVOICE
INDEX RANGE SCAN IX_INVDATE
TABLE ACCESS BY ROWID CUSTOMER
INDEX UNIQUE SCAN PK_CUSTNO

SELECT INVNO, INVDATE, CUSTNAME FROM CUSTOMER Y, INVOICE X

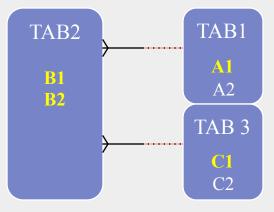
WHERE RTRIM(X.CUSTNO) = RTRIM(Y.CUSTNO)

AND X.INVDATE = '20151013'

AND Y.NATION = 'USA'

rows, sec 7

JOIN: Deciding Access Path



```
SELECT A1,A2,.., B1,B2,..,C1,C2,..
FROM TAB1 t1, TAB2 t2, TAB3 t3
WHERE t1.A1 = t2.B1
  AND t3.C1 = t2.B2
  AND t1.A2 = '10'
  AND t2.B2 LIKE 'WB%'
```

Driving	Access Path	Index Strategy
TAB1	A2 = '10' B1 = A1 and B2 like 'WB%' C1 = B2	A2B1 + B2C1
TAB2	B2 like 'WB%' C1 = B2 A1 = B1 and A2 = '10'	B2C1A1 + A2
TAB3	Full table scan B2 = C1 and B2 like 'WB%' A1 = B1 and A2 = '10'	B2A1 + A2

JOIN: Deciding Access Path (cont'd)

```
SELECT A1,A2,.., B1,B2,..,C1,C2,..

FROM TAB1 t1, TAB2 t2, TAB3 t3

WHERE t1.A1 = t2.B1

AND t3.C1 = t2.B2

AND t1.A2 = '10'

AND t2.B2 LIKE 'WB%'
```

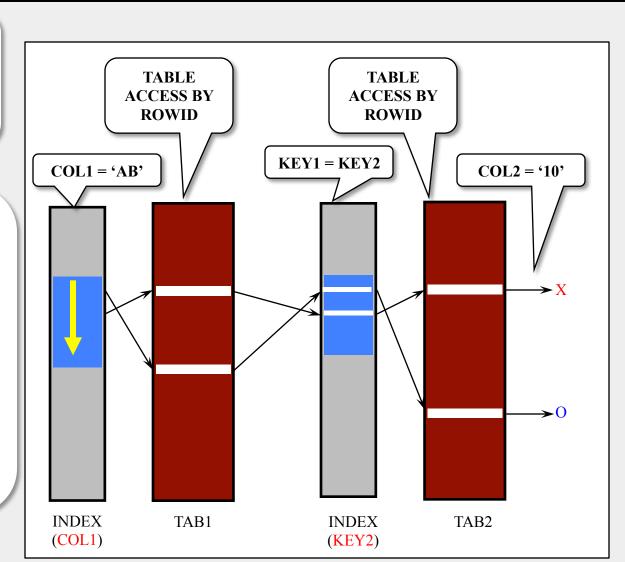
```
SELECT A1,A2,.., B1,B2,..,C1,C2,..
FROM TAB1 t1, TAB2 t2, TAB3 t3
WHERE t1.A1 = t2.B1
   AND t3.C1 = t2.B2
   AND t1.A2 = '10'
AND t3.C1 LIKE 'WB%'
```

Driving	Access Path	Index Strategy
TAB3	Full table scan B2 = C1 and B2 like 'WB%' A1 = B1 and A2 = '10'	• B2 • A1 + A2
TAB3	C1 like 'WB%' C1 = B2 A1 = B1 and A2 = '10'	C1B2A1 + A2

SQL Optimization for DW JOIN — Nested Loop JOIN

```
SELECT a.COL1, b.COL2
FROM TAB1 a, TAB2 b
WHERE a.KEY1 = b.KEY2
AND a.COL1 = 'AB'
AND b.COL2 = '10';
```

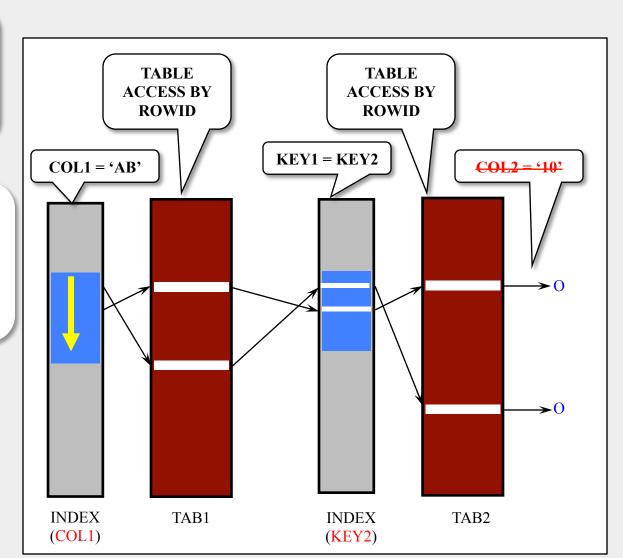
- Sequential
- Partial Range Scan
- Dependency
- Random Table I/O
- Fit for small range
- Indexing on join is important
- •OLTP



JOIN - Nested Loop JOIN (cont'd)

```
SELECT a.COL1, b.COL2
FROM TAB1 a, TAB2 b
WHERE a.KEY1 = b.KEY2
AND a.COL1 = 'AB'
AND b.COL2 = '10';
```

Eliminating b.col2='10' condition has positive impact to the response time

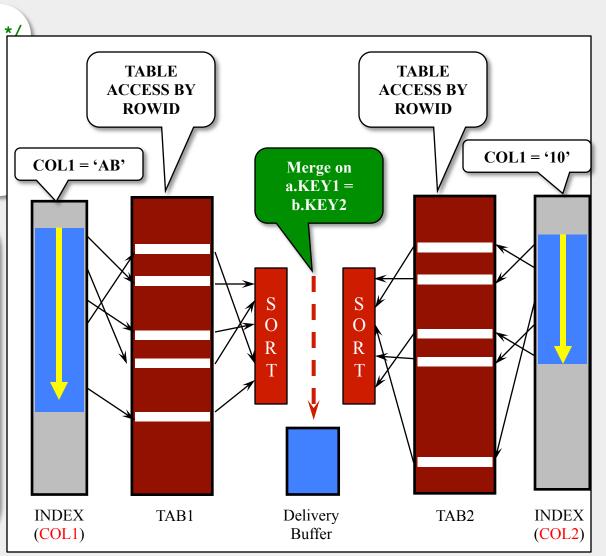


SQL Optimization for DW JOIN — Sort Merge JOIN

```
SELECT /*+ use merge(a b)
a.COL1, b.COL2, ...
FROM TAB1 a, TAB2 b
WHERE a.KEY1 = b.KEY2
AND a.COL1 = 'AB'
AND b.COL2 = '10';
```

- Parallel
- Full Range Scan
- Constrain on both tables needed
- Scan access on sort
- Indexing on join is not important
- Fit for wide range

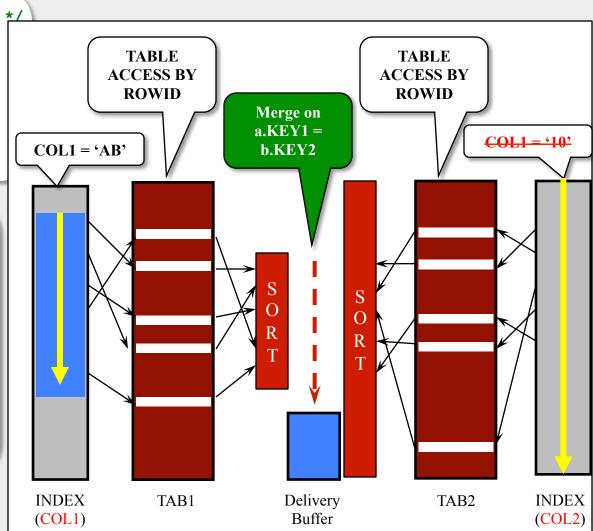
• DW



JOIN - Sort Merge JOIN (cont'd)

```
SELECT /*+ use_merge(a b)
a.COL1, b.COL2, ...
FROM TAB1 a, TAB2 b
WHERE a.KEY1 = b.KEY2
AND a.COL1 = 'AB'
AND b.COL2 = '10';
```

Sort merge becomes ineffective when one of the sort buffer size is much bigger than the other → Hash merge join can be an option



JOIN – Execution Plan Examples

> Throughput

```
SELECT a.COL1, b.COL2

FROM TAB1 a, TAB2 b

WHERE a.KEY1 = b.KEY2

AND a.COL1 = 'AB'

AND b.COL2 = '10';
```

MERGE JOIN (SORT)

TABLE ACCESS BY ROWID TAB1
INDEX RANGE SCAN IX_COL1
TABLE ACCESS BY ROWID TAB2
INDEX RANGE SCAN IX_COL2

> Response time

```
SELECT --+ RULE
a.COL1, b.COL2
FROM TAB1 a, TAB2 b
WHERE a.KEY1 = b.KEY2
AND a.COL1 = 'AB'
AND b.COL2 = '10';
```

NESTED LOOPS

TABLE ACCESS BY ROWID TAB1
INDEX RANGE SCAN IX_COL1
TABLE ACCESS BY ROWID TAB2
INDEX UNIQUE SCAN IX_KEY2

SQL Optimization for DW (cont'd) JOIN – Execution Plan Examples

> Throughput

```
SELECT b.COL3, SUM(...)

FROM TAB1 a, TAB2 b

WHERE a.KEY1 = b.KEY2

GROUP BY b.COL3;

SORT GROUP BY

MERGE JOIN

TABLE ACCESS FULL TAB1

SORT JOIN

TABLE ACCESS FULL TAB2
```

> Response time

```
SELECT --+ RULE

b.COL3, SUM(...)

FROM TAB1 a, TAB2 b

WHERE a.KEY1 = b.KEY2

GROUP BY b.COL3;

SORT GROUP BY

NESTED LOOPS

TABLE ACCESS FULL TAB1

TABLE ACCESS BY ROWID TAB2

INDEX RANGE SCAN KEY2
```