

FONDIND

1. PRELIMINARY

Definition 1.1. A *Reedy category* \mathbb{A} consists of the following data:

- A small category \mathbb{A} .
- A strict factorisation system $(\mathbb{A}_-, \mathbb{A}_+)$. We consider \mathbb{A}_- and \mathbb{A}_+ as wide subcategories of \mathbb{A} .
- A function $d: \mathbf{Obj} \mathbb{A} \rightarrow \omega$ that extends to identity-reflecting functors $d: \mathbb{A}_-^{\text{op}} \rightarrow \omega$ and $d: \mathbb{A}_+ \rightarrow \omega$, where the set of natural numbers ω is seen as a linearly ordered set. \blacklozenge

Notation 1.2. Let \mathbb{A} be a Reedy category and $X \in \mathbf{Psh}[\mathbb{A}]$ be a presheaf. We view the category \mathbb{A} as a full subcategory of $\mathbf{Psh}[\mathbb{A}]$.

- We mean by \twoheadrightarrow a morphism in \mathbb{A}_- , while by \twoheadrightarrow , we mean a morphism in \mathbb{A}_+ .
- A *cell* is a morphism $a \rightarrow X$ in $\mathbf{Psh}[\mathbb{A}]$ whose domain a is in \mathbb{A} .
- A cell $x: a \rightarrow X$ is *degenerate* if there is a factorisation $a \xrightarrow{\sigma} b \xrightarrow{y} X$ of x such that σ is non-identity morphism in \mathbb{A}_- .
- For each cell $x: a \rightarrow X$, a *face* of x is a cell $y: b \rightarrow X$ that factors as $b \xrightarrow{\iota} a \xrightarrow{x} X$ with ι in \mathbb{A}_+ , while a *degeneracy* of x is a cell $y: b \rightarrow X$ that factors as $b \xrightarrow{\sigma} a \xrightarrow{x} X$ with σ in \mathbb{A}_- . \blacklozenge

Lemma 1.3. Let \mathbb{A} be a Reedy category. For each presheaf $X \in \mathbf{Psh}[\mathbb{A}]$ and any cell $x: a \rightarrow X$, there exists a factorisation $a \xrightarrow{\sigma} b \xrightarrow{y} X$ of x such that y is non-degenerate.

Proof. Consider the set of factorisations of x by morphisms in \mathbb{A}_- . This is not empty because (id_a, x) gives one such factorisation. Therefore, there exists a factorisation (σ, y) whose degree of the domain of y is minimum among this set. It suffices to show y is non-degenerate. To this end, suppose y factors as $z \circ \tau$ where τ is a morphism in \mathbb{A}_- . Since $(\tau \circ \sigma, z)$ gives a factorisation of x , the degree of the domain of z is not less than that of y . Therefore, τ must be an identity because $d: \mathbb{A}_-^{\text{op}} \rightarrow \omega$ reflects identity. \square

Proposition 1.4. For a Reedy category \mathbb{A} , the following are equivalent.

- For each presheaf $X \in \mathbf{Psh}[\mathbb{A}]$, any cell $x: a \rightarrow X$ factors uniquely as $a \xrightarrow{\sigma} b \xrightarrow{y} X$ where σ is in \mathbb{A}_- and y is non-degenerate.
- Any span in \mathbb{A}_- admits a pushout in \mathbb{A} that is preserved by the yoneda $\mathbb{A} \hookrightarrow \mathbf{Psh}[\mathbb{A}]$.
- Any span in \mathbb{A}_- admits an absolute pushout in \mathbb{A} .

Proof. The equivalence of **ii** and **iii** follows from a folklore result: a small colimit is absolute if and only if it is preserved by the cocompletion.

[**i** \Rightarrow **ii**] Suppose **i** and that we are given a span $b \xleftarrow{\sigma} a \xrightarrow{\sigma'} b'$ in \mathbb{A}_- , and take a pushout in $\mathbf{Psh}[\mathbb{A}]$ as in the diagram below. Since the yoneda embedding is fully faithful, it suffices to show the pushout is representable. Let us factorise x and x' by **i** as follows so that the cells y and y' are non-degenerate.

$$\begin{array}{ccccc}
 a & \xrightarrow{\sigma'} & b' & & \\
 \sigma \downarrow & & \downarrow x' & \searrow \tau' & \\
 & & & & c' \\
 & \nearrow y' & & & \\
 b & \xrightarrow{x} & X & & \\
 \tau \searrow & & \nearrow y & & \\
 & & c & &
 \end{array}$$

By the uniqueness of the factorisation of $x \circ \sigma$, we have $(\tau \circ \sigma, y) = (\tau' \circ \sigma', y')$. By the universality of the pushout applying to the square $\tau \circ \sigma = \tau' \circ \sigma'$, we obtain a section u of $y = y'$. Therefore,

we have a factorisation $(u \circ y, y)$ of y . Since y is non-degenerate, through factorising $u \circ y$ by the factorisation system, we conclude $u \circ y$ is an endo-morphism in \mathbb{A}_+ , which shows $u \circ y = \text{id}_c$ and that X is representable.

[ii \Rightarrow i] By [Lemma 1.3](#), it suffices to show the uniqueness of such factorisation (σ, y) for any cell x . Suppose ii and that we are given a cell x and its two factorisations (σ, y) and (σ', y') where y and y' are non-degenerate. By taking the pushout of the span formed by σ and σ' , we obtain a square $\tau \circ \sigma = \tau' \circ \sigma'$ in \mathbb{A} and factorisations (τ, z) and (τ', z') of y and y' respectively. Since the left class of an orthogonal factorisation system must be closed under pushout, τ and τ' are in \mathbb{A}_- . Moreover, since y and y' are non-degenerate, τ and τ' are identities, and this shows $(\sigma, y) = (\sigma', y')$. \square

Definition 1.5. A Reedy category is *elegant* if the equivalent conditions in [Proposition 1.4](#) hold. \blacklozenge

Proposition 1.6. For an elegant Reedy category \mathbb{A} , the left class \mathbb{A}_- coincides with the wide subcategory of split epimorphisms.

Proof. \square

2. FONDIND

Let \mathbb{A} be a Reedy category. The following definition is after Yuki Maehara.

Definition 2.1. A presheaf X is *fondind* if any face of a non-degenerate cell is non-degenerate: i.e., X is fondind if for each factorisation $b \xrightarrow{\iota} a \xrightarrow{x} X$ of a cell $y: b \rightarrow X$ with ι in \mathbb{A}_+ , the cell y is non-degenerate whenever x is as well. \blacklozenge

Proposition 2.2. The class of fondind presheaves is closed under limits and subobjects.

Proof. It suffices to show the class is closed under products and subobjects. [straightforward](#) \square

Note that for each presheaf X , the category of elements \mathbb{A}/X has as its objects cells of X . We write $(\mathbb{A}/X)_{\text{nd}}$ for the full subcategory of \mathbb{A}/X consisting of all non-degenerate cells.

Lemma 2.3. Let X be a presheaf. Consider the following commutative diagram in $\mathbf{Psh}[\mathbb{A}]$.

$$\begin{array}{ccc} \cdot & \xrightarrow{\sigma} & \cdot \\ \sigma' \downarrow & \nearrow f & \downarrow x \\ \cdot & \xrightarrow{x'} & X \end{array}$$

Suppose that x is non-degenerate. Then we have $\sigma = \sigma'$ and $x \circ i = x' \circ i'$.

Proof. Considering the factorisation of the morphism f in \mathbb{A} , we can conclude f is a morphism in \mathbb{A}_+ since x is non-degenerate. Now since both (σ', i') and $(\sigma, f \circ i)$ give the $(\mathbb{A}_-, \mathbb{A}_+)$ -factorisation of the same morphism, they coincide. This implies $\sigma = \sigma'$ and $x \circ i = x' \circ f \circ i' = x' \circ i'$. \square

Proposition 2.4. A presheaf $X \in \mathbf{Psh}[\mathbb{A}]$ is fondind if $(\mathbb{A}/X)_{\text{nd}}$ is a final full subcategory of \mathbb{A}/X .

Proof. Suppose $(\mathbb{A}/X)_{\text{nd}}$ is a final full subcategory of \mathbb{A}/X . To show X is fondind, take a factorisation $a \xrightarrow{i} b \xrightarrow{x} X$ of a cell $a \xrightarrow{y} X$ by a morphism i in \mathbb{A}_+ and a non-degenerate cell x . We will show y is non-degenerate. To this end, consider another factorisation $a \xrightarrow{\sigma} c \xrightarrow{z} X$ of y obtained by [Lemma 1.3](#). Since z is non-degenerate, it suffices to show $\sigma = \text{id}$. Since $(\mathbb{A}/X)_{\text{nd}}$ is final in \mathbb{A}/X , the comma category $y \downarrow (\mathbb{A}/X)_{\text{nd}}$ is connected. In particular, (σ, z) and (i, x) are connected by a zigzag in this comma category. Now, iterated use of [Lemma 2.3](#) to this zigzag shows $\sigma = \text{id}$. \square

Proposition 2.5. The converse of [Proposition 2.4](#) is true when \mathbb{A} is elegant; i.e., if \mathbb{A} is elegant, a presheaf X is fondind if and only if $(\mathbb{A}/X)_{\text{nd}}$ is final in \mathbb{A}/X .

Proof. Suppose we are given a fondind presheaf X and a cell $a \xrightarrow{x} X$. It suffices to show the comma category $x \downarrow (\mathbb{A}/X)_{\text{nd}}$ is connected. Since it is inhabited by [Lemma 2.3](#), it suffices to show that any two factorisations of x by non-degenerate cells are connected in this category. To this end, take two factorisations $a \xrightarrow{f} y \rightarrow X$ and $a \xrightarrow{f'} y' \rightarrow X$ of x where y and y' are non-degenerate. Moreover, let us

decompose f and f' by the factorisation system as $f = i \circ \sigma$ and $f' = i' \circ \sigma'$. Since X is fondind, $y \circ i$ and $y' \circ i'$ are non-degenerate, and hence $(\sigma, y \circ i)$ and $(\sigma', y' \circ i')$ become objects of the comma category. Moreover, since \mathbb{A} is elegant, those two factorisations coincide, and now we have a span $(f, y) \xleftarrow{i} (\sigma, y \circ i) = (\sigma', y' \circ i') \xrightarrow{i'} (f', y')$ in $x \downarrow (\mathbb{A}/X)_{\text{nd}}$. \square