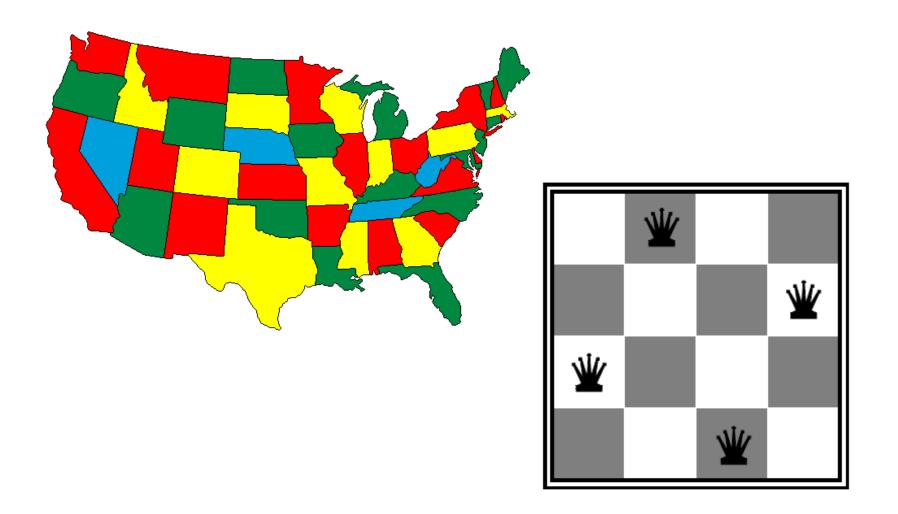
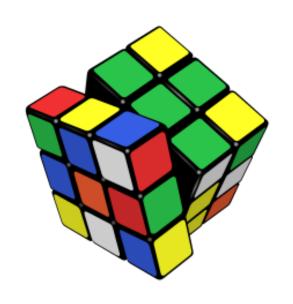
Constraint Satisfaction Problems



Sanja Lazarova-Molnar

CSP Definitions

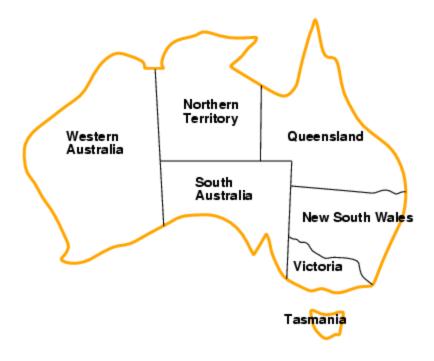
- Variables X₁, X₂,..., X_n each X_i
 having a non-empty domain D_i of
 possible values.
- Constraints C₁, C₂,..., C_m consisting of some subset of variables and specifies allowable combinations of values for that subset.
- State defined by an assignment of values to some or all variables
 (X₁=v₁, X_j=v_j, ...)
- Consistent assignment that does not violate any constraints.
- Solution complete assignment that satisfies all constraints



CSP Formulation

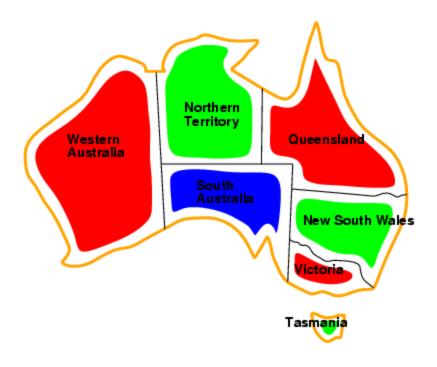
- Initial state empty assignment: all variables are unassigned.
- Successor function assigns value to an unassigned variable that does not conflict with previously assigned variables
- Goal test complete current assignment
- Path cost constant cost per step

Example: Map Coloring



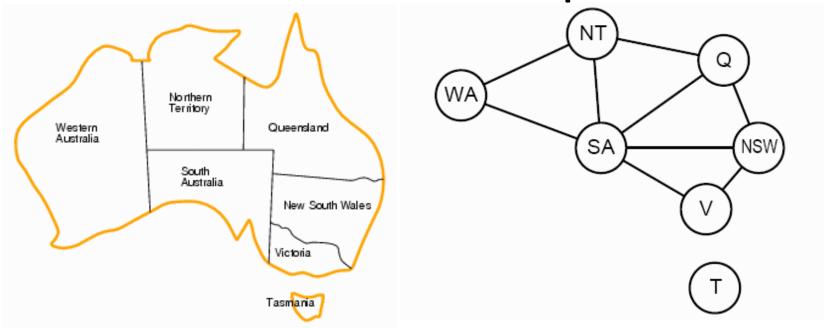
- Variables: WA, NT, Q, NSW, V, SA, T
- Domains: {red, green, blue}
- Constraints: adjacent regions must have different colors e.g., WA ≠ NT, or (WA, NT) in {(red, green), (red, blue), (green, red), (green, blue), (blue, red), (blue, green)}

Example: Map Coloring



- State one of many but not a solution, e.g. WA = red, NT = red, Q = red, NSW = red, V = red, SA = red, T = red
- Solutions are complete and consistent assignments, e.g., WA = red,
 NT = green, Q = red, NSW = green, V = red, SA = blue, T = green

Constraint Graph

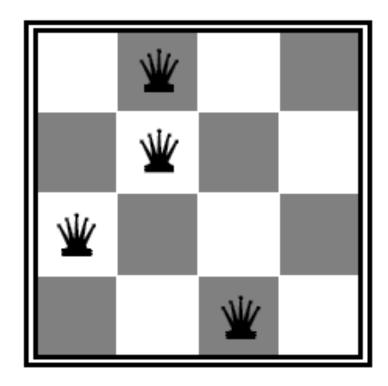


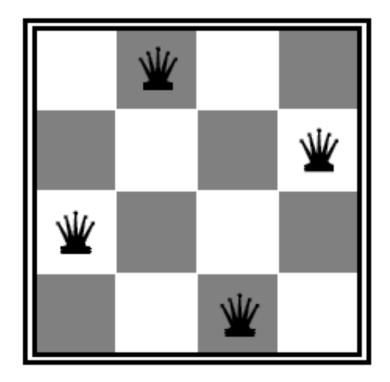
Nodes are variables, arcs show constraints.

- The constraint a != b in map coloring means that no adjacent Australian states are the same color.
- What is meaning of arcs from SA under the a != b constraint?
- What about the fact that there are no arcs to T.
- What is the implication of the arcs between WA, NT and SA?

Example: n-queens problem

• Put n queens on an $n \times n$ board with no two queens on the same row, column, or diagonal





Example: N-Queens

- Variables: X_{ij}
- **Domains:** {0, 1}
- Constraints:

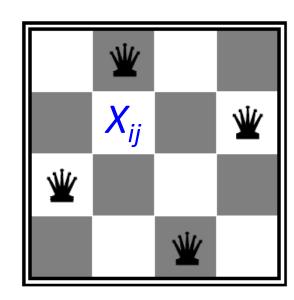
$$\Sigma_{i,j} X_{ij} = N$$

$$(X_{ij}, X_{ik}) \in \{(0, 0), (0, 1), (1, 0)\}$$

$$(X_{ij}, X_{kj}) \in \{(0, 0), (0, 1), (1, 0)\}$$

$$(X_{ij}, X_{i+k, j+k}) \in \{(0, 0), (0, 1), (1, 0)\}$$

$$(X_{ij}, X_{i+k, j-k}) \in \{(0, 0), (0, 1), (1, 0)\}$$



Example: Cryptarithmetic

• Variables: T, W, O, F, U, R

$$X_1, X_2$$

- **Domains**: {0, 1, 2, ..., 9}
- Constraints:

$$O + O = R + 10 * X_1$$
 $W + W + X_1 = U + 10 * X_2$
 $T + T + X_2 = O + 10 * F$
Alldiff(T, W, O, F, U, R)
 $T \neq 0, F \neq 0$

Example: Sudoku

- Variables: X_{ij}
- **Domains:** {1, 2, ..., 9}
- Constraints:

Alldiff(X_{ii} in the same *unit*)

					0			4
_			_		8	_		4
	8	4		1	6		,	
			5			1	0,	
1		3	8			9		
6		8		X _{ij}		4		3
		2			9	5		1
		7			2			
			7	8		2	6	
2			3					

Real-world CSPs

- Assignment problems
 - e.g., who teaches what class
- Timetable problems
 - e.g., which class is offered when and where?
- Transportation scheduling
- Factory scheduling

More examples of CSPs: http://www.csplib.org/

Standard search formulation (incremental)

States:

Variables and values assigned so far

Initial state:

The empty assignment

Action:

- Choose any unassigned variable and assign to it a value that does not violate any constraints
 - Fail if no legal assignments

Goal test:

The current assignment is complete and satisfies all constraints

Standard search formulation (incremental)

- What is the depth of any solution (assuming n variables)?
 n (this is good)
- Given that there are m possible values for any variable, how many paths are there in the search tree?

```
n! \cdot m^n (this is bad)
```

Backtracking search

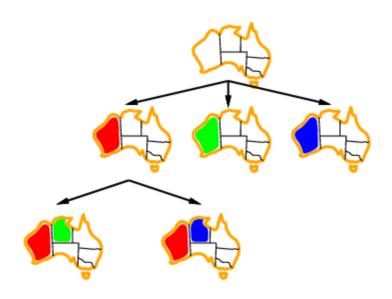
- In CSP's, variable assignments are commutative
 - For example, $[WA = red \ then \ NT = green]$ is the same as $[NT = green \ then \ WA = red]$
- We only need to consider assignments to a single variable at each level (i.e., we fix the order of assignments)
 - Then there are only m^n leaves (n number of variables and m number of values)
- Depth-first search for CSPs with single-variable assignments is called backtracking search



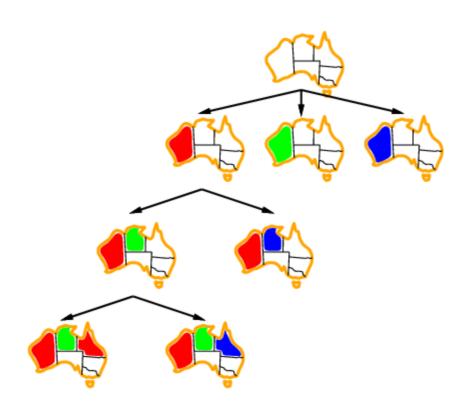








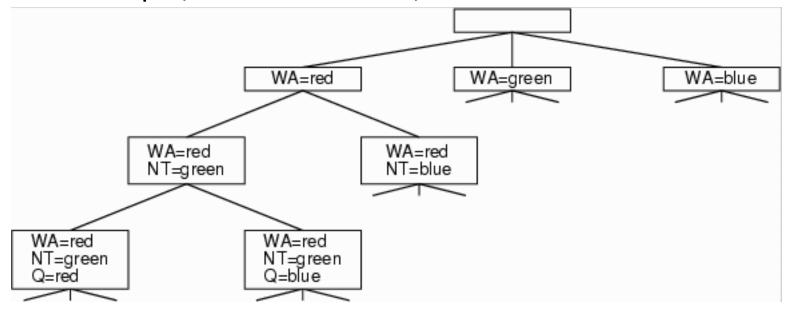


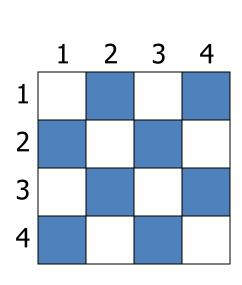


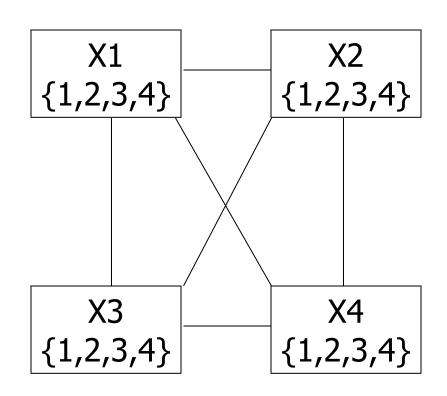


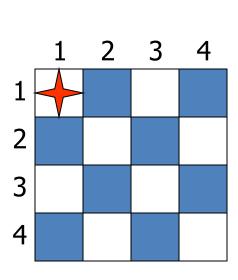
Backtracking

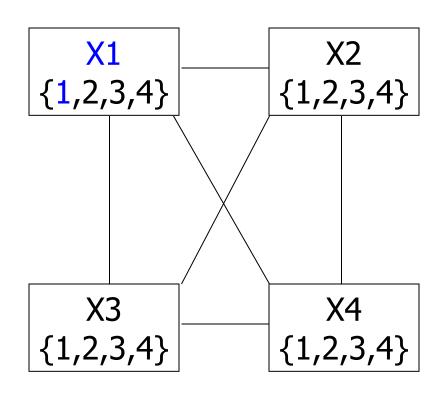
- Constraints on SA will eventually cause failure when WA !=
 Q. When not the same color (bottom right), SA cannot be
 assigned.
- The algorithm will backtrack to a node with unexplored states.
- For example, such as WA=red, NT=blue.

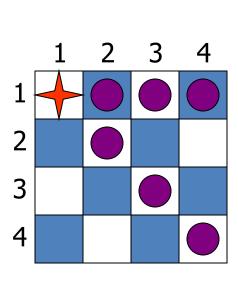


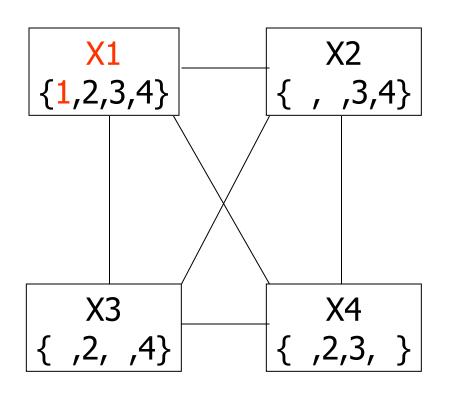


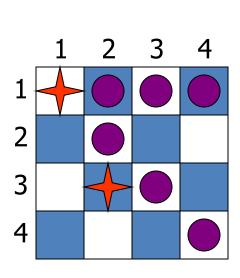


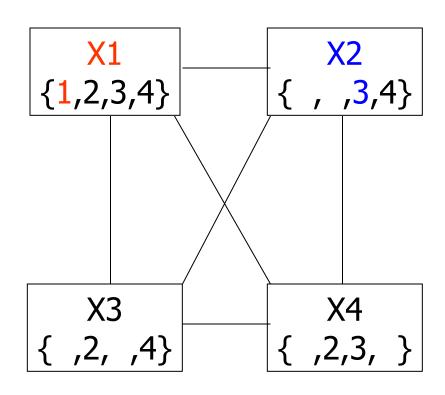


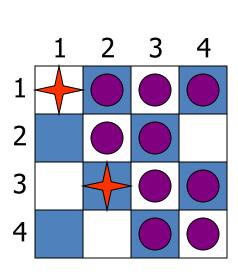


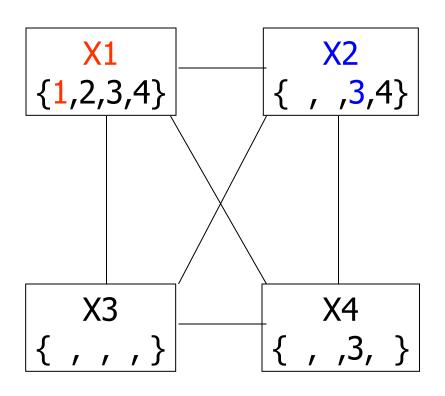


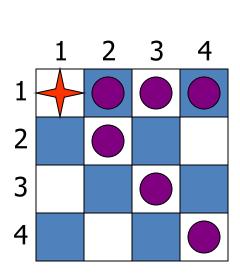


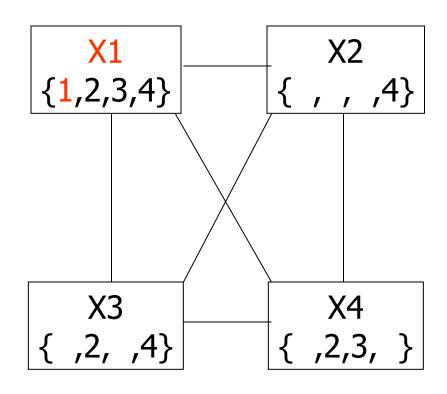


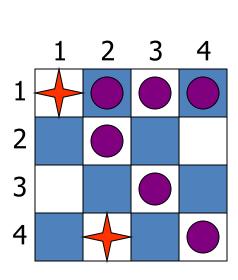


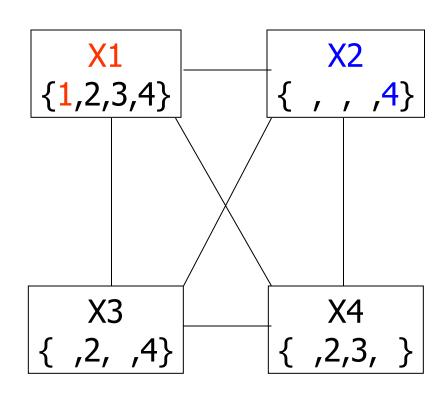


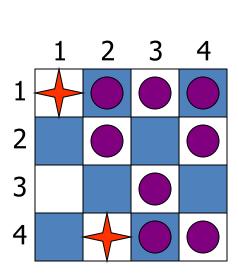


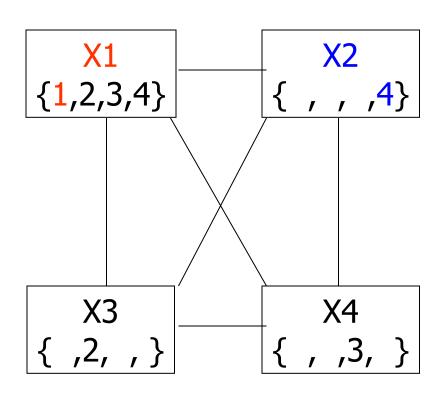


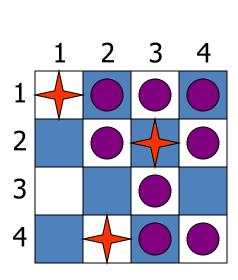


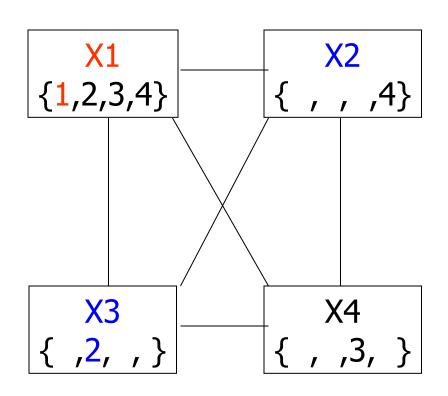


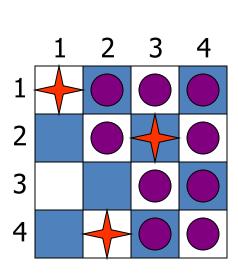


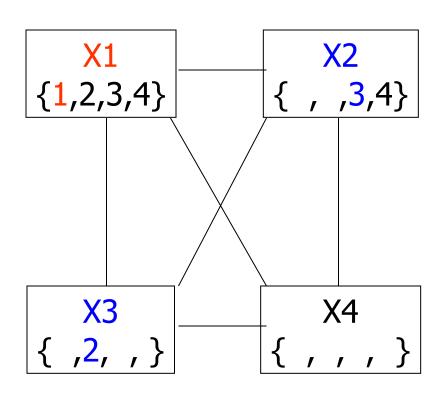












Most constrained variable:

- Choose the variable with the fewest legal values
- A.k.a. minimum remaining values (MRV) heuristic

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- Choose the variable that imposes the most constraints on the remaining variables
- Tie-breaker among most constrained variables

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N.B. Among the variables with the smallest remaining domains, select the one that appears in the largest number of constraints on variables not in the current assignment



Given a variable, in which order should its values be tried?

- Choose the least constraining value:
 - The value that rules out the fewest values in the remaining variables

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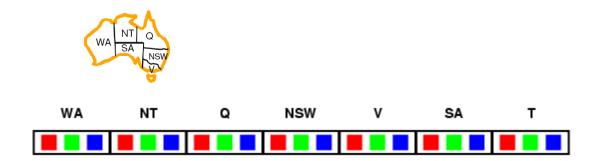
Which assignment for Q should we choose?



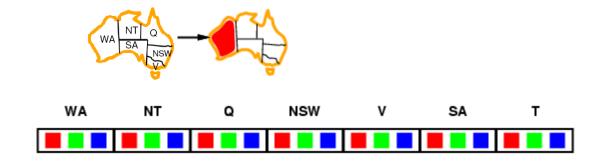
Early detection of failure: Forward checking

- Keep track of remaining legal values for unassigned variables
- Terminate search when any variable has no legal values

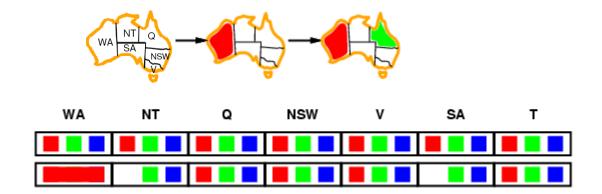
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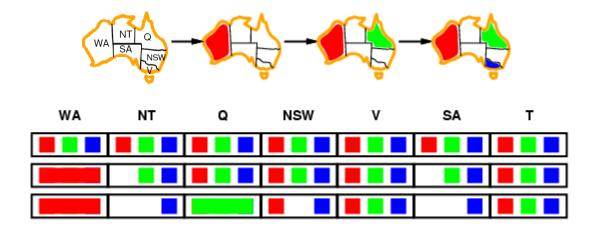
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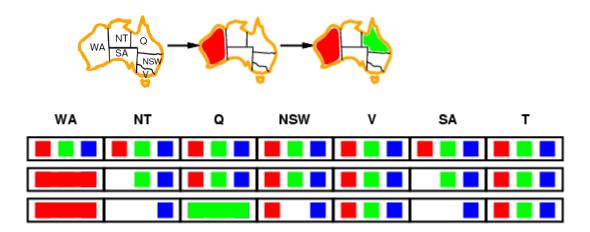


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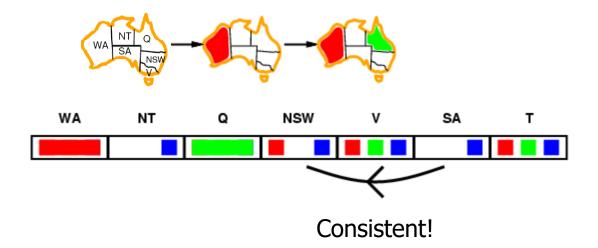
Constraint propagation

 Forward checking propagates information from assigned to unassigned variables, but doesn't provide early detection for all failures

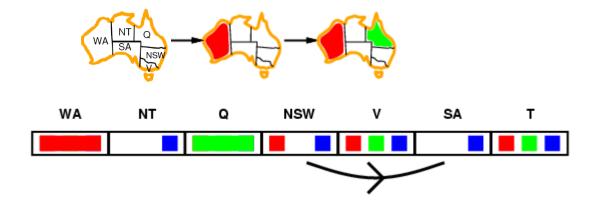


- NT and SA cannot both be blue!
- Constraint propagation repeatedly enforces constraints locally

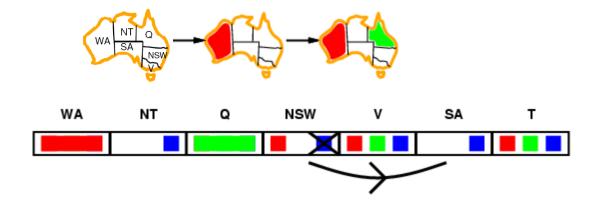
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 - $-X \rightarrow Y$ is consistent iff for every value of X there is some allowed value of Y



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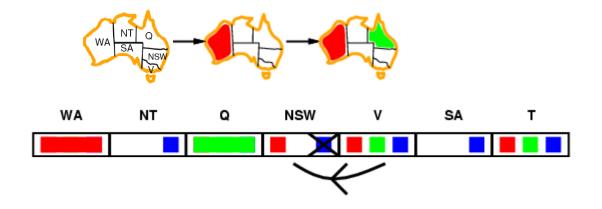


- Simplest form of propagation makes each pair of variables consistent:
 - $-X \rightarrow Y$ is consistent iff for every value of X there is some allowed value of Y
 - When checking $X \rightarrow Y$, throw out any values of X for which there isn't an allowed value of Y



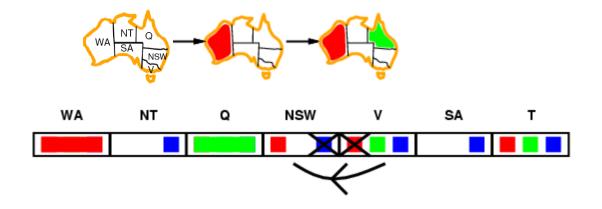
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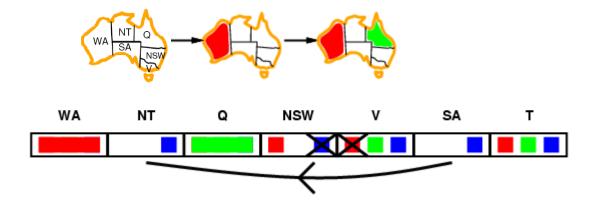
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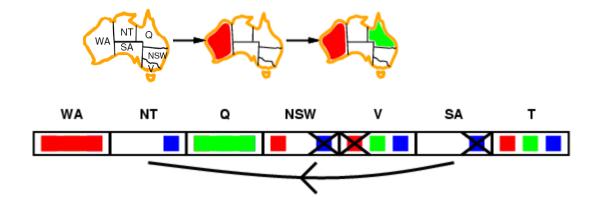


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- Arc consistency detects failure earlier than forward checking
- Can be run before or after each assignment

Summary

- CSPs are a special kind of search problem:
 - States defined by values of a fixed set of variables
 - Goal test defined by constraints on variable values
- Backtracking = depth-first search where successor states are generated by considering assignments to a single variable
 - Variable ordering and value selection heuristics can help significantly
 - Forward checking prevents assignments that guarantee later failure
 - Constraint propagation (e.g., arc consistency) does additional work to constrain values and detect inconsistencies
- Complexity of CSPs
 - NP-complete in general (exponential worst-case running time)