

Rodney Wotton

I pledge my honor that I have abided by the Stevens Honor System - Rodney Wotton

1.

- a. There are .04167 FLOPS/byte. We get this by seeing that the code has two loads and store them in one number. One operand means that we only have one FLOP.  $1/((2+1)*8) = .04167$
- b. There are .04167 FLOPS/byte. We get this by seeing that the code has two loads and store them in one number. One operand means that we only have one FLOP.  $1/((2+1)*8) = .04167$
- c. We get .0625 FLOPS/byte. We get this by seeing that the code has one load and store them in one number. One operand means that we only have one FLOP.  $1/((1+1)*8) = .0625$
- d. We get .0625 FLOPS/byte. We get this by seeing that the code has three loads and store them in one number. We have two operands in his equation which means that we have 2 FLOPs.  $2/((3+1)*8) = .0625$
- e. We get .0833 FLOPS/byte. We get this by seeing that the code has two loads and store them in one number. We have two operands in his equation which means that we have 2 FLOPs.  $2/((2+1)*8) = .0833$

2.

- a. The Peak is  $2.5 \times 8 \times 8 \times 32$  which = 640 GFLOPS. We get this by multiplying the clockrate \* #SIMD \* FP
- b. New peak is  $2.5 \times 8 \times 32 \times 2 = 1280$ . The throughput is calculated by comparing (2 load + 1 store) \* 4 bytes = 12 bytes/2 flops or 6 bytes/flop  
 $6 \text{ bytes}/1 \text{ FLOP} \times 1280 \text{ GFLOP/sec} = 7680 \text{ GB/s} > 112 \text{ GB/s}$ .  
NOT SUSTAINABLE

3. To get this answer we must first use the equation  $1/((1-F) + F/\text{processor}) = \text{speedup}$ . In this instance, F is parallelizable parts and 1-F is sequential. So we get  $1/((1-F)+F/2000) = 100$   
 $F = .99$ .

Therefore we get the Sequential Parts = .95%

Parallelizable Parts= 99.05%

4.

- a. 1.  $P * B = 16$   
2.  $2 * (P - \text{sqrt}(P)) = 24$   
3.  $(P * (P - 1) / 2) * B = 120$
- b. 1.  $2*B=2$   
2.  $\text{sqrt}(P) * B = 4$   
3.  $(P/2)^2 * B=64$
- c. 1 link

2. 9 links
3. 105 links