CryptDB: Processing Queries on an Encrypted Database

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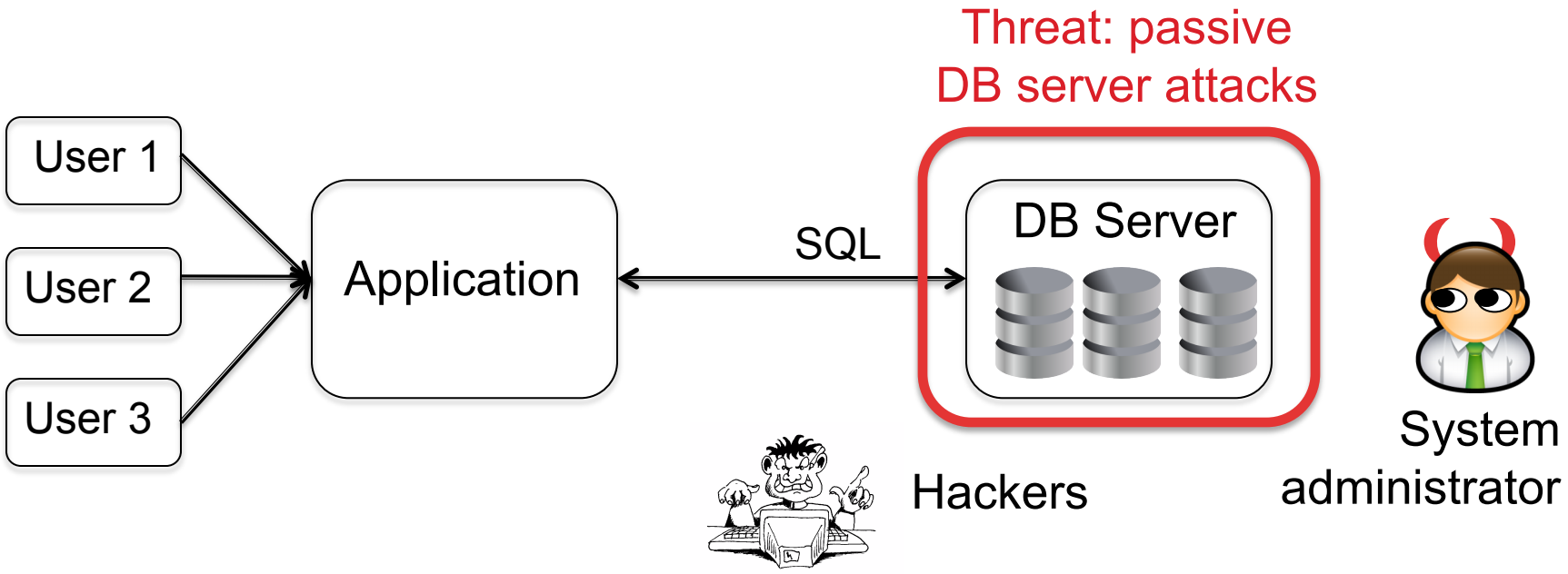
Problem

- Confidential data leaks from databases (DB)
  - 2012: hackers extracted 6.5 million hashed passwords from the DB of LinkedIn



- Process SQL queries on encrypted data

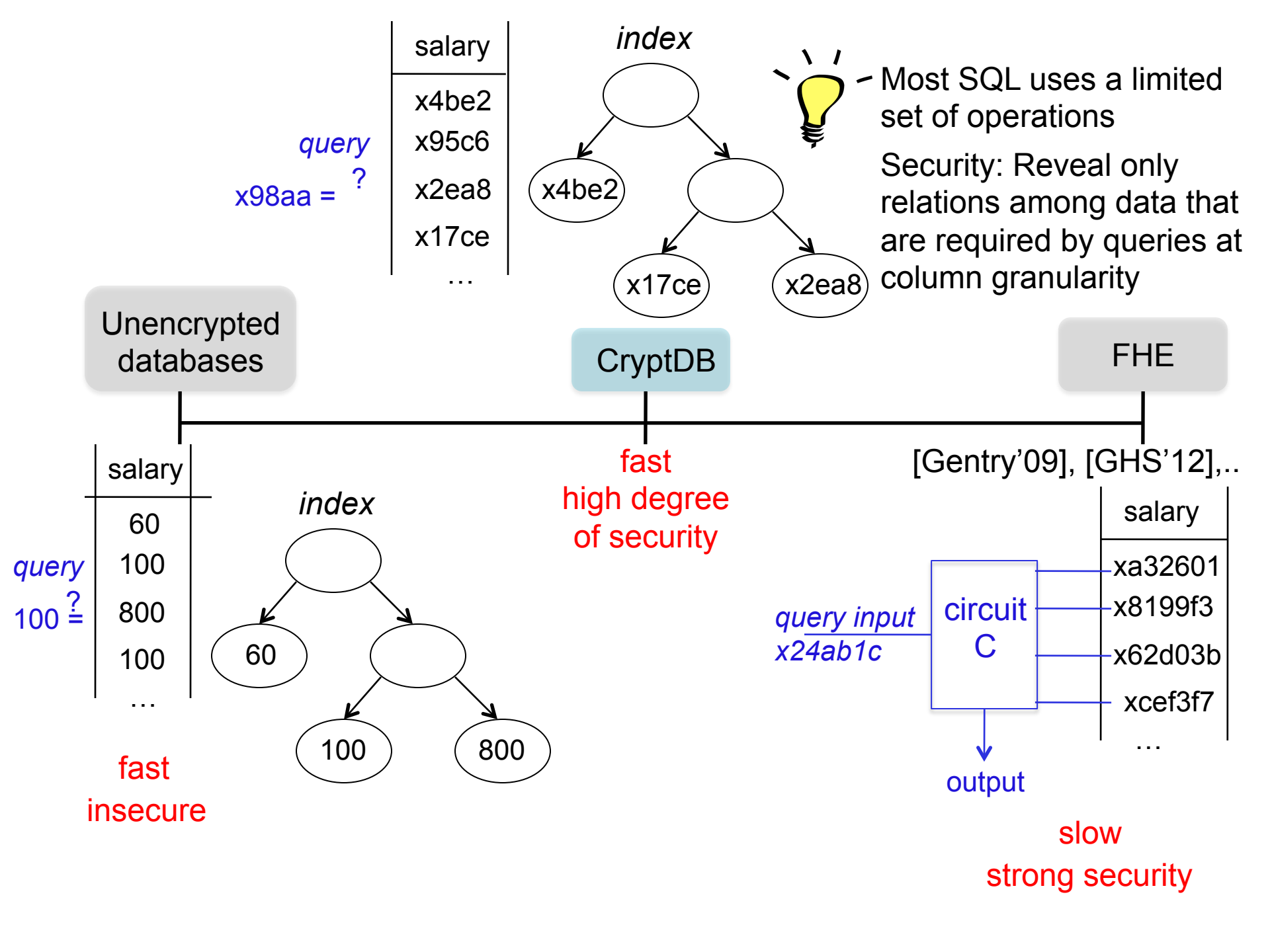
Contributions

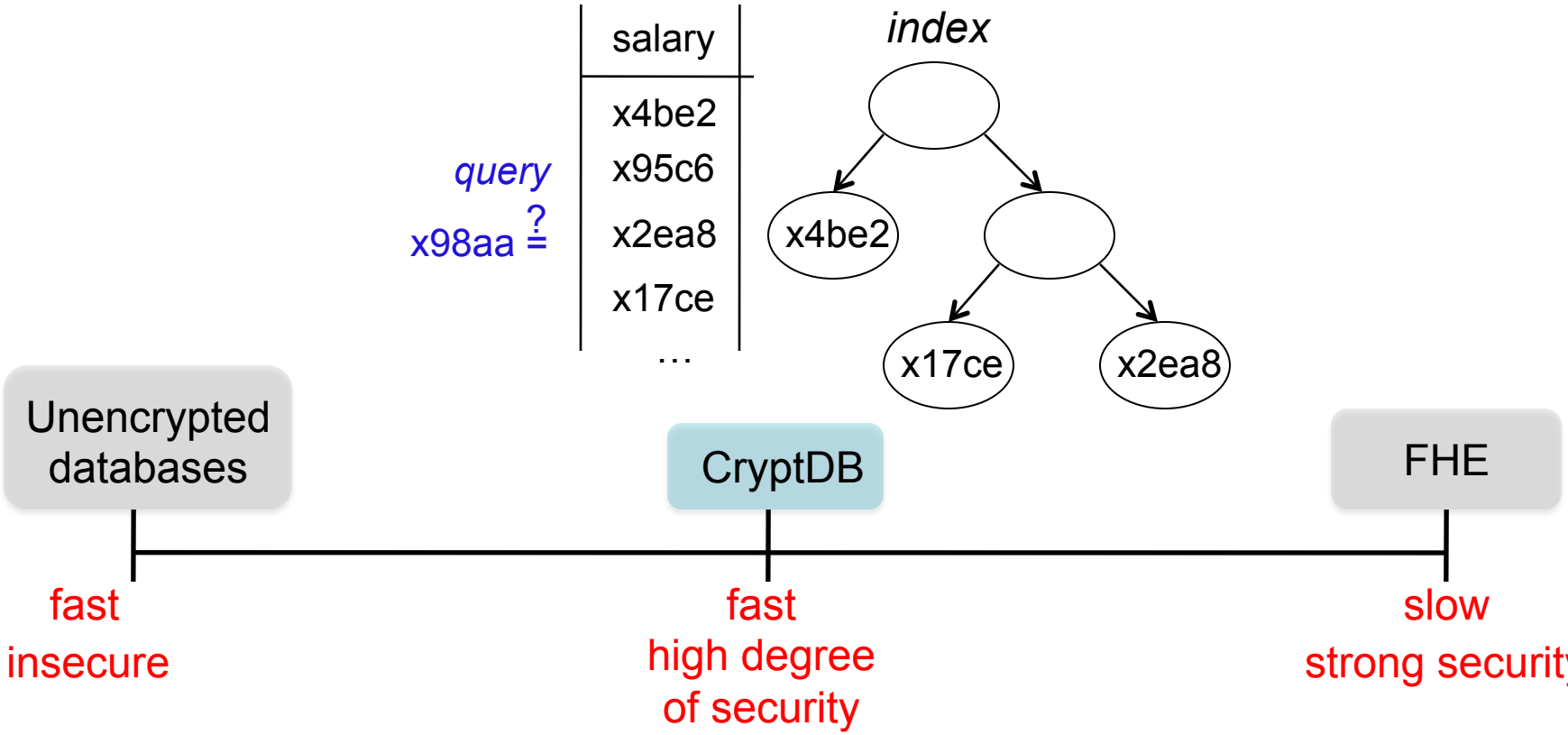
1. First *practical* DBMS to process most SQL queries on encrypted data

   ⇒ Hide DB from sys. admins., outsource DB to the cloud

2. Modest overhead: 26% throughput loss for TPC-C

3. No changes to DBMS (e.g., Postgres, MySQL) and no changes to applications

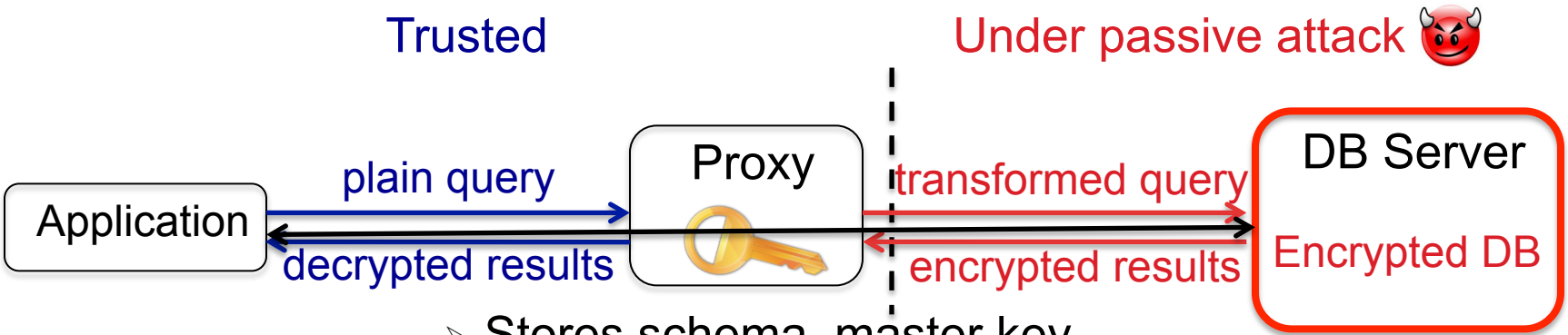

salary
x4be2
x95c6
x2ea8
x17ce
…
query
x98aa = ?
index
x4be2
x17ce
x2ea8
Most SQL uses a limited set of operations
Security: Reveal only relations among data that are required by queries at column granularity
Unencrypted databases
CryptDB
FHE
fast
high degree of security
[Gentry'09], [GHS'12],..
salary
60
100
800
100
…
query
100 = ?
index
60
100
800
fast
insecure
query input
x24ab1c
circuit C
output
salary
xa32601
x8199f3
x62d03b
xcef3f7
…
slow
strong security



Other work: weaker security, functionality, and/or efficiency:

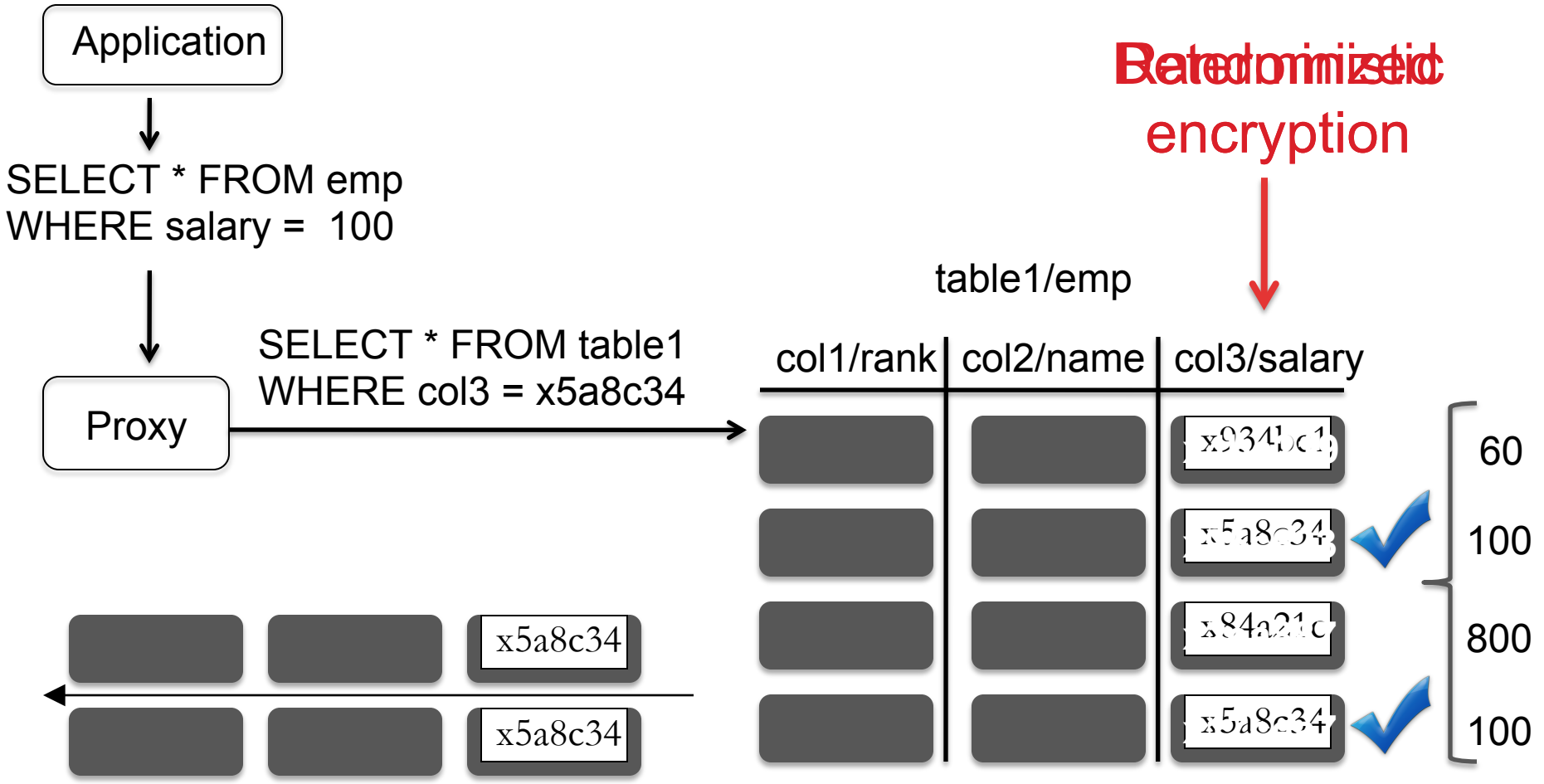
- Search on encrypted data (e.g., [Song et al.,'00])
- Systems proposals (e.g., [Hacigumus et al.,'02])
  - Require significant client-side processing

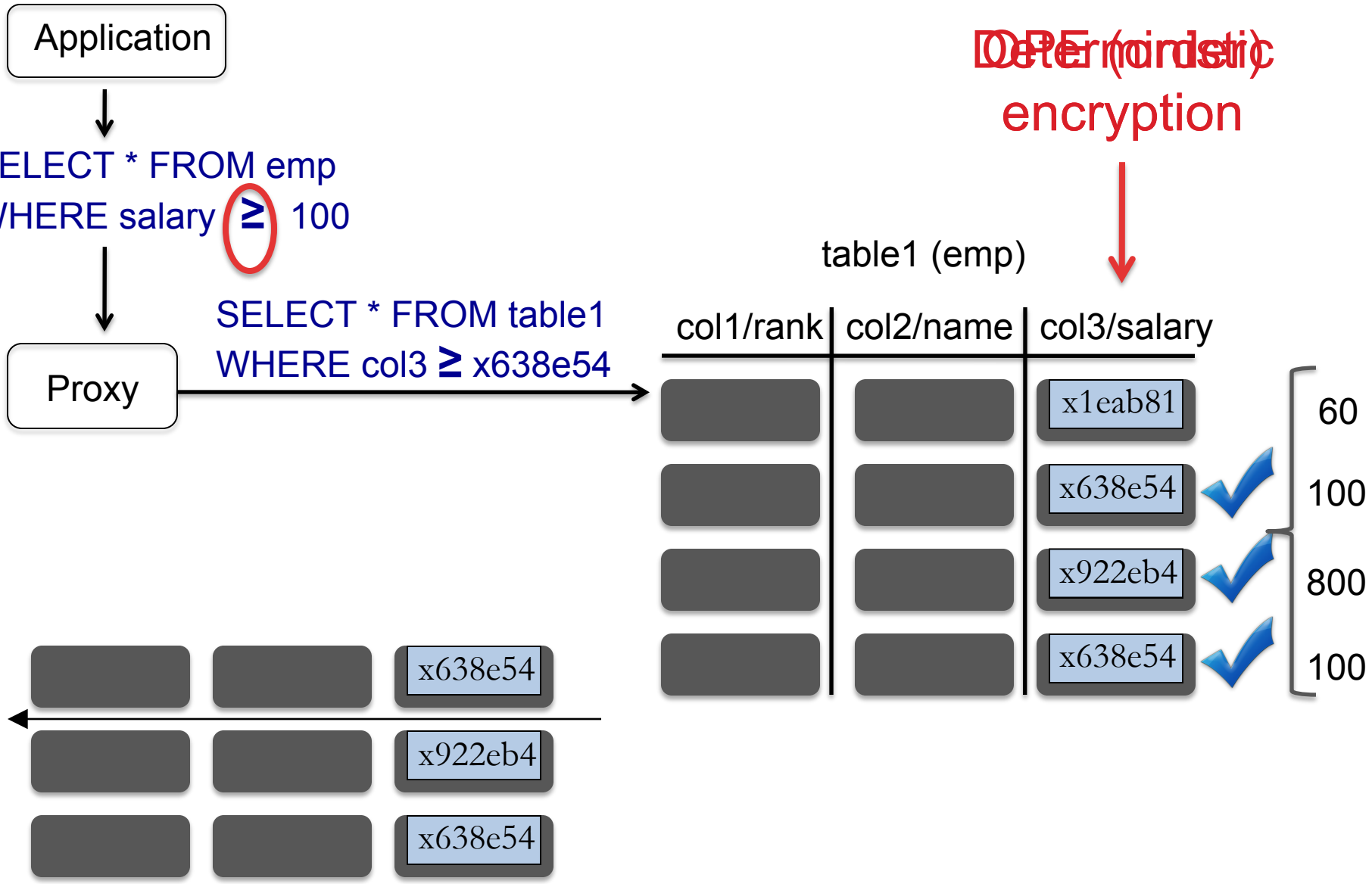
System Setup



- Stores schema, master key
- No data storage
- No query execution

- Process queries completely at the DBMS, on encrypted database


Application
SELECT * FROM emp
WHERE salary = 100
Proxy
SELECT * FROM table1
WHERE col3 = x5a8c34
x5a8c34
x5a8c34
Deterministic
encryption
table1/emp
col1/rank
col2/name
col3/salary
60
100
800
100

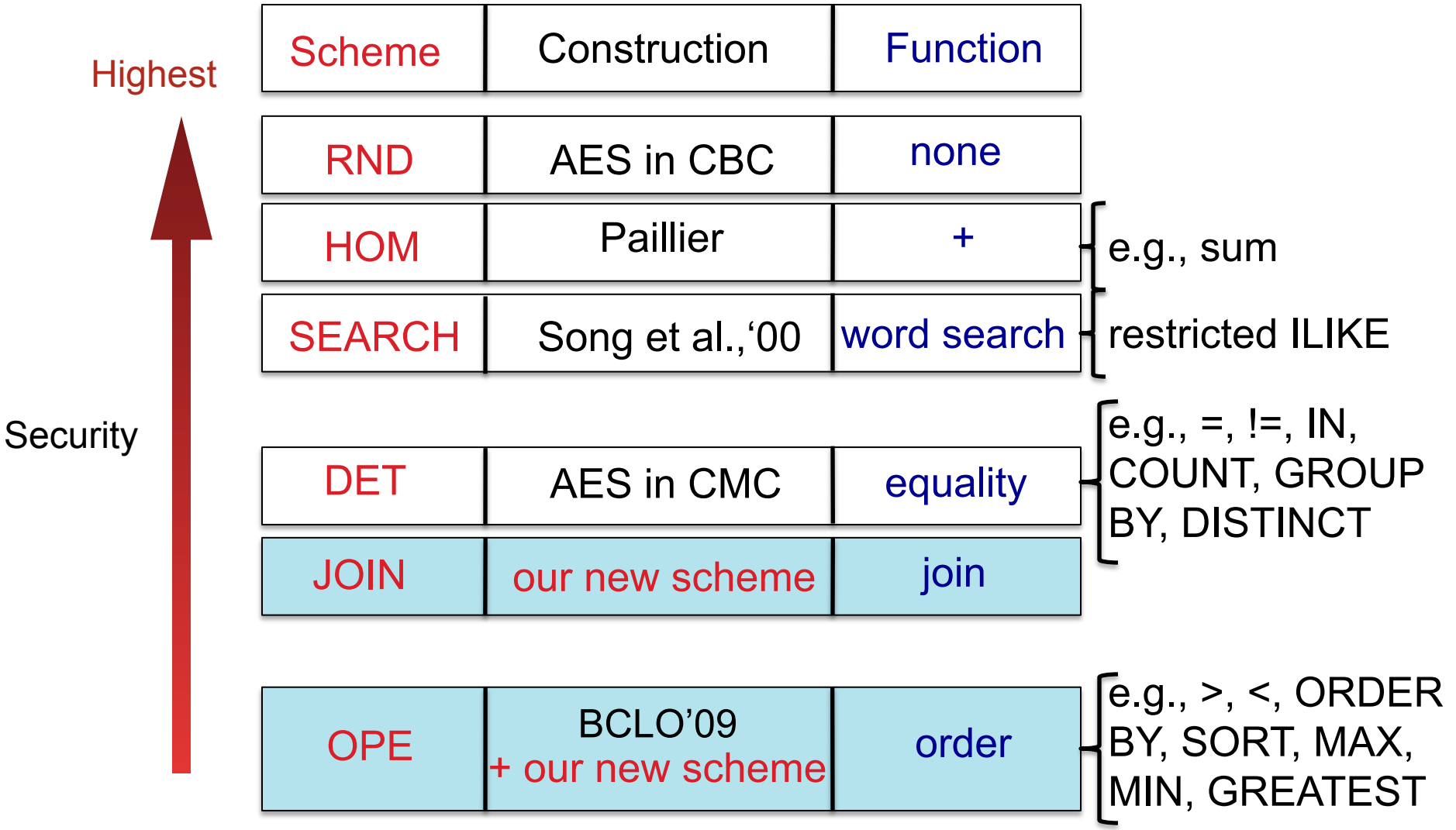

Application
SELECT * FROM emp
WHERE salary ≥ 100
Proxy
SELECT * FROM table1
WHERE col3 ≥ x638e54
Deterministic
encryption
table1 (emp)
col1/rank
col2/name
col3/salary
x1eab81
x638e54
x922eb4
x638e54
60
100
800
100
x638e54
x922eb4
x638e54

Two techniques

1. Use SQL-aware set of encryption schemes

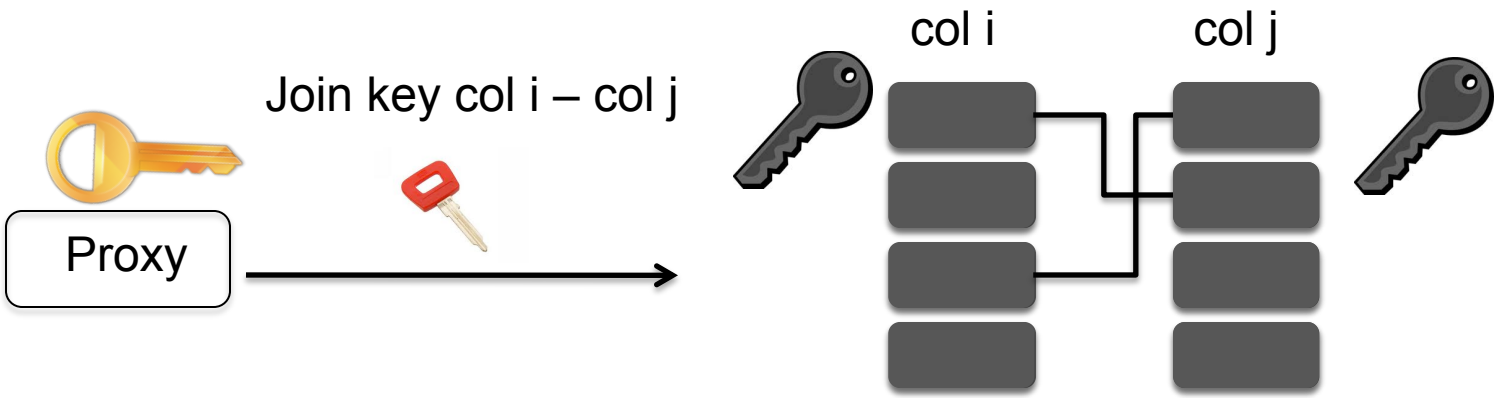
2. Adjust encryption of database based on queries

Encryption schemes



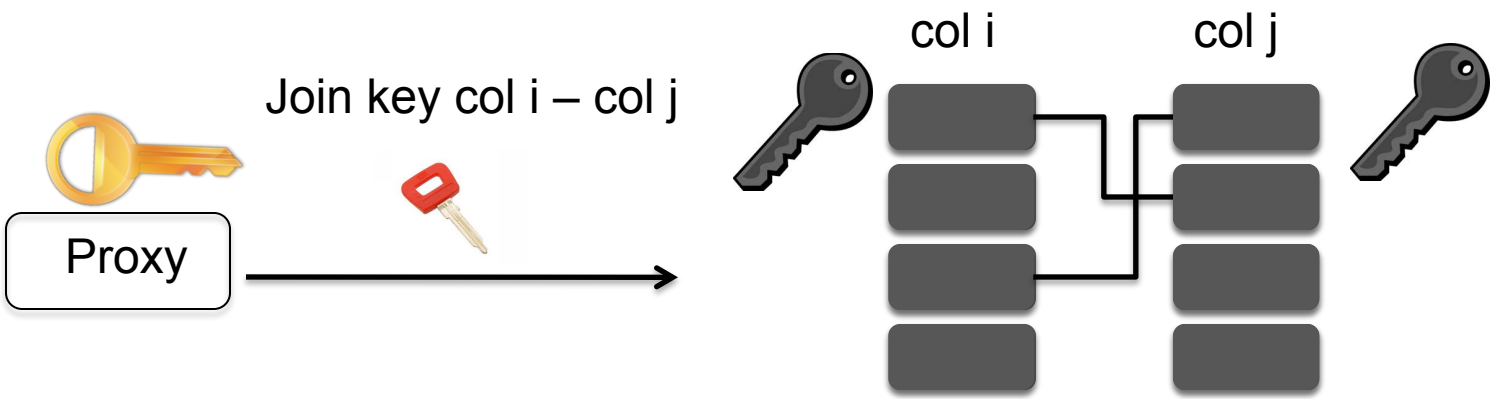
JOIN

- Do not know columns to be joined a priori!



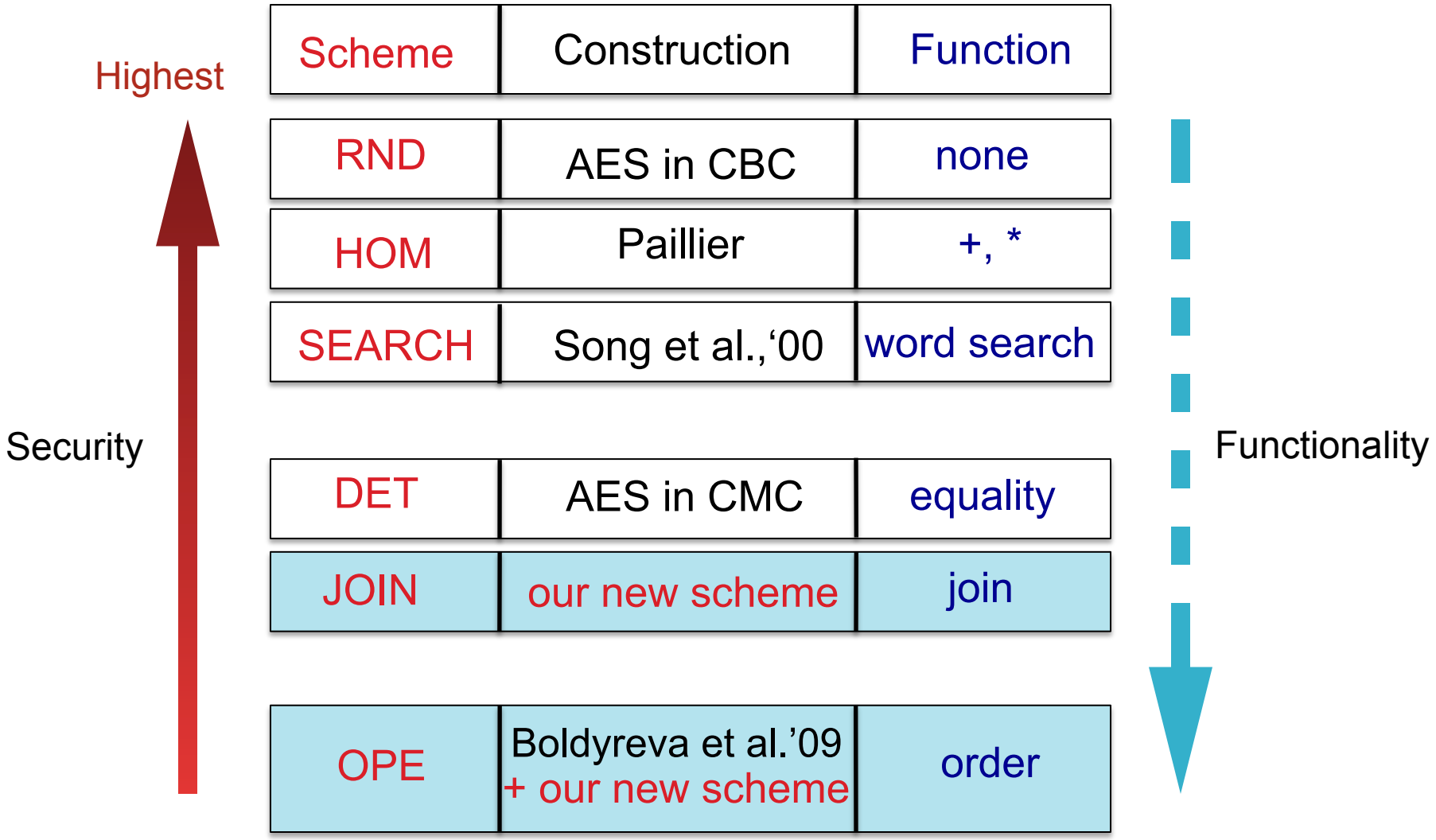
- *KeyGen* (sec. param): SK
- *Encrypt* (SK, m, col i): C_m^i (with) - deterministic
- *Token* (SK, col i, col j): (t_i, t_j)
- *Adjust* (t_i,C_m^i): C_m (with)

JOIN (cont'd)



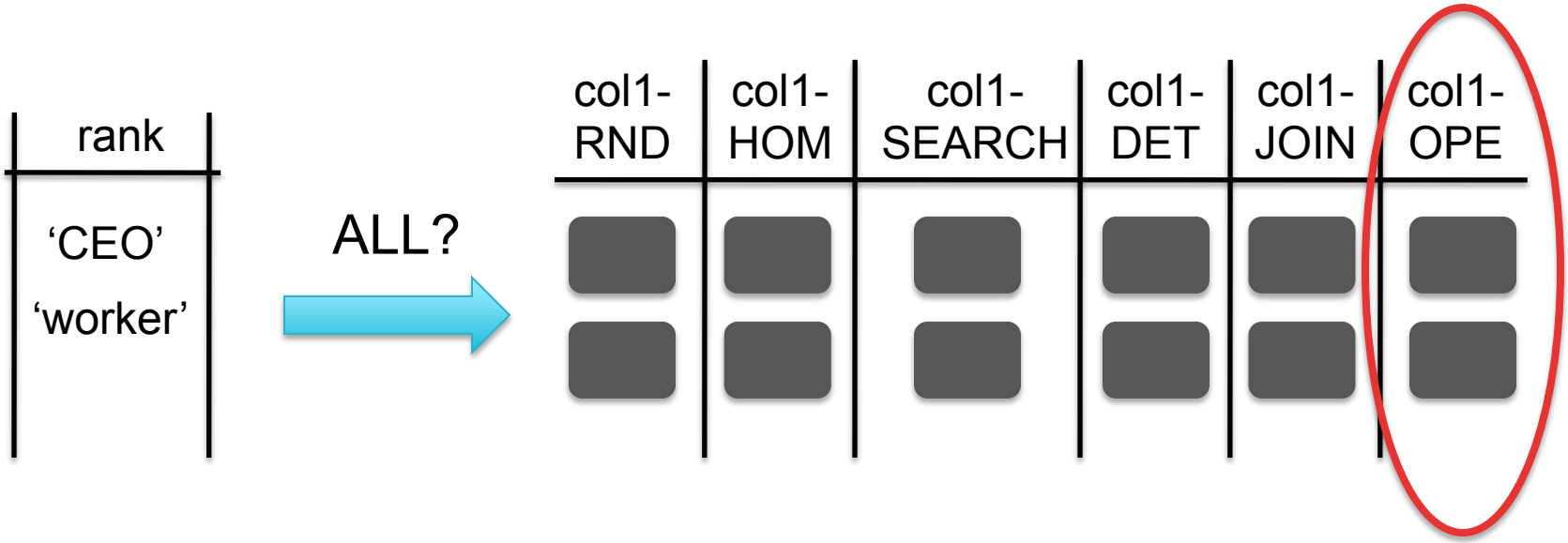
- Security: do not learn join relations without token
- Implementation:
  - 192 bits long, 0.52 ms encrypt, 0.56 ms adjust

Encryption schemes

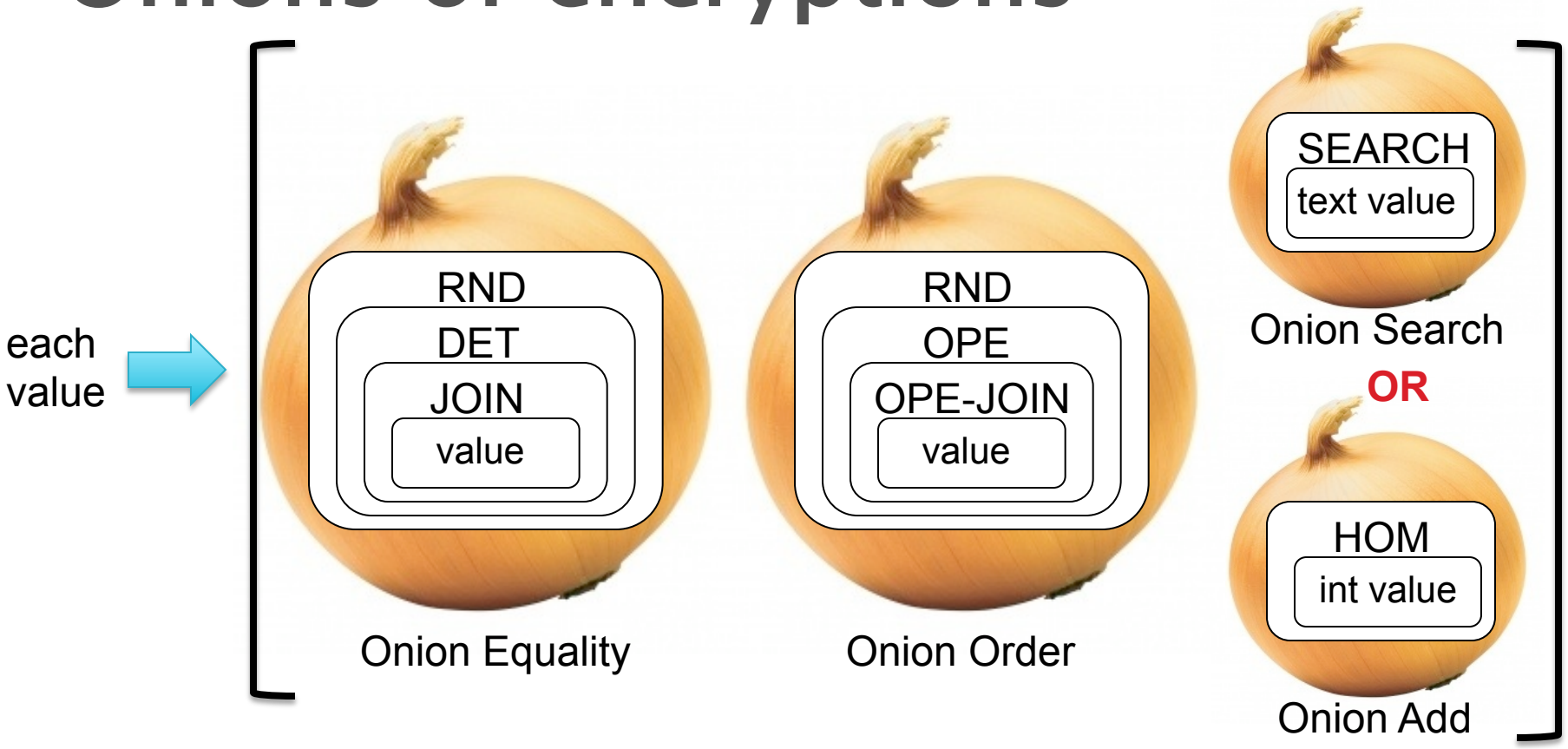


How to encrypt each data item?

- Encryption schemes needed depend on queries
- May not know queries ahead of time



Onions of encryptions

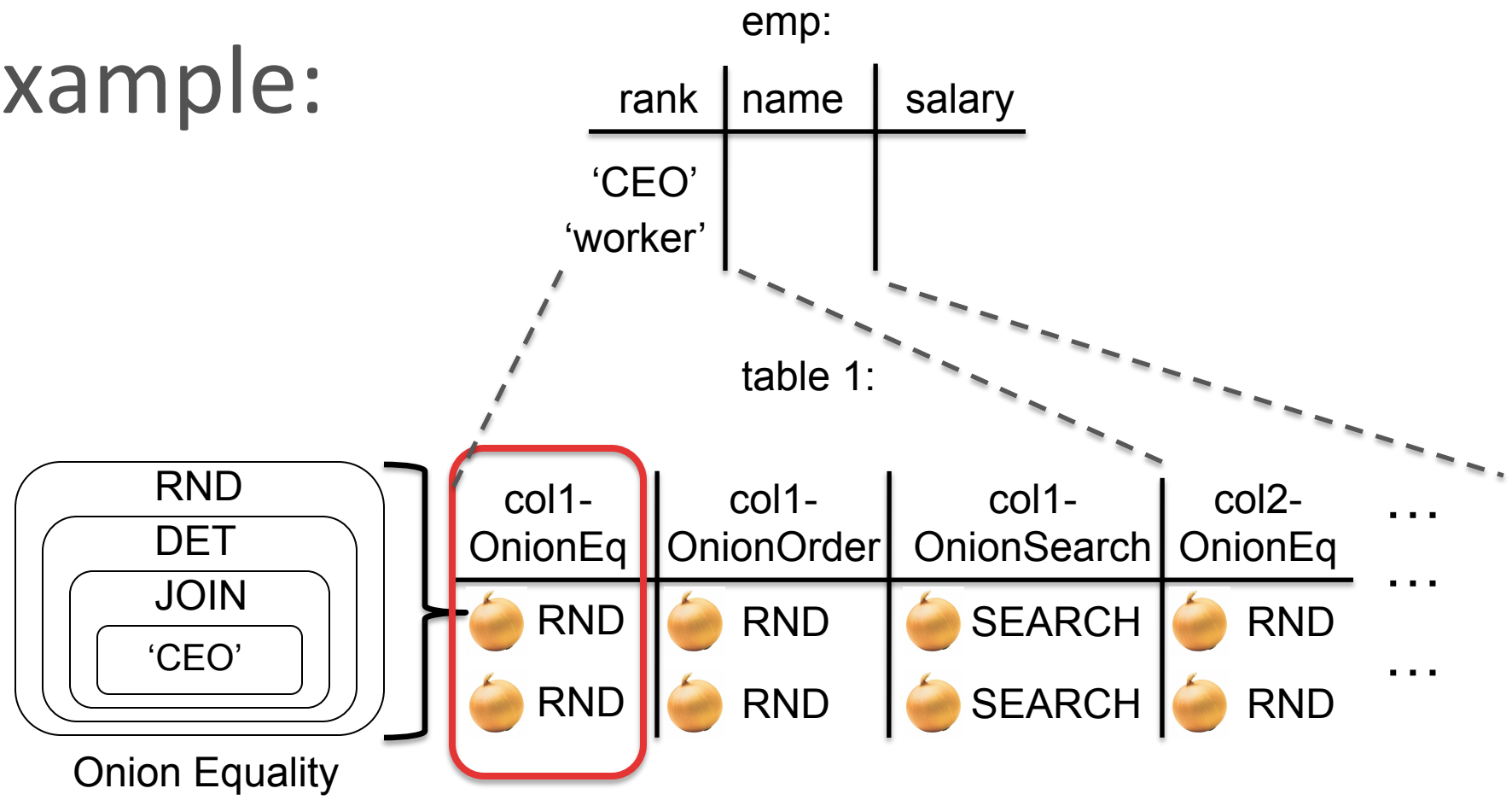


- Same key for all items in a column for same onion layer
- Start out the database with the most secure encryption scheme

Adjust encryption

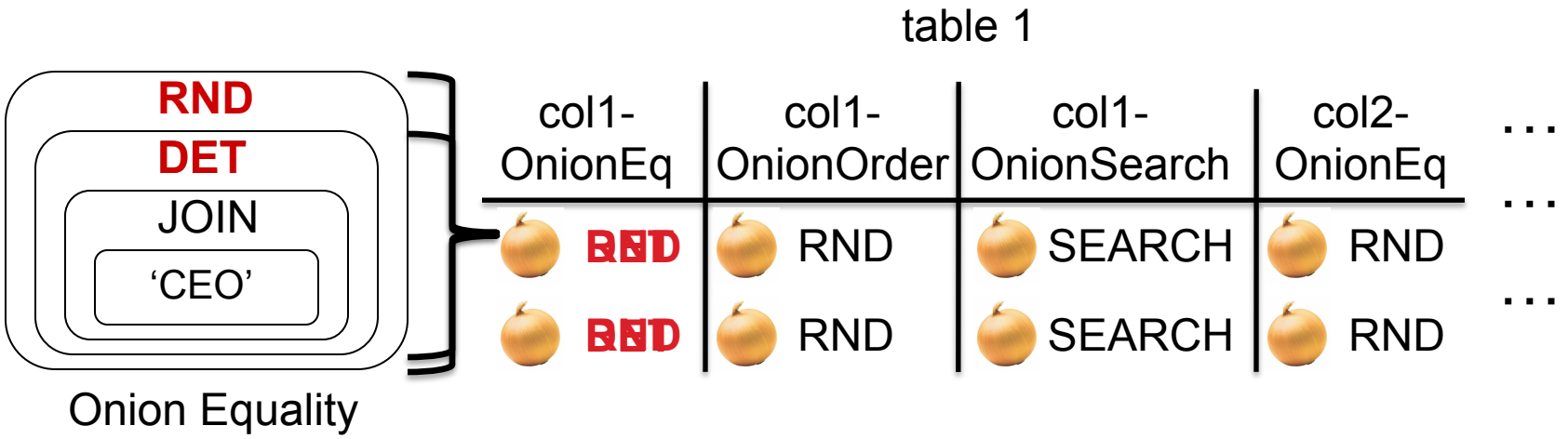
- Strip off layers of the onions
  - Proxy gives keys to server using a SQL UDF ("user-defined function")
  - Proxy remembers onion layer for columns
- Do not put back onion layer

Example:



SELECT * FROM emp WHERE rank = 'CEO';

Example (cont'd)



SELECT * FROM emp WHERE rank = 'CEO';

UPDATE table1 SET col1-OnionEq =
Decrypt_RND(key, col1-OnionEq);

SELECT * FROM table1 WHERE col1-OnionEq = xda5c0407;

Security guarantees

Queries ➡ encryption schemes ➡ leakage

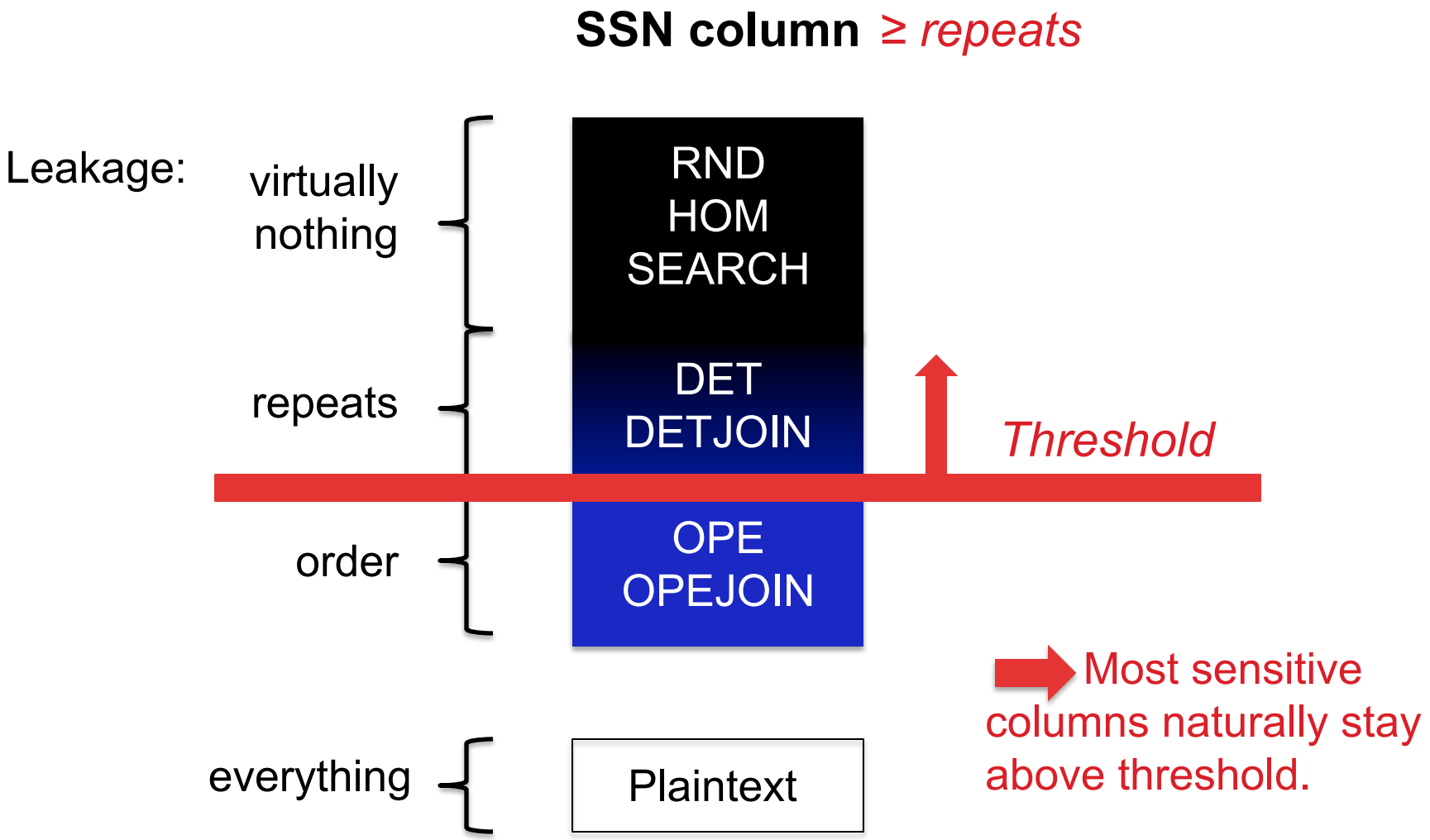
- Encryption schemes exposed for each column are the most secure enabling queries
- Overall: Reveal only data relations needed for query type, at column granularity

• equality predicate on a column	➡ DET	➡ repeats
• aggregation on a column	➡ HOM	➡ nothing
• no filter on a column	➡ RND	➡ nothing

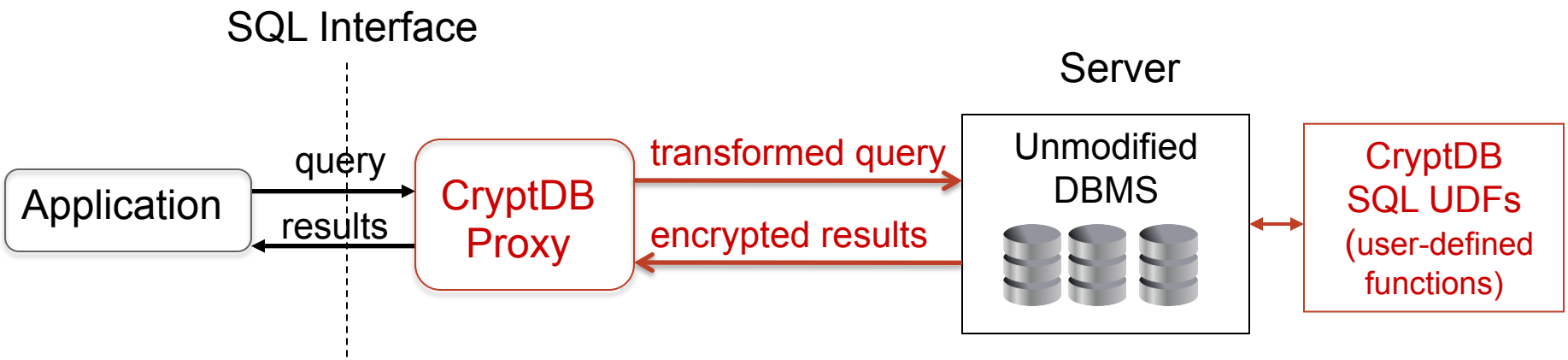
common in practice

➡ Never reveals plaintext

Security threshold



Implementation



- No change to the DBMS
- Portable: from Postgres to MySQL with 86 lines
- No change to applications

Evaluation

1. Does it support real queries/applications?
2. What is the resulting confidentiality?
3. What is the performance overhead?

Queries not supported

- More complex operators, e.g., trigonometry
- Operations that require combining encryption schemes
  - e.g., T1.a + T1.b > T2.c

Extensions: split queries, precompute columns, use FHE or other encryption schemes

Real queries/applications

Application	Total columns	Encrypted columns	# cols not supported
phpBB	563	23	0
HotCRP	204	22	0
grad-apply	706	103	0
TPC-C	92	92	0
sql.mit.edu	128,840	128,840	1,094

SELECT 1/log(series_no+1.2) …

… WHERE sin(latitude + PI()) …

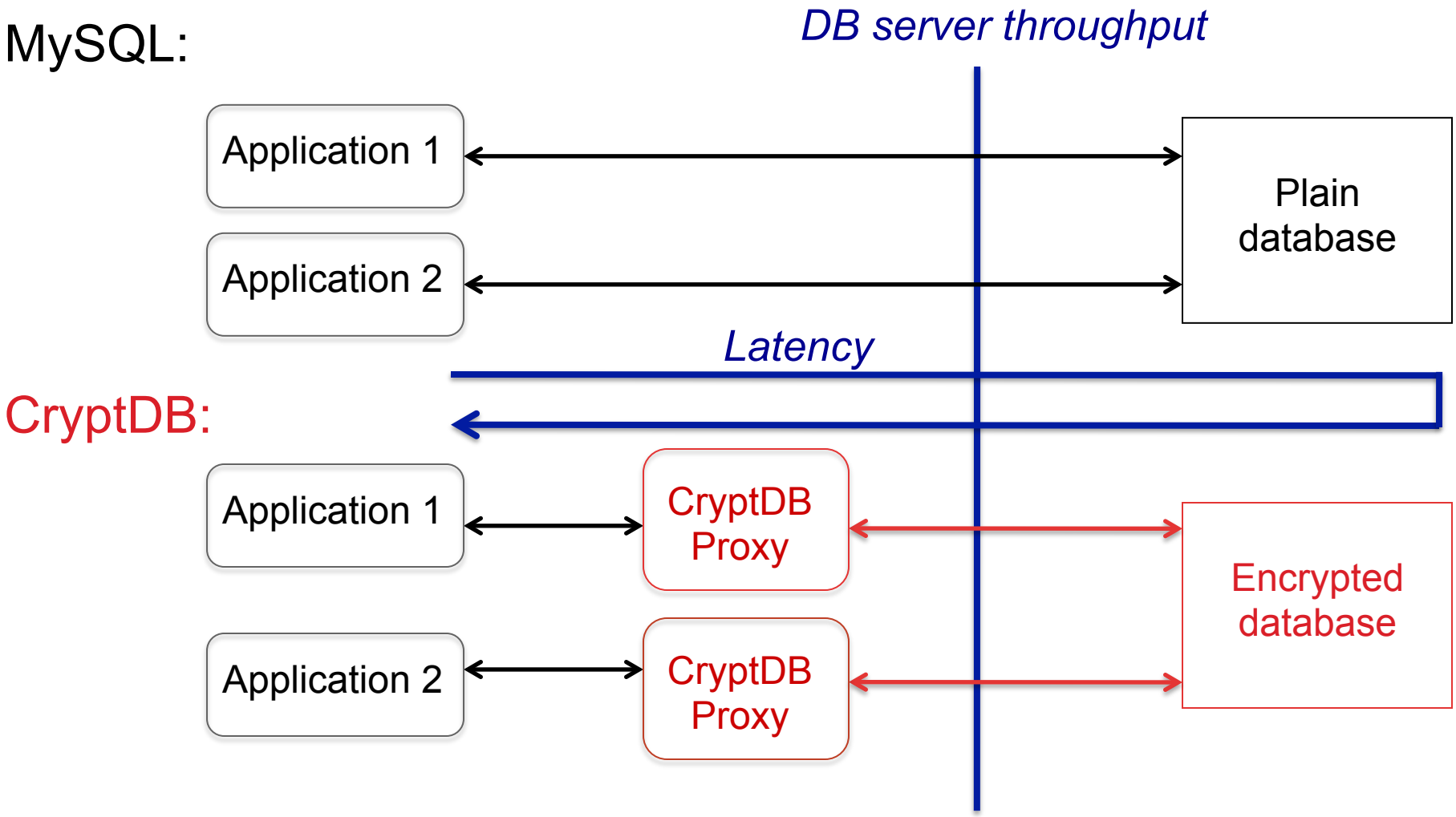
Resulting confidentiality

Application	Total columns	Encrypted columns	Min level is RND	Min level is DET	Min level is OPE
phpBB	563	23	21	1	1
HotCRP	204	22	18	1	2
grad-apply	706	103	95	6	2
TPC-C	92	92	65	19	8
sql.mit.edu	128,840	128,840	80,053	34,212	13,131

Most columns at RND

Most columns at OPE analyzed were less sensitive

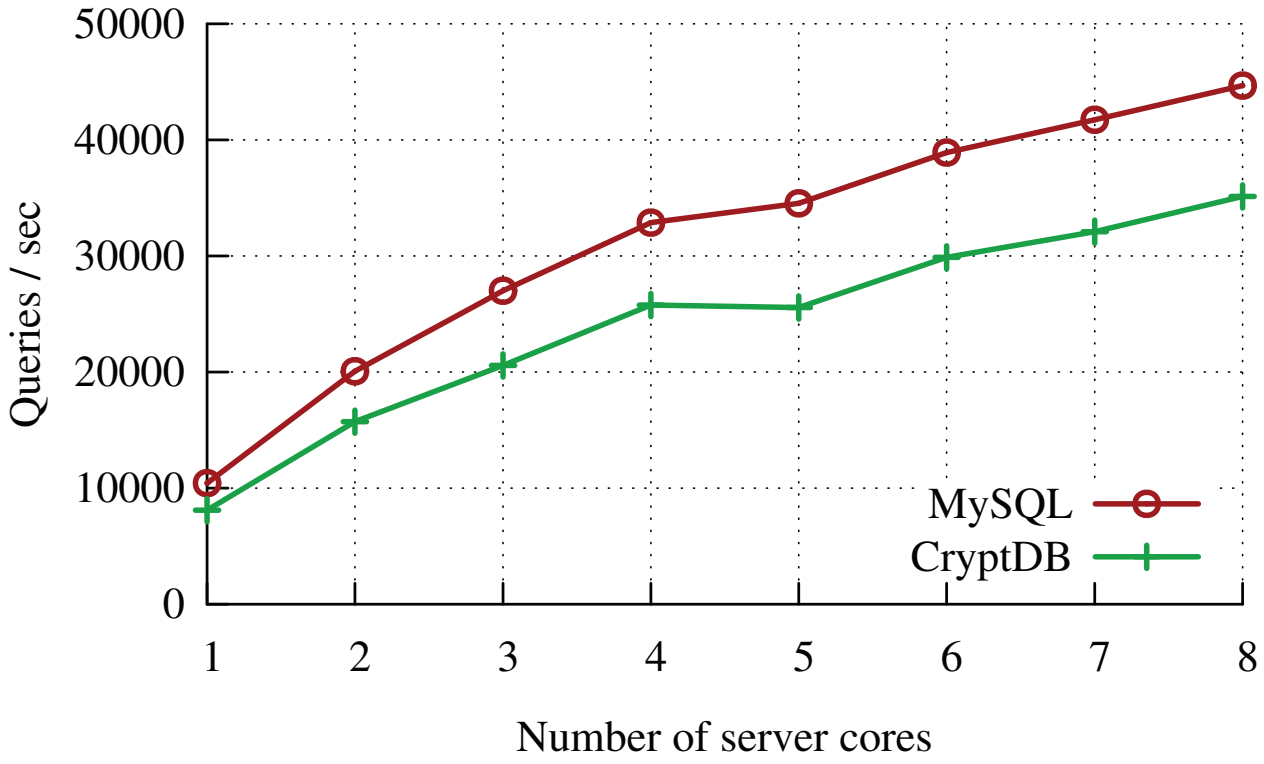
Performance



- Hardware: 2.4 GHz Intel Xeon E5620 – 8 cores, 12 GB RAM

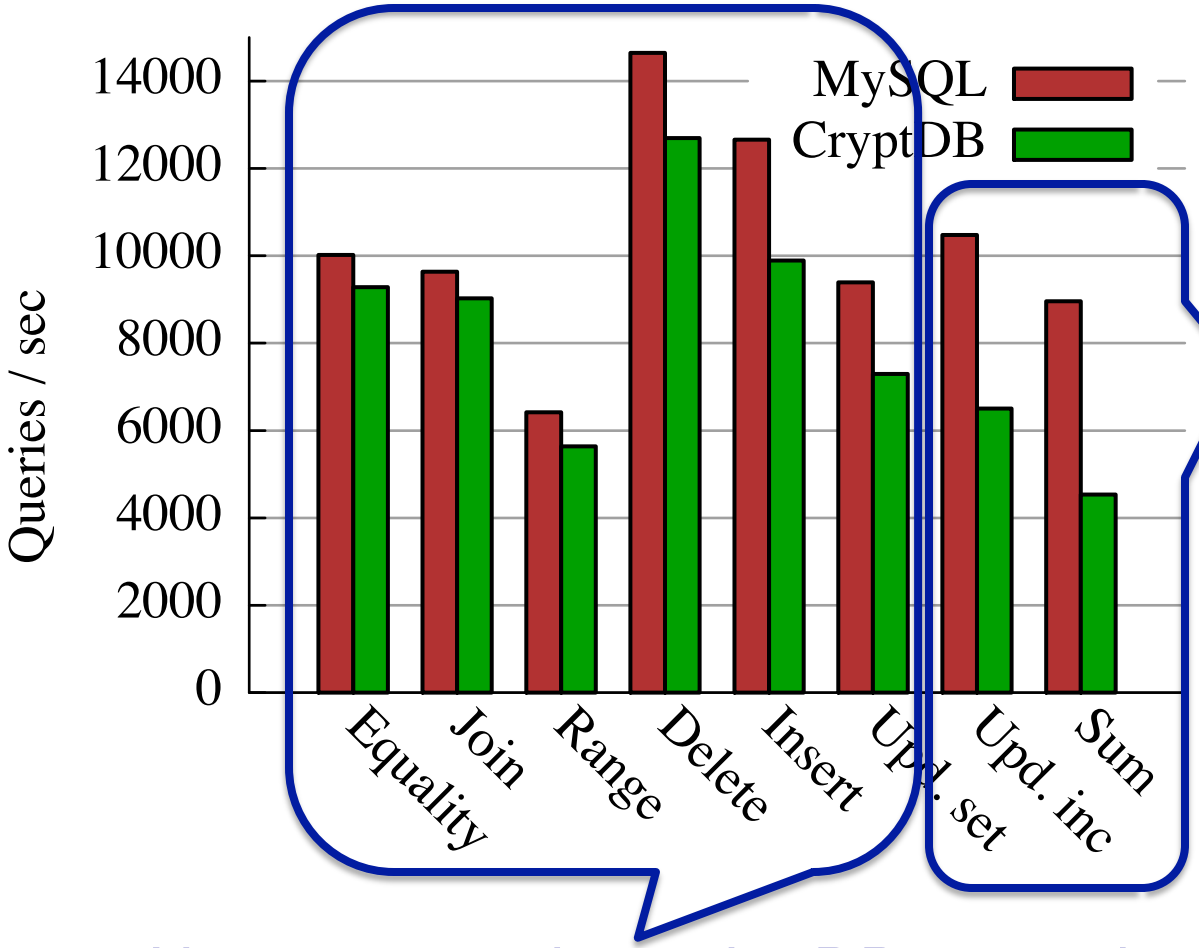
TPC-C performance

- Latency (ms/query): 0.10 MySQL vs. 0.72 CryptDB



Throughput loss 26%

TPC-C microbenchmarks



Homomorphic addition

No cryptography at the DB server in the steady state!

CryptDB is practical

Demo

Conclusions

CryptDB:

1. The first practical DBMS for running most standard queries on encrypted data
2. Modest overhead and no changes to DBMS

Website: http://css.csail.mit.edu/cryptdb/

Thanks!