

Least squares for programmers

— with color plates —

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3 Minimization of quadratic functions

4 Least squares through examples

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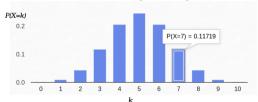
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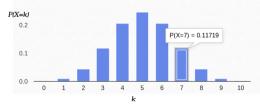
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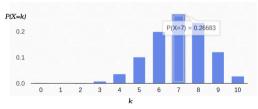
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a biased coin (p = 7/10)

Coin toss: the likelihood function

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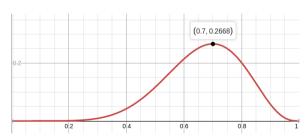
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N.B. the function is continuous!

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That is, the maximum likelihood (about 27%) is reached at the point p = 7/10. Just in case, let us check the second derivative:

$$\frac{d^2 \log \mathcal{L}}{dp^2} = -\frac{7}{p^2} - \frac{3}{(1-p)^2}$$

At the point p = 7/10 it is negative, therefore this point is indeed a maximum of the function \mathcal{L} :

$$\frac{d^2\log\mathcal{L}}{dp^2}(0.7)\approx -48<0$$

Let us measure a constant value; all measurements are inherently noisy.

For example, if we measure the battery voltage *N* times, we get *N* different measurements:

$$\{U_j\}_{j=1}^N$$

Suppose that each measurement U_j is i.i.d. and subject to a Gaussian noise, e.g. it is equal to the real value plus the Gaussian noise. The probability density can be expressed as follows:

$$p(U_j) = rac{1}{\sqrt{2\pi}\sigma} \exp\left(-rac{(U_j - U)^2}{2\sigma^2}
ight),$$

where U is the (unknown) value and σ is the noise amplitude (can be unknown).

$$\log \mathcal{L}(U,\sigma) = \log \left(\prod_{j=1}^{N} \frac{1}{\sqrt{2\pi}\sigma} \exp \left(-\frac{(U_j - U)^2}{2\sigma^2} \right) \right)$$

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Under Gaussian noise

$$\underset{U}{\text{arg max}} \log \mathcal{L} = \underset{U}{\text{arg min}} \sum_{j=1}^{N} (U_j - U)^2$$

$$\frac{\partial \log \mathcal{L}}{\partial U} = -\frac{1}{\sigma^2} \sum_{i=1}^{N} (U_i - U) = 0$$

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The most plausible estimation of the unknown value U is the simple average of all measurements:

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$$\sigma = \sqrt{\frac{\sum_{j=1}^{N} (U_j - U)^2}{N}}$$

Such a convoluted way to obtain a simple average of all measurements...

It is much harder for less trivial examples. Suppose we have N measurements $\{x_j, y_j\}_{j=1}^N$, and we want to fit a straight line onto it.

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$$\log \mathcal{L}(a, b, \sigma) = \log \left(\prod_{j=1}^{N} \frac{1}{\sqrt{2\pi}\sigma} \exp \left(-\frac{(y_j - ax_j - b)^2}{2\sigma^2} \right) \right) =$$

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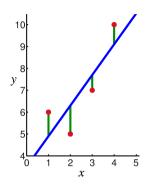
As before, $\underset{a,b}{\operatorname{arg\,max}} \log \mathcal{L} = \underset{a,b}{\operatorname{arg\,min}} S(a,b).$

Linear regression
$$S(a,b) := \sum_{j=1}^{N} (y_j - ax_j - b)^2$$

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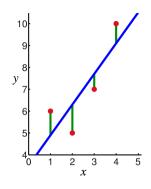


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$$a = \frac{N \sum_{j=1}^{N} x_j y_j - \sum_{j=1}^{N} x_j \sum_{j=1}^{N} y_j}{N \sum_{j=1}^{N} x_j^2 - \left(\sum_{j=1}^{N} x_j\right)^2}$$



$$b = \frac{1}{N} \left(\sum_{j=1}^{N} y_j - a \sum_{j=1}^{N} x_j \right)$$

The takeaway message

The least squares method is a particular case of maximizing likelihood in cases where the probability density is Gaussian.

The more we parameters we have, the more cumbersome the analytical solutions are. Fortunately, we are not living in XVIII century anymore, we have computers!

Next we will try to build a geometric intuition on least squares, and see how can least squares problems be efficiently implemented.

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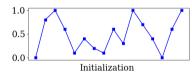
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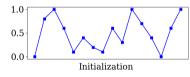


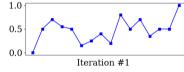
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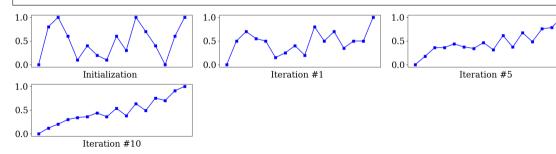


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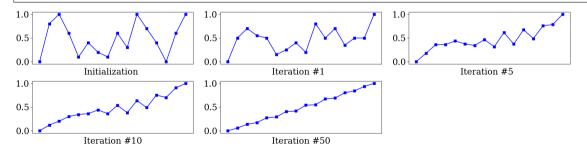
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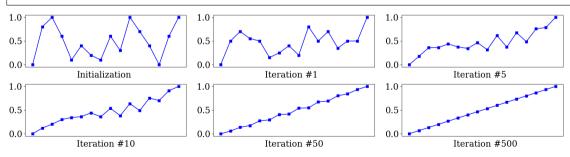
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Given an ordinary system of linear equations:

$$\begin{cases} a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n = b_1 \\ a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n = b_2 \\ \vdots \\ a_{n1}x_1 + a_{n2}x_2 + \cdots + a_{nn}x_n = b_n \end{cases}$$

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Let us rewrite it as follows:

$$x_{1} = \frac{1}{a_{11}}(b_{1} - a_{12}x_{2} - a_{13}x_{3} - \dots - a_{1n}x_{n})$$

$$x_{2} = \frac{1}{a_{22}}(b_{2} - a_{21}x_{1} - a_{23}x_{3} - \dots - a_{2n}x_{n})$$

$$\vdots$$

$$x_{n} = \frac{1}{a_{nn}}(b_{n} - a_{n1}x_{1} - a_{n2}x_{2} - \dots - a_{n,n-1}x_{n-1})$$

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Repeating the process k times, the solution can be approximated by the vector $\vec{x}^{(k)} = (x_1^{(k)}, x_2^{(k)}, \dots, x_n^{(k)}).$

Back to the array smoothing

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The Gauß-Seidel iterative method

Jacobi:

$$x_i^{(k)} = \frac{1}{a_{ii}} \left(b_i - \sum_{j=1, j \neq i}^n a_{ij} x_j^{(k-1)} \right), \quad \text{for } i = 1, 2, \dots, n$$

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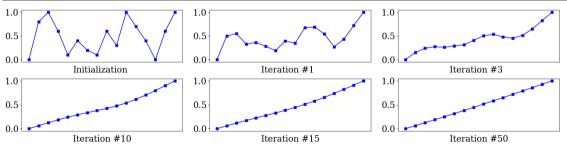
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Smooth an array: Gauß-Seidel

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Equality of derivatives vs zero curvature

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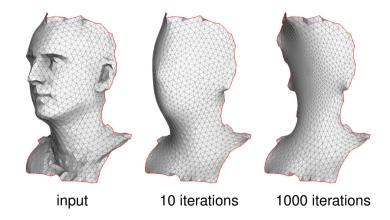
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\vdots
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```

It also works for 3d surfaces

```
#include "model.h"
int main(void) {
   Model m("../input.obj"); // parse the input mesh
    // smooth the surface through Gauss-Seidel iterations
    for (int it=0; it<1000; it++)
        for (int v=0; v<m.nverts(); v++) // for all vertices</pre>
            if (!m.is boundary vert(v)) // fix the boundary
                m.point(v) = m.one ring barycenter(v);
    std::cout << m; // drop the result
    return 0;
```

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Prescribe the right hand side

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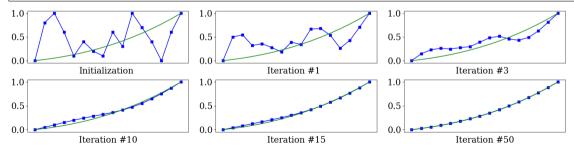


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Is it a function $f(x) = ax : \mathbb{R} \to \mathbb{R}$?

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What is a number a?

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	return a*x;
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float f(float x) {
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The same goes for matrices, what is a matrix
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Is it f(x) = Ax : \mathbb{R}^2 \to \mathbb{R}^2?

vector<float> f(vector<float> x) {

return vector<float>{al1*x[0] + al2*x[1], a21*x[0] + a22*x[1]};
}
```

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Is it f(x) = Ax : \mathbb{R}^2 \to \mathbb{R}^2?
vector<float> f(vector<float> x) {
   return vector < float> {a11*x[0] + a12*x[1], a21*x[0] + a22*x[1]};
Or f(x) = x^{\top} A x = \sum_{i} \sum_{i} a_{ij} x_i x_j : \mathbb{R}^2 \to \mathbb{R}?
float f(vector<float> x) {
      return x[0]*a11*x[0] + x[0]*a12*x[1] +
                 x[1]*a21*x[0] + x[1]*a22*x[1];
```

What is a positive number?

We have a great tool called the predicate "greater than" >.

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Definition

The real number a is positive if and only if for all non-zero real $x \in \mathbb{R}, \ x \neq 0$ the condition $ax^2 > 0$ is satisfied.

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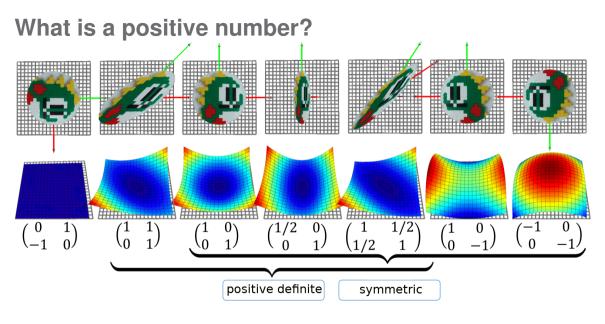
Definition

The real number a is positive if and only if for all non-zero real $x \in \mathbb{R}, x \neq 0$ the condition $ax^2 > 0$ is satisfied.

This definition looks pretty awkward, but it applies perfectly to matrices:

Definition

The square matrix A is called positive definite if for any non-zero x the condition $x^{T}Ax > 0$ is met, i.e. the corresponding quadratic form is strictly positive everywhere except at the origin.



Minimizing a 1d quadratic function

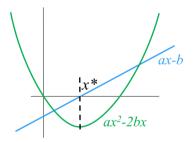
Let us find the minimum of the function $f(x) = ax^2 - 2bx$ (with a positive).

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In 1d, the solution x^* of the equation ax - b = 0 solves the minimization problem $arg min(ax^2 - 2bx)$ as well.

Differentiating matrix expressions

The first theorem states that 1×1 matrices are invariant w.r.t the transposition:

Theorem

$$x \in \mathbb{R} \Rightarrow x^{\top} = x$$

The proof is left as an exercise.

Differentiating matrix expressions

For a 1d function bx we know that $\frac{d}{dx}(bx) = b$, but what happens in the case of a real function of n variables?

Theorem

$$\nabla b^{\mathsf{T}} x = \nabla x^{\mathsf{T}} b = b$$

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Theorem

$$\nabla b^{\mathsf{T}} x = \nabla x^{\mathsf{T}} b = b$$

$$\nabla(b^{\top}x) = \begin{bmatrix} \frac{\partial(b^{\top}x)}{\partial x_1} \\ \vdots \\ \frac{\partial(b^{\top}x)}{\partial x_n} \end{bmatrix} = \begin{bmatrix} \frac{\partial(b_1x_1 + \dots + b_nx_n)}{\partial x_1} \\ \vdots \\ \frac{\partial(b_1x_1 + \dots + b_nx_n)}{\partial x_n} \end{bmatrix} = \begin{bmatrix} b_1 \\ \vdots \\ b_n \end{bmatrix} = b$$

For a 1d function ax^2 we know that $\frac{d}{dx}(ax^2) = 2ax$, but what about quadratic forms?

Theorem

$$\nabla(\mathbf{X}^{\top}\mathbf{A}\mathbf{X}) = (\mathbf{A} + \mathbf{A}^{\top})\mathbf{X}$$

Note that if *A* is symmetric, then $\nabla(x^{\top}Ax) = 2Ax$.

The proof is straightforward, let us express the quadratic form as a double sum:

$$x^{\top}Ax = \sum_{i} \sum_{j} a_{ij} x_{i} x_{j}$$

$$\frac{\partial (\mathbf{x}^{\top} \mathbf{A} \mathbf{x})}{\partial \mathbf{x}_{i}} = \frac{\partial}{\partial \mathbf{x}_{i}} \left(\sum_{k_{1}} \sum_{k_{2}} a_{k_{1} k_{2}} \mathbf{x}_{k_{1}} \mathbf{x}_{k_{2}} \right) =$$

$$\frac{\partial(\mathbf{x}^{\top} A \mathbf{x})}{\partial \mathbf{x}_{i}} = \frac{\partial}{\partial \mathbf{x}_{i}} \left(\sum_{\mathbf{k}_{1}} \sum_{\mathbf{k}_{2}} \mathbf{a}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{2}} \mathbf{x}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{2}} \right) =$$

$$= \frac{\partial}{\partial \mathbf{x}_{i}} \left(\underbrace{\sum_{\mathbf{k}_{1} \neq i} \sum_{\mathbf{k}_{2} \neq i} \mathbf{a}_{i\mathbf{k}_{2}} \mathbf{x}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{2}}}_{\mathbf{k}_{1} \neq i, \mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{2} \neq i} \mathbf{a}_{i\mathbf{k}_{2}} \mathbf{x}_{i} \mathbf{x}_{\mathbf{k}_{2}}}_{\mathbf{k}_{1} \neq i, \mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{1}} \mathbf{x}_{i}}_{\mathbf{k}_{1} \neq i, \mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{1}} \mathbf{x}_{i}}_{\mathbf{k}_{1} \neq i, \mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{1}} \mathbf{x}_{i}}_{\mathbf{k}_{1} \neq i, \mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{1}} \mathbf{x}_{i}}_{\mathbf{k}_{1} \neq i, \mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{1}} \mathbf{x}_{i}}_{\mathbf{k}_{1} \neq i, \mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{2}} \mathbf{x}_{i}}_{\mathbf{k}_{1} \neq i, \mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{2}} \mathbf{x}_{i}}_{\mathbf{k}_{1} \neq i, \mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{2}} \mathbf{x}_{i}}_{\mathbf{k}_{1} \neq i, \mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{2}} \mathbf{x}_{i}}_{\mathbf{k}_{1} \neq i, \mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{2}} \mathbf{x}_{i}}_{\mathbf{k}_{1} \neq i, \mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{1}} \mathbf{x}_{\mathbf{k}_{2}} \mathbf{x}_{i}}_{\mathbf{k}_{1} \neq i, \mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{1}} \mathbf{x}_{i}}_{\mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{2}} \mathbf{x}_{i}}_{\mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{2}} \mathbf{x}_{i}}_{\mathbf{k}_{2}} \mathbf{x}_{i}}_{\mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{2}} \mathbf{x}_{i}}_{\mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{2}} \mathbf{x}_{i}}_{\mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{1} \neq i} \mathbf{a}_{\mathbf{k}_{2}} \mathbf{x}_{i}}_{\mathbf{k}_{2} \neq i} + \underbrace{\sum_{\mathbf{k}_{2} \neq i} \mathbf{a}_{\mathbf{k}_{2}} \mathbf{x}_{i}}_{\mathbf{k}_{2} \neq i} + \underbrace{$$

$$\frac{\partial(x^{\top}Ax)}{\partial x_{i}} = \frac{\partial}{\partial x_{i}} \left(\sum_{k_{1}} \sum_{k_{2}} a_{k_{1}k_{2}} x_{k_{1}} x_{k_{2}} \right) =
= \frac{\partial}{\partial x_{i}} \left(\sum_{k_{1} \neq i} \sum_{k_{2} \neq i} a_{ik_{2}} x_{k_{1}} x_{k_{2}} + \sum_{k_{2} \neq i} a_{ik_{2}} x_{i} x_{k_{2}} + \sum_{k_{1} \neq i} a_{k_{1}i} x_{k_{1}} x_{i} + \underbrace{a_{ii} x_{i}^{2}}_{k_{1} \neq i, k_{2} \neq i} \right) =
= \sum_{k_{2} \neq i} a_{ik_{2}} x_{k_{2}} + \sum_{k_{1} \neq i} a_{k_{1}i} x_{k_{1}} + 2a_{ii} x_{i} =$$

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Recall that for a > 0 solving the equation ax = b is equivalent to the quadratic function $arg min(ax^2 - 2bx)$ minimization.

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To minimize a quadratic form $\underset{x \in \mathbb{R}^n}{\arg\min(x^{\top}Ax - 2b^{\top}x)}$ with a symmetric positive definite matrix A, equate the derivative to zero:

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Recall that A is symmetric: Ax = b.

Given two points (x_1, y_1) and (x_2, y_2) , find the line that passes through: $y = \alpha x + \beta$.

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$$\underbrace{\begin{bmatrix} x_1 & 1 \\ x_2 & 1 \end{bmatrix}}_{\text{A}} \underbrace{\begin{bmatrix} \alpha \\ \beta \end{bmatrix}}_{\text{B}} = \underbrace{\begin{bmatrix} y_1 \\ y_2 \end{bmatrix}}_{\text{A}} \qquad \Rightarrow x^* = A^{-1}b$$

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$$\underbrace{\begin{bmatrix} x_1 & 1 \\ x_2 & 1 \\ x_3 & 1 \end{bmatrix}}_{:=A(3\times2)} \underbrace{\begin{bmatrix} \alpha \\ \beta \end{bmatrix}}_{:=x(2\times1)} = \underbrace{\begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix}}_{:=b(3\times1)}$$

A is rectangular, and thus it is not invertible. Oops!

No biggie, let us rewrite the system:

$$\alpha \underbrace{\begin{bmatrix} x_1 & x_2 & x_3 \end{bmatrix}^\top}_{:=\vec{i}} + \beta \underbrace{\begin{bmatrix} 1 & 1 & 1 \end{bmatrix}^\top}_{:=\vec{i}} = \begin{bmatrix} y_1 & y_2 & y_3 \end{bmatrix}^\top$$

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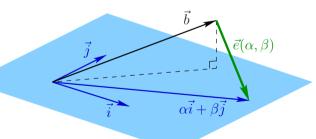
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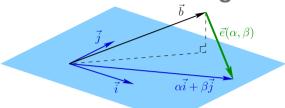
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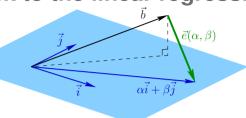
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$$\alpha \vec{i} + \beta \vec{j} = \vec{b}.$$

Solve for $\underset{\alpha,\beta}{\arg\min} \|\vec{e}(\alpha,\beta)\|$, where $\vec{e}(\alpha,\beta) := \alpha\vec{i} + \beta\vec{j} - \vec{b}$:

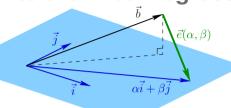






The $\|\vec{e}(\alpha, \beta)\|$ is minimized when $\vec{e}(\alpha, \beta) \perp \text{span}\{\vec{i}, \vec{j}\}$:

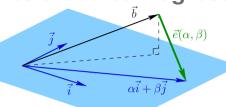
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$$\begin{bmatrix} x_1 & x_2 & x_3 \\ 1 & 1 & 1 \end{bmatrix} \left(\alpha \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} + \beta \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} - \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix} \right) = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

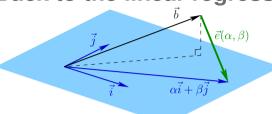


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$$A^{\top}(Ax-b)=\begin{bmatrix}0\\0\end{bmatrix}$$

In a general case the matrix $A^{T}A$ can be invertible!

$$A^{\top}Ax = A^{\top}b.$$

Theorem

 $A^{\top}A$ is symmetric.

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Theorem

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It follows from the fact that $x^{\top}A^{\top}Ax = (Ax)^{\top}Ax > 0$. Moreover, $A^{\top}A$ is positive definite in the case where A has linearly independent columns (rank A is equal to the number of the variables in the system).

The same reasoning applies, here is an algebraic way to show it:

 $arg min ||Ax - b||^2$

$$arg min ||Ax - b||^2 = arg min(Ax - b)^{\top}(Ax - b) =$$

$$\arg\min \|Ax - b\|^2 = \arg\min (Ax - b)^\top (Ax - b) = \arg\min (x^\top A^\top - b^\top) (Ax - b) =$$

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$$\arg\min \|Ax - b\|^2 = \arg\min(Ax - b)^{\top}(Ax - b) = \arg\min(x^{\top}A^{\top} - b^{\top})(Ax - b) =$$

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$$= \arg\min(x^{\top}A^{\top}Ax - 2b^{\top}Ax) =$$

$$\arg\min \|Ax - b\|^2 = \arg\min(Ax - b)^{\top}(Ax - b) = \arg\min(x^{\top}A^{\top} - b^{\top})(Ax - b) =$$

$$= \arg\min(x^{\top}A^{\top}Ax - b^{\top}Ax - x^{\top}A^{\top}b + \underbrace{b^{\top}b}_{\text{const}}) =$$

$$= \arg\min(x^{\top}A^{\top}Ax - 2b^{\top}Ax) = \arg\min(x^{\top}\underbrace{(A^{\top}A)}_{\text{indices}}x - 2\underbrace{(A^{\top}b)}_{\text{indices}}^{\top}x)$$

The same reasoning applies, here is an algebraic way to show it:

$$\arg\min \|Ax - b\|^2 = \arg\min(Ax - b)^{\top}(Ax - b) = \arg\min(x^{\top}A^{\top} - b^{\top})(Ax - b) =$$

$$= \arg\min(x^{\top}A^{\top}Ax - b^{\top}Ax - x^{\top}A^{\top}b + \underbrace{b^{\top}b}_{\text{const}}) =$$

$$= \arg\min(x^{\top}A^{\top}Ax - 2b^{\top}Ax) = \arg\min(x^{\top}\underbrace{(A^{\top}A)}_{:=A'}x - 2\underbrace{(A^{\top}b)}_{:=b'}^{\top}x)$$

The takeaway message

The least squares problem $\arg\min \|Ax - b\|^2$ is equivalent to minimizing the quadratic function $\arg\min \left(x^\top A'x - 2b'^\top x\right)$ with (in general) a symmetric positive definite matrix A'. This can be done by solving a linear system A'x = b'.

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Linear-quadratic regulator

Imagine that we have a car going at 0.5m/s. The goal is to accelerate to 2.3m/s. We can not control the speed directly, but we can act on the acceleration via the gas pedal:

$$v_{k+1} = v_k + u_k$$

where the signals are sampled every 1 second, v_k is the car speed and u_k is the acceleration of the car.

We have half a minute maximum to reach the goal, i.e, $v_0 = 0.5$ m/s, $v_n = 2.3$ m/s, n = 30s. So, we need to find $\{u_i\}_{i=0}^{n-1}$ that optimizes some quality criterion $J(\vec{v}, \vec{u})$:

$$\min J(\vec{v}, \vec{u})$$
 s.t. $v_{i+1} = v_i + u_i \ \forall i \in 0..n-1$



Dmitry Sokolov

Least squares for programmers

— with color plates —