

# 1 Model Framework

We are interested in models with asymptotically power-law decay of the mutual information measure with respect to the distance between the tokens in the sequence. So far so good. But what does it *actually* mean?

The tokens, represented by random variables  $X_t$ , are elements of a finite alphabet  $\Sigma$ . The distance between  $X_t$  and  $X_{t+\tau}$  is  $\tau$ , and for every  $t$  and every  $\tau$  we want to bound

$$I(X_t, X_{t+\tau}) \in \Omega(\tau^{-\alpha}), \quad I(X_t, X_{t+\tau}) \in \mathcal{O}(\tau^{-\beta}) \quad ,$$

for some fixed  $\alpha, \beta \in \mathbb{R}_{>0}$ . The first condition is the important one, while the latter ensures that  $I(X_t, X_{t+\tau}) \xrightarrow{\tau \rightarrow \infty} 0$ . We also may replace the latter condition by this one.

This was straight forward. The challenging part is to define what a model is. In the case of Markov chains this seems trivial: We define a finite set of parameters (the transition probabilities), and we get a model over  $\Sigma^*$ , that is for every  $n \in \mathbb{N}$  the model defines a probability measure over  $\Sigma^n$ . Thus, the first conclusion is every model  $S$  must define a probability measure over  $\Sigma^n$  for every  $n \in \mathbb{N}$ .

As a first formalization,  $S$  is a function  $S : (n, w) \mapsto [0, 1]$ , for  $n \in \mathbb{N}$ ,  $w \in \Sigma^n$  s.t.  $\sum_{w \in \Sigma^n} S(n, w) = 1$ .

But really, we want to restrain  $S$  in order to have reasonable time and space complexity, and to ensure the model is *reasonable*, which means that the language of  $S_n(w)$  should look *similar* to  $S_{n+d}(w)$ , whatever this might mean, where we used the notation  $S_n(w) \equiv S(n, w)$ . We also write  $w_i$  for  $X_i$ . Really,  $w$  is a 1-indexed String of  $X_i$ .

We present one strict definition for this *similarity* in the following definition:

**Definition 1.1.** We say  $S$  is *well behaved* iff for every  $n \in \mathbb{N}$ ,  $w \in \Sigma^{n+1}$  it holds true that

$$\sum_{w_{n+1} \in \Sigma} S_{n+1}(w) = S_n(w_{-(n+1)}) \quad .$$

**Lemma 1.1.** For every  $d \in \mathbb{N}$ , let  $I := [n+d] \setminus [n] = \{n+1, \dots, n+d\}$ . Then, if  $S$  is well behaved, we have for every  $w \in \Sigma^{n+d}$ :

$$\sum_{w_I \in \Sigma^d} S_{n+d}(w) = S_n(w_{-I}) \quad .$$

*Proof.* We use induction over  $d$ . The base case follows directly from the definition of  $S$  being well behaved. Thus, assume the claim holds for some  $d := k$ .

Then we have

$$\begin{aligned}
\sum_{w_I \in \Sigma^{k+1}} S_{n+k+1}(w) &= \sum_{w_{I \setminus \{k+1\}} \in \Sigma^k} \sum_{w_{k+1} \in \Sigma} S_{n+k+1}(w) \\
&\stackrel{S \text{ well defined}}{=} \sum_{w_{I \setminus \{k+1\}} \in \Sigma^k} S_{n+k}(w_{-\{k+1\}}) \\
&\stackrel{\text{induction hypothesis}}{=} S_n(w_{-I}) \quad ,
\end{aligned}$$

which concludes the induction step.  $\square$

**Definition 1.2.** From any model  $S$ , we can construct the *induced well behaved* model  $S^*$  by defining  $S_n^*$  recursively as

- $S_1^* := S_1$  ,
- $S_{n+1}^*(w) := S_n^*(w_{-\{n+1\}}) \frac{S_{n+1}(w)}{\sum_{w_{n+1} \in \Sigma} S_{n+1}(w)}$  .

**Lemma 1.2.** *The induced well behaved model  $S^*$  is indeed well behaved.*

*Proof.* We have:

$$\begin{aligned}
\sum_{w_{n+1} \in \Sigma} S_{n+1}^*(w) &= \sum_{w_{n+1} \in \Sigma} S_n^*(w_{-\{n+1\}}) \frac{S_{n+1}(w)}{\sum_{w_{n+1} \in \Sigma} S_{n+1}(w)} \\
&= \frac{S_n^*(w_{-\{n+1\}})}{\sum_{w_{n+1} \in \Sigma} S_{n+1}(w)} \sum_{w_{n+1} \in \Sigma} S_{n+1}(w) \\
&\stackrel{\check{}}{=} S_n^*(w_{-\{n+1\}}) \quad .
\end{aligned}$$

$\square$

Now, we want to look at how we might restrict our model  $(S_n)_{n \in \mathbb{N}} \equiv S$ . One approach might be to define a model structure for every  $n \in \mathbb{N}$  with finite parameters  $\theta_n$ , thus  $S_n \in \{S_n(\theta_n) : \theta_n \in \Theta_n\} =: \mathcal{S}_n$ . We write  $S_{n, \theta_n}$  for  $S_n$  with parameters  $\theta_n$ . Hence,  $(S_n)_{n \in \mathbb{N}}$  is completely defined by  $(S_n, \theta_n)_{n \in \mathbb{N}}$ . We call  $\mathcal{S}(n) \equiv \mathcal{S}_n$  the *model space*.

**Definition 1.3.** We say  $(S_n)_{n \in \mathbb{N}}$  is a family of models over the model space  $\mathcal{S}$  iff  $S_n \in \mathcal{S}_n$  for every  $n \in \mathbb{N}$ . As a shorthand, we write  $S \in \mathcal{S}$ .

For our model  $S$ , we want the mutual information between *any* two variables  $X_t, X_{t+\tau}$  to be  $I(X_t, X_{t+\tau}) \in \Omega(\tau^{-\alpha})$ , i.e. it has to hold for every  $t$  and *every*  $S_n$ . Thus, let us define the following:

**Definition 1.4.** We define  $i_{S_n}(\tau)$  (,  $I_{S_n}(\tau)$ ) to be the minimal (, maximal) mutual information between any two variables of  $S_n$  with distance  $\tau$ . Formally, let  $X_t, X_{t+\tau}$  be random variables with distributions defined by  $S_n$ . Then:

- $i_{S_n}(\tau) := \min_{t \in [n-\tau]} I(X_t, X_{t+\tau})$  ,
- $I_{S_n}(\tau) := \max_{t \in [n-\tau]} I(X_t, X_{t+\tau})$  .

**Definition 1.5.** A model  $S$  has *lower bound power-law behavior* iff there exist constants  $c, \alpha \in \mathbb{R}_{>0}$  s.t. for every  $n \in \mathbb{N}$  it holds true that  $i_{S_n}(\tau) \geq c\tau^{-\alpha}$ . Furthermore,  $S$  has *decaying behavior* iff for every  $n \in \mathbb{N}$  we have  $I_{S_{n+\tau}}(\tau) \xrightarrow{\tau \rightarrow \infty} 0$ . Lastly,  $S$  has *power-law behavior* iff it has lower bound power-law behavior and decaying behavior.