

$$\mathbf{G} = \begin{bmatrix} 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 0 \end{bmatrix}$$

$$\mathbf{c} = [1 \quad 0 \quad 1]$$

$$\mathbf{d} = \mathbf{c} \cdot \mathbf{G} = [1 \quad 0 \quad 1] \cdot \begin{bmatrix} 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 0 \end{bmatrix} = [\underbrace{1 \quad 0 \quad 1}_{\mathbf{c}} \quad \underbrace{0 \quad 1 \quad 1}_{\mathbf{w}}]$$

$$\mathbf{G} \cdot \mathbf{H}^T = \mathbf{0}$$

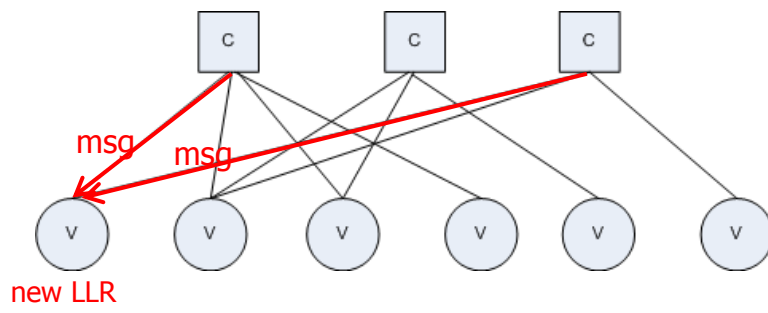
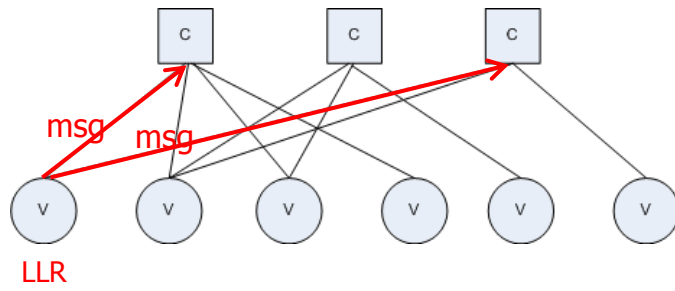
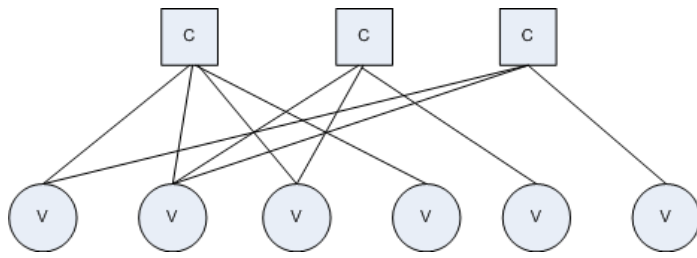
$$\mathbf{H}^T \cdot \mathbf{c} = \mathbf{0}$$

$$\mathbf{H} = \begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}$$

$$c_0 \oplus c_1 \oplus c_2 \oplus c_3 = 0$$

$$c_1 \oplus c_2 \oplus c_4 = 0$$

$$c_0 \oplus c_1 \oplus c_5 = 0$$



(1) Initialization: For  $i \in \{1, \dots, M\}, j \in \{1, \dots, N\}$

$$\alpha_j^0 = \ln \frac{P(VN_j = 0|y_j)}{P(VN_j = 1|y_j)} = \frac{2y_j}{\sigma^2} \quad (1)$$

$$\alpha_{i,j}^0 = \alpha_j^0 \quad (2)$$

(2) Update CN:  $\forall CN_i, i \in \{1, \dots, M\}$  do

$$\beta_{i,j}^n = \text{sgn}(\beta_{i,j}^n) \cdot |\beta_{i,j}^n| \quad (3)$$

where

$$\begin{aligned} \text{sgn}(\beta_{i,j}^n) &= \prod_{j' \in \mathcal{N}(i) \setminus j} \text{sgn}(\alpha_{i,j'}^{n-1}) \\ |\beta_{i,j}^n| &= \bigwedge_{j' \neq j} (\alpha_{i,j'}^{n-1}) \end{aligned}$$

(3) Update VN:  $\forall VN_j, j \in \{1, \dots, N\}$

$$\alpha_{i,j}^n = \alpha_{i,j}^0 + \sum_{i' \in \mathcal{M}(j) \setminus i} \beta_{i',j}^n \quad (4)$$

(4) Update a posteriori LLR:

$$\alpha_j^n = \alpha_j^0 + \sum_{i' \in \mathcal{M}(j)} \beta_{i',j}^n \quad (5)$$

(5) Repeat from (2) to (4) until the maximum number of iteration is reached, the estimated codeword is then given by:

$$\hat{c}_j = \begin{cases} 0, & \text{if } \alpha_j^n > 0 \\ 1, & \text{else} \end{cases} \quad (6)$$

$$\alpha_{i,j}^n = \alpha_j^n - \beta_{i,j}^n$$

(1) Initialization: For  $i \in \{1, \dots, M\}, j \in \{1, \dots, N\}$

$$\alpha_j^0 = \ln \frac{P(VN_j = 0|y_j)}{P(VN_j = 1|y_j)} = \frac{2y_j}{\sigma^2} \quad (7)$$

$$\alpha_{i,j}^0 = \alpha_j^0 \quad (8)$$

$$\beta_{i,j}^0 = 0 \quad (9)$$

(2)  $\forall CN_i, i \in \{1, \dots, M\}$  do

$$\beta_{i,j}^n = \text{sgn}(\beta_{i,j}^n) \cdot |\beta_{i,j}^n| \quad (10)$$

and

$$\alpha_j^n = (\alpha_j^{n-1} - \beta_{i,j}^{n-1}) + \beta_{i,j}^n \quad (11)$$

where

$$\text{sgn}(\beta_{i,j}^n) = \prod_{j' \in \mathcal{N}(i) \setminus j} \text{sgn}(\alpha_j^{n-1} - \beta_{i,j'}^{n-1})$$

$$|\beta_{i,j}^n| = \lambda \cdot \min_{j' \in \mathcal{N}(i) \setminus j} |\alpha_j^{n-1} - \beta_{i,j'}^{n-1}|$$

(3) Repeat (2) until the maximum number of iteration is reached, the estimated codeword is then given by (6)

Algorithm	$\bigwedge_{j' \neq j} (\alpha_{i,j'}^{n-1})$
SP	$\Phi \left( \sum_{j' \in \mathcal{N}(i) \setminus j} \Phi( \alpha_{i,j'}^{n-1} ) \right)$ $\Phi(x) = -\log \left( \tanh \left( \frac{x}{2} \right) \right)$
MS	$\min_{j' \in \mathcal{N}(i) \setminus j} \{ \alpha_{i,j'}^{n-1} \}$
OMS	$\max\{\min_{j' \in \mathcal{N}(i) \setminus j} \{ \alpha_{i,j'}^{n-1} \} - \beta, 0\}$ $\beta \geq 0$
NMS	$\lambda \cdot \min_{j' \in \mathcal{N}(i) \setminus j} \{ \alpha_{i,j'}^{n-1} \}$ $\lambda < 1$