

# Contents

- Event Groups
- Task Notifications

# Event Groups



Barry, R. (2016). *Mastering the freertos real time kernel. Pre-release 161204 Edition*. Real Time Engineers Ltd.

Chapter 8

# Event driven design

Real-time embedded systems have to take actions in response to events

Events must be communicated to tasks, for example with semaphores and queues

Properties of this type of communication are:

- Allows a task to wait in the Blocked state for a **single event** to occur
- Unblock a **single task** (the one with the highest priority) when the event occurs

# Event driven design

Event groups are another FreeRTOS feature

Properties of event groups are:

- Allows a task to wait in the Blocked state for **a combination of one of more events** to occur.
- Unblock **all the tasks** that were waiting for the same event, or combination of events, when the event occurs

# Event driven design

This makes event groups useful for :

- synchronizing multiple tasks
- broadcasting events
- allowing a task to wait in the Blocked state for any one of a set of events to occur
- allowing a task to wait in the Blocked state for multiple actions to complete
- possibly reduce RAM usage, because they may replace many binary semaphores

# Characteristics of an Event Group

An event 'flag' or 'bit' is a Boolean (1 or 0) stored in a single bit

- 0: event has not occurred
- 1: event has occurred

An event 'group' is a set of event flags: EventBits\_t

An application designer assigns the meaning to individual bits

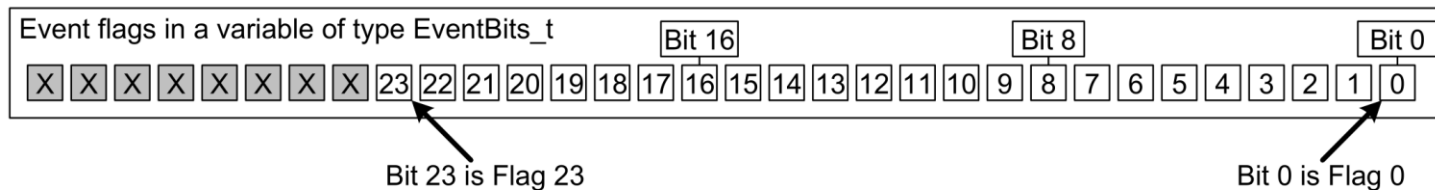


Figure 71 Event flag to bit number mapping in a variable of type EventBits\_t

# Characteristics of an Event Group

The number of event bits in an event group is dependent on configUSE\_16\_BIT\_TICKS

- If configUSE\_16\_BIT\_TICKS is 1, then each event group contains **8** usable event bits
- If configUSE\_16\_BIT\_TICKS is 0, then each event group contains **24** usable event bits

# Event Management Using Event Groups

Event groups are referenced using variables of type `EventGroupHandle_t`.

---

```
EventGroupHandle_t xEventGroupCreate( void );
```

---

**Listing 132. The `xEventGroupCreate()` API function prototype**

If NULL is returned, then the event group cannot be created.

A non-NULL value being returned indicates that the event group has been created successfully. The returned value is the handle to the created event group.



# Event Management Using Event Groups

Notify one or more tasks that the events represented by the bit, or bits, being set has occurred.

---

```
EventBits_t xEventGroupSetBits( EventGroupHandle_t xEventGroup,  
                                const EventBits_t uxBitsToSet );
```

---

**Listing 133. The xEventGroupSetBits() API function prototype**

The handle of the event group in which bits are being set.

# Event Management Using Event Groups

Notify one or more tasks that the events represented by the bit, or bits, being set has occurred.

---

```
EventBits_t xEventGroupSetBits( EventGroupHandle_t xEventGroup,  
                                const EventBits_t uxBitsToSet );
```

---

Listing 133. The xEventGroupSetBits() API function prototype

A bit mask that specifies the event bit, or event bits, to set to 1 in the event group. The value of the event group is updated by **bitwise ORing** the event group's existing value with the value passed in uxBitsToSet.

# Event Management Using Event Groups

Notify one or more tasks that the events represented by the bit, or bits, being set has occurred.

---

```
EventBits_t xEventGroupSetBits( EventGroupHandle_t xEventGroup,  
                                const EventBits_t uxBitsToSet );
```

---

Listing 133. The xEventGroupSetBits() API function prototype

The value of the event group at the time the call to xEventGroupSetBits() returned.

# Event Management Using Event Groups

Read the value of an event group, and optionally wait in the Blocked state for one or more event bits in the event group to become set.

The **unblock condition** is specified by two variables

- *uxBitsToWaitFor* specifies which event bits in the event group to test
- *xWaitForAllBits* specifies if a single of those bits must be set (bitwise OR test), or if all those bits must be set (bitwise AND test)

A task will not enter the Blocked state if its unblock condition is met at the time `xEventGroupWaitBits()` is called.

# Event Management Using Event Groups

## Example unblock conditions

Event group value	uxBitsToWaitFor	xWaitForAllBits	Resultant Behavior
0b00000000	0b00000101	pdFALSE	The calling task will enter the Blocked state because neither of bit 0 or bit 2 are set.
0b00000100	0b00000101	pdTRUE	The calling task will enter the Blocked state because bit 0 and bit 2 are not both set
0b00000100	0b00000101	pdFALSE	The calling task will not enter the Blocked state because bit 2 is already set
0b00000111	0b00000101	pdTRUE	The calling task will not enter the Blocked state because both bits are already set

# Event Management Using Event Groups

Read the value of an event group, and optionally wait in the Blocked state for one or more event bits in the event group to become set, if the event bits are not already set.

---

```
EventBits_t xEventGroupWaitBits( const EventGroupHandle_t xEventGroup,  
                                const EventBits_t uxBitsToWaitFor,  
                                const BaseType_t xClearOnExit,  
                                const BaseType_t xWaitForAllBits,  
                                TickType_t xTicksToWait );
```

---

Listing 135. The xEventGroupWaitBits() API function prototype

The handle of the event group in which bits are being set.

# Event Management Using Event Groups

Read the value of an event group, and optionally wait in the Blocked state for one or more event bits in the event group to become set, if the event bits are not already set.

---

```
EventBits_t xEventGroupWaitBits( const EventGroupHandle_t xEventGroup,  
                                const EventBits_t uxBitsToWaitFor,  
                                const BaseType_t xClearOnExit,  
                                const BaseType_t xWaitForAllBits,  
                                TickType_t xTicksToWait );
```

---

Listing 135. The xEventGroupWaitBits() API function prototype

A bit mask that specifies the event bit, or event bits, to test in the event group.

# Event Management Using Event Groups

Read the value of an event group, and optionally wait in the Blocked state for one or more event bits in the event group to become set, if the event bits are not already set.

---

```
EventBits_t xEventGroupWaitBits( const EventGroupHandle_t xEventGroup,  
                                const EventBits_t uxBitsToWaitFor,  
                                const BaseType_t xClearOnExit,  
                                const BaseType_t xWaitForAllBits,  
                                TickType_t xTicksToWait );
```

---

Listing 135. The xEventGroupWaitBits() API function prototype

If xWaitForAllBits is set to **pdFALSE**, then a task will leave the Blocked state when **any** of the bits specified by uxBitsToWaitFor become set.

If xWaitForAllBits is set to **pdTRUE**, then a task will leave the Blocked state when **all** of the bits specified by uxBitsToWaitFor become set.



# Event Management Using Event Groups

Read the value of an event group, and optionally wait in the Blocked state for one or more event bits in the event group to become set, if the event bits are not already set.

---

```
EventBits_t xEventGroupWaitBits( const EventGroupHandle_t xEventGroup,  
                                const EventBits_t uxBitsToWaitFor,  
                                const BaseType_t xClearOnExit,  
                                const BaseType_t xWaitForAllBits,  
                                TickType_t xTicksToWait );
```

---

**Listing 135. The xEventGroupWaitBits() API function prototype**

If the calling task's unblock condition has been met, and xClearOnExit is set to pdTRUE, then the event bits specified by uxBitsToWaitFor will be cleared back to 0 in the event group.

# Event Management Using Event Groups

Read the value of an event group, and optionally wait in the Blocked state for one or more event bits in the event group to become set, if the event bits are not already set.

---

```
EventBits_t xEventGroupWaitBits( const EventGroupHandle_t xEventGroup,  
                                const EventBits_t uxBitsToWaitFor,  
                                const BaseType_t xClearOnExit,  
                                const BaseType_t xWaitForAllBits,  
                                TickType_t xTicksToWait );
```

---

**Listing 135. The xEventGroupWaitBits() API function prototype**

The maximum amount of time the task should remain in the Blocked state to wait for its unblock condition to be met.

# Event Management Using Event Groups

Read the value of an event group, and optionally wait in the Blocked state for one or more event bits in the event group to become set, if the event bits are not already set.

---

```
EventBits_t xEventGroupWaitBits( const EventGroupHandle_t xEventGroup,  
                                const EventBits_t uxBitsToWaitFor,  
                                const BaseType_t xClearOnExit,  
                                const BaseType_t xWaitForAllBits,  
                                TickType_t xTicksToWait );
```

---

**Listing 135. The xEventGroupWaitBits() API function prototype**

The value of the event group at the time the calling task's unblock condition was met or the value at the time the block time expired.

# Event Groups Example

```
/*-----*/  
// Local variables  
/*-----*/  
static EventGroupHandle_t xEventGroup;  
  
const EventBits_t xEventBits[3] =  
{  
    0b00000000000000000000000000000001, // Is set by vTask1  
    0b00000000000000000000000000000010, // Is set by vTask2  
    0b00000000000000000000000000000100, // Is set by vTask3  
};  
  
const TickType_t xDelaysMs[3] =  
{  
    pdMS_TO_TICKS(1000), // Delay for vTask1  
    pdMS_TO_TICKS(5000), // Delay for vTask2  
    pdMS_TO_TICKS(10000), // Delay for vTask3  
};
```

# Event Groups Example

```
int main(void)
{
    rgb_init();
    xSerialPortInit(921600, 128);

    vSerialPutString("\r\nFRDM-KL25Z FreeRTOS demo Week 6 - Example 01\r\n");
    vSerialPutString("By Hugo Arends\r\n\r\n");

    xTaskCreate(vTask,          "vTask1",          configMINIMAL_STACK_SIZE, (void *)1, 3, NULL );
    xTaskCreate(vTask,          "vTask2",          configMINIMAL_STACK_SIZE, (void *)2, 2, NULL );
    xTaskCreate(vTask,          "vTask3",          configMINIMAL_STACK_SIZE, (void *)3, 1, NULL );
    xTaskCreate(vEventReader, "vEventReader", configMINIMAL_STACK_SIZE, NULL,      4, NULL );

    // Create the event group
    xEventGroup = xEventGroupCreate();

    /* Start the scheduler so the tasks start executing. */
    vTaskStartScheduler();

    /* If all is well then main() will never reach here as the scheduler will
    now be running the tasks. If main() does reach here then it is likely that
    there was insufficient heap memory available for the idle task to be created.
    Chapter 2 provides more information on heap memory management. */
    for( ;; );
}
```

# Event Groups Example

```
static void vTask(void *parameters)
{
    char str[48];

    // Set task specific values.
    int32_t n = (int32_t)parameters;
    TickType_t xLastWakeTime = xTaskGetTickCount();
    TickType_t xDelayMs = xDelaysMs[n-1];

    sprintf(str, "[% 7d - Task %d      ] Created\r\n", xLastWakeTime, n);
    vSerialPutString(str);

    // As per most tasks, this task is implemented in an infinite loop.
    for( ;; )
    {
        // Wait
        vTaskDelayUntil(&xLastWakeTime, xDelayMs);

        // Print info
        sprintf(str, "[% 7d - Task %d      ] Set bit %d\r\n", xLastWakeTime, n, (n-1));
        vSerialPutString(str);

        // Set the bit in the event group
        xEventGroupSetBits(xEventGroup, xEventBits[n-1]);
    }
}
```

# Event Groups Example

```
static void vEventReader(void *parameters)
{
    char str[48];

    EventBits_t xEventGroupValue;
    const EventBits_t xBitsToWaitFor = xEventBits[0] | xEventBits[1] | xEventBits[2];

    const BaseType_t xWaitForAllBits = pdTRUE; // pdFALSE

    sprintf(str, "[% 7d - EventReader] Created\r\n", xTaskGetTickCount());
    vSerialPutString(str);
}
```

# Event Groups Example

```
// As per most tasks, this task is implemented in an infinite loop.
for( ;; )
{
    // Block to wait for event bits to become set within the event group.
    xEventGroupValue = xEventGroupWaitBits(xEventGroup,
                                           xBitsToWaitFor,
                                           pdTRUE,
                                           xWaitForAllBits,
                                           portMAX_DELAY);

    // Prepare the string
    sprintf(str, "[% 7d - EventReader] EventBits: 000..0", xTaskGetTickCount());

    // Add the value of the last three bits
    strcat(str, ((xEventGroupValue & xEventBits[2]) != 0) ? "1" : "0");
    strcat(str, ((xEventGroupValue & xEventBits[1]) != 0) ? "1" : "0");
    strcat(str, ((xEventGroupValue & xEventBits[0]) != 0) ? "1" : "0");

    // Finalize string and print it
    strcat(str, "\r\n");
    vSerialPutString(str);
}
}
```



# Event Groups Example

**xWaitForAllBits = pdFALSE;**

```
[ 0 - EventReader] Created
[ 1 - Task 1      ] Created
[ 2 - Task 2      ] Created
[ 3 - Task 3      ] Created
[ 1001 - Task 1   ] Set bit 0
[ 1002 - EventReader] EventBits: 000..0001
[ 2001 - Task 1   ] Set bit 0
[ 2002 - EventReader] EventBits: 000..0001
[ 3001 - Task 1   ] Set bit 0
[ 3002 - EventReader] EventBits: 000..0001
[ 4001 - Task 1   ] Set bit 0
[ 4002 - EventReader] EventBits: 000..0001
[ 5001 - Task 1   ] Set bit 0
[ 5002 - EventReader] EventBits: 000..0001
[ 5002 - Task 2   ] Set bit 1
[ 5004 - EventReader] EventBits: 000..0010
[ 6001 - Task 1   ] Set bit 0
[ 6002 - EventReader] EventBits: 000..0001
[ 7001 - Task 1   ] Set bit 0
[ 7002 - EventReader] EventBits: 000..0001
[ 8001 - Task 1   ] Set bit 0
[ 8002 - EventReader] EventBits: 000..0001
[ 9001 - Task 1   ] Set bit 0
[ 9002 - EventReader] EventBits: 000..0001
[ 10001 - Task 1   ] Set bit 0
[ 10002 - EventReader] EventBits: 000..0001
[ 10002 - Task 2   ] Set bit 1
[ 10004 - EventReader] EventBits: 000..0010
[ 10003 - Task 3   ] Set bit 2
[ 10006 - EventReader] EventBits: 000..0100
[ 11001 - Task 1   ] Set bit 0
[ 11002 - EventReader] EventBits: 000..0001
```

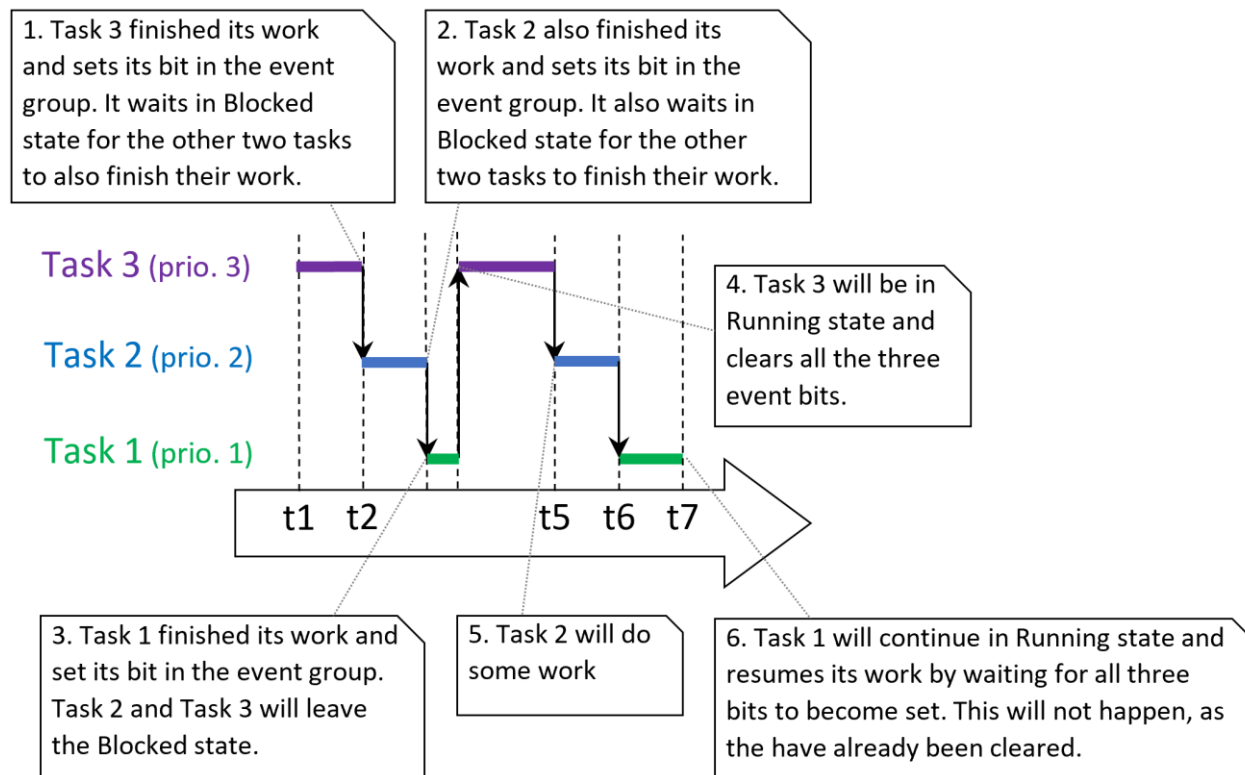
**xWaitForAllBits = pdTRUE;**

```
[ 0 - EventReader] Created
[ 1 - Task 1      ] Created
[ 2 - Task 2      ] Created
[ 3 - Task 3      ] Created
[ 1001 - Task 1   ] Set bit 0
[ 2001 - Task 1   ] Set bit 0
[ 3001 - Task 1   ] Set bit 0
[ 4001 - Task 1   ] Set bit 0
[ 5001 - Task 1   ] Set bit 0
[ 5002 - Task 2   ] Set bit 1
[ 6001 - Task 1   ] Set bit 0
[ 7001 - Task 1   ] Set bit 0
[ 8001 - Task 1   ] Set bit 0
[ 9001 - Task 1   ] Set bit 0
[ 10001 - Task 1   ] Set bit 0
[ 10002 - Task 2   ] Set bit 1
[ 10003 - Task 3   ] Set bit 2
[ 10004 - EventReader] EventBits: 000..0111
[ 11001 - Task 1   ] Set bit 0
[ 12001 - Task 1   ] Set bit 0
```

# Task Synchronization Using an Event Group

Goal: to synchronize three tasks with event groups

Problem example: the bits in the event group are already cleared!



# Task Synchronization Using an Event Group

`xEventGroupSync()` is a special function for task synchronisation with event groups

- It allows a task to set one or more event bits in an event group
- It then waits for a combination of event bits to become set in the same event group **as a single un interruptable operation.**

---

```
EventBits_t xEventGroupSync( EventGroupHandle_t xEventGroup,  
                             const EventBits_t uxBitsToSet,  
                             const EventBits_t uxBitsToWaitFor,  
                             TickType_t xTicksToWait );
```

---

Listing 142. The `xEventGroupSync()` API function prototype

# Summary

- Practical uses for event groups.
- The advantages and disadvantages of event groups relative to other FreeRTOS features.
- How to set bits in an event group.
- How to wait in the Blocked state for bits to become set in an event group.
- How to use an event group to synchronize a set of tasks.

# Task Notifications



Barry, R. (2016). *Mastering the freertos real time kernel. Pre-release 161204 Edition*. Real Time Engineers Ltd.

Chapter 9

# FreeRTOS communication

Set of independent tasks that:

- Communicate with each other
- Collectively implement system functionality

So far we have used **communication objects**

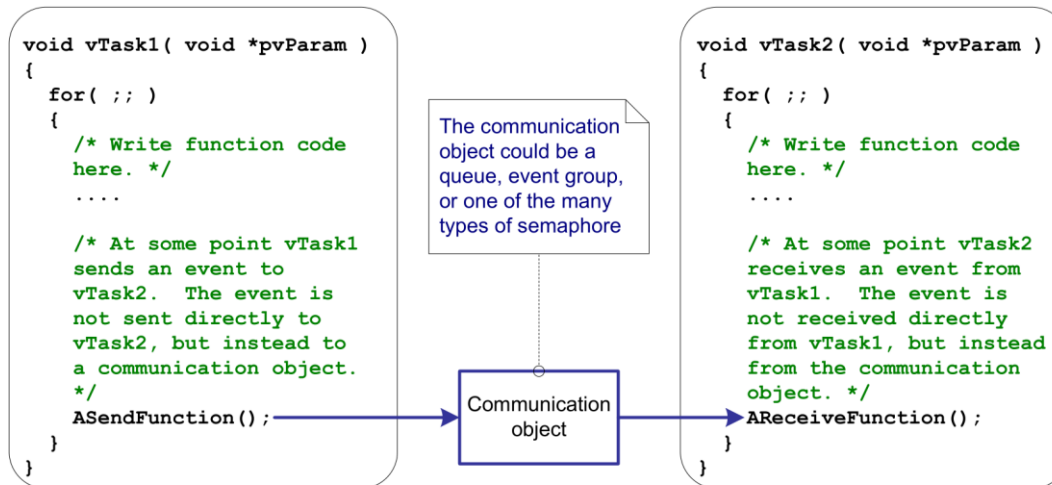


Figure 76 A communication object being used to send an event from one task to another

# FreeRTOS communication

We can use task notifications instead, to send an event directly

Enabled by setting configUSE\_TASK\_NOTIFICATIONS to 1

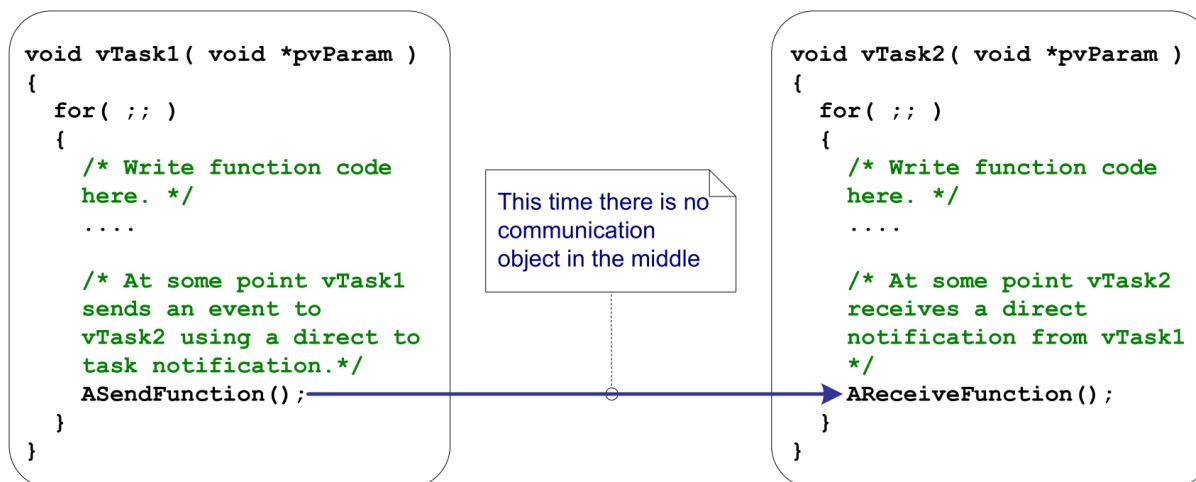


Figure 77 A task notification used to send an event directly from one task to another

# Task notification characteristics

Each task has a **Notification State**:

- *Pending*  
When a task receives a notification, its notification state is set to pending.
- *Not-Pending*  
When a task reads its notification value, its notification state is set to not-pending.

Each task has a **Notification Value**:

- A 32-bit unsigned integer

A task can wait in the Blocked state, with an optional time out, for its notification state to become pending.



# Benefits

When compared to a queue, semaphore or event group to perform an equivalent operation

- Performance - significantly faster
- RAM usage - significantly less RAM

# Limitations

Scenarios in which a task notification cannot be used

- Sending an event or data to an ISR
- Buffering multiple data items
- Broadcasting to more than one task
- Waiting in the blocked state for a send to complete

*If a task attempts to send a task notification to a task that already has a notification pending, then it is not possible for the sending task to wait in the Blocked state for the receiving task to reset its notification state.*

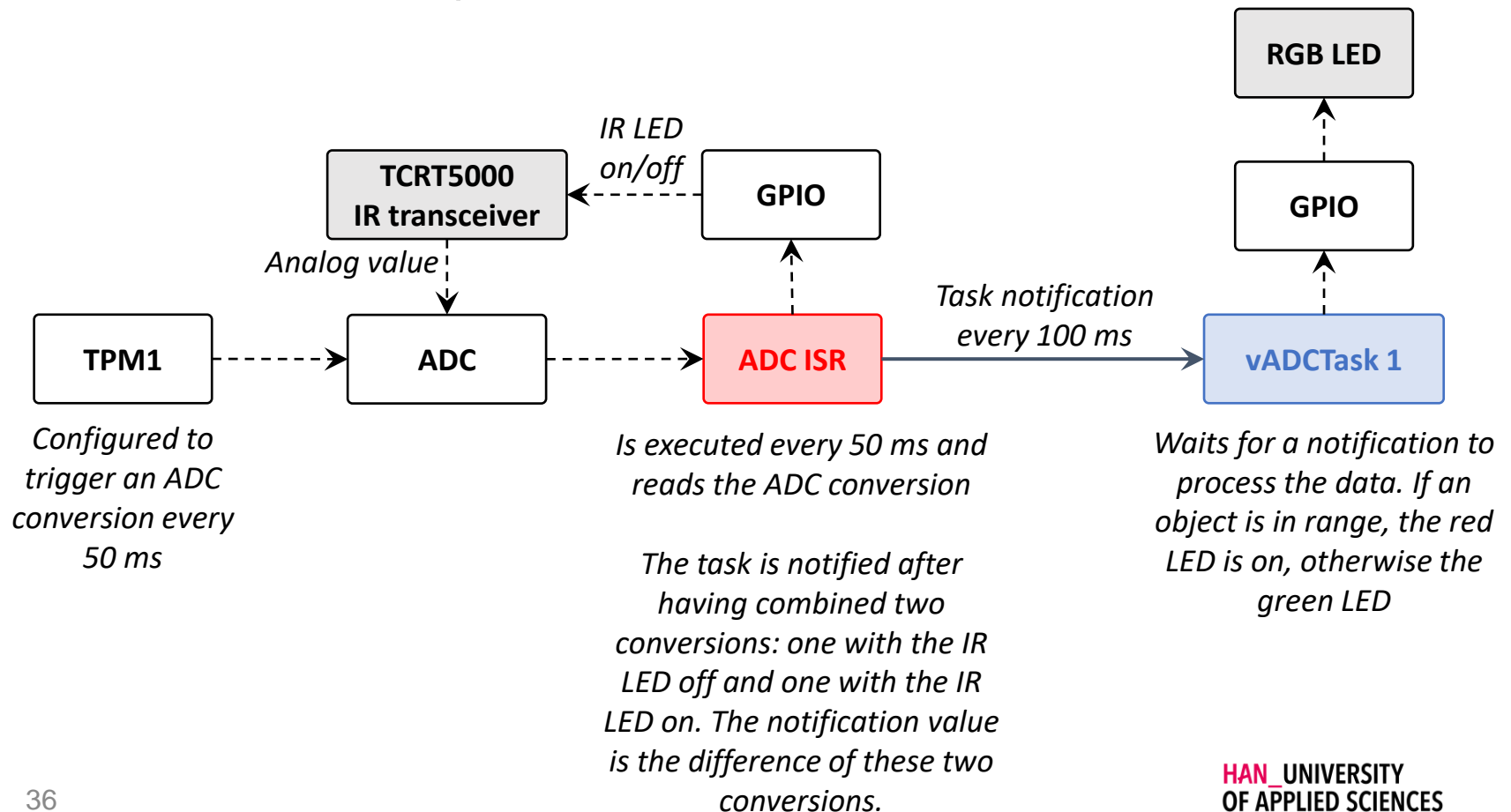
# Using Task Notifications

## API functions overview

	<i><u>Lightest weight and fastest alternative to a binary or counting semaphore</u></i>	<i><u>Lighter weight and faster alternative to a binary or counting semaphore, or event group, or queue of length one</u></i>
<i>Sending</i>	<b>xTaskNotifyGive(...)</b> <b>xTaskNotifyGiveFromISR(...)</b> <ul style="list-style-type: none"><li>• Sets notification state to pending</li><li>• Notification value: +1</li></ul>	<b>xTaskNotify(...)</b> <b>xTaskNotifyFromISR(...)</b> <ul style="list-style-type: none"><li>• Sets notification state to pending</li><li>• Notification value: +1 or set bits or a completely new number</li></ul>
<i>Receiving</i>	<b>ulTaskNotifyTake(...)</b> <ul style="list-style-type: none"><li>• Task optionally waits in Blocked state for notification state to become pending</li><li>• Notification value: -1 or cleared to 0</li></ul>	<b>xTaskNotifyWait(...)</b> <ul style="list-style-type: none"><li>• Task optionally waits in Blocked state for notification state to become pending</li><li>• Notification value: bits to clear on entry <i>and</i> bits to clear on exit</li></ul>

# Task Notifications Example

See Week 6 – Example 03



# Summary

- A task's notification state and notification value.
- How and when a task notification can be used in place of a communication object, such as a semaphore.
- The advantages of using a task notification in place of a communication object.