# 3007 Final Exam Review

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## 1 Definitions

Define the following terms and provide examples or sample code as appropriate.

## 1.1 Imperative vs Declarative

## 1.1.1 Imperative

- Series of instructions
- Iterative functions
- Command driven, statement oriented
- Procedural
  - C
  - Pascal
  - Assembly
- · Object oriented
  - C++
  - Java

#### 1.1.2 Declarative

- No side effects
- Focus on relations
- "What to get" instead of "How to get"
- $\bullet$  Order of statements *shouldn't* matter
- Examples:
  - SQL
  - Prolog
  - Regex

## 1.2 Scope vs Visibility

## 1.2.1 Scope

- The set of expressions for which the variable exists
- In lexical scoping
  - variables in the scope we were defined in
  - and local variables
  - who uses this?
    - \* C-family languages
    - \* Scheme
    - \* Algol
- In dynamic scoping
  - variables in the scope we were *called* in
  - and local variables
  - who uses this?
    - \* early LISP
    - \* APL
    - \* BASH

#### 1.2.2 Visibility

- The set of expressions for which the variable can be reached
- If we declare a local variable with the same name as a variable in enclosing scope
  - that enclosing scope variable is now hidden
  - all references to name are to our locally scoped variable instead

## 1.3 Lexical Scope vs Dynamic Scope

#### 1.3.1 Lexical

- Function scope is enclosed in the scope which defined us
  - if you can't find a binding, recursively search in the function that defined you

## 1.3.2 Dynamic

- Function scope is enclosed in the scope which called us
  - if you can't find a binding, recursively search in the function that called you

#### 1.4 Free Variables

- Used locally but bound in an enclosing scope
- In the following example:

```
(define (f x y)
  (define z 2)
  (define (g)
        (* x y z)
  )
)
```

- x,y,z are free variables in (g)

## 1.5 Applicative Order Evaluation vs Normal Order Evaluation

## 1.5.1 Applicative Order Evaluation

- Strict evaluation
- Evaluate an expression before it is passed in as an argument
  - go as deep as you can until you hit primitives, then evaluate and go back
  - as deep into the nest as possible and work backwards

```
- e.g.,
  (double (* (+ 1 3) 4))
  (double (* 4 4))
  (double 16)
  (* 16 2)
  32
```

#### 1.5.2 Normal Order Evaluation

- Lazy evaluation
- Evaluate an expression only when its value is needed
  - first expand, then reduce
     e.g.,
     (double (\* (+ 1 3) 4))
     (\* (\* (+ 1 3) 4) 2)
     (\* (\* 4 4) 2)
     (\* 16 2)
     32

## 1.6 Special Forms

- Exceptions to the usual evaluation order
  - they have their own evaluation rules
  - e.g., take the first argument without evaluating right away, evaluate the second symbol right away
- Use constructs like (delay foo), (force foo) behind the scenes

## 1.7 Tail Recursion

- Linear iterative processes in Scheme
- No deferred operations
  - recusrive call is the last operation of the procedure
- In Scheme, recursion is tail optimized
  - this means that it will run in constant space
  - number of steps will **grow linearly**, but memory will **remain constant**
- Even though the *program* is still recursive, the *process* is linear iterative because of tail-recursion optimization
- E.g., to compute a factorial using tail recursion, we do the following:

```
(define (factorial x)
  (define (iter prod i)
    (if (> i x)
        prod
        (iter (* i prod) (+ i 1))
    )
    (iter 1 1)
)
```

• To compute a factorial using normal recursion, we would do the following instead:

```
(define (factorial x)
  (if (= x 1)
      x
      (* x (factorial - x 1))
  )
)
```