

# Got Python?

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- python.org and the “cheese shop”
  - [www.python.org](http://www.python.org)
  - <http://pypi.python.org/pypi>
- Norm Matloff (UC-Davis)
  - [heather.cs.ucdavis.edu/~matloff/python.html](http://heather.cs.ucdavis.edu/~matloff/python.html)
- Dive into Python
  - [diveintopython.net](http://diveintopython.net)
- Lutz's “Programming Python”
  - 3/e uses Python 2.x      4/e uses Python 3.x
  - available from O'Reilly

# Other built-in Python modules to investigate (1)

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- math/cmath
  - higher-order math functions
- argparse
  - powerful command line option parser
- CSV
  - “Comma-Separated Value” file reading and writing
- gzip/zipfile/tarfile/bz2
  - front-door to reading/writing compressed files
- thread/threading
  - low-level and high-level thread control/mgmt.
- os and os.xxxx
  - platform-independent functions ( **VERY USEFUL!** )
- popen2
  - subprocess control/access/mgmt.

# Other built-in Python modules to investigate (2)

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- unittest
  - unit testing framework and test automation
- re
  - regular expression support ( ***VERY USEFUL!*** )
- socket
  - module access to BSD sockets
- time
  - time-related functions and manipulations
- pdb
  - Python debugger (can be called programmatically)

# 3rd party Python resources

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- **ipython**
  - feature-rich interactive Python shell
  - supports “parallel” python computing
  - I use this instead of “python”
- **wxPython/pyQT**
  - Python bindings for GUI widgets libraries
- **Stani's Python editor (SPE) -- or -- Spyder**
  - SPE: not-too-big IDE for Python – Spyder: everything IDE
- **scipy**
  - scientific numerical routines for Python
- **others?**
  - visit [http://www.scipy.org/Topical\\_Software](http://www.scipy.org/Topical_Software)

# **ECE 4723/6723**

# **Embedded Systems**

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## **Lecture**

Embedded Systems Operating System  
(ESOS)

## **Reading:**

# Why do we need an OS?

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- Embedded systems are usually reactive
  - Lots of flow control, handshaking, and conditional execution
- Embedded systems usually do more than one thing “at a time”
  - Most practical embedded system “reactions” have more than one causal condition
    - read sensors and buttons, process HMI, control motors, flash LEDs, send telemetry over comm. links, etc.

***You can write your own application execution framework.  
But WHY would you WANT to ????***

# Why do we need an OS?

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- Lots of choices for embedded OS designer
  - Real-time: soft- vs. hard- vs. non-?
  - Multitasking? Multiprocessing?
  - Threaded vs event-driven?
  - Memory protection? Memory mgmt? Virtual mem?
  - Provided services, libraries, etc.
- Many commercial and free OS-es available
  - Wind River/VxWorks, LynxOS, uCOS-ii/-iii, linux, android
  - Windows CE/Mobile/Phone/Phone7/Embedded/RT\*
  - \*\*or whatever Microsoft is calling it today

***FACT: Any embedded system you will work with that has any complexity to it whatsoever will use an OS!***

# Multitasking contexts

## Roll-your-own

```
while (1) {  
    do audio functions 1;  
    do GUI functions 1;  
    do I/O functions 1;  
    do audio functions 2;  
    do GUI functions 2;  
    do I/O functions 2;  
  
    :  
  
    do audio functions N;  
    do GUI functions N;  
    do I/O functions N;  
}
```

## OS

```
while (1) {  
    call tasks periodically;  
}
```

## audio task

```
while (1) {  
    do audio functions 1;  
    do audio functions 2;  
  
    :  
  
    do audio functions Y;  
}
```

## I/O task

```
while (1) {  
    do I/O functions 1;  
    do I/O functions 2;  
  
    :  
  
    do I/O functions X;  
}
```

## GUI task

```
while (1) {  
    do GUI functions 1;  
    do GUI functions 2;  
  
    :  
  
    do GUI functions Z;  
}
```



# Multitasking

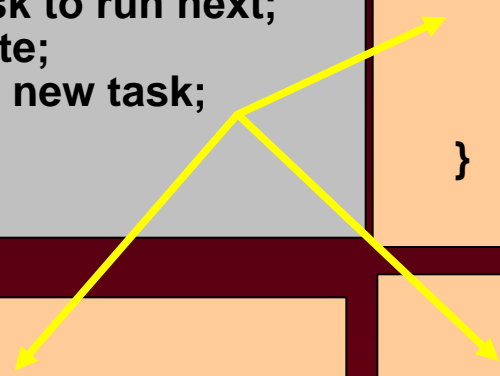
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- Multitask programs, threads, or tasks
- Cooperative multitasking
  - task must choose to give up focus/control
  - Dangerous???
    - one task can kill everything
- Preemptive multitasking
  - some “higher” force can wrestle control away from your task
  - Safer??
    - theoretically, a watchdog could shut down errant task

# Preemptive Multitasking

## Preemptive Multitasking OS

```
OS task timer ISR(void) {  
    save current task's state;  
    determine "best" task to run next;  
    restore new task state;  
    resume execution in new task;  
}
```



## audio task

```
while (1) {  
    copy compressed data from MP3 file;  
    uncompress music data to a buffer;  
    copy buffer to D/A converter;  
    compute elapsed time values;  
    while (D/A converter is busy);  
}
```

## I/O task

```
while (1) {  
    read button/scroll wheel states;  
    update status LEDs;  
    delay (50ms);  
}
```

## GUI task

```
while (1) {  
    update song title on screen;  
    update artist/album on screen;  
    while (elapsed time data is stale);  
    update song's elapsed time;  
}
```

# Cooperative Multitasking

## Cooperative Multitasking OS

```
while(1) {  
    determine "best" task to run next;  
    call that task;  
}
```

## audio task

```
while (1) {  
    copy compressed data from MP3 file;  
    uncompress music data to a buffer;  
    copy buffer to D/A converter;  
    compute elapsed time values;  
    yield until D/A converter is ready for  
    more data;  
}
```

## I/O task

```
while (1) {  
    read button/scroll wheel states;  
    update status LEDs;  
    yield;  
}
```

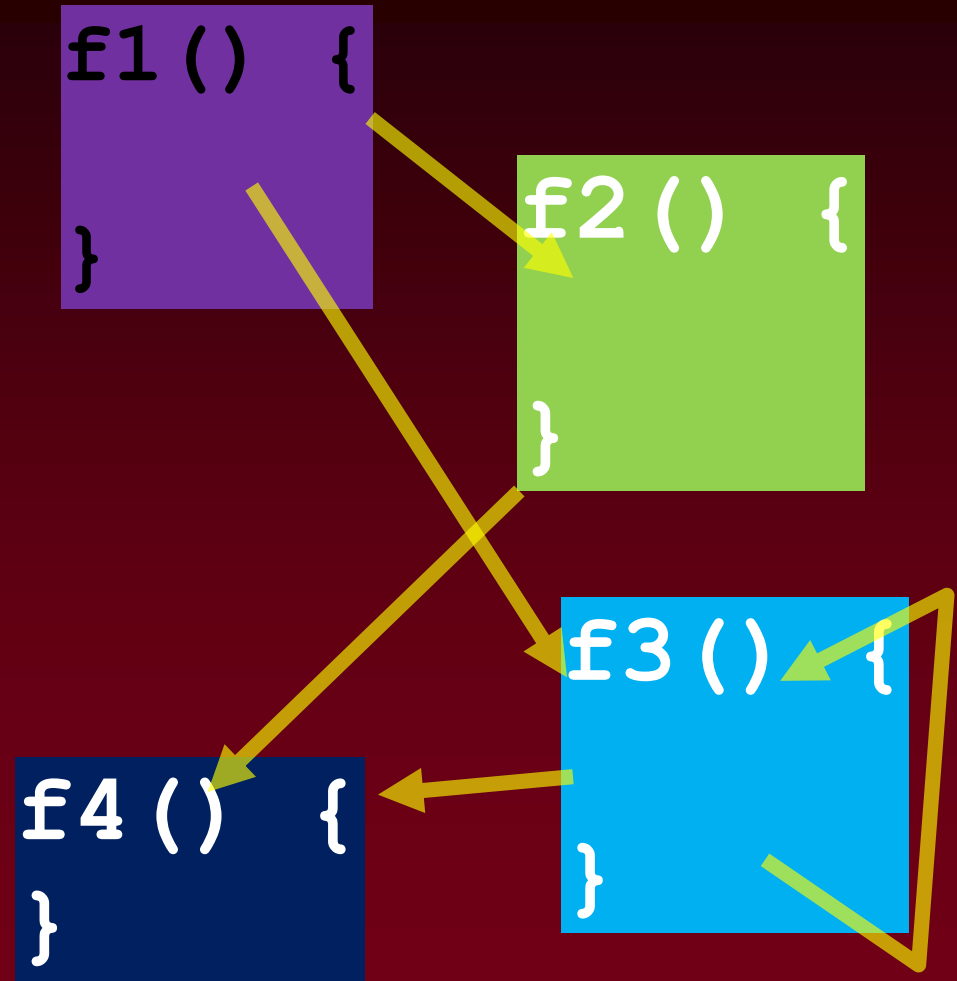
## GUI task

```
while (1) {  
    update song title on screen;  
    update artist/album on screen;  
    yield until elapsed time is updated;  
    update song's elapsed time;  
}
```

# Tasks vs. event-driven function

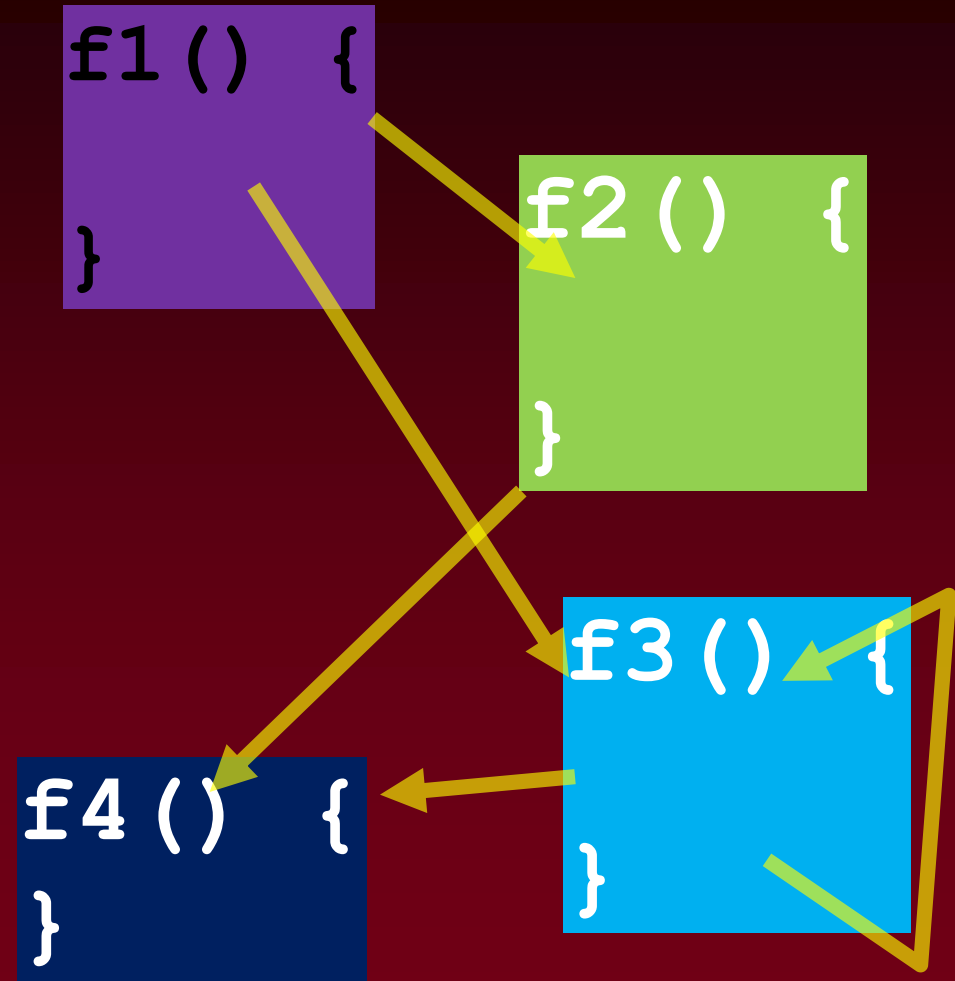
```
TaskX() {  
  
    if() {  
        YIELD();  
    } else {  
        YIELD();  
        while() {  
            YIELD();  
        }  
    }  
}
```

where X is 1, 2, 3, 4



# Event-driven: Problem with explicit state machine flow control

- State machine keeps track of flow control
- Control flow is difficult to follow (at best)
  - Flow is not evident from reading code
- Code is hard to write, understand, debug, and maintain



# Event-driven programs

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- Programs are not not sequential
  - code is often logically “muddled”
- Programs are event-handlers
  - An event-handler is a C function
  - “calling” function is not clear
  - An event-handler must **always** return;
  - An event-handler cannot block/wait for something to happen
    - Handler functions must be “atomic”

# Background:

## Event-driven programming

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- Why use event-driven programming?
  - Often used in memory-constrained systems
  - Threading libraries require too much memory and overhead
- “Backwards” to “normal” programming
  - Code is executed only when events occur
    - Incoming packets, sensor input, mouse clicks, time-outs, HMI
  - Nothing happens without an “event”
    - Programs have no “idle loop”
    - Events are called/triggered by some other entity (OS, GUI, etc.)

# Example of event-driven code

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```
void handle_sensor_input(int data) {  
    if(check_sensor_data(data)) {  
        send_radio_packet(data);  
    }  
    return;  
}
```

```
void handle_incoming_packet(char *packet) {  
    if(packet_should_be_forwarded(packet)) {  
        send_radio_packet(packet);  
    }  
    return;  
}
```



# Threads/Tasks:

## No explicit state machine needed

```
THREAD (f()) {  
  
    if() {  
        THREAD_WAIT_UNTIL();  
    } else {  
        THREAD_WAIT_UNTIL();  
        while() {  
            THREAD_WAIT_UNTIL();  
        }  
    }  
}
```

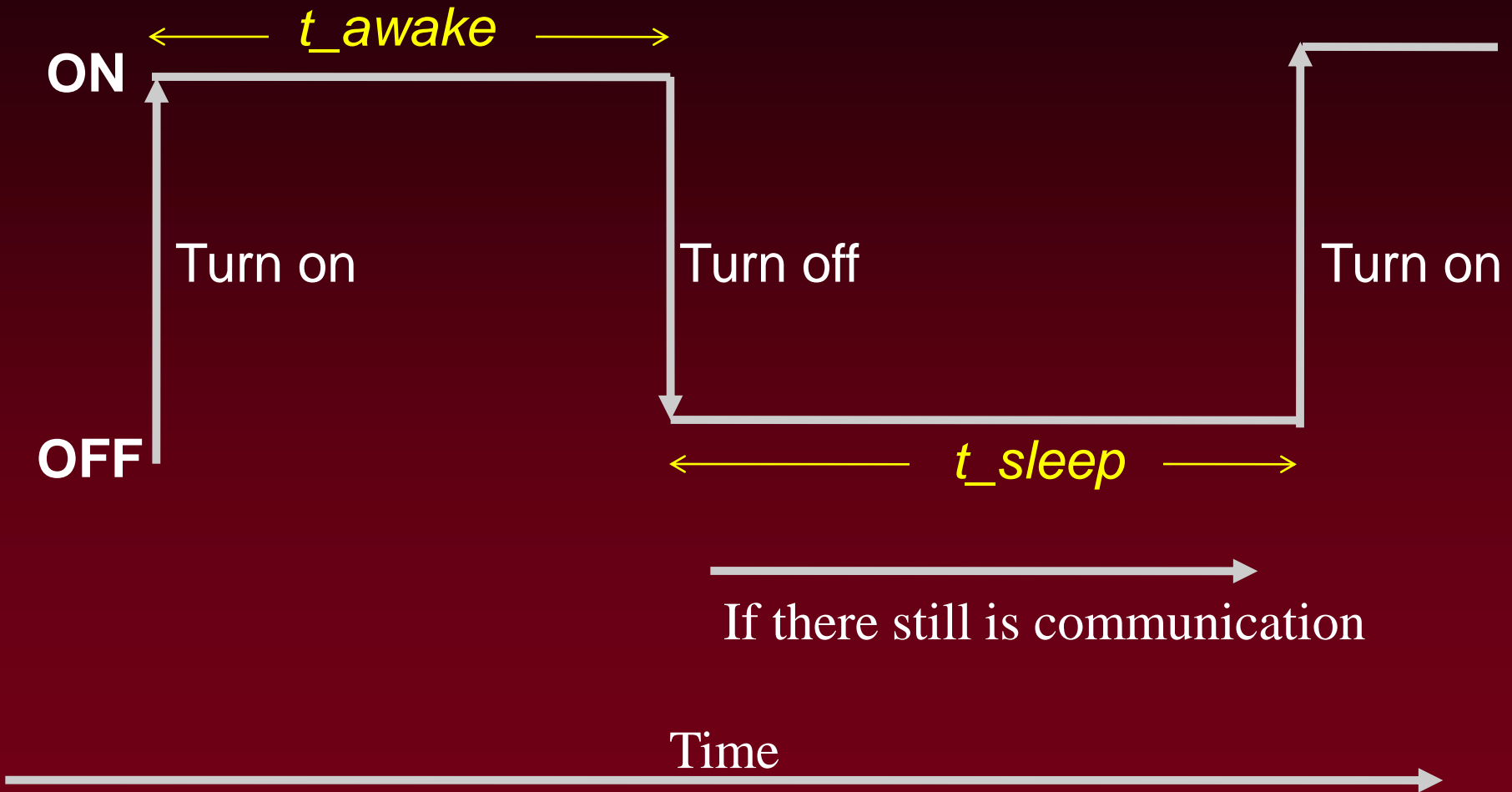
- With threads, code looks “sequential”
- Flow control is apparent
- Code is easier to understand, debug, and maintain
- Easy to write once you start thinking “threaded”

# The STACK!

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- Threading
  - every thread requires a stack
  - stacks contain dead space
  - may use large amounts of the available memory
  - thread “manager” required
- Event-driven
  - only one stack is needed.
    - more efficient use of memory
  - Stack rewound on each event
  - Event-handlers are normal functions that explicitly return

# Example: radio sleep cycle



# ***Example: radio sleep cycle***

## ***6-step informal specification***

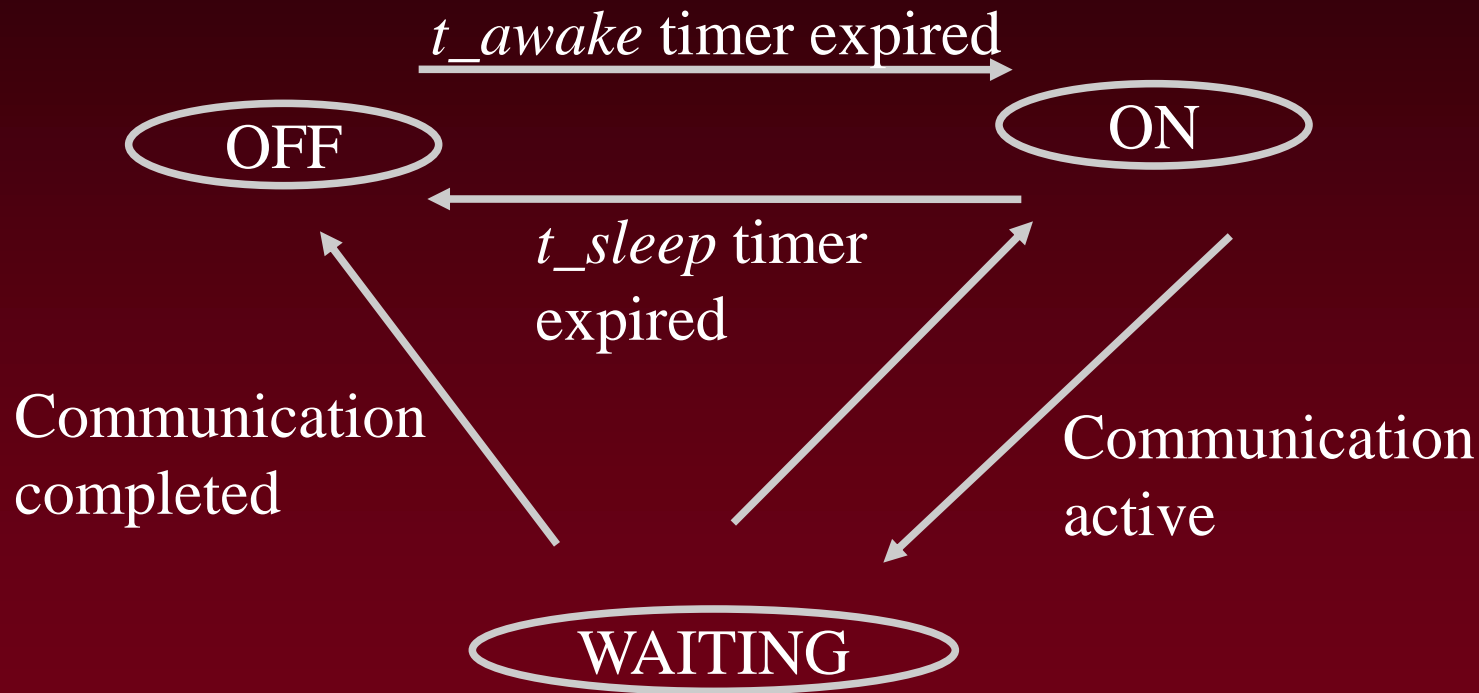
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1. Turn radio on.
2. Wait for  $t\_awake$  milliseconds.
3. Turn radio off, but only if all communication has completed.
4. If communication has not completed, wait until it has completed.  
Then turn off the radio.
5. Wait for  $t\_sleep$  milliseconds. If the radio could not be turned off before  $t\_sleep$  milliseconds because of remaining communication, do not turn the radio off at all.
6. Repeat from step 1.

***Problem with events:***  
***We can't write this as a 6-step program!***

# State machine implementation

With events, we must use an explicit state machine!



# Our state machine coded in C

```
enum {  
    ON,  
    WAITING,  
    OFF  
} state;
```

*Called by somebody...*

```
void radio_wake_eventhandler() {  
    switch(state) {  
        case OFF:  
            if(timer_expired(&timer)) {  
                radio_on();  
                state = ON;  
                timer_set(&timer, T_AWAKE);  
            }  
            break;
```

***Quite complex!***  
***And not even correct!***

```
        case ON:  
            if(timer_expired(&timer)) {  
                timer_set(&timer, T_SLEEP);  
                if(!communication_complete()) {  
                    state = WAITING;  
                } else {  
                    radio_off();  
                    state = OFF;  
                }  
            }  
            break;  
        case WAITING:  
            if(communication_complete() || \  
                timer_expired(&timer)) {  
                state = ON;  
                timer_set(&timer, T_AWAKE);  
            } else {  
                radio_off();  
                state = OFF;  
            }  
            break;  
    }  
}
```

# Have our cake and eat it too!

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**We want the memory  
efficiency of events**

**AND**

**the code clarity and  
maintainability of  
isolated threads/tasks!**

# Our cake: MSU's *Embedded Systems Operating System* (ESOS)

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- ESOS “threading” is based on protothreads
  - ANSI C code from Adam Dunkels (*www.sics.se*)
- Protothreads are a mixture of the event-driven and “true” threads
  - stackless, non-preemptive threading
  - can be driven by an event-handler
  - With protothreads, we can write blocking waits, inside an event-handler
    - `ESOS_TASK_WAIT_UNTIL()` – conditional blocking
    - `ESOS_TASK_WAIT_WHILE()` – conditional blocking



# ESOS implementation

```
ESOS_USER_TASK( radio_wake_thread ) {  
    ESOS_TASK_BEGIN();  
    while( TRUE ) {  
1      radio_on();  
        timer_set(&timer, T_AWAKE);  
2      ESOS_TASK_WAIT_UNTIL(timer_expired(&timer));  
        timer_set(&timer, T_SLEEP);  
        if(!communication_complete()) {  
3, 4      ESOS_TASK_WAIT_UNTIL(  
            communication_complete() || timer_expired(&timer));  
        }  
        if(!timer_expired(&timer)) {  
5          radio_off();  
            ESOS_TASK_WAIT_UNTIL( timer_expired(&timer) );  
        }  
6    }  
    ESOS_TASK_END();  
}
```

# ESOS implementation

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- Provides sequential code flow
  - The 6-step informal specification visible in the code
- Possible to use C control structures
  - if(), while(), for(), etc.
- Debug stack retained
  - Possible to follow calls
- Implicit blocking calls evident

# Limitations in ESOS

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- “Automatic” variables (stack variables) are **NOT** saved across blocking waits
  - Programmer must manually save automatic variables
  - However, **static** local variables still work as expected
- Standard C compiler may issue some warnings
- C language “switch” statements can not be used across “waits”
  - Not a big deal really
  - use *if-else if-else* construct instead (for short switch-es)
  - Refactor (for large switch statements)

# ESOS thread functions (1)

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- **ESOS\_TASK\_WAIT\_UNTIL( cond )**
  - blocks current task until condition is TRUE
- **ESOS\_TASK\_WAIT\_WHILE( cond )**
  - blocks current task while condition is TRUE
- **ESOS\_TASK\_YIELD()**
  - blocks current task until its next execution opportunity
- **ESOS\_TASK\_RESTART()**
  - blocks and re-inits the current task. Task will start executing at its beginning.
- **ESOS\_TASK\_EXIT()**
  - causes the current task to exit/end. If the ESOS task was spawned by a parent task, the parent task will become unblocked and resume execution

# ESOS thread functions (2)

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- **ESOS\_TASK\_SLEEP()**
  - current task goes to sleep and will be blocked until some other task wakes it up
- **ESOS\_TASK\_WAKE( pstTask )**
  - wakes up task pstTask. Task pstTask executes at next opportunity
- **ESOS\_TASK\_SEM\_INIT( s, v )**
  - creates a semaphore s with initial value v
- **ESOS\_TASK\_WAIT\_SEMAPHORE( s, val )**
  - Task will wait until semaphore s signaled val times
- **ESOS\_SIGNAL\_SEMAPHORE( s, val )**
  - signal (increment) the s semaphore val times

# ESOS thread functions (3)

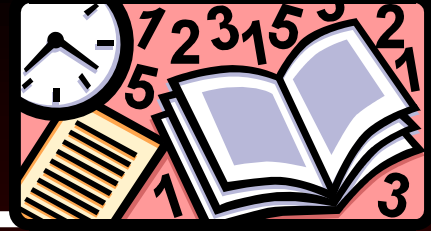
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- **ESOS\_ALLOCATE\_CHILD\_TASK( thChild )**
  - Allocates/initializes the task handle thChild
- **ESOS\_TASK\_SPAWN\_AND\_WAIT( thChild, pfnChild )**
  - Current task waits on the completion (exit/death) of child task thChild that runs function pfnChild.

***Next time, we will look at the***

- ***structure of an ESOS application***
- ***other ESOS services***
  - ***communications***
  - ***interrupt control***
  - ***timer services***

# References



New PIC24/dsPIC33 users:

Read Chapters 1-9 in R/B/J

Read Chapter 14 in R/B/J

*You may want to build a few of the examples from Chapter 14.*

Little Book of Semaphores

*Reference 74 in R/B/J bibliography*

Adam Dunkel's Protothreads

*Reference 75 in R/B/J bibliography*

Read ECE4723 C language coding conventions

Ganssle Chapt. 1-3

*If you bought this nice little book*