

1. Answer:

We will use UDP. Assume file transmission time is the same, for TCP, we need extra 2RTTs (One for the TCP connection set-up, the other for client to send request to sever and sever to send back). But for UDP, we only need RTT(the client sends the transaction request into a UDP socket, and the sever sends the reply back to the client' s UDP socket). So the UDP is faster.

2. Answer:

Transport Layer provides services such as connection-oriented communication, reliability, flow control, and multiplexing. The two most important protocols in the Transport Layer are Transmission Control Protocol (TCP) and User Datagram Protocol (UDP). Both protocols deliver data between the Application Layer and the Internet Layer.

TCP provides reliable data delivery service with end-to-end error detection and correction.

UDP provides low-overhead, connectionless datagram delivery service.

But both of them don' t provide minimum throughput guarantee or delay.

3. Answer:

```
[dyn-129-236-236-89:~ zhanghaopeng$ telnet www.ee.columbia.edu 80
Trying 128.59.64.28...
Connected to broadband.ee.columbia.edu.
Escape character is '^]'.
GET /~jghaderi/index.html HTTP/1.1
Host: www.ee.columbia.edu
If-Modified-Since: Wed, 10 Jan 2019 20:33:23 GMT

HTTP/1.1 304 Not Modified
Date: Wed, 13 Feb 2019 19:43:55 GMT
Server: Apache/2.4.29 (Ubuntu)
ETag: "1796-57ea7ca58aac0"

Connection closed by foreign host. _
```

Cannot identify my IP address.

4. Answer:

The user will get first chunk as a result of being selected by one of his neighbors as a result of an “optimistically unchoke” for sending out chunks to him. Every 30 secs, a peer will select one of its neighbors randomly as a peer for uploading no matter whether this neighbor is uploading data to it or not.

5. Answer:

Every time, we need 2RTT to setup TCP connection, request and receive HTML file.

a. Non-persistent HTTP requires 2 RTTs per object.

So we need $2RTT + 8 * 2RTT = 18RTT$

b. We need $2RTT + (8 // 5) * 2RTT = 6RTT$

c. Server leaves connection open after sending response so that subsequent HTTP message between same client/server can be sent over open connection, and client will send request as soon as it encounters the reference objects.

So we need $2RTT + RTT = 3RTT$

6. Answer:

```
zhanghaopengdeMacBook-Air:~ zhanghaopeng$ nslookup
> set type=ns
> .
Server:      209.18.47.62
Address:     209.18.47.62#53

Non-authoritative answer:
.    nameserver = a.root-servers.net.
.    nameserver = b.root-servers.net.
.    nameserver = c.root-servers.net.
.    nameserver = d.root-servers.net.
.    nameserver = e.root-servers.net.
.    nameserver = f.root-servers.net.
.    nameserver = g.root-servers.net.
.    nameserver = h.root-servers.net.
.    nameserver = i.root-servers.net.
.    nameserver = j.root-servers.net.
.    nameserver = k.root-servers.net.
.    nameserver = l.root-servers.net.
.    nameserver = m.root-servers.net.

Authoritative answers can be found from:
```

There are 13 root DNS servers.

```
[> edu.
Server:      209.18.47.62
Address:     209.18.47.62#53

Non-authoritative answer:
edu      nameserver = f.edu-servers.net.
edu      nameserver = i.edu-servers.net.
edu      nameserver = b.edu-servers.net.
edu      nameserver = l.edu-servers.net.
edu      nameserver = e.edu-servers.net.
edu      nameserver = h.edu-servers.net.
edu      nameserver = g.edu-servers.net.
edu      nameserver = k.edu-servers.net.
edu      nameserver = j.edu-servers.net.
edu      nameserver = m.edu-servers.net.
edu      nameserver = d.edu-servers.net.
edu      nameserver = a.edu-servers.net.
edu      nameserver = c.edu-servers.net.
```

Authoritative answers can be found from:

There are 13 TLD DNS servers for edu

```
[> columbia.edu
Server:      209.18.47.62
Address:     209.18.47.62#53

Non-authoritative answer:
columbia.edu  nameserver = sns-pb.isc.org.
columbia.edu  nameserver = ns1.lse.ac.uk.
columbia.edu  nameserver = sdns1.berkeley.edu.
columbia.edu  nameserver = ext-ns1.columbia.edu.
columbia.edu  nameserver = dns2.itd.umich.edu.

Authoritative answers can be found from:
```

Only 5 authoritative DNS servers for that organization.

7. Answer:

a. In this distribution scheme, the server will send the file to N clients at a rate of

$\frac{us}{N}$, which is larger than client's download rate, so the receiving rate will also

be $\frac{us}{N}$. Because $\frac{us}{N} \leq d_{min}$, so $\frac{NF}{us} \geq \frac{F}{d_{min}}$. The distribution time will be:

$$\max\left(\frac{NF}{us}, \frac{F}{d_{min}}\right) = \frac{NF}{us}$$

b. In this distribution scheme, the server will send the file to N clients at a rate of

d_{min} , and the total rate is $Nd_{min} \leq us$. The client's receiving rate can be

d_{min} . Because $\frac{us}{N} \geq d_{min}$, so $\frac{NF}{us} \leq \frac{F}{d_{min}}$. The distribution time will be:

$$\max\left(\frac{NF}{us}, \frac{F}{d_{min}}\right) = \frac{F}{d_{min}}$$

c. Distribution time(DT) $\geq \max\left(\frac{NF}{us}, \frac{F}{d_{min}}\right)$

If $\frac{us}{N} \leq d_{min}$, we get DT $\geq \frac{NF}{us}$

From part a, we know that DT = $\frac{NF}{us}$

so DT = $\frac{NF}{us}$

If $\frac{us}{N} \geq d_{min}$, we get DT $\geq \frac{F}{d_{min}}$

From part b, we know that DT = $\frac{F}{d_{min}}$

so DT = $\frac{F}{d_{min}}$

Above all, DT = $\max\left(\frac{NF}{us}, \frac{F}{d_{min}}\right)$