

# RiNNAL+: a Riemannian ALM Solver for SDP-RLT Relaxations of Mixed-Binary QP

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# Outline

- 1 Problem: SDP-RLT relaxations of mixed-binary QP
- 2 RiNNAL+: Riemannian ALM solver
- 3 Numerical experiments

# Mixed-binary QP

$$\begin{array}{ll} \min_x & x^\top Qx + 2c^\top x \\ \text{s.t.} & Ax = b, \ Bx \leq d, \ x \in \{0,1\}^n. \end{array} \quad (\text{MBQP})$$

- **Applications:** Max-Cut, clustering, portfolio optimization ...
- **Generalization:** mixed-binary, QCQP.
- **Global optimality:** solve **convex relaxations** to obtain lower bound.
- **Challenge:** tradeoff between **tightness** and **computational cost**.

# SDP Relaxation

$$\begin{aligned} \min_x \quad & x^\top Q x + 2c^\top x \\ \text{s.t.} \quad & Ax = b, \quad Bx \leq d, \quad x \in \{0, 1\}^n. \end{aligned} \quad (\text{MBQP})$$

$$\begin{aligned} \min_{X, x} \quad & \langle Q, X \rangle + 2c^\top x \\ \text{s.t.} \quad & Ax = b, \quad Bx \leq d, \quad \text{diag}(X) = x, \\ & \begin{bmatrix} 1 & x^\top \\ x & X \end{bmatrix} \succeq 0. \end{aligned} \quad (\text{SDP})$$

Relaxation	Size	Bound	Speed
SDP	$\Omega(n^2)$	Loose	Fast

# DNN Relaxation

$$\begin{aligned} \min_x \quad & x^\top Qx + 2c^\top x \\ \text{s.t.} \quad & Ax = b, \quad Bx + s = d, \quad x \in \{0, 1\}^n, \quad (x, s) \geq 0. \end{aligned} \quad (\text{MBQP})$$

$$\begin{aligned} \min_{X, x} \quad & \langle Q, X \rangle + 2c^\top x \\ \text{s.t.} \quad & \begin{bmatrix} -b & A & 0 \\ -d & B & I \end{bmatrix} \hat{X} = 0, \quad \text{diag}(X) = x, \\ & \hat{X} = \begin{bmatrix} 1 & x^\top & s^\top \\ x & X & Z^\top \\ s & Z & W \end{bmatrix} \succeq 0 \text{ and } \geq 0. \end{aligned} \quad (\text{DNN})$$

Relaxation	Size	Bound	Speed
SDP	$\Omega(n^2)$	Loose	Fast
DNN <sup>1</sup>	$\Omega((n + \ell)^2)$	Tight	Slow

<sup>1</sup>Burer, 2010.

# SDP-RLT Relaxation

$$\min_{X, x} \quad \langle Q, X \rangle + 2c^\top x$$

$$\text{s.t.} \quad Ax = b, \quad Bx \leq d, \quad \text{diag}(X) = x,$$

$$AX - bx^\top = 0, \quad X \geq 0,$$

$$dx^\top - BX \geq 0,$$

$$BXB^\top - Bxd^\top - dx^\top B^\top + dd^\top \geq 0,$$

$$\begin{bmatrix} 1 & x^\top \\ x & X \end{bmatrix} \succeq 0.$$

(SDP-RLT)

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<sup>2</sup>Sherali and Adams, 1990.

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## Theorem

*The DNN and SDP-RLT relaxations of (MBQP) yield the same optimal value.*

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## Tightening Strategy:

$$\min \left\{ x^T Q x + 2c^T x : x \in \{0, 1\}^n \right\} \quad (\text{BIQ})$$

**Strengthen** by adding redundant bound constraint:

$$\min \left\{ x^T Q x + 2c^T x : \mathbf{x} \leq \mathbf{e}, x \in \{0, 1\}^n \right\} \quad (\text{BIQ-S})$$

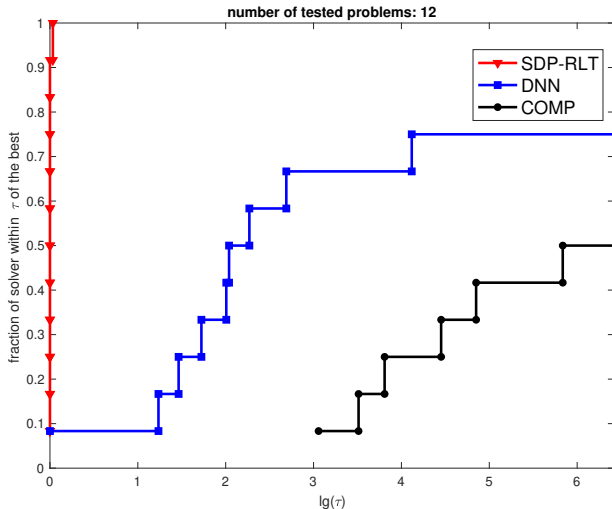
Data: bqp100.1	(BIQ)		(BIQ-S)	
UB: -7.97e3	lower bound	% gap	lower bound	% gap
(SDP)	-8.72e3	9.424	-8.72e3	9.424
(DNN)	-8.38e3	5.149	-8.04e3	0.836
(SDP-RLT)	-8.38e3	5.149	-8.04e3	0.836

↓  
**Tighter**



# Time Comparison

- Solving SDP-RLT is **faster** due to lower dimensionality.



# How to Solve SDP-RLT?

**Conclusion:** SDP-RLT delivers the **same relaxation bound** as the DNN approach but with substantially **lower computational cost**.

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- $\Omega(n^2)$  matrix variables.
- $\Omega(nm)$  equality constraints from  $Ax = b$ .
- $\Omega(n\ell)$  inequality constraints from  $Bx \leq d$ .
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## Limitations of existing solvers:

- Interior-point methods:  $\mathcal{O}(n^6)$  per iteration.
- ADMM/ALM: costly eigenvalue decompositions.
- RiNNAL: applicable to DNN, not to SDP-RLT.

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## Contribution:

- **RiNNAL+**: a two-phase Riemannian ALM method.
- Exploits **low-rank structure** and handles **many constraints** efficiently.

# SDP-RLT Relaxation

## Compact form

$$\begin{array}{ll} \min_Y & \langle C, Y \rangle \\ \text{s.t.} & Y \in \mathcal{I} \cap \mathcal{E} \cap \mathbb{S}_+^{n+1}. \end{array} \quad (\text{SDP-RLT})$$

**Variables and data:**

$$C := \begin{bmatrix} 0 & c^\top \\ c & Q \end{bmatrix}, \quad Y = \begin{bmatrix} z & x^\top \\ x & X \end{bmatrix}.$$

**Feasible set:**

$$(\text{Equalities}) \quad \mathcal{E} := \left\{ Y \in \mathbb{S}^{n+1} : \mathcal{A}(X) = g \right\},$$

$$(\text{Inequalities}) \quad \mathcal{I} := \left\{ Y \in \mathbb{S}^{n+1} : \mathcal{B}(X) \geq h \right\}.$$

# Augmented Lagrangian Method for (SDP-RLT)

Reformulation:

$$\begin{aligned} \min_Y \quad & \langle C, Y \rangle + \delta_{\mathcal{E} \cap \mathbb{S}_+^{n+1}}(Y) \\ \text{s.t.} \quad & \mathcal{B}(Y) \geq h. \end{aligned}$$

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Augmented Lagrangian function:

$$L_\sigma(Y; \lambda) := \langle C, Y \rangle + \frac{\sigma}{2} \|\sigma^{-1} \lambda - (\mathcal{B}(Y) - h)\|^2.$$



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## ALM Iteration

$$\begin{aligned} Y^{k+1} &= \arg \min \left\{ \phi(Y) := L_{\sigma_k}(Y; \lambda^k) \mid Y \in \mathcal{E} \cap \mathbb{S}_+^{n+1} \right\}, & (\text{ALM-sub}) \\ \lambda^{k+1} &= \lambda^k - \sigma_k (\mathcal{B}(Y^{k+1}) - g). \end{aligned}$$

**Challenge:** How to solve (ALM-sub) efficiently?

# A Hybrid Method to Solve ALM Subproblem

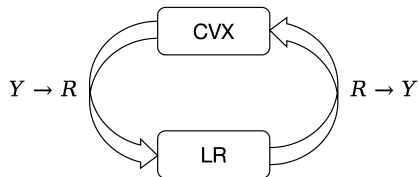


Figure: Two-phase transition.

A diagram illustrating low-rank factorization. A large square box labeled  $Y$  is shown on the left, followed by an equals sign. To the right of the equals sign are two terms. The first term is a vertical yellow rectangle divided into two parts: a small top section labeled  $e_1^T$  and a larger bottom section labeled  $R$ . Below this rectangle is a curly brace labeled  $\hat{R}$ . The second term is a horizontal yellow rectangle divided into two parts: a small left section labeled  $e_1$  and a larger right section labeled  $R^T$ . Below this rectangle is a curly brace labeled  $\hat{R}^T$ .

Figure: Low-Rank factorization.

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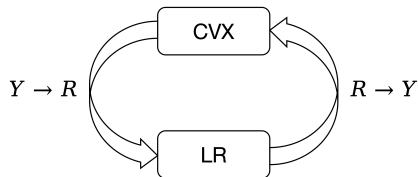


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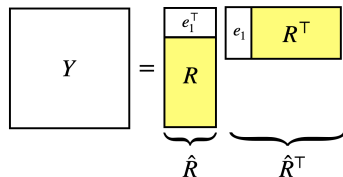


Figure: Low-Rank factorization.

## CVX phase

Projected gradient descent on  $Y$

$$\begin{aligned} \min_Y \quad & \phi(Y) \\ \text{s.t.} \quad & Y \in \mathcal{E} \cap \mathbb{S}_+^{n+1}. \end{aligned}$$

- Ensures **global convergence**.
- Automatically adjusts the **rank**.

## LR phase

Riemannian gradient descent on  $R$

$$\begin{aligned} \min_R \quad & \phi(\hat{R} \hat{R}^T) \\ \text{s.t.} \quad & \hat{R} \hat{R}^T \in \mathcal{E}. \end{aligned}$$

- Reduces **dim** and **#constraints**.
- Reduces the **objective value**.

# CVX Phase: Computation

**Perform one PG step:** (Lee et al., 2024)

$$G = Y - t \nabla \phi(Y),$$

$$Y^+ = \Pi_{\mathcal{E} \cap \mathbb{S}_+^{n+1}}(G).$$

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**Projection subproblem:** Generalized nearest correlation matrix problem:

$$Y^+ = \operatorname{argmin} \left\{ \frac{1}{2} \|Y - G\|^2 : \operatorname{diag}(Y) = e, \langle NN^\top, Y \rangle = 0, Y \succeq 0 \right\}, \quad (\text{Proj-sub})$$

efficiently solved via [Semismooth Newton-CG method](#) (Qi and Sun, 2006).

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**Acceleration:** (10+ times faster)

- **Preprocessing:** transform  $\operatorname{diag}(X) = x$  to  $\operatorname{diag}(Y) = e$ .
- **Warm-start:** recover dual variables of (Proj-sub) from LR phase.

# CVX Phase: Auto-rank

## Rank-adaptive strategy:

- Rank Increase: Escape saddle points by following the eigenvector direction with the smallest eigenvalue. (Journée et al., 2010)
- Rank Decrease: Reduce computational cost by truncating small singular values. (Gao and Absil, 2022)
- Frequently used: (Boumal, 2015; J. Wang and Hu, 2025; Y. Wang et al., 2023)

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- Requires careful parameter tuning (step sizes, thresholds).
- Truncating small singular values may increase the function value.



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- Truncating small singular values may increase the function value.

## PG Step:

- Automatically adjust the rank.
- parameter-free.
- monotonically decreasing.

# CVX Phase: Auto-rank

## BIQ experiment:

- Data: bqp500.1.
- observation: PG identifies the correct rank rapidly in both scenarios.

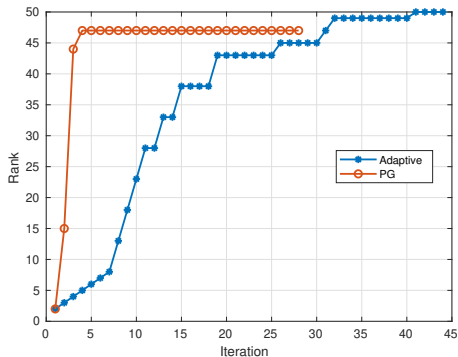


Figure: Initial rank:  $r = 2$ .

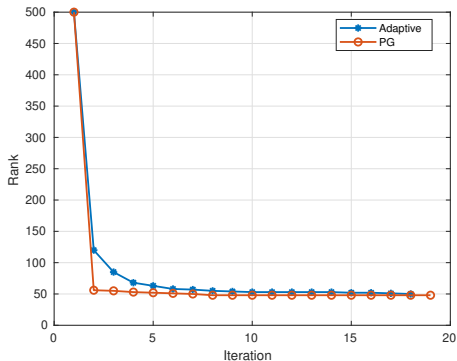


Figure: Initial rank:  $r = n$ .

# LR Phase: RiNNAL Framework

**Manifold:**

$$Y \in \mathcal{E} \cap \mathbb{S}_+^{n+1} \iff R \in \mathcal{M}_r,$$

$$\mathcal{M}_r := \left\{ R \in \mathbb{R}^{n \times r} : AR = be_1^\top, \text{diag}(RR^\top) = Re_1 \right\}.$$

The diagram shows a square matrix  $Y$  on the left, followed by an equals sign. To the right of the equals sign is a vertical yellow rectangle labeled  $R$  with a bracket underneath labeled  $\hat{R}$ . To the right of  $R$  is a horizontal yellow rectangle labeled  $R^\top$  with a bracket underneath labeled  $\hat{R}^\top$ . Above the  $R$  block is a small white box containing  $e_1^\top$ , and above the  $R^\top$  block is a small white box containing  $e_1$ . The entire right-hand side represents the product  $Y = R R^\top$ .

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- **Dim:** Reduce  $\mathcal{O}(n^2) \rightarrow \mathcal{O}(nr)$ .
- **Constraints:** Reduce  $\Omega(mn) \rightarrow \Omega(mr)$ .

$$Y = \underbrace{\begin{bmatrix} e_1^\top \\ R \end{bmatrix}}_{\hat{R}} \underbrace{\begin{bmatrix} e_1 & R^\top \end{bmatrix}}_{\hat{R}^\top}$$

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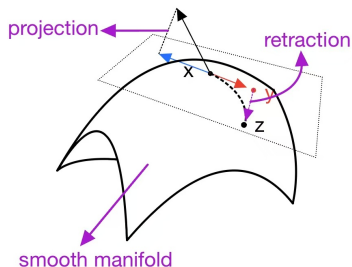
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## Properties:

- **Smoothness:** Random perturbation.
- **Projection:** Schur-based reduction.
- **Retraction:** Convex reformulation (solved by Weiszfeld algorithm).
- **Dual recovery:**  
(LR) KKT  $\Rightarrow$  partial (CVX) KKT.



# LR Phase: Smoothness

**Spherical formulation of  $\mathcal{M}_r$ :**

$$\mathcal{M}_r = \left\{ R \in \mathbb{R}^{n \times r} : AR = be_1^\top, \text{diag}_B((2R - ee_1^\top)(2R - ee_1^\top)^\top) = e \right\}.$$

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**Perturbed manifold  $\mathcal{M}_{r,v}$ :**

$$\mathcal{M}_{r,v} = \left\{ R \in \mathbb{R}^{n \times r} : AR = be_1^\top, \text{diag}_B((2R - ee_1^\top)(2R - ee_1^\top)^\top) = e + v \right\},$$

where  $v \in \mathbb{R}^p$  is random with  $\|v\| = \varepsilon > 0$ .

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where  $v \in \mathbb{R}^p$  is random with  $\|v\| = \varepsilon > 0$ .

**Lemma (Tang and Toh, 2024b)**

*For a generic  $v$ , every point of  $\mathcal{M}_{r,v}$  satisfies LICQ.*

**Implementation:** switch to  $\mathcal{M}_{r,v}$  if encounter nonsmooth point.



# Projection

To compute  $\text{grad} f_r(R)$ , we need to compute **projection onto  $T_R \mathcal{M}_r$** , i.e.

$$\text{grad} f_r(R) = \text{Proj}_{T_R \mathcal{M}_r}(\nabla f_r(R)) = \nabla f_r(R) - h_R^*(\lambda, \mu),$$

where  $T_R \mathcal{M}_r$  is the tangent space,  $h_R : \mathbb{R}^{n \times r} \rightarrow \mathbb{R}^{m \times r} \times \mathbb{R}^{|B|}$  is the Fréchet differential mapping of the constraints defining  $\mathcal{M}_r$  at  $R$ :

$$h_R(H) := (AH; 2 \text{diag}_B(HR^\top) - H_B e_1),$$

and  $(\lambda, \mu) \in \mathbb{R}^{m \times r} \times \mathbb{R}^{|B|}$  is the solution to the following linear system:

$$h_R(h_R^*(\lambda, \mu)) = h_R(\nabla f_r(R)). \quad (\text{Proj-sub})$$

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By leveraging the special structure of (Proj-sub) with Schur complement and SMW formula, we reduce the complexity for solving (Proj-sub) to

$$\mathcal{O}\left(\min\left\{|B|^3 + m^2 r + mr|B|, (mr)^2|B| + (mr)^3\right\}\right),$$

which is smaller than vanilla  $\mathcal{O}((|B| + mr)^3)$  when either  $|B|$  or  $mr$  is small.

# Retraction

We use the following **projection retraction** to compute the **projection onto  $\mathcal{M}_r$** :

$$\text{Rtr}_{\bar{R}}(H) := \text{Proj}_{\mathcal{M}_r}(V) = \arg \min \left\{ \|R - V\|_F^2 : R \in \mathcal{M}_r \right\},$$

where  $\bar{R} \in \mathcal{M}_r$ ,  $H \in \mathcal{T}_{\bar{R}}\mathcal{M}_r$  and  $V := \bar{R} + H$ . This problem is **nonconvex**, but can be **transformed into a convex problem**.

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$$\mathcal{B}_r := \left\{ R \in \mathbb{R}^{n \times r} : \text{diag}_B(RR^\top) - R_B e_1 = 0 \right\},$$

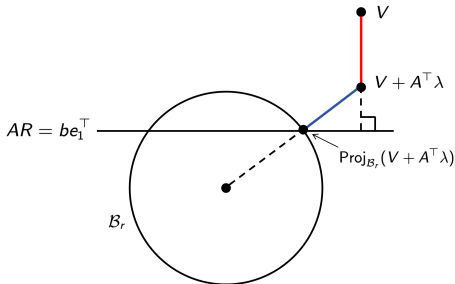
$$\mathcal{M}_r := \left\{ R \in \mathbb{R}^{n \times r} : AR = be_1^\top, \text{diag}_B(RR^\top) - R_B e_1 = 0 \right\}.$$

**Lemma 1.** For any  $V \in \mathbb{R}^{n \times r}$ , if there exists  $\lambda \in \mathbb{R}^{m \times r}$  such that

$$A \text{Proj}_{\mathcal{B}_r}(V + A^\top \lambda) = be_1^\top,$$

then

$$\text{Proj}_{\mathcal{M}_r}(V) = \text{Proj}_{\mathcal{B}_r}(V + A^\top \lambda).$$



# Retraction Subproblem

$\lambda$  is solution to the following unconstrained nonsmooth **convex** problem:

$$\min_{\lambda \in \mathbb{R}^{m \times r}} \sum_{i \in B} \|(V' + A^\top \lambda)_i\| + \sum_{i \in [n] \setminus B} \|(V' + A^\top \lambda)_i\|^2 + \langle b', A^\top \lambda \rangle,$$

where  $V' = V - ee_1^\top/2$ ,  $b' \in \mathbb{R}^{m \times r}$  such that  $(Ae - 2b)e_1^\top = Ab'$ , and  $R_i$  denotes the  $i$ -th row of  $R$ . This is a **generalization of the geometric median problem**, which can be solved by the Weiszfeld algorithm (Weiszfeld 1937).

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**Generalized Weiszfeld algorithm:** assume every row of  $(V' + A^\top \lambda^k)_B$  is non-zero, we design the following algorithm with convergence guarantee:

$$\lambda^{k+1} = T(\lambda^k) := -(A \operatorname{diag}(\mu^k) A^\top)^{-1} A(b' + \operatorname{diag}(\mu^k) V'),$$

where  $\mu^k \in \mathbb{R}^n$  is the vector such that  $\mu_i^k = 1/\|(V' + A^\top \lambda^k)_i\|$  for any  $i \in [B]$ , and  $\mu_i^k = 2$  otherwise.

We also use the **Newton method** to accelerate the convergence.

# Numerical Experiments

Machine: All experiments were conducted using MATLAB R2023b on a workstation with an Intel Xeon E5-2680 v3 @ 2.50GHz CPU and 96GB RAM.

Tolerance for relative KKT residue:  $= 10^{-6}$ ; Maxtime:  $= 3600$ s.

## Solvers comparison:

Solver	LR Factorization	Handle Ineq
ManiSDP (J. Wang and Hu, 2025)	Yes	No
HALLaR (Monteiro, Sujanani, and Cifuentes, 2024)	Yes	No
LoRADS (Han et al., 2024)	Yes	No
SDPF (Tang and Toh, 2024a)	Yes	No
RiNNAL (Hou, Tang, and Toh, 2025)	Yes	No
SketchyCGAL (Yurtsever et al., 2021)	No	Yes
SDPNAL+ (Yang, Sun, and Toh, 2015)	No	Yes
SDPLR (Burer and Monteiro, 2003)	Yes	Yes
SDPDAL (Y. Wang et al., 2023)	Yes	Yes
RiNNAL+ (ours)	Yes	Yes

- Not comparing SDPDAL: code unavailable.
- Not Comparing SDPLR/SketchyCGAL: slower than SDPNAL+.

# Binary Integer Nonconvex Quadratic Programming

$$\min \left\{ x^\top Qx + 2c^\top x : x \leq e, x \in \{0, 1\}^n \right\}$$

Problem	Algorithm	Iteration	Rank	$R_{\max}$	Objective	Time	TPG
$n = 500$	RiNNAL+	16, 850, 3	50	8.70e-07	-1.2259545e+05	8.0	0.9
	SDPNAL+	44, 97, 1535	76	9.98e-07	-1.2259531e+05	613.1	–
$n = 1000$	RiNNAL+	12, 650, 3	79	9.58e-07	-3.8983061e+05	19.3	2.3
	SDPNAL+	–	–	–	–	–	–
$n = 2500$	RiNNAL+	11, 600, 2	134	7.28e-07	-1.6096512e+06	133.3	15.9
	SDPNAL+	–	–	–	–	–	–
$n = 5000$	RiNNAL+	12, 650, 3	177	8.61e-07	-4.6838013e+06	1103.1	254.8
	SDPNAL+	–	–	–	–	–	–

- PG step is efficient due to preprocessing and warm-start.



# Maximum Stable Set Problems

$$\max \left\{ x^\top x : x_i x_j = 0 \ \forall (i, j) \in E, \ x \leq e, \ x \in \{0, 1\}^n \right\}$$

Problem	Algorithm	Iteration	Rank	$R_{\max}$	Objective	Time	TPG
G11 $n = 800$	RiNNAL+	56, 2850, 11	8	6.41e-07	-4.0000014e+02	47.7	6.4
	SDPNAL+	482, 722, 9358	2	2.46e-07	-4.0000005e+02	3008.3	–
G18 $n = 800$	RiNNAL+	34, 1750, 7	75	9.08e-07	-2.7900052e+02	34.0	5.5
	SDPNAL+	–	–	–	–	–	–
G54 $n = 1000$	RiNNAL+	46, 2350, 10	168	9.01e-07	-3.4100141e+02	68.4	11.3
	SDPNAL+	–	–	–	–	–	–
G30 $n = 2000$	RiNNAL+	13, 700, 3	98	9.00e-07	-5.7703872e+02	87.7	10.4
	SDPNAL+	–	–	–	–	–	–
G50 $n = 3000$	RiNNAL+	24, 1250, 5	134	9.39e-07	-1.4940617e+03	754.8	149.3
	SDPNAL+	–	–	–	–	–	–
1tc.1024 $n = 1024$	RiNNAL+	116, 5850, 24	288	9.38e-07	-2.0420520e+02	192.4	28.1
	SDPNAL+	–	–	–	–	–	–
1tc.2048 $n = 2048$	RiNNAL+	94, 4750, 26	520	9.85e-07	-3.7049061e+02	820.7	162.8
	SDPNAL+	–	–	–	–	–	–

- Efficient even on problems with **high-rank** solutions.

# Quadratic Knapsack Problems

$$\max \left\{ x^T Q x : a^T x \leq \tau, x \leq e, x \in \{0, 1\}^n \right\}$$

$n$	Iteration	Rank	$R_{\max}$	Objective	Time	TPG
500	87, 4400, 18	26	9.91e-07	-1.1420448e+06	40.5	4.3
1000	27, 1400, 6	37	9.29e-07	-4.5986965e+06	53.1	5.5
2000	15, 800, 3	42	6.84e-07	-1.8316334e+07	123.7	22.7
5000	18, 950, 4	77	7.89e-07	-1.1399134e+08	1761.8	517.4

- SDPNAL+ fails to solve all instances.
- Projection and retraction nontrivial, but taking only 10–20% of total time.

# Cardinality-Constrained Minimum SOS Clustering

$$\min \left\{ \sum_{j=1}^k \frac{1}{c_j} \sum_{s=1}^m \sum_{t=1}^m d_{st} \pi_j^{(s)} \pi_j^{(t)} : \pi \in \mathcal{S}, \pi \leq \mathbf{e}_{m \times k}, \pi \in \{0, 1\}^{m \times k} \right\}$$

where  $\pi = [\pi^{(1)}, \dots, \pi^{(k)}]$  with  $\pi^{(j)} \in \{0, 1\}^m$ , and

$$\mathcal{S} := \left\{ \pi \in \mathbb{R}^{m \times k} : \sum_{j=1}^k \pi_i^{(j)} = 1 \ \forall i \in [m], \quad \sum_{i=1}^m \pi_i^{(j)} = c_j \ \forall j \in [k] \right\}$$

Problem	Algorithm	Iteration	Rank	$R_{\max}$	Objective	Time	TPG
1 $n = 720$	RiNNAL+	9, 500, 2	5	9.27e-07	7.3254609e+04	14.3	1.4
	SDPNAL+	-	-	-	-	-	-
2 $n = 902$	RiNNAL+	3, 200, 1	7	8.78e-07	4.0536895e+09	12.6	1.1
	SDPNAL+	54, 246, 2312	1	7.91e-07	4.0536886e+09	2313.4	-
3 $n = 1662$	RiNNAL+	12, 511, 3	21	3.90e-07	8.5151301e+02	1500.3	193.9
	SDPNAL+	-	-	-	-	-	-
4 $n = 1800$	RiNNAL+	2, 150, 1	28	5.87e-07	7.3610392e+08	42.2	13.7
	SDPNAL+	-	-	-	-	-	-
5 $n = 2250$	RiNNAL+	51, 2600, 11	61	5.42e-07	1.0500893e+06	2215.9	170.8
	SDPNAL+	-	-	-	-	-	-

- Efficient for **many constraints** (up to 3.7M eq/ineq and 5M nonneg).

# Sparse StQP

$$\min \left\{ x^\top Qx : e^\top x = 1, \|x\|_0 \leq \rho, x \in \mathbb{R}_+^m \right\},$$

Big-M reformulation:

$$\min \left\{ x^\top Qx : e^\top x = 1, e^\top u = \rho, x \leq u \leq e, x \in \mathbb{R}_+^m, u \in \{0, 1\}^m \right\}.$$

Problem	Algorithm	Iteration	Rank	R <sub>max</sub>	Objective	Time	TPG
COP $n = 200$	RiNNAL+	125, 6300, 25	102	6.64e-07	-1.0190589e-01	30.7	2.1
	SDPNAL+	76, 112, 1170	102	8.59e-07	-1.0196083e-01	40.1	-
PSD $n = 200$	RiNNAL+	374, 18750, 75	5	8.59e-07	1.4174794e-02	80.7	14.5
	SDPNAL+	801, 1058, 17306	2	8.50e-07	1.4174999e-02	468.1	-
SPN $n = 200$	RiNNAL+	747, 37400, 150	107	9.31e-07	1.3592384e-02	204.0	20.8
	SDPNAL+	801, 1032, 29187	12	9.94e-07	1.4887842e-02	866.1	-
COP $n = 400$	RiNNAL+	300, 15050, 60	202	9.99e-07	-1.0402158e-01	230.5	15.8
	SDPNAL+	103, 145, 1511	202	9.24e-07	-1.0445084e-01	170.7	-
PSD $n = 400$	RiNNAL+	360, 18050, 72	108	9.50e-07	9.0349366e-03	231.0	19.4
	SDPNAL+	-	-	-	-	-	-
SPN $n = 400$	RiNNAL+	1812, 90650, 363	248	9.97e-07	-1.2530973e-02	1556.1	145.7
	SDPNAL+	-	-	-	-	-	-
PSD $n = 1000$	RiNNAL+	507, 25400, 102	234	9.88e-07	2.9188026e-03	1838.0	152.6
	SDPNAL+	-	-	-	-	-	-

# Quadratic Minimum Spanning Tree Problem

$$\min \left\{ x^\top Qx : x \in \mathcal{T}' \right\}.$$

where

$$\mathcal{T}' := \left\{ x \in \{0, 1\}^m : x \leq e, \sum_{e \in E} x_e = n - 1, \sum_{e \in \partial S} x_e \geq 1, \forall S \subsetneq V, |S| = 1 \right\},$$

Problem	Algorithm	Iteration	Rank	$R_{\max}$	Objective	Time	TPG
vsym $m = 435$	RiNNAL+	29, 1480, 6	5	4.86e-07	7.7804931e+04	12.9	1.7
	SDPNAL+	148, 230, 3713	1	2.91e-07	7.7804694e+04	299.7	-
sym $m = 435$	RiNNAL+	103, 5190, 21	197	7.70e-07	5.3262719e+03	99.3	2.8
	SDPNAL+	96, 109, 1486	195	1.09e-06	5.3262716e+03	130.5	-
esym $m = 435$	RiNNAL+	4912, 254289, 4912	49	9.85e-07	7.4273402e+03	3456.3	931.4
	SDPNAL+	529, 692, 10111	46	9.55e-07	7.4273883e+03	1083.3	-
vsym $m = 1225$	RiNNAL+	96, 4850, 22	8	7.99e-07	1.6444786e+05	289.2	64.3
	SDPNAL+	-	-	-	-	-	-
sym $m = 1225$	RiNNAL+	84, 4275, 17	509	8.07e-07	1.4104759e+04	484.0	21.2
	SDPNAL+	192, 199, 2628	507	1.22e-06	1.4104759e+04	1987.4	-

# Conclusion

- We justify the use of **SDP-RLT** over DNN: same tightness, smaller dimension.
- We propose **RiNNAL+**: a two-phase Riemannian ALM for solving SDP-RLT.
- Efficient for problems with **many constraints** and even **high-rank** solutions.
- The framework is broadly applicable to other structured SDPs.

# Conclusion






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- We propose **RiNNAL+**: a two-phase Riemannian ALM for solving SDP-RLT.
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- The framework is broadly applicable to other structured SDPs.

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**Code:** <https://github.com/HouDiOpt/RiNNALplus>






Thank you

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