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Index of programming contest code	e library	
Howard Cheng		
Arithmetic:		
bigint: Big (signed) integer arithme	tic	
bignumber.java:	e integer arithmetic (BigInteger	-1
cra:	e integer affenmette (biginteger	- /
Chinese remainder theorem diophantine_sys:		
Linear system of diophantine	equations (works for single equ	ation too!)
euclid: Euclidean algorithm		
eulerphi: Computes the Euler phi (totio	ent) function: given a positive	n. return
	en 1 and n relatively prime to n	
Extended Euclidean algorithm exp:		
Fast exponentiation		
expmod: Fast exponentiation mod m		
factor: Integer prime factorization		
factor_large:		
Integer prime factorization fflinsolve:	for larger integers (>= 2^40)	
Fraction-free solution of li- with integer coefficients)	near systems of equations (for s	systems
frac2dec:		
Obtain the decimal representation:	ation of a fraction.	
A rational number class. infix:		
Parses and evaluates infix a	rithmetic expressions.	
<pre>int_mult: Multiply integer factors on factors in the denominator w</pre>	the numerator, divide by the int	eger
linsolve:		
Solves linear systems of equation mult:	ations with LU decomposition.	
denominator without overflow	rator, divide by the factors in	the
ratlinsolve: Rational solution of linear fflinsolve as well).	systems of equations (can be sol	ved by
roman_numerals Converts between Arabic and	Roman numerals.	
Geometric (mostly 2-D):		
areapoly: Computes the signed area of	a simple (no self-intersection)	polvaon.
ccw:	-	
undefined). circle_3pts:	f 3 points (counterclockwise, cl	OCKWISE,
	us of a circle given 3 points.	
Computes the convex hull of dist3D:	a list of points.	
Computes the distance between two line segments, or a point	n two points, a point and a line t and a triangle in 3D. There a	re also
dist_line:	nfinite lines and infinite plane	:5.
Computes the distance of a pagreatcircle:	oint to a line.	
	n two points on a sphere along t	he surface.

Index Jan 16, 14 15:05 Page 2/3 Also has routines to convert between Cartesian coordinates and spherical coordinates. heron: Computes the area of a triangle given the lengths of 3 sides. intersect_circle_circle: Computes the intersection of two circles. intersectTF: Given two line segments, return whether they intersect or not (but doesn't return the point of intersection) intersect line: Given two 2-D line segments, return whether they intersect or not, and return the point of intersection if there is a unique one. intersect iline: Given two 2-D lines (infinite), return whether they intersect or not, and return the point of intersection if there is a unique one. intersect_iline_circle: Given an infinite 2-D line and a circle, return whether they intersect and also the point(s) of intersection. pointpoly: Given a polygon and a point, determines whether the point is in the polygon. The behaviour when the point is on the boundary is left to the user. polygon_inter: Given two convex polygons, compute their intersection as another polygon. Graph: bellmanford: Computes the shortest distance from one vertex to all other vertices. Also computes the paths. It is slow $(O(n^3))$ but handles negative weights. Can also be used to detect negative cycles. bfs_path: Computes the shortest distance from one vertex to all other vertices. Also computes the paths. The edges in the graph must have equal cost. bicomp: Finds the biconnected components and articulation points of a graph. Computes the shortest distance from one vertex to all other vertices. Also computes the paths. dijkstra_sparse: Same as dijkstra but for sparse graphs. Complexity $O((n+m) \log(n+m))$. Determines if there is an Eulerian tour in the graph. If so, find one. flovd: Computes the shortest distance between any two vertices. floyd_path: Like floyd, but also stores the paths. hungarian: Maximum/minimum weight bipartite matching. O(N^3). Compute unweighted matching of bipartite graphs. (Matthew) Compute the minimum spanning tree. mincostmaxflowdense: Compute the minimum cost maximum flow in a network. Good for dense graphs when maximum flow is small. Complexity is $O(n^2 * flow)$. mincostmaxflowsparse: Compute the minimum cost maximum flow in a network. Good for sparse graphs when maximum flow is small. Complexity is $O(m \log(m) * flow)$. networkflow: Compute the maximum flow in a network. Uses Ford-Fulkerson with complexity O(fm) where f is the value of the maximum flow and m is the number of edges. Good for sparse graphs where the maximum flow is small. networkflow2:

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    Compute the maximum flow in a network. Uses relabel-to-front
    with complexity O(n^3). Good for dense (but small) graphs
    where the maximum flow is large.
    Compute the strongly connected components (and possibly the
    compressed graph) of a directed graph.
 top_sort:
    Topological sort on directed acyclic graph (or detect if a
    cycle exists). O(n+m)
Data Structures:
 fenwicktree:
   A data structure that supports the maintainence of cumulative sums
    in an array dynamically. Most operations can be done in O(log N)
   time where N is the number of elements.
   An O(n) algorithm to construct a suffix array (and longest
    common prefix information) from a string.
Miscellaneous:
 asc_subseq:
    Longest (strictly) ascending/decreasing subsequence.
 binsearch:
   Binary search that also returns the position to insert an
    element if it is not found.
 common_subseq:
   Find the longest common subsequence of the two sequences.
 date:
   A class for dealing with dates in the Gregorian calendar.
 dow:
   Computing the day of the week.
 josephus:
   Finding the last survivor and killing order of the Josephus problem.
    Linear time string searching routines.
 int_prog:
    Integer programming.
 simplex:
   Linear programming by simplex algorithm.
 str_rotation_period:
   Computes the lexicographically least rotation of a string, as well
    as its period.
 unionfind:
    Union-find implementation to compute equivalence classes.
 vecsum:
   Find the contiguous subvector that gives the largest sum.
 zero_one:
    Zero-one programming.
```

```
2sat.cc
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                                                                         Page 1/2
// 2SAT solver: returns T/F whether it is satisfiable -- O(n+m)
    - use NOT() to negate a variable (works on negated ones too!)
    - ALWAYS use VAR() to talk about the non-negated version of the var i
    - use add clause to add a clause
    - one possible satisfying assignment is returned in val[], if
      it exists
     - To FORCE i to be true: add clause(G, VAR(i), VAR(i));
     - To implement XOR -- say (i XOR j) :
       add_clause(G, VAR(i), VAR(j)); add_clause(G, NOT(VAR(i)), NOT(VAR(j)));
       NOTE: val[] is indexed by i for var i, not by VAR(i)!!!
#include <iostream>
#include <algorithm>
#include <stack>
#include <cassert>
#include <vector>
using namespace std;
const int MAX VARS = 100;
                                    // maximum number of variables
const int MAX_NODES = 2*MAX_VARS;
struct Graph{
 int numNodes;
 vector<int> adj[MAX_NODES];
 void clear(){
    numNodes = 0;
    for(int i=0;i<MAX_NODES;i++)</pre>
      adi[i].clear();
 void add_edge(int u,int v){
   if(find(adj[u].begin(),adj[u].end(),v) == adj[u].end())
      adj[u].push_back(v);
};
int po[MAX_NODES], comp[MAX_NODES];
int num scc;
void DFS(int v, const Graph& G, int& C, stack<int>& P, stack<int>& S) {
 po[v] = C++;
 S.push(v); P.push(v);
 for (unsigned int i=0;i<G.adj[v].size();i++) {</pre>
    int w = G.adj[v][i];
    if (po[w] == -1) {
     DFS(w,G,C,P,S);
    } else if (comp[w] == -1) {
      while(!P.empty() && (po[P.top()] > po[w]))
       P.pop();
 if(!P.empty() && P.top() == v){
    while(!S.empty()){
      int t = S.top();
     S.pop();
      comp[t] = num_scc;
      if(t == v)
       break;
    P.pop();
    num_scc++;
int SCC(const Graph& G) {
 num\_scc = 0;
 int^{-}C = 1:
 stack<int> P,S;
 fill (po, po+G.numNodes, -1);
```

```
2sat.cc
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  fill(comp,comp+G.numNodes,-1);
  for (int i=0; i < G.numNodes; i++)</pre>
    if(po[i] == -1)
      DFS(i,G,C,P,S);
  return num_scc;
int VAR(int i) { return 2*i; }
int NOT(int i) { return i ^ 1; }
void add_clause(Graph &G, int v, int w) { // adds (v || w)
 if (v == NOT(w)) return;
  G.add_edge(NOT(v), w);
  G.add_edge(NOT(w), v);
bool twoSAT(const Graph &G, bool val[]) { // assumes graph is built
  for (int i = 0; i < G.numNodes; i += 2) {</pre>
    if (comp[i] == comp[i+1]) return false;
    val[i/2] = (comp[i] < comp[i+1]);
 return true;
// Declare this as a global variable if MAX_NODES is large to
    avoid Runtime Error.
Graph G;
int main(){
 int m,n;
  while(cin >> n >> m && (n || m)){
    G.clear();
    G.numNodes = 2*n;
    for (int i = 0; i < m; i++)
      cout << "Enter two variables for clause (1 - " << n
           < "), negative means negated: ";
      int x, y;
      cin >> x >> y;
      int var1 = VAR(abs(x)-1), var2 = VAR(abs(y)-1);
      if (x < 0) var1 = NOT(var1);
      if (y < 0) var2 = NOT (var2);
      add_clause(G, var1, var2);
    bool val[MAX_VARS];
    if (twoSAT(G, val)) {
      for (int i = 0; i < n; i++) {
        cout << val[i] << '';
      cout << endl;
    } else {
      cout << "Impossible" << endl;</pre>
 return 0;
```

```
areapoly.cc
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                                                                         Page 1/1
 * Area of a polygon
 * Author: Howard Cheng
 * Reference:
    http://www.exaflop.org/docs/cgafaq/cga2.html
 ^{\star} This routine returns the SIGNED area of a polygon represented as an
 * array of n points (n \geq 1). The result is positive if the orientation is
 * counterclockwise, and negative otherwise.
 */
#include <iostream>
#include <iomanip>
#include <cmath>
#include <cassert>
using namespace std;
struct Point {
 double x, y;
double area_polygon(Point polygon[], int n)
 double sum = 0.0;
 for (int i = 0; i < n-1; i++) {
   sum += polygon[i].x * polygon[i+1].y - polygon[i].y * polygon[i+1].x;
 sum += polygon[n-1].x * polygon[0].y - polygon[n-1].y * polygon[0].x;
 return sum/2.0;
int main(void)
 Point *polygon;
 int n;
 while (cin >> n && n > 0) {
   polygon = new Point[n];
    assert (polygon);
    for (int i = 0; i < n; i++) {
      cin >> polygon[i].x >> polygon[i].y;
    cout << "Area = " << fixed << setprecision(2)</pre>
         << area_polygon(polygon, n) << endl;
    delete[] polygon;
 return 0;
```

```
asc subseq.cc
 Jan 16, 14 15:08
                                                                        Page 1/3
 * Longest Ascending Subsequence
 * Author: Howard Cheng
 * Reference:
    Gries, D. The Science of Programming
 * Given an array of size n, asc_seq returns the length of the longest
 ^{\star} ascending subsequence, as well as one of the subsequences in S.
 * sasc_seq returns the length of the longest strictly ascending
 * subsequence. It runs in O(n log n) time.
 * Also included are simplified versions when only the length is needed.
 * Note: If we want to find do the same things with descending
 * subsequences, just reverse the array before calling the routines.
 */
#include <iostream>
#include <algorithm>
#include <vector>
#include <cassert>
using namespace std;
int asc_seq(int A[], int n, int S[])
 vector < int > last(n+1), pos(n+1), pred(n);
 if (n == 0)
    return 0;
 int len = 1;
 last[1] = A[pos[1] = 0];
 for (int i = 1; i < n; i++) {
    int j = upper_bound(last.begin()+1, last.begin()+len+1, A[i]) -
     last.begin();
    pred[i] = (j-1 > 0) ? pos[j-1] : -1;
    last[j] = A[pos[j] = i];
   len = max(len, j);
 int start = pos[len];
 for (int i = len-1; i >= 0; i--) {
   S[i] = A[start];
   assert(i == 0 || pred[start] < start);</pre>
   start = pred[start];
 return len;
int asc_seq(int A[], int n)
 vector<int> last(n+1);
 if (n == 0) {
    return 0;
 int len = 1;
 last[1] = A[0];
 for (int i = 1; i < n; i++) {
   int j = upper_bound(last.begin()+1, last.begin()+len+1, A[i]) -
     last.begin();
    last[j] = A[i];
```

```
asc subseq.cc
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                                                                         Page 2/3
    len = max(len, j);
 return len;
int sasc_seq(int A[], int n, int S[])
 vector < int > last(n+1), pos(n+1), pred(n);
  if (n == 0) {
    return 0;
  int len = 1;
 last[1] = A[pos[1] = 0];
  for (int i = 1; i < n; i++) {
    int j = lower_bound(last.begin()+1, last.begin()+len+1, A[i]) -
     last.begin();
   pred[i] = (j-1 > 0) ? pos[j-1] : -1;
    last[j] = A[pos[j] = i];
   len = max(len, j);
  int start = pos[len];
  for (int i = len-1; i >= 0; i--) {
   S[i] = A[start];
   start = pred[start];
 return len;
int sasc_seq(int A[], int n)
 vector<int> last(n+1);
 if (n == 0) {
   return 0;
  int len = 1;
 last[1] = A[0];
  for (int i = 1; i < n; i++) {
   int j = lower_bound(last.begin()+1, last.begin()+len+1, A[i]) -
     last.begin();
    last[j] = A[i];
   len = max(len, j);
 return len;
int main (void)
 int *A, *S, n, i, k;
  while (cin >> n && n > 0) {
   A = new int[n];
    S = new int[n];
   for (i = 0; i < n; i++) {</pre>
     cin >> A[i];
   k = asc\_seq(A, n, S);
    cout << "length = " << k << endl;
    for (i = 0; i < k; i++) {
     cout << S[i] << "";
    cout << endl;
```

```
asc_subseq.cc
                                                                                                                                                 Page 3/3
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     k = sasc_seq(A, n, S);
cout << "length = " << k << endl;
for (i = 0; i < k; i++) {
  cout << S[i] << " ";</pre>
     cout << endl;
delete[] A;
delete[] S;</pre>
 return 0;
```

```
bellmanford.cc
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                                                                        Page 1/3
 * Bellman-Ford Shortest Path Algorithm
 * Author: Howard Cheng
 * Given a weight matrix representing a graph and a source vertex, this
 * algorithm computes the shortest distance, as well as path, to each
 * of the other vertices. The paths are represented by an inverted list,
 * such that if v preceeds immediately before w in a path from the
 * source to vertex w, then the path P[w] is v. The distances from
 * the source to v is given in \overline{\text{D[v]}} (DISCONNECT if not connected).
 * Call get_path to recover the path.
 * Note: the Bellman-Ford algorithm has complexity O(n^3), but it works even
         when edges have negative weights. As long as there are no negative
         cycles the computed results are correct.
         We can make this O(n*m) if we use an adjacency list representation.
         This works for directed graphs too.
         You can use this to detect negative cycles too. See code.
 */
#include <iostream>
#include <climits>
#include <cassert>
using namespace std;
const int MAX_NODES = 20;
const int DISCONNECT = INT MAX;
/* assume that D and P have been allocated */
void bellmanford(int graph[MAX_NODES][MAX_NODES], int n, int src,
                 int D[], int P[])
 int v, w, k;
 for (v = 0; v < n; v++) {
   D[v] = INT MAX;
   P[v] = -1;
 D[src] = 0;
 for (k = 0; k < n-1; k++)
   for (v = 0; v < n; v++) {
     for (w = 0; w < n; w++) {
       if (graph[v][w] != DISCONNECT && D[v] != INT_MAX) {
          if (D[w] == INT_MAX \mid \mid D[w] > D[v] + graph[v][w]) {
           D[w] = D[v] + graph[v][w];
           P[w] = v;
          } else if (D[w] == D[v] + graph[v][w]) {
            /* do some tie-breaking here */
  /* the following loop is used only to detect negative cycles, not */
  /* needed if you don't care about this
 for (v = 0; v < n; v++) {
   for (w = 0; w < n; w++)
     if (graph[v][w] != DISCONNECT && D[v] != INT_MAX) {
        if (D[w] == INT_MAX \mid D[w] > D[v] + graph[v][w]) {
          /* if we get here then there is a negative cycle somewhere */
```

```
bellmanford.cc
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          /* on the path from src to
int get_path(int v, int P[], int path[])
 int A[MAX_NODES];
 int i, k;
 k = 0;
 A[k++] = v;
 while (P[v] != -1) {
   v = P[v];
   A[k++] = v;
 for (i = k-1; i >= 0; i--) {
   path[k-1-i] = A[i];
 return k;
int main(void)
 int m, w, num;
 int i, j;
 int graph[MAX_NODES][MAX_NODES];
 int P[MAX_NODES][MAX_NODES], D[MAX_NODES][MAX_NODES];
 int path[MAX_NODES];
  /* clear graph */
 for (i = 0; i < MAX NODES; i++) {
   for (j = 0; j < MAX_NODES; j++) {</pre>
     graph[i][j] = DISCONNECT;
  /* read graph */
 cin >> i >> j >> w;
 while (!(i = -1 & k & j = -1)) {
   assert(0 <= i && i < MAX_NODES && 0 <= j && j < MAX_NODES);
   graph[i][j] = graph[j][i] = w;
   cin >> i >> j >> w;
 for (i = 0; i < MAX NODES; i++) {</pre>
   bellmanford(graph, MAX_NODES, i, D[i], P[i]);
  /* do queries */
 cin >> i >> j;
 while (!(i = -1 \&\& j = -1)) {
   assert(0 <= i && i < MAX_NODES && 0 <= j && j < MAX_NODES);
   cout << i << "" << j << ":" << D[i][j] << endl;
   for (m = j; m != -1; m = P[i][m]) {
     cout << " " << m;
   cout << endl;
   num = get_path(j, P[i], path);
   for (m = 0; m < num; m++) {</pre>
     cout << " " << path[m];
   cout << endl;
   cin >> i >> j;
```

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return 0;		
,		

```
bfs path.cc
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                                                                          Page 1/3
 * Shortest Path with BFS
 * Author: Howard Cheng
 * Given a graph represented by an adjacency list, this algorithm uses
 * BFS to find the shortest path from a source vertex to each of the
 ^{\star} other vertices. The distances from the source to v is given in D[v], and
 * D[v] is set to -1 if the source vertex is not connected to w. Also,
 * the shortest path tree is stored in the array P.
 * Call get_path to recover the path.
 \mbox{\ensuremath{^{\star}}} Note: All edges must have the same cost for this to work.
         This algorithm has complexity O(n+m).
#include <iostream>
#include <cassert>
#include <algorithm>
#include <queue>
using namespace std;
const int MAX_NODES = 100;
struct Node {
                      /* number of outgoing edges */
 int deg;
 int adj[MAX_NODES];
 /* the following is not necessary, but useful in many situations */
 int cost[MAX_NODES];
};
void BFS_shortest_path(Node graph[], int n, int src, int D[], int P[])
 char used[MAX NODES];
 queue<int> q;
 int i, v, w;
 fill(used, used+MAX_NODES, 0);
 q.push(src);
 used[src] = 1;
 for (i = 0; i < MAX_NODES; i++) {</pre>
   D[i] = -1;
   P[i] = -1;
 D[src] = 0;
 while (!q.empty()) {
   v = q.front();
    q.pop();
    for (i = 0; i < graph[v].deg; i++) {</pre>
      w = graph[v].adj[i];
      if (!used[w]) {
       D[w] = D[v] + 1;
        P[w] = v;
        q.push(w);
       used[w] = 1;
      } else if (D[v] + 1 == D[w]) {
        /* put tie-breaker here */
        /* eg. find largest path in lexicographic order, when the path */
              is considered in REVERSE!
        P[w] = max(P[w], v);
```

```
bfs path.cc
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                                                                         Page 2/3
void clear(Node graph[], int n)
 int i;
 for (i = 0; i < n; i++) {</pre>
   graph[i].deg = 0;
void add_edge(Node graph[], int v, int w, int cost)
 int i;
  /* make sure that we have no duplicate edges */
 for (i = 0; i < graph[v].deg; i++) {</pre>
   if (graph[v].adj[i] == w) {
     assert(0);
 }
 graph[v].cost[graph[v].deg] = cost;
 graph[v].adj[graph[v].deg] = w;
 graph[v].deg++;
int get_path(int v, int P[], int path[])
 int A[MAX_NODES];
 int i, k;
 k = 0;
 A[k++] = v;
 while (P[v] != -1) {
   v = P[v];
   A[k++] = v;
 for (i = k-1; i >= 0; i--) {
   path[k-1-i] = A[i];
 return k;
int main (void)
 int v, w, num;
 int i;
 Node graph[MAX_NODES];
 int P[MAX_NODES][MAX_NODES], D[MAX_NODES][MAX_NODES];
 int path[MAX_NODES];
 clear(graph, MAX_NODES);
 while (cin >> v >> w && v >= 0 && w >= 0) {
   add_edge(graph, v, w, 1);
 for (i = 0; i < MAX_NODES; i++) {</pre>
   BFS_shortest_path(graph, MAX_NODES, i, D[i], P[i]);
 while (cin >> v >> w && v >= 0 && w >= 0) {
   cout << v << "" << w << ":" << D[v][w] << endl;
   num = get_path(w, P[v], path);
   assert (D[v][w] == -1 || num == D[v][w]+1);
   for (i = 0; i < num; i++) {
     cout << " " << path[i];
   cout << endl;
```

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return 0;		
}		

```
bicomp.cc
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                                                                        Page 1/2
 * Biconnected Components
 * Author: Howard Cheng
 * Date: Oct 15, 2004
 * The routine bicomp() uses DFS to find the biconnected components in
 * a graph. The graph is stored as an adjacency list. Use clear_graph()
 * and add_edge() to build the graph.
 * Note: This works only on connected graphs. See comment below in code.
 ^{\star} The code simply prints the biconnected components and the articulation
 * points. Replace the printing code to do whatever is appropriate.
 * NOTE: some articulation points may be printed multiple times.
#include <iostream>
#include <stack>
#include <algorithm>
#include <cassert>
using namespace std;
/* maximum number of nodes, maximum degree, and maximum number of edges */
const int MAX_N = 1000;
const int MAX_DEG = 4;
struct Node {
 int deg;
 int nbrs[MAX_DEG];
 int dfs, back;
int dfn;
void clear_graph(Node G[], int n)
 int i;
 for (i = 0; i < n; i++) {
   G[i].deq = 0;
void add_edge(Node G[], int u, int v)
 G[u].nbrs[G[u].deg++] = v;
 G[v].nbrs[G[v].deq++] = u;
void do_dfs(Node G[], int v, int pred, stack<int> &v_stack,
            stack<int> &w_stack)
 int i, w, child = 0;
 G[v].dfs = G[v].back = ++dfn;
 for (i = 0; i < G[v].deg; i++) {</pre>
   w = G[v].nbrs[i];
    if (G[w].dfs < G[v].dfs && w != pred) {
     /* back edge or unexamined forward edge */
     v_stack.push(v);
     w_stack.push(w);
    if (!G[w].dfs) {
     do_dfs(G, w, v, v_stack, w_stack);
      child++;
```

```
bicomp.cc
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                                                                               Page 2/2
       /* back up from recursion */
      if (G[w].back >= G[v].dfs) {
        /* new bicomponent */
        cout << "edges in new biconnected component:" << endl;</pre>
        while (v_stack.top() != v || w_stack.top() != w) {
  cout << v_stack.top() << " " << w_stack.top() << endl;</pre>
           v_stack.pop();
          w_stack.pop();
        cout << v_stack.top() << " " << w_stack.top() << endl;</pre>
        v_stack.pop();
        w_stack.pop();
        if (pred != −1)
          cout << "articulation point: " << v << endl;</pre>
      } else
        G[v].back = min(G[v].back, G[w].back);
    } else
      /* w has been examined already */
      G[v].back = min(G[v].back, G[w].dfs);
  if (pred == -1 && child > 1) {
    cout << "articulation point: " << v << endl;</pre>
void bicomp(Node G[], int n)
  int i;
 stack<int> v_stack, w_stack;
  dfn = 0;
  for (i = 0; i < n; i++) {</pre>
   G[i].dfs = 0;
  do_dfs(G, 0, -1, v_stack, w_stack);
  // NOTE: if you wish to process all connected components, you can simply
  // run the following code instead of the line above:
  // for (int i = 0; i < n; i++) {
  // if (G[i].dfs == 0) {
         do_dfs(G, i, -1, v_stack, w_stack);
  //
  // }
int main (void)
 Node G[MAX_N];
 int n, m, i, u, v;
  cin >> n;
  clear_graph(G, n);
  cin >> m;
  for (i = 0; i < m; i++) {</pre>
    cin >> u >> v;
    add_edge(G, u-1, v-1);
 bicomp(G, n);
  return 0;
```

```
bigint.cc
 Sep 29, 16 13:22
                                                                        Page 1/11
 * Big integer implementation
 * Author: Howard Cheng
 * Each digit in our representation represents LOG_BASE decimal digits
#include <vector>
#include <string>
#include <cstdio>
#include <cctype>
#include <iostream>
#include <algorithm>
#include <utility>
#include <cassert>
using namespace std;
using namespace std::rel_ops;
typedef long long Digit;
#define BASE 1000000000
#define LOG_BASE 9
#define FMT_STR "%lld"
#define FMT_STR0 "%09lld"
class BigInteger {
private:
                      // +1 = positive, 0 = zero, -1 = negative
 int sign;
 vector<Digit> mag; // magnitude
 void normalize();
public:
 BigInteger(Digit n = 0);
 BigInteger(const string &s);
                                  // no error checking
 long long toLongLong() const;
                                  // convert to long long (assumes no overflow)
 string toString() const;
                                   // convert to string
 void clear(); // set to zero
  // comparison
 bool operator < (const BigInteger &a) const;
 bool operator == (const BigInteger &a) const;
 bool isZero() const;
  // arithmetic
 BigInteger &operator+=(const BigInteger &a);
 BigInteger & operator -= (const BigInteger &a);
 BigInteger & operator *= (const BigInteger &a);
 BigInteger & operator *= (Digit a);
 BigInteger & operator <<= (Digit a);</pre>
 BigInteger & operator/=(const BigInteger &a);
 BigInteger & operator/= (Digit a);
 BigInteger & operator % = (const BigInteger &a);
 friend Digit operator% (const BigInteger &a, Digit b);
 // we have *this = b * q + r
 // r is such that 0 <= r < |b|
 void divide(const BigInteger &b, BigInteger &q, BigInteger &r) const;
 void divide (Digit b, BigInteger &g, Digit &r) const;
  // root = floor(sqrt(a)). Returns 1 if a is a perfect square, 0 otherwise.
 // assume >= 0
 int sqrt (BigInteger &root) const;
```

```
bigint.cc
 Sep 29, 16 13:22
                                                                        Page 2/11
BigInteger operator+(const BigInteger &a, const BigInteger &b);
BigInteger operator-(const BigInteger &a, const BigInteger &b);
BigInteger operator* (const BigInteger &a, const BigInteger &b);
BigInteger operator* (const BigInteger &a, Digit b);
BigInteger operator << (const BigInteger &a, Digit b);
BigInteger operator/(const BigInteger &a, const BigInteger &b);
BigInteger operator/(const BigInteger &a, Digit b);
BigInteger operator% (const BigInteger &a, const BigInteger &b);
Digit operator% (const BigInteger &a, Digit b);
BigInteger power(BigInteger x, Digit y);
istream & operator >> (istream & is, BigInteger & a);
ostream & operator << (ostream & os, const BigInteger & a);
void BigInteger::normalize()
 if (mag.size() == 0) {
    return;
  vector<Digit>::iterator p = mag.end();
   if (*(--p) != 0) break;
  } while (p != mag.begin());
 if (p == mag.begin() && *p == 0) {
    clear();
    sign = 0;
 } else {
    mag.erase(++p, mag.end());
BigInteger::BigInteger(Digit n)
  if (n == 0) {
    sign = 0;
    return:
  if (n < 0)
   sign = -1:
    n = -n;
  } else {
   sign = 1;
  while (n > 0) {
    mag.push_back(n % BASE);
    n /= BASE;
BigInteger::BigInteger(const string &s)
  int 1 = 0;
 bool zero = true;
 bool neg = false;
  clear();
  sign = 1;
  if (s[1] == '-') {
   neg = true;
    1++;
  for (; 1 < s.length(); l++) {</pre>
    *this *= 10;
    *this += s[1] - '0';
    zero &= s[1] == '0';
```

```
bigint.cc
 Sep 29, 16 13:22
 if (zero) {
   clear();
 if (neg) {
   sign = -1;
long long BigInteger::toLongLong() const
 long long a = 0;
 for (int i = mag.size()-1; i >= 0; i--) {
   a *= BASE;
   a += mag[i];
 return sign * a;
string BigInteger::toString() const
 char buffer[LOG_BASE+1];
 string s;
 if (isZero()) {
    return "0";
  else
    if (sign < 0) {
     s += "-";
    for (int i = mag.size()-1; i >= 0; i--) {
      if (i == (int) (mag.size()-1))
       sprintf(buffer, FMT_STR, mag[i]);
      } else {
       sprintf(buffer, FMT_STR0, mag[i]);
     s += buffer;
    return s;
void BigInteger::clear()
 sign = 0;
 mag.clear();
bool BigInteger::operator<(const BigInteger &a) const
 if (sign != a.sign) +
   return sign < a.sign;
 } else if (sign == 0) {
   return false;
 } else if (mag.size() < a.mag.size()) {</pre>
    return sign > 0;
 } else if (mag.size() > a.mag.size()) {
    return sign < 0;
  } else {
    for (int i = mag.size()-1; i >= 0; i--) {
      if (mag[i] < a.mag[i]) {
       return sign > 0;
      } else if (mag[i] > a.mag[i]) {
       return sign < 0;</pre>
    return false;
```

```
bigint.cc
 Sep 29, 16 13:22
                                                                         Page 4/11
bool BigInteger::operator == (const BigInteger &a) const
  return sign == a.sign && mag == a.mag;
bool BigInteger::isZero() const
 return sign == 0;
BigInteger &BigInteger::operator+=(const BigInteger &a)
  if (a.sign == 0) {
    return *this;
  } else if (sign == 0) {
    sign = a.sign;
    mag = a.mag;
    return *this;
  } else if (sign < 0 && a.sign > 0) {
    BigInteger b(a);
    sign = 1;
    b -= *this;
    return *this = b;
  } else if (sign > 0 && a.sign < 0) {
    BigInteger b(a);
    b.sign = 1;
    return (*this) -= b;
    else {
    Digit carry = 0;
    unsigned int limit = max(mag.size(), a.mag.size());
    for (unsigned int i = 0; i < limit; i++) {</pre>
      Digit s\tilde{1} = (i < mag.size()) ? mag[i] : 0;
      Digit s2 = (i < a.mag.size()) ? a.mag[i] : 0;
      Digit sum = s1 + s2 + carry;
      Digit result = (sum < BASE) ? sum : sum - BASE;
      carry = (sum >= BASE);
      if (i < mag.size())
        mag[i] = result;
      } else {
        mag.push_back(result);
    if (carry) {
      mag.push_back(carry);
    return *this;
BigInteger &BigInteger::operator == (const BigInteger &a)
  if (a.sign == 0) {
    return *this;
  } else if (sign == 0) {
    sign = -a.sign;
    mag = a.mag;
    return *this;
  } else if (sign != a.sign) {
    BigInteger b(a);
    b.sign *=-1;
    return *this += b;
  } else if (sign < 0)
    BigInteger b(a);
    b.sign \dot{*}=-1;
    sign^* = -1;
    b -= *this;
    return *this = b;
```

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```
bigint.cc
 Sep 29, 16 13:22
                                                                        Page 5/11
 } else
    if (*this == a) {
      clear();
      return *this;
    } else if (*this < a) {</pre>
     BigInteger b(a);
      b -= *this;
     b.sign *=-1;
     return *this = b;
    } else {
      // we know that *this > a
      unsigned int limit = mag.size();
      Digit borrow = 0;
      for (unsigned int i = 0; i < limit; i++) {</pre>
       Digit s1 = mag[i];
       Digit s2 = (i < a.mag.size()) ? a.mag[i] : 0;
       Digit diff = s1 - s2 - borrow;
       mag[i] = (diff >= 0) ? diff : diff + BASE;
       borrow = (diff < 0);
     normalize();
     return *this;
BigInteger &BigInteger::operator*=(const BigInteger &a)
 BigInteger temp(*this);
 BigInteger c;
 if (this == &a) {
                    // make a copy to prevent clobbering it
 const BigInteger &b = (this == &a) ? c : a;
 clear();
 if (b.sign) {
    for (unsigned int i = 0; i < b.mag.size(); i++) {</pre>
     if (b.mag[i] != 0) {
        *this += (temp * b.mag[i]);
     temp <<= 1;
    sign *= b.sign;
 return *this;
BigInteger &BigInteger::operator*=(Digit a)
 if (a <= -BASE | | a >= BASE) {
   BigInteger b(a);
   return (*this *= b);
 if (isZero()) {
    return *this;
  } else if (a == 0) {
   clear();
    return *this;
  } else if (a < 0) {
   sign *= -1;
    a = -a;
 Digit carry = 0;
 for (unsigned int i = 0; i < mag.size(); i++) {
```

```
bigint.cc
 Sep 29, 16 13:22
                                                                       Page 6/11
    Digit prod = a * mag[i];
    mag[i] = (carry + prod) % BASE;
    carry = (carry + prod) / BASE;
  if (carry) {
    mag.push_back(carry);
  return *this;
BigInteger &BigInteger::operator<<=(Digit a)</pre>
  assert(a >= 0);
  if (sign) {
   while (a-- > 0)
     mag.insert(mag.begin(), 0);
  return *this;
BigInteger &BigInteger::operator/=(const BigInteger &a)
  BigInteger temp(*this), r;
 temp.divide(a, *this, r);
  return *this;
BigInteger &BigInteger::operator/=(Digit a)
  BigInteger temp(*this);
  Digit r;
  temp.divide(a, *this, r);
  return *this;
BigInteger &BigInteger::operator%=(const BigInteger &a)
  BigInteger temp(*this), q;
  temp.divide(a, q, *this);
  return *this;
void BigInteger::divide(const BigInteger &b, BigInteger &g,
                        BigInteger &r) const
  // reference Knuth v.2 Algorithm D
  assert(!b.isZero());
  if (b.mag.size() == 1) {
   Digit r2;
    divide(b.sign*b.mag[0], q, r2);
    r = r2;
    return;
  r = *this:
  if (r.sign < 0) {
    r.sign = 1;
  q.clear();
  int n = b.mag.size();
  int m = mag.size() - n;
  if (m >= 0) {
    BigInteger v(b);
    q.maq.resize(m+1);
    q.sign = 1;
```

```
bigint.cc
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                                                                     Page 7/11
  // D1: normalize
  Digit d = BASE / (v.mag[n-1] + 1); // Book is wrong. See errata on web
  r *= d;
  v *= d;
  while ((int)r.mag.size() < m+n+1) {
    r.mag.push_back(0);
  // loop
  for (int j = m; j >= 0; j--) {
    // D3: calculate q2
    Digit t = r.mag[j+n] * BASE + r.mag[j+n-1];
    Digit q2 = t / v.mag[n-1];
    Digit r2 = t - q2 * v.mag[n-1];
    if (q2 == BASE | | q2 * v.mag[n-2] > BASE * r2 + r.mag[j+n-2]) {
      q2--;
      r2 += v.mag[n-1];
      if (r2 < BASE &&
           (q2 == BASE \mid | q2 * v.mag[n-2] > BASE * r2 + r.mag[j+n-2])) {
        q2--;
        r2 += v.mag[n-1];
    // D4: multiply and subtract
    Digit carry, borrow, diff;
    carry = borrow = 0;
    for (int i = 0; i <= n; i++)
      t = q2 * ((i < n) ? v.mag[i] : 0) + carry;
      carry = t / BASE;
      t %= BASE;
      diff = r.mag[j+i] - t - borrow;
      r.mag[j+i] = (diff >= 0 || i == n) ? diff : diff + BASE;
      borrow = (diff < 0);
    // D5: test remainder
    q.mag[j] = q2;
    if (r.mag[n+j] < 0) {
      // D6: add back
      q.mag[j]--;
      carry = 0;
      for (int i = 0; i < n; i++) {
        t = r.mag[j+i] + v.mag[i] + carry;
        r.mag[j+i] = (t < BASE) ? t : t - BASE;
        carry = (t >= BASE);
      r.mag[j+n] += carry;
  q.normalize();
  r.normalize();
  // D8: unnormalize
  r /= d;
 // normalize
if (sign < 0 && b.sign > 0) {
  q.sign *=-1;
  r *= -1;
  if (!r.isZero()) {
    r += b;
    q -= 1;
} else if (sign > 0 && b.sign < 0) {
```

```
bigint.cc
 Sep 29, 16 13:22
                                                                      Page 8/11
    g.sign *=-1;
  } else if (sign < 0 && b.sign < 0 && !r.isZero()) {
   r += b;
   r *= -1;
   q += 1;
void BigInteger::divide(Digit b, BigInteger &q, Digit &r) const
 if (b <= -BASE | | b >= BASE) {
   BigInteger bb(b), rr;
   divide(bb, q, rr);
   r = rr.toLongLong();
   return;
  int bsign = 1;
  if (b < 0) {
   b *= -1;
   bsign = -1;
  q.clear();
  r = 0;
  for (int i = mag.size()-1; i >= 0; i--) {
   Digit t = r * BASE + mag[i];
   if (t / b > 0) {
     q.sign = 1;
   q.mag.insert(q.mag.begin(), t / b);
   r = t - q.mag[0] * b;
  // normalize
 q.normalize();
 if (sign < 0 && bsign > 0) {
   q.sign *=-1;
    r *= -1;
   if (r)
     r += b;
     q -= 1;
 } else if (sign > 0 && bsign < 0) {
   a.sign *= -1;
 } else if (sign < 0 && bsign < 0 && r) {
   r = b - r;
   q += 1;
int BigInteger::sqrt(BigInteger &root) const
 assert(sign >= 0);
 root.clear();
 if (sign == 0) {
   return 1;
 // figure out how many digits there are
 BigInteger x, r, t2;
 r.sign = 1;
 int d = mag.size();
  int root_d = (d % 2) ? (d+1)/2 : d / 2;
  if (d % 2) {
   r.mag.resize(1);
```

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```
bigint.cc
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                                                                       Page 9/11
    r.mag[0] = mag[--d];
 } else {
   r.mag.resize(2);
   r.mag[1] = mag[--d];
   r.mag[0] = mag[--d];
 root.sign = 1;
  // figure out one digit at a time
 for (int k = root_d - 1; k >= 0; k--) {
    // invariant: result is the sqrt (integer part) of the digits processed
    // so far
    // look for next digit in result by binary search
   x = root * 2;
    x <<= 1;
    Digit t;
    Digit lo = 0, hi = BASE;
    while (hi - lo > 1) {
     Digit mid = (lo + hi) / 2;
     x.mag[0] = t = mid;
     t2 = x * t;
     if (t2 < r || t2 == r) {
       lo = mid;
     } else {
       hi = mid;
   root <<= 1;
    root.mag[0] = lo;
    // form the next r
    x.mag[0] = t = lo;
   t2 = x * t;
   r -= t2;
   r <<= 1;
   r += (d > 0) ? mag[--d] : 0;
   r <<= 1;
    r += (d > 0) ? mag[--d] : 0;
 return r.isZero();
BigInteger operator+(const BigInteger &a, const BigInteger &b)
 BigInteger r(a);
 r += b;
 return r;
BigInteger operator-(const BigInteger &a, const BigInteger &b)
 BigInteger r(a);
 r -= b;
 return r;
BigInteger operator* (const BigInteger &a, const BigInteger &b)
 BigInteger r(a);
 r *= b;
 return r;
BigInteger operator*(const BigInteger &a, Digit b)
```

```
bigint.cc
 Sep 29, 16 13:22
                                                                      Page 10/11
  BigInteger r(a);
  r *= b;
  return r;
BigInteger operator << (const BigInteger &a, Digit b)
  BigInteger r(a);
 r <<= b;
  return r;
BigInteger operator/(const BigInteger &a, const BigInteger &b)
  BigInteger r(a);
 r /= b;
  return r;
BigInteger operator/(const BigInteger &a, Digit b)
 BigInteger r(a);
 r /= b;
 return r;
BigInteger operator%(const BigInteger &a, const BigInteger &b)
 BigInteger r(a);
 r %= b;
 return r;
Digit operator% (const BigInteger &a, Digit b)
  Digit r;
  if (b > 0 && b < BASE) {
    r = 0;
    for (int i = a.mag.size()-1; i >= 0; i--) {
     r = ((r * BASE) + a.mag[i]) % b;
    if (a.sign < 0) {
     r = (b - r) % b;
    return r;
  BigInteger q;
  a.divide(b, q, r);
  return r;
BigInteger power (BigInteger x, Digit y)
  BigInteger result(1), sx(x);
  assert (y >= 0);
  while (y > 0)
    if (y & 0x01) {
      result *= sx;
    } else {
      sx *= sx;
      y >>= 1;
  return result;
```

src/bigint.cc

```
bigint.cc
 Sep 29, 16 13:22
                                                                        Page 11/11
istream & operator>>(istream &is, BigInteger &a)
 string s;
char c = '';
 is.get(c);
 while (!is.eof() && isspace(c)) {
    is.get(c);
 if (is.eof()) {
    if (isdigit(c)) {
      a = (int)(c - '0');
      is.clear();
    return is;
 if (c == '-') {
   s = "-";
 } else {
    is.unget();
    if (!isdigit(c)) {
      return is;
 is.get(c);
 while (!is.eof() && isdigit(c)) {
   s += c;
    is.get(c);
 if (!is.eof()) {
    is.unget();
 a = s;
 is.clear();
 return is;
ostream &operator<<(ostream &os, const BigInteger &a)
 return (os << a.toString());</pre>
int main()
 BigInteger a, b;
 while (cin >> a >> b && (!(a == 0) || !(b == 0))) {
    cout << "a = " << a << endl;
    cout << "b = " << b << endl;
    if (!(a < 0)) {
      if (a.sqrt(b)) {
        cout << "perfect square" << endl;</pre>
      cout << "sqrt(a) = " << b << endl;
 return 0;
```

```
bignumber.java
 Jan 16, 14 15:11
                                                                        Page 1/2
// Java template for using BigInteger class.
// Note that there is also a similar BigDecimal class which may be useful.
// Name of the file must be NAME.java where NAME is the class name.
// To compile:
  javac NAME.java
// To run:
// java NAME
// Note: in Java, all non-native types (including arrays) need to be
// allocated by new. Multidimensional arrays can be allocated in one
// call. See below.
// for importing IO routines
import java.io.*;
import java.util.Scanner;
import java.math.BigInteger;
class bignumber {
  // this is main
 public static void main(String argv[])
    // A scanner can be used to read many different types
    Scanner sc = new Scanner(System.in);
    \ensuremath{//} checking whether there is a next token can be done before
    // actually reading it and fail (vs. I/O model in C++)
    while (sc.hasNextInt()) {
     int N = sc.nextInt();
     int K = sc.nextInt();
      // here is how to allocate two dimensional arrays
     BigInteger binom[][] = new BigInteger[N+1][N+1];
     for (int n = 0; n <= N; n++)
       // here is how you construct from an integer.
       binom[n][0] = binom[n][n] = BigInteger.valueOf(1);
       for (int k = 1; k < n; k++) {
          binom[n][k] = binom[n-1][k-1].add(binom[n-1][k]);
      // to print something, use System.out.println().
      // Arguments are strings (in double quotes), and most data types
      // can be converted to strings and concatenated.
      // Call it with no argument to produce a blank line, or use print()
      // to print without a trailing end-of-line
      System.out.println("C(" + N + "," + K + ") = " + binom[N][K]);
    Here are a bunch of other things you can do with BigIntegers.
    Assuming a, b, c, d are BigIntegers, n is an int
    a = BigInteger.ZERO;
                                       a = 0
    a = BigInteger.ONE;
                                       a = 1
    a = new BigInteger("FF", 16);
                                       a = 255
    a = sc.nextBigInteger();
                                       cin >> a
    s = a.toString();
                                       convert to string representation
```

```
bignumber.java
Jan 16, 14 15:11
                                                                        Page 2/2
   s = a.toString(base);
                                       convert to string representation in
                                       given base
  x = a.intValue();
                                       convert to smaller types, but may
  x = a.longValue();
                                       lose precision
  x = a.floatValue();
  x = a.doubleValue();
                                       a = |b|
  a = b.abs():
                                      n = 0, +1, -1 depending on sign of a
  n = a.signum();
                                      a = -b;
  a = b.negate();
                                      a = b + c
  a = b.add(c);
  a = b.subtract(c);
                                      a = b-c
                                      a = b*c
  a = b.multiply(c);
  a = b.divide(c);
                                      a = b/c
  a = b.remainder(c);
                                      a = b %c
                                      a = b%c, but c must be positive
  a = b.mod(c);
                                                and a >= 0
  if (a.compareTo(b))
                                      -1 if a < b
                                       0 \text{ if } a == b
                                      1 \text{ if } a > b
  if (a.equals(b))
                                      true iff a == b
  a = b.min(c);
                                      a = min(b, c)
  a = b.max(c);
                                      a = max(b, c)
  a = b.pow(n);
                                      a = pow(b, n)
  a = b.modpow(n, c);
                                      a = pow(b, n) \mod c, c > 0
                                      a = gcd(|b|, |c|)
  a = b.gcd(c);
  a = b.modInverse(c);
                                       a = b^{(-1)} \mod c
  n = a.bitLength();
                                       number of bits in 2's complement
                                       representation, minus the sign bit
  if (a.isProbablePrime(n))
                                       whether a is prime, with error
                                       of (1/2)^n
  System.exit(0);
```

```
binsearch.cc
 Jan 16, 14 15:08
                                                                         Page 1/2
 * Binary Search
 * Author: Howard Cheng
 * Note: you may wish to use the STL functions lower_bound and upper_bound
 * instead.
 ^{\star} Given a sorted array A of size n, it tries to find an item x in the
 * the array using binary search. The function returns non-zero if
 * x is found, and zero otherwise. Furthermore, if it is found, then
 * A[index] = x. If it is not found, then index is the place x should
 * be inserted into A.
 * ie. A[i] \le x
                            for 0 \le i \le index
                x < A[i] for index \leq i < n
 ^{\star} This routine is written for integer arrays, but can be adapted to
 ^{\star} other types by changing the comparison operator.
 * There is also an insert routine here that will insert the element into
 * the right place after the array has been reallocated (if necessary) to
 * store n+1 elements.
 */
#include <iostream>
#include <cassert>
using namespace std;
bool bin_search(const int A[], int n, int x, int &index)
  int 1, u, m;
  if (n \le 0 \mid | x \le A[0]) { // check the first element, but only if it exists
      index = 0;
      return false;
  if (A[n-1] < x)
      index = n;
      return false;
  if (x == A[n-1]) {
      index = n-1;
      return true:
  1 = 0;
  u = n-1:
  while (1+1 < u) {
     assert(A[1] \leq x && x \leq A[u]);
     m = (1+u)/2;
     if (A[m] <= x) {
        1 = m;
      } else {
        u = m;
  if (A[1] == x) {
      index = 1;
      return true;
  } else {
      index = u;
      return false;
void insert(int A[], int n, int x, int index)
```

```
Printed by Howard Cheng
                                     binsearch.cc
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                                                                         Page 2/2
  for (i = n-1; i >= index+1; i--) {
     A[i] = A[i-1];
  A[index] = x;
int main (void)
   int A[10000];
  int n, i, x, index;
  // implements binary insertion sort, but only keeps the unique elements
  n = 0;
  while (cin >> x && n < 10000)</pre>
     if (!bin_search(A, n, x, index)) {
        insert(A, n, x, index);
     cout << "List:";
      for (i = 0; i < n; i++) {
        cout << " " << A[i];
       if (i == index) {
         cout << "*";
                             // show which one is just inserted
     cout << endl;
  return 0;
```

```
Jan 16, 14 15:08
                                         ccw.cc
                                                                          Page 1/2
 * Orientation analysis
 * Author: Howard Cheng
 * Reference:
    http://wilma.cs.brown.edu/courses/cs016/packet/node18.html
 * Given three points a, b, c, it returns whether the path from a to b to c\,
 * is counterclockwise, clockwise, or undefined.
 ^{\star} Undefined is returned if the 3 points are colinear, and c is between
 * a and b.
 */
#include <iostream>
#include <cmath>
using namespace std;
/* how close to call equal */
const double EPSILON = 1E-8;
struct Point {
 double x, y;
};
/* counterclockwise, clockwise, or undefined */
enum Orientation {CCW, CW, CNEITHER};
Orientation ccw(Point a, Point b, Point c)
 double dx1 = b.x - a.x;
 double dx2 = c.x - b.x;
 double dy1 = b.y - a.y;
 double dy2 = c.y - b.y;
double t1 = dy2 * dx1;
 double t2 = dy1 * dx2;
 if (fabs(t1 - t2) < EPSILON)</pre>
    if (dx1 * dx2 < 0 | | dy1 * dy2 < 0) {
      if (dx1*dx1 + dy1*dy1) >= dx2*dx2 + dy2*dy2 - EPSILON) {
        return CNEITHER;
      } else {
        return CW;
    } else {
      return CCW;
 } else if (t1 > t2) {
    return CCW;
 } else {
    return CW;
int main (void)
 Point a, b, c;
 Orientation res;
 while (cin >> a.x >> a.y >> b.x >> b.y >> c.x >> c.y) {
    res = ccw(a,b,c);
    if (res == CW) {
      cout << "CW" << endl;
    } else if (res == CCW)
      cout << "CCW" << endl;
    } else if (res == CNEITHER)
      cout << "CNEITHER" << endl;
    } else {
```

```
Printed by Howard Cheng
                                          ccw.cc
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                                                                             Page 2/2
     printf("Help, I am in trouble!\n");
     exit(1);
return 0;
```

```
circle 3pts.cc
 Jan 16, 14 15:08
                                                                         Page 1/1
 * Parameters of circle from 3 points
 * Author: Howard Cheng
 * Reference:
    http://www.exaflop.org/docs/cgafaq/
 ^{\star} This routine computes the parameters of a circle (center and radius)
 * from 3 points. Returns non-zero if successful, zero if the three
 * points are colinear.
 */
#include <iostream>
#include <iomanip>
#include <cmath>
#include <cassert>
using namespace std;
/* how close to call equal */
const double EPSILON = 1E-8;
struct Point {
 double x, y;
int circle(Point p1, Point p2, Point p3, Point &center, double &r)
 double a,b,c,d,e,f,g;
 a = p2.x - p1.x;
 b = p2.y - p1.y;
 c = p3.x - p1.x;
 d = p3.y - p1.y;
 e = a*(p1.x + p2.x) + b*(p1.y + p2.y);
 f = c*(p1.x + p3.x) + d*(p1.y + p3.y);
 q = 2.0*(a*(p3.y - p2.y) - b*(p3.x - p2.x));
 if (fabs(g) < EPSILON) {</pre>
   return 0;
 center.x = (d*e - b*f) / g;
 center.y = (a*f - c*e) / g;
 r = sqrt((p1.x-center.x)*(p1.x-center.x) + (p1.y-center.y)*(p1.y-center.y));
 return 1;
int main (void)
 Point a, b, c, center;
 double r;
 while (cin >> a.x >> a.y >> b.x >> b.y >> c.x >> c.y) {
   if (circle(a, b, c, center, r)) {
     cout << fixed << setprecision(3);</pre>
      cout << "center = (" << center.x << ", " << center.y << ")" << endl;</pre>
      cout << "radius = " << r << endl;
   } else {
      cout << "colinear" << endl;
 return 0;
```

common subseq.cc Jan 16, 14 15:08 Page 1/2 * Longest common subsequence * Author: Howard Cheng * Reference: http://www.ics.uci.edu/~eppstein/161/960229.html * Given two arrays A and B with sizes n and m respectively, compute the * length of the longest common subsequence. It also returns in s a longest * common subsequence (it may not be unique). One can specify which one * to choose when multiple longest common subsequences exist. * Running time and space requirement is O(mn). */ #include <iostream> #include <algorithm> #include <cassert> using namespace std; const int MAX_LEN = 20; int LCS(int A[], int n, int B[], int m, int s[]) int L[MAX_LEN+1][MAX_LEN+1]; int i, j, k; for (i = n; i >= 0; i--) { for (j = m; j >= 0; j--) { **if** (i == n | | j == m) { L[i][j] = 0;} else if (A[i] == B[j]) + L[i][j] = 1 + L[i+1][j+1];} else L[i][j] = max(L[i+1][j], L[i][j+1]);/* the following is not needed if you are not interested in the sequence */ k = 0;i = j = 0;**while** (i < n && j < m) { **if** (A[i] == B[j]) { s[k++] = A[i];i++; j++; } else if (L[i+1][j] > L[i][j+1]) { } else if (L[i+1][j] < L[i][j+1]) {</pre> } else { /* put tie-breaking conditions here */ /* eg. pick the one that starts at the first one the earliest */ j++; **return** L[0][0]; int main(void) int A[MAX_LEN], B[MAX_LEN], s[MAX_LEN]; *int* m, n, i, 1; **while** (cin >> n >> m && 1 <= n && 1 <= m &&

```
Printed by Howard Cheng
                              common subseq.cc
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                                                                     Page 2/2
       n <= MAX LEN && m <= MAX LEN)
  for (i = 0; i < n; i++) {
    cin >> A[i];
  for (i = 0; i < m; i++) {
    cin >> B[i];
  1 = LCS(A, n, B, m, s);
  for (i = 0; i < 1; i++) {
    cout << s[i] << "";
  cout << endl << "Len = " << 1 << endl;
return 0;
```

```
convex hull.cc
 Jan 16, 14 15:08
                                                                         Page 1/3
 * Convex hull
 * Author: Howard Cheng
 * Reference:
    http://wilma.cs.brown.edu/courses/cs016/packet/node25.html
 * Given a list of n (n \geq 1) points in an array, it returns the vertices of
 * the convex hull in counterclockwise order. Also returns the number of
 * vertices in the convex hull. Assumes that the hull array has been
 ^{\star} allocated to store the right number of elements (n elements is safe).
 * The points in the original polygon will be re-ordered.
 * Note: The hull contains a maximum number of points. ie. all colinear
         points and non-distinct points are included in the hull.
#include <iostream>
#include <iomanip>
#include <cmath>
#include <algorithm>
#include <cassert>
using namespace std;
/* how close to call equal */
const double EPSILON = 1E-8;
struct Point {
 double x, y;
 bool operator < (const Point &p) const {
    return y < p.y || (y == p.y && x < p.x);
};
/* counterclockwise, clockwise, or undefined */
enum Orientation {CCW, CW, CNEITHER};
/* Global point for computing convex hull */
Point start_p, max_p;
bool colinear (Point a, Point b, Point c)
 double dx1 = b.x - a.x;
 double dx2 = c.x - b.x;
 double dy1 = b.y - a.y;
 double dy2 = c.y - b.y;
 double t1 = dy2 * dx1;
 double t2 = dy1 * dx2;
 return fabs(t1 - t2) < EPSILON;
Orientation ccw(Point a, Point b, Point c)
 double dx1 = b.x - a.x;
 double dx2 = c.x - b.x;
 double dy1 = b.y - a.y;
 double dy2 = c.y - b.y;
 double t1 = dy2 * dx1;
 double t2 = dy1 * dx2;
 if (fabs(t1 - t2) < EPSILON)</pre>
    if (dx1 * dx2 < 0 | | dy1 * dy2 < 0) {
      if (dx1*dx1 + dy1*dy1) >= dx2*dx2 + dy2*dy2 - EPSILON) {
        return CNEITHER;
      } else {
```

```
convex hull.cc
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                                                                         Page 2/3
        return CW;
    } else {
      return CCW;
  } else if (t1 > t2) {
    return CCW;
  } else {
    return CW;
bool ccw_cmp(const Point &a, const Point &b)
  return ccw(start_p, a, b) == CCW;
bool sort_cmp(const Point &a, const Point &b)
  if (colinear(start_p, a, max_p) && colinear(start_p, b, max_p)) {
    double dx1 = abs(start_p.x - a.x);
    double dx2 = abs(start_p.x - b.x);
    double dy1 = abs(start_p.y - a.y);
    double dy2 = abs(start_p.y - b.y);
    return dx1 > dx2 \mid \mid (dx1 == dx2 \&\& dy1 > dy2);
    return ccw(start_p, a, b) == CCW;
int convex_hull(Point polygon[], int n, Point hull[]) {
 int count, best_i, i;
  sort(polygon, polygon+n);
  for (int i = n-1; i >= 1; i--) {
    if (fabs(polygon[i].x - polygon[i-1].x) < EPSILON &&</pre>
        fabs(polygon[i].y - polygon[i-1].y) < EPSILON) {</pre>
      for (int j = i; j < n-1; j++) {
        polygon[j] = polygon[j+1];
      n--;
  assert (n > 0);
  if (n == 1) {
   hull[0] = polygon[0];
    return 1:
  /* find the first point: min y, and then min x */
  best_i = min_element(polygon, polygon+n) - polygon;
  swap(polygon[0], polygon[best_i]);
  start_p = polygon[0];
  /* find the maximum angle wrt start_p and positive x-axis */
 best_i = 1;
  for (i = 2; i < n; i++) {
    if (ccw_cmp(polygon[best_i], polygon[i])) {
      best_i = i;
  max_p = polygon[best_i];
  /* get simple closed polygon */
  sort(polygon+1, polygon+n, sort_cmp);
  /* do convex hull */
```

```
convex hull.cc
                                                                           Page 3/3
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 count = 0;
 hull[count] = polygon[count]; count++;
 hull[count] = polygon[count]; count++;
 for (i = 2; i < n; i++) {
    while (count > 1 &&
          ccw(hull[count-2], hull[count-1], polygon[i]) == CW) {
      /* pop point */
      count--;
    hull[count++] = polygon[i];
 return count;
int main(void)
 Point *polygon, *hull;
 int n, hull_size;
 int i;
 while (cin >> n && n > 0) {
    polygon = new Point[n];
    hull = new Point[n];
    assert (polygon && hull);
    for (i = 0; i < n; i++) {
      cin >> polygon[i].x >> polygon[i].y;
    hull_size = convex_hull(polygon, n, hull);
    cout << "Sorted:" << endl;
    for (i = 0; i < n; i++) {
      cout << fixed << setprecision(2);</pre>
      cout << "(" << polygon[i].x << "," << polygon[i].y << ")" << endl;</pre>
    cout << endl;
    cout << "Hull size = " << hull_size << endl;</pre>
    for (i = 0; i < hull_size; i++) {</pre>
      cout << "(" << hull[i].x << "," << hull[i].y << ")" << endl;</pre>
    cout << endl;
    delete[] polygon;
delete[] hull;
 return 0;
```

```
Jan 16, 14 15:08
                                        cra.cc
                                                                        Page 1/2
 * Chinese Remainder Theorem
 * Author: Howard Cheng
 * Reference:
    Geddes, K.O., Czapor, S.R., and Labahn, G. Algorithms for Computer
    Algebra, Kluwer Academic Publishers, 1992, p. 180
 * Given n relatively prime modular in m[0], ..., m[n-1], and right-hand
 * sides a[0], ..., a[n-1], the routine solves for the unique solution
 * in the range 0 <= x < m[0]*m[1]*...*m[n-1] such that x = a[i] mod m[i]
 * for all 0 \le i \le n. The algorithm used is Garner's algorithm, which
 * is not the same as the one usually used in number theory textbooks.
* It is assumed that m[i] are positive and pairwise relatively prime.
* a[i] can be any integer.
 */
#include <iostream>
#include <cassert>
using namespace std;
int gcd(int a, int b, int &s, int &t)
 int r, r1, r2, a1, a2, b1, b2, q;
 int A = a;
 int B = b;
 a1 = b2 = 1;
 a2 = b1 = 0;
 while (b) {
   assert (a1*A + a2*B == a);
    q = a / b;
   r = a % b;
   r1 = a1 - q*b1;
   r2 = a2 - q*b2;
    a = b;
    a1 = b1;
   a2 = b2;
   b = r;
   b1 = r1;
   b2 = r2;
 s = a1;
 t = a2;
 assert (a >= 0);
 return a;
int cra(int n, int m[], int a[])
 int x, i, k, prod, temp;
 int *gamma, *v;
 gamma = new int[n];
 v = new int[n];
 assert (gamma && v);
  /* compute inverses */
 for (k = 1; k < n; k++) {
   prod = m[0] % m[k];
    for (i = 1; i < k; i++) {
     prod = (prod * m[i]) % m[k];
    gcd(prod, m[k], gamma[k], temp);
```

```
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                                         cra.cc
                                                                          Page 2/2
    gamma[k] %= m[k];
    if (gamma[k] < 0) {
     gamma[k] += m[k];
 /* compute coefficients */
 v[0] = a[0];
 for (k = 1; k < n; k++) {
   temp = v[k-1];
    for (i = k-2; i >= 0; i--) {
     temp = (temp * m[i] + v[i]) % m[k];
     if (temp < 0) {
        temp += m[k];
    v[k] = ((a[k] - temp) * gamma[k]) % m[k];
   if (v[k] < 0) {
     v[k] += m[k];
 /* convert from mixed-radix representation */
 x = v[n-1];
 for (k = n-2; k \ge 0; k--) {
   x = x * m[k] + v[k];
 delete[] gamma;
 delete[] v;
 return x;
int main (void)
 int n, *m, *a, i, x;
 while (cin >> n && n > 0) {
   m = new int[n];
   a = new int[n];
   assert (m && a);
   cout << "Enter moduli:" << endl;</pre>
   for (i = 0; i < n; i++) {
    cin >> m[i];
    cout << "Enter right-hand side:" << endl;
   for (i = 0; i < n; i++) {
     cin >> a[i];
   x = cra(n, m, a);
   cout << "x = " << x << endl;
    for (i = 0; i < n; i++) {
     assert ((x-a[i]) % m[i] == 0);
   delete[] m;
   delete[] a;
 return 0;
```

```
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                                        date.cc
                                                                        Page 1/3
// Date class
// This is an implementation of some common functionalities for dates.
// It can represent dates from Jan 1, 1753 to after (dates before that
// time are complicated...).
#include <iostream>
#include <string>
#include <utility>
#include <iomanip>
#include <cctype>
using namespace std;
using namespace std::rel_ops;
struct Date {
 int yyyy;
 int mm;
 int dd;
 // no dates before 1753
 static int const BASE YEAR = 1753;
 // Enumerated type for names of the days of the week
 enum dayName {SUN, MON, TUE, WED, THU, FRI, SAT};
  // Is a date valid
 static bool validDate(int yr, int mon, int day)
    return yr >= BASE_YEAR && mon >= 1 && mon <= 12 &&
     day > 0 && day <= daysIn(mon, yr);</pre>
 bool isValid() const
    return validDate(yyyy, mm, dd);
  // Constructor to create a specific date. If the date is invalid,
  // the behaviour is undefined
 Date(int yr = 1970, int mon = 1, int day = 1)
   yyyy = yr;
   mm = mon;
   dd = day;
  // Returns the day of the week for this
 dayName dayOfWeek() const
    int a = (14 - mm) / 12;
   int y = yyyy - a;
    int m = mm + 12 * a - 2;
    return (dayName) ((dd + y + y/4 - y/100 + y/400 + 31 * m / 12) % 7);
  // comparison operators
 bool operator == (const Date &d) const
    return dd == d.dd && mm == d.mm && yyyy == d.yyyy;
 bool operator<(const Date &d) const
    return yyyy < d.yyyy || (yyyy == d.yyyy && mm < d.mm) ||
      (yyyy == d.yyyy && mm == d.mm && dd < d.dd);
```

```
date.cc
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                                                                      Page 2/3
// Returns true if yr is a leap year
static bool leapYear(int v)
  return (y % 400 ==0 || (y % 4 == 0 && y % 100 != 0));
 // number of days in this month
static int daysIn(int m, int y)
  switch (m) {
  case 4 :
  case 6 :
  case 9 :
  case 11 :
    return 30;
  case 2 :
    if (leapYear(y)) {
      return 29;
    else {
      return 28;
  default :
    return 31;
 // increment by day, month, or year
// Use negative argument to decrement
// If adding a month/year results in a date before BASE_YEAR, the result
// is undefined.
// If adding a month/year results in an invalid date (Feb 29 on a non-leap
// year, Feb 31, Jun 31, etc.), the results are automatically "rounded down"
 // to the last valid date
 // add n days to the date: complexity is about n/30 iterations
void addDay(int n = 1)
  dd += n;
  while (dd > daysIn(mm,yyyy)) {
    dd -= daysIn(mm, yyyy);
    if (++mm > 12) {
      mm = 1;
      уууу++;
  while (dd < 1) {
    if (--mm < 1) {
      mm = 12;
      уууу--;
    dd += daysIn(mm, yyyy);
// add n months to the date: complexity is about n/12 iterations
void addMonth(int n = 1)
  mm += n;
  while (mm > 12) {
    mm -= 12:
    уууу++;
```

```
date.cc
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                                                                         Page 3/3
    while (mm < 1)
     mm += 12;
     уууу--;
    if (dd > daysIn(mm, yyyy)) {
      dd = daysIn(mm, yyyy);
  // add n years to the date
  void addYear(int n = 1)
    yyyy += n;
   if (!leapYear(yyyy) && mm == 2 && dd == 29) {
     dd = 28;
  // number of days since 1753/01/01, including the current date
 int daysFromStart() const
    int c = 0;
    Date d(BASE_YEAR, 1, 1);
    Date d2(d);
    d2.addYear(1);
    while (d2 < *this)
     c += leapYear(d.yyyy) ? 366 : 365;
     d = d2;
      d2.addYear(1);
    d2 = d;
    d2.addMonth(1);
    while (d2 < *this) {</pre>
     c += daysIn(d.mm, d.yyyy);
     d = d2;
     d2.addMonth(1);
    while (d <= *this) {</pre>
     d.addDay();
    return c;
};
// Reads a date in yyyy/mm/dd format, assumes date is valid and in the
// right format
istream & operator >> (istream & is, Date &d)
 return is >> d.yyyy >> c >> d.mm >> c >> d.dd;
// print date in yyyy/mm/dd format
ostream& operator<< (ostream &os, const Date &d) {
 char t = os.fill('0');
 os << d.yyyy << '/' << setw(2) << d.mm << '/' << setw(2) << d.dd;
 os.fill(t);
 return os;
```

```
dijkstra.cc
 Jan 16, 14 15:08
                                                                         Page 1/3
 * Dijkstra's Algorithm
 * Author: Howard Cheng
 * Reference:
    Ian Parberry's "Problems on Algorithms", page 102.
 * Given a weight matrix representing a graph and a source vertex, this
 * algorithm computes the shortest distance, as well as path, to each
 * of the other vertices. The paths are represented by an inverted list,
 ^{\star} such that if v preceeds immediately before w in a path from the
 * source to vertex w, then the path P[w] is v. The distances from
 * the source to v is given in D[v] (DISCONNECT if not connected).
 * Call get_path to recover the path.
 * Note: Dijkstra's algorithm only works if all weight edges are
         non-negative.
 */
#include <iostream>
#include <algorithm>
#include <cassert>
using namespace std;
const int MAX NODES = 10;
const int DISCONNECT = -1;
/* assume that D and P have been allocated */
void dijkstra(int graph[MAX_NODES][MAX_NODES], int n, int src, int D[],
              int P[])
 char used[MAX NODES];
 int fringe[MAX_NODES];
 int f_size;
 int v, w, j, wj;
 int best, best_init;
  f_size = 0;
 \overline{\text{for}} (v = 0; v < n; v++) {
   if (graph[src][v] != DISCONNECT && src != v) {
     D[v] = graph[src][v];
     P[v] = src;
      fringe[f_size++] = v;
      used[v] = 1;
    } else {
     D[v] = DISCONNECT;
     P[v] = -1;
     used[v] = 0;
 D[src] = 0;
 P[src] = -1;
 used[src] = 1;
 best_init = 1;
 while (best_init)
    /* find unused vertex with smallest D */
    best_init = 0;
    for (j = 0; j < f_size; j++) {</pre>
      v = fringe[j];
      assert(D[v] != DISCONNECT);
      if (!best_init || D[v] < best) {</pre>
       best = D[v];
       w = v;
       w j = j;
       best_init = 1;
```

```
dijkstra.cc
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                                                                         Page 2/3
    if (best init) {
      assert (D[w] != DISCONNECT);
      assert(fringe[wj] == w);
      /* get rid of w from fringe */
      f_size--;
      for (j = wj; j < f_size; j++) {</pre>
       fringe[j] = fringe[j+1];
      /* update distances and add new vertices to fringe */
      for (v = 0; v < n; v++) {
        if (v != src && graph[w][v] != DISCONNECT)
          if (D[v] == DISCONNECT \mid\mid D[w] + graph[w][v] < D[v]) 
            D[v] = D[w] + graph[w][v];
            P[v] = w;
          } else if (D[w] + graph[w][v] == D[v]) {
            /* put tie-breaker here */
          if (!used[v]) {
            used[v] = 1;
            fringe[f_size++] = v;
 D[src] = 0;
int get_path(int v, int P[], int path[])
 int A[MAX_NODES];
 int i, k;
  k = 0;
 A[k++] = v;
 while (P[v] != -1) {
   v = P[v];
   A[k++] = v;
 for (i = k-1; i >= 0; i--) {
   path[k-1-i] = A[i];
 return k;
int main (void)
 int m, w, num;
  int i, j;
  int graph[MAX_NODES][MAX_NODES];
  int P[MAX_NODES][MAX_NODES], D[MAX_NODES][MAX_NODES];
 int path[MAX_NODES];
  /* clear graph */
 for (i = 0; i < MAX_NODES; i++) {
    for (j = 0; j < MAX_NODES; j++) {</pre>
      graph[i][j] = DISCONNECT;
  /* read graph */
  cin >> i >> j >> w;
 while (!(i = -1 \&\& j = -1)) {
    assert(0 <= i && i < MAX_NODES && 0 <= j && j < MAX_NODES);
```

```
dijkstra_sparse.cc
 Jan 16, 14 15:08
                                                                        Page 1/3
 * Dijkstra's Algorithm for sparse graphs
 * Author: Howard Cheng
 * Given a weight matrix representing a graph and a source vertex, this
 * algorithm computes the shortest distance, as well as path, to each
 * of the other vertices. The paths are represented by an inverted list,
 ^{\star} such that if v preceeds immediately before w in a path from the
 * source to vertex w, then the path P[w] is v. The distances from
 * the source to v is given in D[v] (-1 if not connected).
 * Call get_path to recover the path.
 ^{\star} Note: Dijkstra's algorithm only works if all weight edges are
        non-negative.
 * This version works well if the graph is not dense. The complexity
 * is O((n + m) \log (n + m)) where n is the number of vertices and
* m is the number of edges.
 */
#include <iostream>
#include <algorithm>
#include <vector>
#include <cassert>
#include <queue>
using namespace std;
struct Edge {
 int to;
                    // can be double or other numeric type
 int weight;
 Edge(int t, int w)
   : to(t), weight(w) { }
};
typedef vector<Edge>::iterator EdgeIter;
struct Graph {
 vector<Edge> *nbr;
 int num_nodes;
 Graph (int n)
   : num_nodes(n)
    nbr = new vector<Edge>[num_nodes];
 ~Graph()
    delete[] nbr;
 // note: There is no check on duplicate edge, so it is possible to
 // add multiple edges between two vertices
 // If this is an undirected graph, be sure to add an edge both
 void add_edge(int u, int v, int weight)
    nbr[u].push_back(Edge(v, weight));
};
/* assume that D and P have been allocated */
void dijkstra(const Graph &G, int src, vector<int> &D, vector<int> &P)
```

```
dijkstra sparse.cc
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                                                                         Page 2/3
 typedef pair<int, int> pii;
 int n = G.num_nodes;
 vector<bool> used(n, false);
 priority_queue<pii, vector<pii>, greater<pii> > fringe;
 D.resize(n);
 P.resize(n);
  fill(D.begin(), D.end(), -1);
 fill(P.begin(), P.end(), -1);
 D[src] = 0;
 fringe.push(make_pair(D[src], src));
 while (!fringe.empty())
   pii next = fringe.top();
   fringe.pop();
   int u = next.second;
   if (used[u]) continue;
   used[u] = true;
   for (EdgeIter it = G.nbr[u].begin(); it != G.nbr[u].end(); ++it) {
      int v = it -> to;
     int weight = it->weight + next.first;
     if (used[v]) continue;
      if (D[v] == -1 \mid | weight < D[v]) +
       D[v] = weight;
       P[v] = u;
        fringe.push(make_pair(D[v], v));
int get_path(int v, const vector<int> &P, vector<int> &path)
 path.clear();
 path.push back(v);
 while (P[v] != -1) {
   v = P[v];
   path.push_back(v);
 reverse(path.begin(), path.end());
 return path.size();
int main(void)
 int n:
 while (cin >> n && n > 0) {
   Graph G(n);
   int u, v, w;
   while (cin >> u >> v >> w && !(u == -1 && v == -1 && w == -1)) {
     G.add_edge(u, v, w);
   while (cin >> u >> v && !(u == -1 && v == -1)) {
     vector<int> D, P, path;
     dijkstra(G, u, D, P);
     get_path(v, P, path);
     cout << "distance = " << D[v] << endl;
     cout << "path = ";
     for (unsigned int i = 0; i < path.size(); i++) {</pre>
       cout << path[i] << '';
     cout << endl;
```

Jan 16, 14 15:08	dijkstra_sparse.cc	Page 3/3
} return 0;		
}		

```
diophantine sys.c
 Jan 16, 14 15:11
                                                                         Page 1/3
 * Solution of system of linear diophantine equations
 * Author: Howard Cheng
 * Date: Nov 25, 2000
 * Reference:
 * http://scicomp.ewha.ac.kr/netlib/tomspdf/
 * Look at Algorithms 287 (sort of) and 288.
  Given a system of m linear diophantine equations in n unknowns,
 * this algorithm finds a particular solution as well as a basis for
 * the solution space of the homogeneous system, if they exist. The
 * system is represented in matrix form as Ax = b where all entries
 * are integers.
 ^{\star} Function: diophantine_linsolve
 * Input:
 * A: an m x n matrix specifying the coefficients of each equation in
    each row (it is okay to have zero rows, or even have A = 0)
 ^{\star} b: an m-dimensional vector specifying the right-hand side of the system
 * m: number of equations in the system
 * n: number of unknowns in the system
 * Output:
 ^{\star} The function returns the dimension of the solution space of the
 * homogeneous system Ax = 0 (hom_dim) if it has a solution.
 * Otherwise, it returns -1.
 * Other results returned in the parameters are:
  xp: an n-dimensional vector giving a particular solution
  hom_basis: an n x n matrix whose first hom_dim columns form a basis
              of the solution space of the homogeneous system Ax = 0
 * All solutions to Ax = b can be obtained by adding integer multiples
 * of the first hom_dim columns of hom_basis to xp.
 * Note:
 ^{\star} The contents of A and b are not changed by this function.
#include <stdio.h>
#include <stdlib.h>
#define MAX N 50
#define MAX_M 50
int triangulate(int A[MAX_N+1][MAX_M+MAX_N+1], int m, int n, int cols)
 int ri, ci, i, j, k, pi, t;
 div_t d;
 ri = ci = 0;
 while (ri < m && ci < cols) {
    /* find smallest non-zero pivot */
    pi = -1:
    for (i = ri; i < m; i++) {
      if (A[i][ci] \&\& (pi == -1 || abs(A[i][ci]) < abs(A[pi][ci]))) {
       pi = i;
```

```
diophantine sys.c
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                                                                         Page 2/3
   if (pi == -1)
      /* the entire column is 0, skip it */
   } else {
     k = 0;
     for (i = ri; i < m; i++) {</pre>
       if (i != pi)
         d = div(A[i][ci], A[pi][ci]);
         if (d.quot) {
           k++;
           for (j = ci; j < n; j++)
              A[i][j] -= d.quot*A[pi][j];
     if (!k)
       ^{\prime *} swap the row to make it triangular...Alg 287 also switches the ^{*\prime}
       /* sign, probably to preserve the sign of the minors. I don't
       /* think this is necessary for our purpose.
       for (i = ci; i < n && ri != pi; i++) {
         t = A[ri][i];
         A[ri][i] = A[pi][i];
         A[pi][i] = t;
       ri++;
       ci++;
 return ri;
int diophantine_linsolve(int A[MAX_M][MAX_N], int b[MAX_M], int m, int n,
                         int xp[MAX_N], int hom_basis[MAX_N][MAX_N])
 int mat[MAX_N+1][MAX_M+MAX_N+1];
 int i, j, rank, d;
  /st form the work matrix st/
 for (i = 0; i < m; i++) {
   mat[0][i] = -b[i];
 for (i = 0; i < m; i++) {
   for (j = 0; j < n; j++) {
     mat[j+1][i] = A[i][j];
 for (i = 0; i < n+1; i++) {
   for (j = 0; j < n+1; j++) {
     mat[i][j+m] = (i == j);
 /* triangluate the first n+1 x m+1 submatrix */
 rank = triangulate(mat, n+1, m+n+1, m+1);
 d = mat[rank-1][m];
  /* check for no solutions */
 if (d != 1 && d != −1) {
   /* no integer solutions */
   return -1;
  /* check for inconsistent system */
 for (i = 0; i < m; i++) {
   if (mat[rank-1][i]) {
     return -1;
```

```
diophantine_sys.c
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                                                                                   Page 3/3
  /* there is a solution, copy it to the result */
 for (i = 0; i < n; i++) {
    xp[i] = d*mat[rank-1][m+1+i];
    for (j = 0; j < n+1-rank; j++) {
  hom_basis[i][j] = mat[rank+j][m+1+i];</pre>
 return n+1-rank;
int main (void)
 int A[MAX_M][MAX_N], b[MAX_M], m, n, xp[MAX_N], hom_basis[MAX_N][MAX_N];
 int i, j, hom_dim;
 while (scanf("%d%d", &m, &n) == 2 && m > 0 && n > 0) {
    for (i = 0; i < m; i++) {
      printf("Enter equation %d:\n", i+1);
      for (j = 0; j < n; j++) {
  scanf("%d", &A[i][j]);</pre>
      scanf("%d", &b[i]);
    if ((hom_dim = diophantine_linsolve(A, b, m, n, xp, hom_basis)) >= 0) {
      printf("Particular solution:\n");
      for (i = 0; i < n; i++) {
        printf("%d", xp[i]);
      printf("\n");
      printf("hom_dim = %d\n", hom_dim);
      printf("Basis for Ax = 0:\n");
      for (j = 0; j < hom_dim; j++) {
  for (i = 0; i < n; i++) {</pre>
           printf("%d", hom_basis[i][j]);
        printf("\n");
    } else {
      printf("No solution.\n");
 return 0;
```

```
dist3D.cc
 Jan 16, 14 15:11
                                                                        Page 1/5
// 3-D distances between point to point, point to line segment,
// line segment to line segment, and point to triangle.
// There are corresponding versions of the same code for distances
// between point to infinite lines, infinite line to infinite line,
// and point to infinite plane.
// It is assumed that segments/lines/triangles/plane are defined by
// distinct points (so the objects are not degenerate).
// They can be used for 2-D objects as well by setting the z coordinates
// to 0.
//
// Author: Howard Cheng
#include <iostream>
#include <iomanip>
#include <cmath>
#include <algorithm>
#include <cassert>
using namespace std;
const double PI = acos(-1.0);
struct Vector {
 double x, y, z;
 Vector (double xx = 0, double yy = 0, double zz = 0)
   : x(xx), y(yy), z(zz) \{ \}
 Vector (const Vector &p1, const Vector &p2)
   : x(p2.x - p1.x), y(p2.y - p1.y), z(p2.z - p1.z) { }
 Vector (const Vector &p1, const Vector &p2, double t)
   : x(p1.x + t*p2.x), y(p1.y + t*p2.y), z(p1.z + t*p2.z) { }
 double norm() const
    return sqrt(x*x + y*y + z*z);
};
istream & operator >> (istream & is, Vector &p)
 return is >> p.x >> p.y >> p.z;
ostream & operator << (ostream & os, const Vector &p)
 return os << "(" << p.x << "," << p.y << "," << p.z << ")";
double dot (const Vector &p1, const Vector &p2)
 return p1.x * p2.x + p1.y * p2.y + p1.z * p2.z;
Vector cross (const Vector &p1, const Vector &p2)
 Vector v(p1.y*p2.z - p2.y*p1.z,
           p2.x*p1.z - p1.x*p2.z,
           p1.x*p2.y - p2.x*p1.y);
 return v;
// distance between two points
```

```
dist3D.cc
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                                                                        Page 2/5
double dist point to point (const Vector &p1, const Vector &p2)
  Vector p(p1, p2);
  return p.norm();
// angle between two vectors (in radians)
double angle (const Vector &p1, const Vector &p2)
 return acos(dot(p1, p2)/p1.norm()/p2.norm());
// distance from p to the line segment from a to b
double dist_point_to_segment(const Vector &p, const Vector &a,
                             const Vector &b)
 Vector u(a, p), v(a, b);
 double s = dot(u,v) / dot(v,v);
 if (s < 0 || s > 1) {
   return min(dist_point_to_point(p, a), dist_point_to_point(p, b));
 Vector proj(a, v, s);
 return dist_point_to_point(proj, p);
// distance from p to the infinite line defined by a and b
double dist_point_to_line(const Vector &p, const Vector &a,
                          const Vector &b)
 Vector u(a, p), v(a, b);
  double s = dot(u,v) / dot(v,v);
 Vector proj(a, v, s);
 return dist_point_to_point(proj, p);
// distance from p to the triangle defined by a, b, c
double dist_point_to_triangle(const Vector &p, const Vector &a,
                              const Vector &b, const Vector &c)
 Vector u(a, p), v1(a, b), v2(a, c);
 Vector normal = cross(v1, v2);
  double s = dot(u, normal) / (normal.norm() * normal.norm());
 Vector proj(p, normal, -s);
  // check projection: inside if sum of angles is 2*pi
 Vector wa(proj, a), wb(proj, b), wc(proj, c);
  double a1 = angle(wa, wb);
  double a2 = angle(wa, wc);
  double a3 = angle(wb, wc);
  if (fabs(a1 + a2 + a3 - 2*PI) < 1e-8) {
    return dist_point_to_point(proj, p);
  } else {
   return min(dist_point_to_segment(p, a, b),
               min(dist_point_to_segment(p, a, c),
                   dist_point_to_segment(p, b, c)));
// distance from p to the infinite plane defined by a, b, c
double dist_point_to_plane(const Vector &p, const Vector &a,
                           const Vector &b, const Vector &c)
 Vector u(a, p), v1(a, b), v2(a, c);
 Vector normal = cross(v1, v2);
  double s = dot(u, normal) / (normal.norm() * normal.norm());
 Vector proj(p, normal, -s);
 return dist_point_to_point(proj, p);
```

```
dist3D.cc
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                                                                         Page 3/5
// distance from segment p1->q1 to p2->q2
double dist_segment_to_segment (const Vector &p1, const Vector &q1,
                               const Vector &p2, const Vector &q2)
 //
  // the points on the 1st line are p1 + t * v1
 // the points on the 2nd line are p2 + s * v2
 //
                                 0 <= s, t <= 1
 //
 // squared distance is
 //S = (p1.x - p2.x + t * v1.x - s * v2.x)^2 +
        (p1.y - p2.y + t * v1.y - s * v2.y)^2 + (p1.z - p2.z + t * v1.z - s * v2.z)^2
 //
 11
 //
 // deriviative wrt t and s are:
  // 1/2 \, ds/dt = norm(v1)^2 * t - dot(v1, v2) * s + dot(v1, p1) - dot(v1, p2)
 // 1/2 dS/ds = -dot(v1, v2) * t + norm(v2)^2 * s - dot(v2, p1) + dot(v2, p2)
 // solving for s and t with both derivatives = 0:
 Vector v1(p1, q1), v2(p2, q2);
 Vector rhs(dot(v1, p2) - dot(v1, p1), dot(v2, p1) - dot(v2, p2));
 double det = v1.norm()*v1.norm()*v2.norm()*v2.norm() -
   dot(v1, v2)*dot(v1, v2);
 if (det < 1e-8) {
    // parallel lines (if v1 and v2 != 0)
    goto degenerate;
  } else {
    double t = (rhs.x*v2.norm()*v2.norm() + rhs.y*dot(v1, v2)) / det;
    double s = (v1.norm()*v1.norm()*rhs.y + dot(v1, v2) * rhs.x) / det;
    if (0 <= s && s <= 1 && 0 <= t && t <= 1) {
      Vector pp1(p1, v1, t), pp2(p2, v2, s);
      return dist_point_to_point(pp1, pp2);
 degenerate:
 return min(min(dist_point_to_segment(p1, p2, q2),
                 dist_point_to_segment(q1, p2, q2)),
             min(dist_point_to_segment(p2, p1, q1),
                 dist_point_to_segment(q2, p1, q1)));
// distance from infinite lines defined by p1->q1 and p2->q2
double dist_line_to_line(const Vector &p1, const Vector &q1,
                         const Vector &p2, const Vector &q2)
 //
 // the points on the 1st line are p1 + t * v1
 // the points on the 2nd line are p2 + s * v2
 //
 //
                                 0 <= s, t <= 1
 // squared distance is
  //S = (p1.x - p2.x + t * v1.x - s * v2.x)^2 +
         (p1.y - p2.y + t * v1.y - s * v2.y)^2 +
  //
         (p1.z - p2.z + t * v1.z - s * v2.z)^2
 // deriviative wrt t and s are:
  // 1/2 dS/dt = norm(v1)^2 * t - dot(v1, v2) * s + dot(v1, p1) - dot(v1, p2)
```

```
dist3D.cc
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                                                                 Page 4/5
 // 1/2 dS/ds = -dot(v1, v2) * t + norm(v2)^2 * s - dot(v2, p1) + dot(v2, p2)
 // solving for s and t with both derivatives = 0:
 Vector v1(p1, q1), v2(p2, q2);
 Vector rhs(dot(v1, p2) - dot(v1, p1), dot(v2, p1) - dot(v2, p2));
 double det = v1.norm()*v1.norm()*v2.norm()*v2.norm() -
   dot(v1, v2)*dot(v1, v2);
 if (det < 1e-8) {
   // parallel lines (if v1 and v2 != 0)
   return dist_point_to_line(p1, p2, q2);
   double t = (rhs.x*v2.norm()*v2.norm() + rhs.y* dot(v1, v2)) / det;
   double s = (v1.norm()*v1.norm()*rhs.y + dot(v1, v2) * rhs.x) / det;
   Vector pp1(p1, v1, t), pp2(p2, v2, s);
   return dist_point_to_point(pp1, pp2);
// This is the solution to 11836 (Star War)
void do case()
 Vector t1[4], t2[4];
 for (int i = 0; i < 4; i++) {
   cin >> t1[i];
 for (int i = 0; i < 4; i++) {
   cin >> t2[i];
 double best = dist_point_to_point(t1[0], t2[0]);
  // vertex-face distance
 for (int i1 = 0; i1 < 4; i1++) {</pre>
   for (int j1 = 0; j1 < 4; j1++) {
     best = min(best, dist_point_to_triangle(t1[i1], t2[j1], t2[(j1+1)%4],
                                          t2[(j1+2)%4]));
     best = min(best, dist_point_to_triangle(t2[i1], t1[j1], t1[(j1+1)%4],
                                          t1[(j1+2)%4]));
  // edge-edge distance
 for (int i1 = 0; i1 < 4; i1++) {
   for (int i2 = i1+1; i2 < 4; i2++) {
     for (int j1 = 0; j1 < 4; j1++) {
       for (int j2 = j1+1; j2 < 4; j2++) {
        best = min(best, dist_segment_to_segment(t1[i1], t1[i2],
                                               t2[j1], t2[j2]));
 cout << setprecision(2) << fixed << best << endl;</pre>
int main (void)
 int T;
 cin >> T;
 while (T-- > 0) {
```

```
dist3D.cc
                                                                    Page 5/5
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  do_case();
return 0;
```

```
dist line.cc
 Jan 16, 14 15:11
                                                                           Page 1/1
 ^{\star} Distance from a point to a line.
 * Author: Howard Cheng
 * Reference:
    http://www.exaflop.org/docs/cgafaq/cgal.html
 \mbox{\scriptsize {\tt *}} This routine computes the shortest distance from a point to a line.
 * ie. distance from point to its orthogonal projection onto the line.
 * Works even if the projection is not on the line.
 */
#include <iostream>
#include <iomanip>
#include <cmath>
#include <cassert>
using namespace std;
struct Point {
 double x, y;
/* computes the distance from "c" to the line defined by "a" and "b" */
double dist_line(Point a, Point b, Point c)
 double L2, s;
 L2 = (b.x-a.x)*(b.x-a.x)+(b.y-a.y)*(b.y-a.y);
 assert (L2 > 0);
 s = ((a.y-c.y)*(b.x-a.x)-(a.x-c.x)*(b.y-a.y)) / L2;
 return fabs(s*sqrt(L2));
int main(void)
 Point a, b, c;
 while (cin >> a.x >> a.y >> b.x >> b.y >> c.x >> c.y) {
    cout << "distance = " << fixed << setprecision(2) << dist_line(a, b, c)</pre>
         << endl;
 return 0;
```

```
dow.c
 Jan 16, 14 15:11
                                                                         Page 1/1
 ^{\star} Computing the Day of the Week
 * Author: Howard Cheng
 * This routine computes the day of the week (Sunday = 0, Saturday = 6)
 * from the year, month, and day.
unsigned DOW(unsigned y, unsigned m, unsigned d)
      if (m < 3)
            m += 13;
            y--;
      else m++;
      return (d + 26 * m / 10 + y + y / 4 - y / 100 + y / 400 + 6) % 7;
#include <stdio.h>
#include <stdlib.h>
int main(int argc, char *argv[])
      int Day;
      void usage(void);
      unsigned d, m, y, days[] = {31, 29, 31, 30, 31, 30,
                                   31, 31, 30, 31, 30, 31};
      char *day[2][7] = {
            ("Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"}, {"Sun", "Mon", "Tue", "Wed", "Thu", "Fri", "Sat"}
      if (4 > argc)
            usage();
      y = atoi(argv[1]);
      \bar{m} = atoi(argv[2]);
      d = atoi(argv[3]);
      if (!m || m > 12)
            usage();
      if (!d || d > days[m - 1])
            usage();
      if (y < 100)
            y += 1900;
      Day = DOW(y, m, d);
      printf ("DOW returned %d, so %d %s %d is a %s\n",
             Day, d, month[m - 1], y, day[6 - 5][Day]);
      return 0;
void usage (void)
      puts ("Usage: DOW yy[yy] mm dd");
      exit(-1);
```

```
euclid.cc
 Jan 16, 14 15:11
                                                                         Page 1/1
 * Euclidean Algorithm
 * Author: Howard Cheng
 ^{\star} Given two integers, return their gcd.
#include <iostream>
#include <cassert>
using namespace std;
int gcd(int a, int b)
 int r;
  /* unnecessary if a, b >= 0 */
 if (a < 0) {
   a = -a;
 if (b < 0) {
   b = -b;
 while (b) {
   r = a % b;
   a = b;
   b = r;
 assert (a >= 0);
 return a;
int main(void)
 int a, b;
 while (cin >> a >> b) {
   cout << gcd(a, b) << endl;</pre>
 return 0;
```

```
eulerphi.cc
 Jan 16, 14 15:11
                                                                        Page 1/2
 * Euler's Phi function:
 * Author: Ethan Kim
 * Complexity: O(sqrt(n))
 * Computes Euler's Phi(Totient) function; Given a positive n, computes
 * the number of positive integers that are <= n and relatively prime to n.
 * For prime n, it is easy to see that phi(n)=n-1.
 * For powers of prime, phi(p^k)=p^k(k-1) * (p-1).
 * Also, phi is multiplicative, so phi(pq)=phi(p)*phi(q), if p and q are
 * relatively prime.
 */
#include <iostream>
#include <cassert>
using namespace std;
int fast_exp(int b, int n)
 int res = 1;
 int x = b;
 while (n > 0) {
   if (n & 0x01) {
     n--;
     res *= x;
    } else {
     n >>= 1;
     x *= x;
 return res;
int phi(int n) {
 int k, res;
 long long p;
 assert (n > 0);
 res=1;
 for (k = 0; n % 2 == 0; k++) {
   n /= 2;
 if (k)
    res *= fast_exp(2, k-1);
 for (p = 3; p*p \le n; p += 2) {
   for (k = 0; n % p == 0; k++) {
     n /= p;
    if (k) {
     res *= fast_exp(p, k-1) * (p-1);
 if (n > 1) {
    res *= n-1;
 return res;
int main(void) {
 while(cin >> p && p) {
```

```
Printed by Howard Cheng
                                    eulerphi.cc
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                                                                       Page 2/2
  cout << phi(p) << endl;
return 0;
```

```
eulertour.cc
 Jan 16, 14 15:11
                                                                         Page 1/4
 * Finding an Eulerian Tour
 * Author: Howard Cheng
 * The routine eulerian() determines if a graph has an Eulerian tour.
 * That is, it checks that it is connected and all vertices have even
 * degree. We assume that the graph is represented as an adjacency matrix
 * and the an auxillary array called "deg" gives the degree of the vertex.
 * The routine eulerian_tour() returns one (arbitrary) Eulerian tour.
 * The tour is stored in an array of the vertices visited in the tour,
 * and the first and last vertex is the same.
 * WARNING: eulerian_tour() destroys the graph as it uses edges. If
           you need the graph back then you should save a copy.
 * NOTE: converting this code for directed graphs should not be that much
         work. You should also be able to convert this code for Eulerian
 */
#include <iostream>
#include <algorithm>
#include <cassert>
using namespace std;
const int NUM_VERTICES = 50;
const int NUM_EDGES = 1000;
                                  /* maximum number of edges in graph */
int graph[NUM_VERTICES+1][NUM_VERTICES+1];
int deg[NUM_VERTICES+1];
void clear_graph(void)
 fill(deg, deg+NUM VERTICES+1, 0);
 for (int i = 1; i <= NUM_VERTICES; i++) {</pre>
   fill(graph[i], graph[i]+NUM_VERTICES+1, 0);
void visit(int v, char visited[])
 int w:
 visited[v] = 1;
 for (w = 1; w <= NUM_VERTICES; w++) {</pre>
   if (!visited[w] && graph[v][w] > 0) {
      visit(w, visited);
 }
int connected (void)
 char visited[NUM_VERTICES+1];
 int i;
 fill(visited, visited+NUM_VERTICES+1, 0);
 for (i = 1; i <= NUM_VERTICES; i++) {</pre>
    if (deg[i] > 0) {
      visit(i, visited);
      break;
 for (i = 1; i <= NUM_VERTICES; i++) {</pre>
    if (deg[i] > 0 && !visited[i]) {
```

```
eulertour.cc
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                                                                          Page 2/4
      return 0;
 return 1;
int eulerian (void)
  int i;
  for (i = 1; i <= NUM_VERTICES; i++) {</pre>
    if (deg[i] % 2 == 1) {
     return 0;
 return connected();
int find_tour(int start, int temp[])
  int len = 0;
  int next;
  temp[len++] = start;
 while (deg[start] > 0) {
    for (next = 1; next <= NUM_VERTICES; next++) {</pre>
      if (graph[start][next] > 0) {
        break;
    temp[len++] = next;
    graph[start][next]--; deg[start]--;
    graph[next][start]--; deg[next]--;
    start = next;
 return len;
int graft_tour(int old[], int old_len, int tour[], int tour_len)
  int pos[NUM VERTICES+1];
 int i, j, p1, p2;
  fill(pos, pos+NUM_VERTICES+1, -1);
  for (i = 0; i < old_len; i++) {</pre>
   pos[old[i]] = i;
 for (i = 0; i < tour_len; i++) {</pre>
    if (pos[tour[i]] >= 0) {
     break:
 assert(i < tour_len);
 p1 = pos[tour[i]];
  p2 = i;
  for (i = old_len-1; i > p1; i--) {
   old[i+tour_len-1] = old[i];
 for (i = p2+1, j = 0; i < tour_len-1; i++, j++) {</pre>
    old[p1+j+1] = tour[i];
 for (i = 0; i <= p2; i++) {
    old[p1+j+1] = tour[i];
 return old_len+tour_len-1;
int eulerian_tour(int tour[])
```

```
eulertour.cc
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                                                                          Page 3/4
 int temp[NUM_EDGES+1];
 int tour_len, temp_len, first_time;
 int i, found;
 tour_len = temp_len = 0;
 first\_time = 1;
 while (1) {
    found = 0;
    if (first_time) {
      for (i = 1; i <= NUM_VERTICES; i++) {</pre>
        if (deg[i] > 0) {
          found = 1;
          break;
    } else {
      /\ast this ensures that we can graft next tour on to existing one \ast/
      for (i = 0; i < tour_len; i++) {</pre>
        if (deg[tour[i]] > 0) {
          found = 1;
          break;
      i = tour[i];
    if (!found) {
     break;
    if (first_time) {
     tour_len = find_tour(i, tour);
     temp_len = find_tour(i, temp);
      tour_len = graft_tour(tour, tour_len, temp, temp_len);
    first_time = 0;
 return tour_len;
int main(void)
 int T, N, i, j, k;
 int u, v;
 int tour[NUM_EDGES+1], tour_len;
 cin >> T;
 for (i = 1; i <= T; i++) {
    clear_graph();
    if (i > 1) {
     cout << endl;
    cout << "Case #" << i << endl;
    cin >> N;
    for (j = 0; j < N; j++) {
     cin >> u >> v;
     graph[u][v]++;
     graph[v][u]++;
     deg[u]++;
     deg[v]++;
    if (eulerian()) {
     tour_len = eulerian_tour(tour);
      for (k = 0; k < tour_len-1; k++) {</pre>
        cout << tour[k] << " " << tour[k+1] << endl;
    } else {
```

```
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                                       eulertour.cc
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                                                                              Page 4/4
     cout << "some beads may be lost" << endl;</pre>
return 0;
```

```
Jan 16, 14 15:08
                                           exp.cc
                                                                              Page 1/1
 * Fast Exponentiation
 * Author: Howard Cheng
 * Given b and n, computes b^n quickly.
#include <iostream>
using namespace std;
int fast_exp(int b, int n)
  int res = 1;
 int x = b;
  while (n > 0) {
    if (n & 0x01) {
     n--;
res *= x;
    } else {
      n >>= 1;
      x *= x;
  return res;
int main(void)
  int b, n;
 while (cin >> b >> n) {
   cout << b << "^" << n << " = " << fast_exp(b, n) << endl;</pre>
 return 0;
```

```
expmod.cc
 Jan 16, 14 15:11
                                                                             Page 1/1
 * Fast Exponentiation mod m
 * Author: Howard Cheng
 * Given b, n, and m, computes b^n mod m quickly.
#include <iostream>
#include <cassert>
using namespace std;
int fast_exp(int b, int n, int m)
 int res = 1;
 long\ long\ x = b;
 while (n > 0) {
   if (n & 0x01) {
      res = (res * x) % m;
    } else {
     n >>= 1;
      x = (x * x) % m;
 return res;
int main(void)
 int b, n, m;
 while (cin >> b >> n >> m) {
   cout << b << "^" << n << " mod " << m << " = " << fast_exp(b, n, m)</pre>
         << endl;
 return 0;
```

```
exteuclid.cc
 Jan 16, 14 15:11
                                                                            Page 1/1
 * Extended Euclidean Algorithm
 * Author: Howard Cheng
 \mbox{\scriptsize *} Given two integers, return their gcd and the cofactors to form the
 * gcd as a linear combination.
 * a*s + b*t = gcd(a,b)
#include <iostream>
#include <cassert>
using namespace std;
int gcd(int a, int b, int &s, int &t)
 int r, r1, r2, a1, a2, b1, b2, q;
 int A = a;
 int B = b;
  /* unnecessary if a, b >= 0 */
 if (a < 0) {
   r = gcd(-a, b, s, t);

s *= -1;
    return r;
 if (b < 0) {
   r = gcd(a, -b, s, t);
   t *= -1;
    return r;
 a1 = b2 = 1;
 a2 = b1 = 0;
 while (b) {
    assert(a1*A + a2*B == a);
    q = a / b;
    r = a % b;
    r1 = a1 - q*b1;
    r2 = a2 - q*b2;
    a = b;
    a1 = b1;
    a2 = b2;
   b = r;
   b1 = r1;
    b2 = r2;
 s = a1;
 t = a2;
 assert (a >= 0);
 return a;
int main (void)
 int a, b, s, t, res;
 while (cin >> a >> b) {
   res = gcd(a, b, s, t);
cout << res << "=" << a << "*" << s << "+"
         << b << "*" << t << endl;
 return 0;
```

```
factor.cc
 Jan 16, 14 15:11
                                                                        Page 1/2
 * Prime Factorization
 * Author: Ethan Kim
 * Complexity: O(sqrt(n))
 ^{\star} Takes an integer and writes out the prime factorization in
 * ascending order. Prints -1 first, when given a negative integer.
 * Note: you can change this code to store the factors in an array or process
 * the factors in other ways.
 * Also, this code works for all integers even on INT_MIN (note that negating
 * INT_MIN does nothing, but it still works because INT_MIN is a power of 2).
 */
#include <iostream>
using namespace std;
void factor(int n) {
 int printed = 0;
 long long p;
 if (n == 0 || n == 1) {
    cout << n << endl;
    return;
 if (n < 0) {
   n *= -1;
   cout << "-1" << endl;
   printed = 1;
 while (n % 2 == 0) {
   n/=2;
    cout << "2" << endl;
    printed = 1;
 for (p = 3; p*p \le n; p += 2) {
   while (n % p == 0) {
     n /= p;
     cout << p << endl;
     printed = 1;
 if(n>1 || !printed)
    cout << n << endl;</pre>
int main(void) {
 int p;
 while(cin >> p && p != 0) {
    factor(p);
 return 0;
```

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Jan 16, 14 15:11	factor.cc	Page 2/2

```
factor large.cc
 Jan 16, 14 15:08
                                                                          Page 1/4
// Gives the prime factorization of natural numbers (Uses probability)
// Author: Darcy Best
// Date : January 7, 2010
// This should be used for factoring large integers. If you're
    dealing with are small integers (N < 2^31), this is going
     overboard. -- The normal sieve of Sieve of Eratosthenes is
    usually good even for values up to 2^40.
// This implementation should only be used if you have numbers
     larger than 2^40 (10^12) to factor.
//
// Notes:
     - You need to handle N < 2 separately.
//
      - Uses Miller-Rabin Primality Test
           - This is a probabilistic test, there is a (1/4) \(^K\)
               probability that a composite will return prime.
               (K = 10 \text{ or } 15 \text{ should be reasonably reliable}).
     - Uses Pollard's Rho algorithm to factor composites.
          - I have also added Brent's improvement
      - This program writes a number as a product of its primes,
           with normal exponents (ie. "60 = 2^2 * 3 * 5")
#include <iostream>
#include <algorithm>
#include <set>
#include <map>
#include <cmath>
#include <ctime>
#include <cstdlib>
#include <vector>
using namespace std;
typedef long long int 11;
const ll MAX_NUM = 1e16;
const 11 CB_RT = 11(pow(1.0*MAX_NUM, 1.0/3)) + 2;
vector<ll> primes;
11 numPrimes;
11 c = 2;
const 11 K = 10;
set<ll> lqPrimes;
map<11,11> semiPrimes;
ll gcd(ll a, ll b) {
 11 r;
 while (b) {
   r = a % b;
    a = b;
   b = r;
 return a;
11 mult_mod(l1 x,l1 y,l1 n) {
 11 \text{ res} = 0;
 while(y) {
    if(y % 2){
     y--;
      res += x;
      if(res >= n)
        res -= n;
    } else {
      x <<= 1;
      v >>= 1;
      if(x >= n)
        x -= n;
```

```
factor large.cc
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                                                                             Page 2/4
  return res;
11 fast_exp_mod(ll b, ll n, ll m) {
 ll res = 1;
  11 x = b;
  while (n)
    if (n % 2) {
      n--:
      res = mult mod(res,x,m);
    } else {
      n >>= 1;
      x = mult_mod(x, x, m);
  return res;
void genSmallPrimes() {
  bool isPrime[CB_RT+7];
  for (int i=3; i < CB_RT; i+=2)</pre>
    isPrime[i] = true;
  primes.clear();
  primes.push_back(2);
  int i;
  for (i=3; i*i<CB_RT; i+=2)</pre>
    if(isPrime[i]){
      primes.push_back(i);
      for (int j=i*i; j<CB_RT; j+=(2*i))</pre>
         isPrime[j] = false;
  for(;i<CB_RT;i+=2)</pre>
    if(isPrime[i])
      primes.push_back(i);
  numPrimes = primes.size();
ll F(ll x,ll n) {
 x = mult_mod(x, x, n);
 return (x < 0 ? x + n : x);
ll pollardRho(ll n){
 ll b,g,x,y,z;
 newC:
  C++;
  q = b = x = 1;
  while(q == 1){
    z = 1;
    y = x;
    for(ll i=0;i<b;i++){
      x = F(x, n);
      z = mult_mod(z, abs(x-y), n);
    g = gcd(z,n);
    b <<= 1;
  if(g == n || g == 0)
    goto newC;
  c = 2;
```

```
factor large.cc
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                                                                             Page 3/4
 return q;
bool miller(ll n) {
 11 d = n-1;
 11 s = 0, a, x;
 while(d % 2 == 0){
    d >>= 1;
    s++;
 for(int i=0; i<K; i++) {
   a = (rand() % (n-2)) + 2; // [2, n-1];
    x = fast_exp_mod(a,d,n);
    if(x == 1 | | x == n-1)
      continue;
    for(ll r=1; r<s; r++) {</pre>
      x = mult_mod(x, x, n);
      if(x == 1)
        return false;
      if(x == n-1)
        goto nextK;
    return false;
 nextK:;
 return true;
void printEntry(bool& printed, ll prime, int ex) {
 if(!printed)
   printed = true;
 else
    cout << " * ";
 cout << prime;</pre>
 if(ex > 1)
    cout << "^" << ex;
void factor(ll x) {
 cout << x << " = ";
 bool printed = false;
 for(int i=0;i<numPrimes;i++)</pre>
    if(x % primes[i] == 0) {
      int ex = 0; // Exponent
      do√
        x \neq primes[i];
        ex++;
      } while (x % primes[i] == 0);
      printEntry(printed, primes[i], ex);
 if(x == 1) {
    cout << endl;
    return;
  // lgPrimes and semiPrimes are useful if there
 // is a lot repetition of large primes/semi-primes // in the test data
 if(lgPrimes.find(x) != lgPrimes.end()){
   printEntry(printed, x, 1);
    cout << endl;
    return;
 if(semiPrimes.find(x) != semiPrimes.end()){
    11 lgFac = semiPrimes[x];
    printEntry(printed, x/lgFac, 1);
    printEntry(printed, lgFac, 1);
```

```
factor large.cc
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                                                                           Page 4/4
    cout << endl;
    return;
  if(miller(x)){ // if x is prime
    printEntry(printed, x, 1);
    cout << endl;
    lgPrimes.insert(x);
    return;
  // Pollard's Rho does not work well with squares,
  // so we'll check for it manually.
  ll \ sqrtX = ll (sqrt(x) + 0.1);
  if(sqrtX*sqrtX == x) {
    printEntry(printed, sqrtX, 2);
    cout << endl;
    return;
  11 smFac = pollardRho(x);
  11 lgFac = x/smFac;
  if(lgFac < smFac)</pre>
    swap(smFac,lgFac);
  printEntry (printed, smFac, 1);
  printEntry (printed, lgFac, 1);
  cout << endl;
  semiPrimes[x] = lgFac;
int main(){
 genSmallPrimes();
  srand((unsigned int) time(NULL));
  11 T,N;
  cin >> T;
  while (T--) {
    cin >> N;
    factor(N);
  return 0:
```

```
fenwicktree.cc
 Sep 29, 14 11:54
                                                                        Page 1/3
 * Fenwick Tree
 * Author: Howard Cheng
  Reference:
    Fenwick, P.M. "A New Data Structure for Cumulative Frequency Tables."
     Software---Practice and Experience, 24(3), 327-336 (March 1994).
 * This code has been tested on UVa 11525 and 11610.
  Fenwick trees are data structures that allows the maintainence of
  cumulative sum tables dynamically. The following operations
  are supported:
 * - Initialize the tree from a list of N integers:
                                                                     0 (N log N)
 * - Read the cumulative sum at index 0 \le k \le N:
                                                                     0(log k)
 * - Read the entry at index 0 <= k < N:
                                                                     0(log N)
 * - Increment/decrement an entry at index 0 \le k \le N in the list: O(\log N)
  - Given a value, find an index such that the cumulative sum at
    that position is the value:
                                                                     0(log N)
 * The space usage is at most 2*N for N input entries.
 * NOTE: it is assumed that all entries are non-negative (even after a
         decrement operation).
 */
#include <vector>
#include <cassert>
using namespace std;
class FenwickTree
public:
 FenwickTree(int n = 0)
   : N(n), tree(n)
   iBM = 1;
    while (iBM < N) {</pre>
     iBM *= 2;
    tree.resize(iBM+1);
    fill(tree.begin(), tree.end(), 0);
  // initialize the tree with the given array of values
 FenwickTree(int val[], int n)
   : N(n)
   iBM = 1;
    while (iBM < N) {
     iBM *= 2;
    tree.resize(iBM+1);
    fill(tree.begin(), tree.end(), 0);
    for (int i = 0; i < n; i++) {
     assert(val[i] >= 0);
     incEntry(i, val[i]);
```

```
fenwicktree.cc
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                                                                       Page 2/3
// increment the entry at position idx by val (use negative val for
// decrement). All affected cumulative sums are updated.
void incEntry(int idx, int val)
  assert(0 <= idx && idx < N);
  if (idx == 0) {
    tree[idx] += val;
  } else {
    do √
      tree[idx] += val;
      idx += idx & (-idx);
    } while (idx < (int)tree.size());</pre>
// return the cumulative sum val[0] + val[1] + ... + val[idx]
int cumulativeSum(int idx) const
  assert(0 <= idx && idx < (int)tree.size());
  int sum = tree[0];
  while (idx > 0)
    sum += tree[idx];
    idx &= idx-1;
  return sum;
// return the entry indexed by idx
int getEntry(int idx) const
  assert(0 <= idx && idx < N);
  int val, parent;
  val = tree[idx];
  if (idx > 0) {
    parent = idx & (idx-1);
    idx--;
    while (parent != idx) {
      val -= tree[idx];
      idx &= idx-1;
  return val:
// return the largest index such that the cumulative frequency is
// what is given, or -1 if it is not found
int getIndex(int sum) const
  int orig = sum;
  if (sum < tree[0]) return -1;</pre>
  sum -= tree[0];
  int idx = 0;
  int bitmask = iBM;
  while (bitmask != 0 \&\& idx < (int)tree.size()-1) {
    int tIdx = idx + bitmask;
    if (sum >= tree[tIdx]) {
      idx = tIdx;
      sum -= tree[tIdx];
    bitmask >>= 1;
  if (sum != 0) {
    return -1;
```

```
fenwicktree.cc
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                                                                                Page 3/3
    idx = min(N-1, idx);
return (cumulativeSum(idx) == orig) ? idx : -1;
private:
  int N, iBM;
 vector<int> tree;
```

```
fflinsolve.cc
 Jan 16, 14 15:08
                                                                         Page 1/2
 * Solution of systems of linear equations over the integers
 * Author: Howard Cheng
 * Reference:
   K.O. Geddes, S.R. Czapor, G. Labahn. "Algorithms for Computer Algebra."
     Kluwer Academic Publishers, 1992, pages 393-399. ISBN 0-7923-9259-0
 * The routine fflinsolve solves the system Ax = b where A is an n x n matrix
 * of integers and b is an n-dimensional vector of integers.
 ^{\star} The inputs to fflinsolve are the matrix A, the dimension n, and an
 * output array to store the solution x_star = det(A) *x. The function
 * also returns the \det(A). In the case that \det(A) = 0, the solution
 * vector is undefined.
 ^{\star} Note that the matrix A and b may be modified.
#include <iostream>
using namespace std;
const int MAX_N = 10;
int fflinsolve(int A[MAX_N][MAX_N], int b[], int x_star[], int n)
 int sign, d, i, j, k, k_c, k_r, pivot, t;
 sign = d = 1;
 for (k_c = k_r = 0; k_c < n; k_c++) {
    /* eliminate column k_c */
    /* find nonzero pivot */
    for (pivot = k_r; pivot < n && !A[pivot][k_r]; pivot++)</pre>
    if (pivot < n) {
      /* swap rows pivot and k_r */
      if (pivot != k_r) {
        for (j = k_c; j < n; j++) {
          t = A[pivot][j];
         A[pivot][j] = A[k_r][j];
         A[k_r][j] = t;
       t = b[pivot];
       b[pivot] = b[k_r];
       b[k_r] = t;
       sign *=-1;
      /* do elimination */
      for (i = k_r+1; i < n; i++) {</pre>
        for (j = k_c+1; j < n; j++) {
         A[i][j] = (A[k_r][k_c]*A[i][j]-A[i][k_c]*A[k_r][j])/d;
       b[i] = (A[k_r][k_c]*b[i]-A[i][k_c]*b[k_r])/d;
       A[i][k_c] = 0;
      if (d) {
       d = A[k_r][k_c];
      k_r++;
    } else {
      /* entire column is 0, det(A) = 0 */
      d = 0;
```

```
fflinsolve.cc
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                                                                           Page 2/2
 if (!d) {
   for (k = k_r; k < n; k++) {
     if (b[k]) {
        /* inconsistent system */
        cout << "Inconsistent system." << endl;</pre>
        return 0;
   /* multiple solutions */
   cout << "More than one solution." << endl;
   return 0;
  /* now backsolve */
 for (k = n-1; k >= 0; k--) {
   x_star[k] = sign*d*b[k];
   for (j = k+1; j < n; j++) {
     x_star[k] -= A[k][j]*x_star[j];
   x_star[k] /= A[k][k];
 return sign*d;
int main(void)
 int A[MAX_N][MAX_N], x_star[MAX_N], b[MAX_N];
 int n, i, j;
 int det;
 while (cin >> n && 0 < n && n <= MAX N) {
   cout << "Enter A:" << endl;
   for (i = 0; i < n; i++)
     for (j = 0; j < n; j++) {
        cin >> A[i][j];
   cout << "Enter b:" << endl;</pre>
   for (i = 0; i < n; i++) {
     cin >> b[i];
   if (det = fflinsolve(A, b, x_star, n)) {
      cout << "det = " << det << endl;
     cout << "x_star = ";
     for (i = 0; i < n; i++) {
        cout << x_star[i] << " ";
     cout << endl;
   } else {
     cout << "A is singular." << endl;
 return 0;
```

```
floyd.cc
 Jan 16, 14 15:08
                                                                             Page 1/2
 * Floyd's Algorithm
 * Author: Howard Cheng
 * The following code takes a graph stored in an adjacency matrix "graph",
 * and returns the shortest distance from node i to node j in dist[i][j].
 ^{\star} We assume that the weights of the edges is not DISCONNECT, and the
 ^{\star} DISCONNECT constant is used to indicate the absence of an edge.
 */
#include <iostream>
#include <cassert>
using namespace std;
const int MAX_NODES = 26;
const int DISCONNECT = -1;
int graph[MAX_NODES][MAX_NODES];
int dist[MAX_NODES][MAX_NODES];
void floyd(void)
 int i, j, k;
 for (i = 0; i < MAX_NODES; i++)</pre>
    for (j = 0; j < MAX_NODES; j++) {</pre>
      dist[i][j] = graph[i][j];
 for (k = 0; k < MAX_NODES; k++) {</pre>
    for (i = 0; i < MAX_NODES; i++) {</pre>
      for (j = 0; j < MAX_NODES; j++) {
  if (dist[i][k] != DISCONNECT && dist[k][j] != DISCONNECT) {</pre>
          int temp = dist[i][k] + dist[k][j];
          if (dist[i][j] == DISCONNECT || dist[i][j] > temp) {
            dist[i][j] = temp;
 for (i = 0; i < MAX_NODES; i++) {</pre>
    dist[i][i] = 0;
int main (void)
 int w;
 int i, j;
  /* clear graph */
 for (i = 0; i < MAX_NODES; i++) {</pre>
    for (j = 0; j < MAX_NODES; j++) {</pre>
      graph[i][j] = DISCONNECT;
  /* read graph */
  cin >> i >> j >> w;
 while (!(i == -1 \&\& j == -1)) {
    assert(0 <= i && i < MAX_NODES && 0 <= j && j < MAX_NODES);
    graph[i][j] = graph[j][i] = w;
    cin >> i >> j >> w;
```

```
Jan 16, 14 15:08 floyd.cc Page 2/2

}
floyd();

/* do queries */
cin >> i >> j;
while (! (i == -1 && j == -1)) {
   assert (0 <= i && i < MAX_NODES && 0 <= j && j < MAX_NODES);
   cout << i << " " << j << ": " << dist[i][j] << endl;
   cin >> i >> j;
}

return 0;
}
```

```
floyd path.cc
 Jan 16, 14 15:08
                                                                          Page 1/2
 * Floyd's Algorithm
 * Author: Howard Cheng
 * The following code takes a graph stored in an adjacency matrix "graph",
 * and returns the shortest distance from node i to node j in dist[i][j].
 * We assume that the weights of the edges is not DISCONNECT, and the
 ^{\star} DISCONNECT constant is used to indicate the absence of an edge.
 * Call extract_path to return the path, as well as its length (in terms
 ^{\star} of vertices). The length is -1 if no such path exists.
 */
#include <iostream>
#include <cassert>
using namespace std;
const int MAX_NODES = 26;
const int DISCONNECT = -1;
int graph [MAX_NODES] [MAX_NODES];
int dist[MAX_NODES][MAX_NODES];
int succ[MAX_NODES][MAX_NODES];
void floyd (void)
 int i, j, k;
 for (i = 0; i < MAX_NODES; i++) {</pre>
    for (j = 0; j < MAX_NODES; j++) {</pre>
      dist[i][j] = graph[i][j];
      if (i == j || graph[i][j] == DISCONNECT) {
        succ[i][j] = -1;
      } else {
        succ[i][j] = j;
 for (k = 0; k < MAX NODES; k++) {
    for (i = 0; i < MAX_NODES; i++) {</pre>
      for (j = 0; j < MAX_NODES; j++) {</pre>
        if (i != k && dist[i][k] != DISCONNECT && dist[k][j] != DISCONNECT) {
          int temp = dist[i][k] + dist[k][j];
          if (dist[i][j] == DISCONNECT || dist[i][j] > temp) {
            dist[i][j] = temp;
            succ[i][j] = succ[i][k];
          } else if (dist[i][j] == temp && succ[i][k] < succ[i][j]) {</pre>
            /* put tie-breaking on paths here */
            /* e.g. the test above chooses lexicographically smallest */
                    paths, but ignores the number of vertices in the */
                    path. To really do lexicographically sorting
                    properly, you also need to have len[i][j] which
                    can be computed easily as well.
            succ[i][j] = succ[i][k];
 for (i = 0; i < MAX_NODES; i++) {</pre>
    dist[i][i] = 0;
```

```
floyd path.cc
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                                                                          Page 2/2
int extract_path(int u, int v, int path[])
 int len = 0;
 if (dist[u][v] == DISCONNECT) {
   return -1;
  path[len++] = u;
  while (u != v)
   u = succ[u][v];
    path[len++] = u;
 return len:
int main (void)
  int m, w, i, j;
 int path [MAX_NODES], len;
  /* clear graph */
  for (i = 0; i < MAX_NODES; i++) {</pre>
    for (j = 0; j < MAX_NODES; j++) {</pre>
      graph[i][j] = DISCONNECT;
  /* read graph */
  cin >> i >> j >> w;
 while (!(i == -1 \&\& j == -1)) {
    assert(0 <= i && i < MAX_NODES && 0 <= j && j < MAX_NODES);
    graph[i][j] = /*graph[j][i] =*/ w;
    cin >> i >> j >> w;
  floyd();
  /* do queries */
  cin >> i >> j;
  while (!(i == -1 \&\& j == -1)) {
    assert(0 <= i && i < MAX_NODES && 0 <= j && j < MAX_NODES);
    cout << i << "" << j << ":" << dist[i][j] << endl;
    len = extract_path(i, j, path);
    for (m = 0; m < len; m++) {
     if (m) {
        cout << " ";
     cout << path[m];
    cout << endl;
    cin >> i >> j;
  return 0;
```

frac2dec.cc Jan 16, 14 15:08 Page 1/2 // Converts a fraction (with integral numerator and denominator) to its decimal expansion. // Author: Darcy Best // Date : August 22, 2010 // Since we are dealing with rational numbers, one of two cases 1. The number will terminate 2. The number will repeat // The algorithm is O(D) where D is the absolute value of the denominator. #include <iostream> #include <string> #include <algorithm> #include <cstdlib> #include <cassert> using namespace std; const int MAX_DENOM = 1001; string itoa(int x) { string ans; while(x) { ans += (x % 10) + '0';x /= 10;reverse(ans.begin(),ans.end()); return (ans.length() ? ans : "0"); int firstSeen[MAX_DENOM]; void frac2dec(int numer, int denom, string& decimal, int& numRepDigs) { assert (denom != 0); // Determine if it is a plus or a minus decimal = ""; **if**(numer < 0 && denom >= 0 || numer >= 0 && denom < 0) { decimal += "-"; } else { decimal += "+"; numer = abs(numer); denom = abs(denom);// Left of the decimal point decimal += itoa(numer / denom); numer %= denom; if(!numer) { numRepDigs = 0;return; // Add the decimal point decimal += '.';// Right of the decimal point fill(firstSeen, firstSeen+denom, -1); int rem = numer; while(rem != 0 && firstSeen[rem] == -1){ firstSeen[rem] = decimal.length(); rem *= 10;decimal += itoa(rem / denom); rem %= denom;

```
fraction.cc
 Jan 16, 14 15:08
                                                                         Page 1/4
// Fraction implementation
// Author: Darcy Best
// Does NOT ever check for division by 0.
// Division by 0 will only cause a runtime error if you use the
    toDouble() function.
#include <iostream>
#include <cstdlib>
using namespace std;
// Change this to whatever integer data type will prevent overflow
// - BigInteger works with this class
typedef long long int dataType;
class Fraction{
public:
 Fraction(dataType num=0, dataType denom=1);
 double toDouble() const;
 void reduce();
 // Changes the fraction itself.
 void selfReciprocal();
  // Returns a new fraction, leaving the original.
 Fraction reciprocal() const;
 Fraction& operator+=(const Fraction& x);
 Fraction& operator = (const Fraction& x);
 Fraction& operator*=(const Fraction& x);
 Fraction& operator/=(const Fraction& x);
 bool operator < (const Fraction & x) const;
 bool operator == (const Fraction& x) const;
 dataType num, denom;
};
Fraction operator+(const Fraction& x, const Fraction& y);
Fraction operator-(const Fraction& x, const Fraction& y);
Fraction operator* (const Fraction& x, const Fraction& y);
Fraction operator/(const Fraction& x, const Fraction& y);
istream& operator>>(istream& is,Fraction& x);
ostream& operator << (ostream& os, const Fraction& x);
Fraction::Fraction(dataType n, dataType d) {
 if(d < 0){
   num = -n;
   denom = -d;
 } else {
   num = n;
   denom = d;
 reduce();
double Fraction::toDouble() const{
 return 1.0*num/denom;
// Howard's GCD function with no checks
dataType gcd(dataType a, dataType b)
```

```
fraction.cc
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                                                                         Page 2/4
  dataType r;
  while (b)
    r = a % b;
    a = b;
    b = r;
  return a;
 void Fraction::reduce(){
  dataType g = gcd(abs(num), denom);
  num /= g;
  denom /= q;
 void Fraction::selfReciprocal(){
  swap (num, denom);
  if (denom < 0) {
    num = -num;
    denom = -denom;
Fraction Fraction::reciprocal() const{
 return Fraction (denom, num);
// Overflow potential in the denominator.
// I've tried to factor out as much as possible before,
// But be careful.
//
// (w)/(a*g) + (z)/(b*g)
// = (w*b)/(a*g*b) + (a*z)/(a*g*b)
// = (w*b + a*z)/(a*g*b)
Fraction& Fraction::operator+=(const Fraction& x) {
  dataType g = gcd(denom, x.denom);
  dataType a = denom / q;
  dataType b = x.denom / g;
  num = num * b + x.num * a;
  denom *= b;
  reduce();
  return (*this);
Fraction& Fraction::operator -= (const Fraction& x) {
  dataType g = gcd(denom, x.denom);
  dataType a = denom / g;
  dataType b = x.denom / g;
  num = num * b - x.num * a;
  denom *= b;
  reduce();
  return (*this);
Fraction& Fraction::operator*=(const Fraction& x) {
 num *= x.num;
  denom *= x.denom;
  reduce();
  return (*this);
Fraction& Fraction::operator/=(const Fraction& x) {
  num *= x.denom;
```

```
fraction.cc
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                                                                         Page 3/4
 denom *= x.num;
 if(denom < 0){
   num = -num;
    denom = -denom;
 reduce();
 return (*this);
// Careful with overflow. If it is an issue, you can compare the
// double values, but you SHOULD check for equality BEFORE converting
bool Fraction::operator<(const Fraction& x) const{</pre>
 return (num*x.denom) < (x.num*denom);</pre>
bool Fraction::operator == (const Fraction& x) const{
 return (num == x.num) && (denom == x.denom);
Fraction operator+(const Fraction& x, const Fraction& y) {
 Fraction a(x);
 a += y;
 return a;
Fraction operator-(const Fraction& x, const Fraction& y) {
 Fraction a(x);
 a -= y;
 return a;
Fraction operator* (const Fraction& x, const Fraction& y) {
 Fraction a(x);
 a *= v;
 return a;
Fraction operator/(const Fraction& x, const Fraction& y) {
 Fraction a(x);
 a /= y;
 return a;
// Note that you can read in Fractions of two forms:
// a/b (With any number of space between a,/,b) - The input "points" to
        the NEXT character after the denom (White space or not)
      (Just an integer - The input "points" to the next NON-WHITE SPACE
       character. Careful when mixing this with getline.)
istream& operator>>(istream& is,Fraction& x) {
 is >> x.num;
 char ch;
 is >> ch;
 if (ch != '/') {
   is.putback(ch);
   x.denom = 1;
 } else {
    is >> x.denom;
    if(x.denom < 0){
     x.num = -x.num;
     x.denom = -x.denom;
   x.reduce();
 return is;
```

```
fraction.cc
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                                                                            Page 4/4
// Will output 5 for 5/1 and 0 for 0/1. If you want always
      fractions, get rid of the if statement
ostream& operator << (ostream& os, const Fraction& x) {
 os << x.num;
 if(x.num != 0 && x.denom != 1)
    os << '/' << x.denom;
  return os;
int main(){
  Fraction x, y;
  while (cin \rightarrow x \rightarrow y) {
    cout << "x: " << x << endl;
    cout << "y: " << y << endl;
    cout << "x+y=" << x+y << endl;
    cout << "x-y=" << x-y << endl;
    cout << "x*y=" << x*y << endl;
    cout << "x/y=" << x/y << endl;
    cout << endl;
  return 0;
```

```
greatcircle.cc
 Jan 16, 14 15:08
                                                                         Page 1/2
// Great Circle distance between two points using Heaverside formula
// Author: Howard Cheng
// Reference: http://mathforum.org/library/drmath/view/51879.html
// Given two points specified by their latitudes and longitudes, as well
// as the radius of the sphere, return the shortest distance between the
// two points along the surface of the sphere.
// latitude should be between -90 to 90 degrees (inclusive), and
// longitude should be between -180 to 180 degrees (inclusive)
// There are also routines that will convert between cartesian coordinates
// (x,y,z) and spherical coordinates (latitude, longitude, radius).
#include <iostream>
#include <iomanip>
#include <cmath>
using namespace std;
const double PI = acos(-1.0);
double greatcircle(double lat1, double long1, double lat2, double long2,
                   double radius)
 lat1 *= PI/180.0;
 lat2 *= PI/180.0;
 long1 *= PI/180.0;
 long2 *= PI/180.0;
 double dlong = long2 - long1;
 double dlat = lat2 - lat1;
 double \ a = \sin(dlat/2) * \sin(dlat/2) +
    cos(lat1)*cos(lat2)*sin(dlong/2)*sin(dlong/2);
 return radius * 2 * atan2(sqrt(a), sqrt(1-a));
void longlat2cart(double lat, double lon, double radius,
                  double &x, double &y, double &z)
 lat *= PI/180.0;
 lon *= PI/180.0;
 x = radius * cos(lat) * cos(lon);
 y = radius * cos(lat) * sin(lon);
 z = radius * sin(lat);
void cart2longlat(double x, double y, double z,
                  double &lat, double &lon, double &radius)
 radius = sqrt(x*x + y*y + z*z);
lat = (PI/2 - acos(z / radius)) * 180.0 / PI;
 lon = atan2(y, x) * 180.0 / PI;
int main(void)
 int T;
 cin >> T;
 while (T-- > 0) {
    const double radius = 6371009;
    double lat1, long1, lat2, long2;
    cin >> lat1 >> long1 >> lat2 >> long2;
    double x1, y1, z1, x2, y2, z2;
```

```
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                                      heron.cc
                                                                        Page 1/1
// Heron's formula
// Computes the area of a triangle given the lengths of the three sides.
// Author: Howard Cheng
#include <iostream>
#include <iomanip>
#include <utility>
#include <cmath>
using namespace std;
// the lengths of the three sides are a, b, and c. The routine returns
// the area of the triangle, or -1 if the three lengths do not make a
// triangle.
double area_heron(double a, double b, double c)
 if (a < b) swap(a, b);
 if (a < c) swap(a, c);
 if (b < c) swap(b, c);
 if (c < a - b) return -1;
 return sqrt((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c))/4.0;
int main(void)
 double a, b, c;
 while (cin >> a >> b >> c) {
   cout << fixed << setprecision(4) << area_heron(a, b, c) << endl;</pre>
 return 0;
```

```
Jan 16, 14 15:08
                                     hungarian.cc
                                                                         Page 1/4
 * Maximum/minimum weight bipartite matching
 * Author: Howard Cheng
 * Reference:
   http://www.topcoder.com/tc?module=Static&d1=tutorials&d2=hungarianAlgorithm
 * This file contains routines for computing the maximum/minimum weight
 * bipartite matching.
 ^{st} It is assumed that each half of the graph has exactly N vertices, labelled
 * 0 to N-1. The weight between vertex i on the left and vertex j on the
 * right is stored in G[i][j]. The cost of the optimal matching is returned,
 * and matching[i] is the vertex on the right that is matched to vertex i
 * on the left.
 \mbox{\ensuremath{^{*}}} If an edge is absent, the corresponding edge weight should be:
     INT_MIN if maximum weight matching is desired
    INT_MAX if minimum weight matching is desired
 * This is an implementation of the Hungarian algorithm. The complexity
 * is O(N^3).
 */
#include <iostream>
#include <algorithm>
#include <queue>
#include <cassert>
#include <climits>
using namespace std;
const int MAX N = 3;
void \ update\_labels(int \ lx[MAX_N], \ int \ ly[MAX_N], \ bool \ S[MAX_N], \ bool \ T[MAX_N],
                   int slack[MAX N], int N)
 int delta;
 bool delta_init = false;
 for (int y = 0; y < N; y++) {
   if (T[y]) continue;
    delta = delta_init ? min(delta, slack[y]) : slack[y];
   delta_init = true;
 for (int x = 0; x < N; x++) {
    if (S[x]) lx[x] -= delta;
 for (int y = 0; y < N; y++) {
   if (T[y]) {
      ly[y] += delta;
    } else {
      slack[y] -= delta;
void add_to_tree(int x, int prevx, int G[MAX_N][MAX_N], bool S[MAX_N],
                 int prev[MAX_N], int lx[MAX_N], int ly[MAX_N],
                 int slack[MAX_N], int slackx[MAX_N], int N)
 S[x] = true;
 prev[x] = prevx;
 for (int y = 0; y < N; y++)
    int temp = (G[x][y] == INT_MIN) ? INT_MAX : lx[x] + ly[y] - G[x][y];
    if (temp < slack[y]) {</pre>
      slack[y] = temp;
```

```
hungarian.cc
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                                                                        Page 2/4
     slackx[v] = x;
 }
int max_weight_matching(int G[MAX_N][MAX_N], int N, int matching[MAX_N])
 int revmatch[MAX_N];
                              // match from right to left
                              // number of vertices in current matching
 int \max_{max_match} = 0;
 fill (matching, matching+N, -1);
 fill(revmatch, revmatch+N, -1);
 // find an initial feasible labelling
 int lx[MAX_N], ly[MAX_N];
 fill(ly, ly+N, 0);
 for (int x = 0; x < N; x++) {
   lx[x] = *max\_element(G[x], G[x]+N);
 // now repeatedly find alternating tree, augment, and relabel
 while (max_match < N) {</pre>
   queue<int> q;
   bool S[MAX_N], T[MAX_N];
   int prev[MAX_N];
   fill(S, S+N, false);
   fill(T, T+N, false);
   fill(prev, prev+N, -1);
   // find root of alternating tree
   int root = find(matching, matching+N, -1) - matching;
   q.push(root);
   prev[root] = -2;
   S[root] = true;
   int slack[MAX_N], slackx[MAX_N];
   for (int y = 0; y < N; y++)
     slack[y] = (G[root][y] == INT_MIN) ? INT_MAX :
       lx[root] + ly[y] - G[root][y];
     slackx[y] = root;
   bool path_found = false;
   int x, y;
   while (!path_found) {
      // build alternating tree with BFS
     while (!path_found && !q.empty()) {
       x = q.front();
       q.pop();
       for (y = 0; y < N; y++) {
         // go through edges in equality graph
         if (G[x][y] == lx[x] + ly[y] && !T[y]) {
           if (revmatch[y] == -1) {
             path_found = true;
             break;
           T[y] = true;
           q.push(revmatch[y]);
           add_to_tree(revmatch[y], x, G, S, prev, lx, ly, slack, slackx, N);
     if (path_found) break;
      // no augmenting path, update the labels
     update_labels(lx, ly, S, T, slack, N);
      while (!q.empty()) {
       q.pop();
```

```
hungarian.cc
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                                                                         Page 3/4
      for (y = 0; y < N; y++) {
       if (!T[y] && slack[y] == 0) {
          if (revmatch[y] == -1) {
            x = slackx[y];
            path_found = true;
            break;
          } else {
            T[y] = true;
            if (!S[revmatch[y]]) {
              q.push(revmatch[y]);
              add_to_tree(revmatch[y], slackx[y], G, S, prev, lx, ly, slack,
                          slackx, N);
    assert (path_found);
   max_match++;
    // augment along the path
    for (int cx = x, cy = y, ty; cx != -2; cx = prev[cx], cy = ty) {
      ty = matching[cx];
      revmatch[cy] = cx;
      matching[cx] = cy;
 // return the final answer
 int weight = 0;
 for (int x = 0; x < N; x++)
   weight += G[x][matching[x]];
 return weight;
int min_weight_matching(int G[MAX_N][MAX_N], int N, int matching[MAX_N])
 int M = INT_MIN;
 for (int i = 0; i < N; i++) {
   for (int j = 0; j < N; j++) {
     if (G[i][j] != INT_MAX) {
       M = max(M, G[i][j]);
 int newG[MAX_N][MAX_N];
 for (int i = 0; i < N; i++) {</pre>
   for (int j = 0; j < N; j++) {
      newG[i][j] = (G[i][j] == INT_MAX) ? INT_MIN : M - G[i][j];
 int weight = max_weight_matching(newG, N, matching);
 return N*M - weight;
int main (void)
 int G[3][3] = \{ \{INT\_MAX, 4, 5\}, \{5, 7, 6\}, \{5, 8, 8\} \};
 int matching[3];
 int w = min_weight_matching(G, 3, matching);
 cout << "weight = " << w << endl;
 for (int i = 0; i < 3; i++) {
    cout << i << "is matched to " << matching[i] << endl;</pre>
```

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<pre>} return 0;</pre>		
,		

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```
infix.cc
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                                                                        Page 1/4
 * Infix expressions evaluation
 * Author: Howard Cheng
 * The evaluate() routine takes a string containing an infix arithmetic
 * expression, and return the numeric result after evaluation. The
 * parameter error indicates whether an error has occurred (syntax
 * error, illegal operation, etc.). If there is an error the result
 * returned is meaningless.
 ^{\star} The routine assumes that the operands in the input are integers
 * with no leading signs. It supports the standard +, -, *, / and
 * parentheses. If you need to support more operators, operand types,
 * etc., you will need to modify the code. See comments below.
#include <iostream>
#include <string>
#include <stack>
#include <cctype>
#include <cstdlib>
using namespace std;
// What is a token? Modify if needed (e.g. to support variables, extra
// operators, etc.)
struct Token
 enum Type {NUMBER, PLUS, MINUS, TIMES, DIVIDE, LEFT PAREN, RIGHT PAREN};
 // priority of the operators: bigger number means higher priority
 // e.g. */ has priority 2, +- has priority 1, ( has priority 0
 int priority[7];
 // is the operator left associative? It's assumed that all operators
 // of the same priority level has the same left/right associative property
 bool left_assoc[7];
 Type type;
 long val;
 Token()
   priority[1] = priority[2] = 1;
   priority[3] = priority[4] = 2;
   priority[5] = 0;
    left_assoc[1] = left_assoc[2] = left_assoc[3] = left_assoc[4] = true;
 int get_priority() {
   return priority[type];
 bool is_left_assoc() {
    return left_assoc[type];
  // returns true if there is a next token
 bool next_token(string &expr, int &start, bool &error)
    int len = expr.length();
    error = false;
    while (start < len && isspace(expr[start])) {</pre>
     start++;
    if (start >= len) {
```

```
infix.cc
 Jan 16, 14 15:08
                                                                         Page 2/4
      return false;
    switch (expr[start]) {
    case '(':
      type = LEFT_PAREN;
      break;
    case ')':
      type = RIGHT_PAREN;
      break;
    case '*'
      type = TIMES;
      break;
    case '/':
      type = DIVIDE;
      break;
    case '+':
      type = PLUS;
      break;
    case '-':
      type = MINUS;
      break;
    default:
      // check for number
      const char *s = expr.c_str() + start;
      val = strtol(s, &p, 10);
      if (s == p) {
       error = true;
        return false;
      type = NUMBER;
      start += (p - s);
    if (type != NUMBER)
      start++;
    return true;
};
// Modify this if you need to support more operators or change their
// meanings.
// returns true if operation is successful
bool apply_op(stack<long> &operands, Token token)
  long a, b;
  if (operands.size() < 2) {</pre>
    return false;
  if (token.type == Token::PLUS)
   b = operands.top(); operands.pop();
    a = operands.top(); operands.pop();
    operands.push(a+b);
  } else if (token.type == Token::MINUS) {
    b = operands.top(); operands.pop();
    a = operands.top(); operands.pop();
    operands.push(a-b);
  } else if (token.type == Token::TIMES) {
    b = operands.top(); operands.pop();
    a = operands.top(); operands.pop();
    operands.push(a*b);
  } else if (token.type == Token::DIVIDE) {
    b = operands.top(); operands.pop();
    a = operands.top(); operands.pop();
    if (b == 0) {
      return false;
```

```
infix.cc
 Jan 16, 14 15:08
                                                                        Page 3/4
    operands.push(a/b);
 } else {
    return false;
 return true;
long evaluate (string expr, bool &error)
 stack<Token> s;
 stack<long> operands;
 int i;
 Token token;
 error = false;
 i = 0;
 while (token.next_token(expr, i, error) && !error) {
    switch (token.type) {
    case Token::NUMBER:
     operands.push(token.val);
     break;
    case Token::LEFT_PAREN:
     s.push(token);
     break;
    case Token::RIGHT_PAREN:
     while (!(error = s.empty()) && s.top().type != Token::LEFT_PAREN) {
       if ((error = !apply_op(operands, s.top()))) {
       s.pop();
     if (!error) {
       s.pop();
     break;
    default:
                  // arithmetic operators
     while (!error && !s.empty() &&
             (token.get_priority() < s.top().get_priority() ||</pre>
              token.get_priority() == s.top().get_priority() &&
              token.is_left_assoc())) {
       error = !apply_op(operands, s.top());
       s.pop();
     if (!error) {
       s.push(token);
    if (error) {
     break;
 while (!error && !s.empty()) {
   error = !apply_op(operands, s.top());
    s.pop();
 error |= (operands.size() != 1);
 if (error) {
   return 0;
 return operands.top();
int main(void)
 int result;
 string expr;
 bool error;
```

```
Printed by Howard Cheng
                                        infix.cc
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                                                                           Page 4/4
getline(cin, expr);
while (!cin.eof())
  result = evaluate(expr, error);
  if (error)
     cout << "Invalid expression" << endl;</pre>
   } else {
    cout << " = " << result << endl;
  getline(cin, expr);
return 0;
```

```
intersect circle circle.cc
 Jan 16, 14 15:08
                                                                         Page 1/2
// Determines the point(s) of intersection if a circle and a circle
// Author: Darcy Best
// Date : October 1, 2010
// Source: http://local.wasp.uwa.edu.au/~pbourke/geometry/2circle/
// Note: A circle of radius 0 must be considered independently.
// See comments in the implementation.
#include <iostream>
#include <iomanip>
#include <cmath>
#include <algorithm>
using namespace std;
\#define SQR(X) ((X) * (X))
// How close to call equal
const double EPS = 1e-4;
bool dEqual (double x, double y) {
 return fabs(x-y) < EPS;
struct Point{
 double x,y;
 bool operator<(const Point& a) const{
   if (dEqual(x,a.x))
     return y < a.y;</pre>
    return x < a.x;
};
// Prints out the ordered pair. This also accounts for the negative 0.
void print (const Point & a) {
 cout << "(";
 if(fabs(a.x) < 1e-4)
    cout << "0.000";
 else
   cout << a.x;
  cout << ",";
 if(fabs(a.y) < 1e-4)
   cout << "0.000";
    cout << a.y;
 cout << ")";
struct Circle{
 double r,x,y;
};
// Input:
   Two circles to intersect
// Output:
// Number of points of intersection points
   If 1 (or 2), then ans1 (and ans2) contain those points.
  If 3, then there are infinitely many. (They're the same circle)
int intersect_circle_circle(Circle c1,Circle c2,Point& ans1,Point& ans2) {
  // If we have two singular points
 if(fabs(c1.r) < EPS && fabs(c2.r) < EPS) {</pre>
   if (dEqual(c1.x,c2.x) && dEqual(c1.y,c2.y)) {
      ans1.x = c1.x;
      ans1.y = c1.y;
      // Here, you need to know what the intersection of two exact points is:
      // "return 1;" - If the points intersect at only 1 point
      // "return 3;" - If the circles are the same
```

```
intersect circle circle.cc
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                                                                         Page 2/2
      // Note that both are true -- It all depends on the problem
      return 1;
     else
      return 0;
 double d = hypot(c1.x-c2.x,c1.y-c2.y);
  // Check if the circles are exactly the same.
 if(dEqual(c1.x,c2.x) && dEqual(c1.y,c2.y) && dEqual(c1.r,c2.r))
   return 3:
  // The circles are disjoint
 if(d > c1.r + c2.r + EPS)
   return 0;
  // One circle is contained inside the other -- No intersection
 if(d < abs(c1.r-c2.r) - EPS)
   return 0;
 double a = (SQR(c1.r) - SQR(c2.r) + SQR(d)) / (2*d);
 double h = sqrt(abs(SQR(c1.r) - SQR(a)));
 Point P:
 P.x = c1.x + a / d * (c2.x - c1.x);
 P.y = c1.y + a / d * (c2.y - c1.y);
 ans1.x = P.x + h / d * (c2.y - c1.y);
 ans1.y = P.y - h / d * (c2.x - c1.x);
 if (fabs(h) < EPS)
   return 1:
 ans2.x = P.x - h / d * (c2.y - c1.y);
 ans2.y = P.y + h / d * (c2.x - c1.x);
 return 2;
int main() {
 cout << fixed << setprecision(3);</pre>
 Circle C1.C2:
 Point al, a2;
 while(cin >> C1.x >> C1.y >> C1.r >> C2.x >> C2.y >> C2.r) {
   int num = intersect_circle_circle(C1,C2,a1,a2);
   switch (num) {
   case 0:
      cout << "NO INTERSECTION" << endl;
     break:
   case 1:
     print(a1); cout << endl;
     break;
   case 2:
     if(a2 < a1)
       swap(a1,a2);
     print(a1);print(a2);cout << endl;</pre>
   case 3:
     cout << "THE CIRCLES ARE THE SAME" << endl;</pre>
     break;
 return 0;
```

```
intersect iline.cc
 Jan 16, 14 15:08
                                                                           Page 1/2
 * 2-D Line Intersection
 * Author: Howard Cheng
 * Reference:
    http://www.exaflop.org/docs/cgafaq/cga1.html
 \mbox{\scriptsize *} This routine takes two infinite lines specified by two points, and
 * determines whether they intersect at one point, infinitely points,
 * or no points. In the first case, the point of intersection is also
 * returned. The points of a line must be different (otherwise,
 ^{\star} the line is not defined).
*/
#include <iostream>
#include <cmath>
#include <cassert>
using namespace std;
/* how close to call equal */
const double EPSILON = 1E-8;
struct Point {
 double x, y;
};
/* returns 1 if intersect at a point, 0 if not, -1 if the lines coincide */
int intersect_iline(Point a, Point b, Point c, Point d, Point &p)
 double denom, num1, num2;
 assert((a.x != b.x || a.y != b.y) && (c.x != d.x || c.y != d.y));
 num1 = (a.y - c.y) * (d.x - c.x) - (a.x - c.x) * (d.y - c.y);
 num2 = (a.y - c.y) * (b.x - a.x) - (a.x - c.x) * (b.y - a.y);
 denom = (b.x - a.x)*(d.y - c.y) - (b.y - a.y)*(d.x - c.x);
 if (fabs(denom) >= EPSILON) {
    r = num1 / denom;
    p.x = a.x + r*(b.x - a.x);
   p.y = a.y + r*(b.y - a.y);
    return 1;
 } else {
    if (fabs(num1) >= EPSILON) {
     return 0;
    } else {
     return -1;
int main (void)
 Point a, b, c, d, p;
 int res;
 while (cin >> a.x >> a.y >> b.x >> b.y >> c.x >> c.y >> d.x >> d.y) {
    res = intersect_iline(a, b, c, d, p);
    if (res == 1) {
      cout << "Intersect at (" << p.x << ", " << p.y << ")" << endl;
    } else if (res == 0)
     cout << "Don't intersect" << endl;
    } else {
      cout << "Infinite number of intersections" << endl;</pre>
```

```
Printed by Howard Cheng
                                 intersect iline.cc
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                                                                        Page 2/2
return 0;
```

intersect iline circle.cc Mar 25, 14 14:08 Page 1/2 // Determines the point(s) of intersection if a circle and a line // Author: Darcy Best // Date : May 1, 2010 // Source: http://mathworld.wolfram.com/Circle-LineIntersection.html #include <iostream> #include <cmath> using namespace std; #define SQR(X) ((X) * (X)) // How close to call equal const double EPS = 1e-7; bool dEqual(double x, double y) { return fabs(x-y) < EPS; struct Point{ double x,y; }; struct Line{ Point p1,p2; }; struct Circle{ Point centre; double radius; }; // Input of: // - 2 distinct points on the line // - The centre of the circle // - The radius of the circle // Output: // Number of points of intersection points // If 1 or 2, then ans1 and ans2 contain those points. int intersect_iline_circle(Line 1, Circle c, Point& ans1, Point& ans2) { Point p1 = 1.p1;Point p2 = 1.p2; Point circCentre = c.centre; double rad = c.radius; p1.x -= circCentre.x; p2.x -= circCentre.x; p1.y -= circCentre.y; p2.y -= circCentre.y; double dx = p2.x - p1.x;double dy = p2.y - p1.y;double dr = SQR(dx) + SQR(dy);double D = p1.x*p2.y - p2.x*p1.y; double desc = SQR(rad)*dr - SQR(D); if (dEqual (desc, 0)) { ans1.x = circCentre.x + (D*dy) / dr;ans1.y = circCentre.y + (-D*dx) / dr; return 1; } **else if**(desc < 0){ return 0; double sgn = (dy < -EPS ? -1 : 1);ans1.x = circCentre.x + (D*dy + sqn*dx*sqrt(desc)) / dr;

```
intersect iline circle.cc
 Mar 25, 14 14:08
                                                                        Page 2/2
 ansl.y = circCentre.y + (-D*dx + abs(dy)*sqrt(desc)) / dr;
  ans2.x = circCentre.x + (D*dy - sgn*dx*sqrt(desc)) / dr;
 ans2.y = circCentre.y + (-D*dx - abs(dy)*sqrt(desc)) / dr;
 return 2:
int main(){
 Line L;
 Circle C;
 Point a1,a2;
 cin >> L.p1.x >> L.p1.y >> L.p2.x >> L.p2.y;
 cin >> C.centre.x >> C.centre.y >> C.radius;
  int num = intersect_iline_circle(L,C,a1,a2);
 if(num == 0)
   cout << "NO INTERSECTION." << endl;
  else if(num == 1)
   cout << "ONE INTERSECTION: (" << a1.x << "," << a1.y << ")" << endl;
 else if(num == 2)
   cout << "TWO INTERSECTIONS: (" << a1.x << "," << a1.y << ")"
         << "(" << a2.x << "," << a2.y << ")" << endl;
 return 0;
```

```
intersect line.cc
 Jan 16, 14 15:11
                                                                        Page 1/2
 * 2-D Line Intersection
 * Author: Howard Cheng
 * Reference:
    http://www.exaflop.org/docs/cgafaq/cgal.html
 ^{\star} This routine takes two line segments specified by endpoints, and
 * determines whether they intersect at one point, infinitely points,
 * or no points. In the first case, the point of intersection is also
 * returned. The endpoints of a line must be different (otherwise,
 * the line is not defined).
*/
#include <iostream>
#include <cmath>
#include <cassert>
using namespace std;
/* how close to call equal */
const double EPSILON = 1E-8;
struct Point {
 double x, y;
};
/* returns 1 if intersect at a point, 0 if not, -1 if the lines coincide */
int intersect_line(Point a, Point b, Point c, Point d, Point &p)
 Point t;
 double r, s;
 double denom, num1, num2;
 assert((a.x != b.x || a.y != b.y) && (c.x != d.x || c.y != d.y));
 num1 = (a.y - c.y)*(d.x - c.x) - (a.x - c.x)*(d.y - c.y);
 num2 = (a.y - c.y) * (b.x - a.x) - (a.x - c.x) * (b.y - a.y);
 denom = (b.x - a.x)*(d.y - c.y) - (b.y - a.y)*(d.x - c.x);
 if (fabs(denom) >= EPSILON) {
   r = num1 / denom;
    s = num2 / denom;
    if (0-EPSILON <= r && r <= 1+EPSILON &&</pre>
       0-EPSILON <= s && s <= 1+EPSILON) {
      /* always do this part if we are interested in lines instead */
     /* of line segments
     p.x = a.x + r^*(b.x - a.x);
     p.y = a.y + r*(b.y - a.y);
     return 1;
    } else {
     return 0;
 } else {
    if (fabs(num1) >= EPSILON) {
     return 0;
    } else {
     /* I am not using "fuzzy comparisons" here, because the comparisons */
      /* are based on the input, not some derived quantities. You may
     /* want to change that if the input points are computed somehow.
      /* two lines are the "same". See if they overlap */
     if (a.x > b.x | | (a.x == b.x && a.y > b.y)) {
       t = a;
       a = b;
       b = t;
      if (c.x > d.x \mid | (c.x == d.x && c.y > d.y)) {
```

```
intersect line.cc
 Jan 16, 14 15:11
                                                                            Page 2/2
        t = c;
        c = d;
        d = t;
      if (a.x == b.x) {
        /* vertical lines */
        if (b.y == c.y) {
          p = b;
          return 1;
        } else if (a.y == d.y) {
          return 1;
        } else if (b.y < c.y || d.y < a.y) {</pre>
          return 0;
        } else
          return -1;
      } else {
        if (b.x == c.x) {
          p = b;
          return 1;
        } else if (a.x == d.x) {
          return 1;
        } else if (b.x < c.x || d.x < a.x) {</pre>
          return 0;
        } else
          return -1;
      return -1;
int main (void)
 Point a, b, c, d, p;
 int res:
  while (cin >> a.x >> a.y >> b.x >> b.y >> c.x >> c.y >> d.x >> d.y) {
   res = intersect_line(a, b, c, d, p);
    if (res == 1) {
      cout << "Intersect at (" << p.x << ", " << p.y << ")" << endl;
    } else if (res == 0)
      cout << "Don't intersect" << endl;
    } else {
      cout << "Infinite number of intersections" << endl;
 }
 return 0;
```

```
intersectTF.cc
 Jan 16, 14 15:08
                                                                        Page 1/2
 * Line Intersection
 * Author: Howard Cheng
 * Reference:
    CLRS, "Introduction to Algorithms", 2nd edition, pages 936-939.
 * Given two lines specified by their endpoints (a1, a2) and (b1, b2),
 * returns true if they intersect, and false otherwise. The intersection
 * point is not known.
 */
#include <iostream>
#include <cmath>
using namespace std;
/* how close to call equal */
const double EPSILON = 1E-8;
struct Point {
 double x, y;
double direction (Point p1, Point p2, Point p3)
 double x1 = p3.x - p1.x;
 double y1 = p3.y - p1.y;
 double x2 = p2.x - p1.x;
 double y2 = p2.y - p1.y;
 return x1*y2 - x2*y1;
int on_segment(Point p1, Point p2, Point p3)
 return ((p1.x <= p3.x && p3.x <= p2.x) || (p2.x <= p3.x && p3.x <= p1.x)) &&
    ((p1.y \le p3.y \&\& p3.y \le p2.y) \mid | (p2.y \le p3.y \&\& p3.y \le p1.y));
int intersect (Point al, Point a2, Point b1, Point b2)
 double d1 = direction(b1, b2, a1);
 double d2 = direction(b1, b2, a2);
 double d3 = direction(a1, a2, b1);
 double d4 = direction(a1, a2, b2);
 if (((d1 > EPSILON && d2 < -EPSILON) || (d1 < -EPSILON && d2 > EPSILON)) &&
     ((d3 > EPSILON && d4 < -EPSILON) || (d3 < -EPSILON && d4 > EPSILON))) {
    return 1;
 } else {
   return (fabs(d1) < EPSILON && on_segment(b1, b2, a1)) | |</pre>
      (fabs(d2) < EPSILON && on_segment(b1, b2, a2)) ||
      (fabs(d3) < EPSILON && on_segment(a1, a2, b1)) ||
      (fabs(d4) < EPSILON && on_segment(a1, a2, b2));
int main (void)
 Point a, b, c, d;
 int a1, a2, a3, a4, a5, a6, a7, a8;
 while (cin >> a1 >> a2 >> a3 >> a4 >> a5 >> a6 >> a7 >> a8) {
   a.x = a1; a.y = a2;
   b.x = a3; b.y = a4;
   c.x = a5; c.y = a6;
    d.x = a7; d.y = a8;
    if (intersect(a, b, c, d)) {
```

```
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                                  intersectTF.cc
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                                                                       Page 2/2
    cout << "Yes" << endl;
  } else
    cout << "No" << endl;
return 0;
```

```
int mult.cc
 Jan 16, 14 15:08
                                                                         Page 1/2
 {\tt *\ Integer\ multiplication/division\ without\ overflow}
 * Author: Howard Cheng
 * Given a list of factors in the numerator (num, size n) and a list \,
 * of factors in the denominator (dem, size m), it returns the product
 ^{\star} of the numerator divided by the denominator. It is assumed that
 ^{\star} the numerator is divisible by the denominator (ie. the result
 * is an integer). Overflow will not occur as long as the final result
* is representable.
*/
#include <iostream>
#include <cassert>
using namespace std;
int gcd(int a, int b)
  int r;
  while (b) {
     r = a % b;
     a = b;
     b = r;
  assert (a >= 0);
  return a;
int mult(int A[], int n, int B[], int m)
  int i, j, prod, d;
  int count = 0;
   /* unnecessary if the two lists are positive */
  for (i = 0; i < n; i++) {
    if (A[i] < 0) {
      A[i] *= -1;
       count++;
  for (i = 0; i < m; i++) {
    if (B[i] < 0) {
      B[i] *= -1;
       count++;
  for (i = 0; i < n; i++) {
      for (j = 0; j < m; j++) {
        d = gcd(A[i], B[j]);
         A[i] /= d;
         B[j] /= d;
  prod = 1;
  for (i = 0; i < n; i++) {
     prod *= A[i];
  for (j = 0; j < m; j++) {
      assert(B[j] == 1);
  return (count % 2 == 0) ? prod : -prod;
int main (void)
```

```
Jan 16, 14 15:08 int_mult.cc Page 2/2

{
    int A[1000], B[1000], n, m, i;

    while (cin >> n >> m && n > 0 && m > 0) {
        for (i = 0; i < n; i++) {
            cin >> A[i];
        }
        for (i = 0; i < m; i++) {
            cin >> B[i];
        }
        cout << "prod = " << mult(A, n, B, m) << endl;
    }

    return 0;
```

```
int prog.c
 Jan 16, 14 15:08
                                                                          Page 1/3
 * All-integer programming
 * Author: Howard Cheng
 * Reference:
    http://www.cs.sunysb.edu/~algorith/implement/syslo/distrib/processed/
 * This algorithm is based on GOMORY cutting plane method.
 * This algorithm solves the following INTEGER LP problem:
                  SUM \quad (A[0][j] * x[j])
                                                  [cost function]
  minimize
              (i=0 \text{ to } n-1)
                  SUM (A[i][j]*x[j]) <= A[i][n]
                                                         1 <= i <= m
       s.t.
              (i=0 \text{ to } n-1)
                  x[i] >= 0
                                     0 <= i <= n-1
       and
 * n = number of variables
  m = number of constraints
 * Input : An input array A with m+n+1 rows and n+1 columns.
           Store the cost function in row 0, and the constraints in rows
             1 to m. Set A[0][n] = 0.
           A vector x allocated for n values to store returned value.
  Output: Returns 1 if a solution is found, 0 if no solution exists.
           The minimum value of the cost function is returned in
             value.
           The variable assignment to x[j] that gives the minimum is given
             in x[j], where 0 \ll j \ll n-1.
  Important Notes:
 * 1. If we want to have constraints that are >=, just multiply all the
    coefficients by -1.
 * 2. If we want to have constraints that are ==, do both >= and <=.
 \star 3. The contents of A is destroyed after this routine.
 * 4. The coefficients in the cost function must be positive. If not,
     make a change of variable x'[i] = m[i] - x[i] where m[i] is the maximum value for variable[i] and adjust all constraints as well
      as the returned optimal value. This is especially useful if you
      wish to maximize the cost function.
      Usually there is some maximum for each variable if you wish to
      maximize the function (or the value could be infinity.
      NOTE: if any coefficient in the objective function is negative or
            0, the routine will crash.
 * 5. If one only wishes to know if there is any variable assignment
      satisfying the constraints, just put 1 in each coefficient
      of the objective function.
#include <stdio.h>
#include <assert.h>
#define MAX_VARS 50
#define MAX_CONS 50
#define MAX_ROWS MAX_VARS+MAX_CONS+1
#define MAX_COLS MAX_VARS+1
int euclid(int u, int v)
 int w = u / v;
 if (w*v > u) {
```

```
int prog.c
Jan 16, 14 15:08
                                                                         Page 2/3
 if ((w+1)*v <= u) {
   w++;
 return w:
int int_prog(int A[MAX_ROWS][MAX_COLS], int n, int m, int *value, int *x)
 int iter, nosol;
 int b, c, i, j, k, l, r, r1, s, t, denom, num;
 for (j = 0; j < n; j++) {
   if (A[0][j] <= 0) {
     // BAD objective function coefficient: make sure it is positive
     assert (false);
 /* set constraints that x[j] >= 0, and clear output */
 for (i = 0; i < n; i++) {
   for (j = 0; j < n+1; j++) {
     A[m+1+i][j] = 0;
   A[m+1+i][i] = -1;
 A[0][n] = 0;
 nosol = 0;
 do {
   r = 0;
   do {
     iter = (A[++r][n] < 0);
    } while (!iter && r != n+m);
   if (iter)
     for (k = iter = 0; k < n && !iter; k++) {</pre>
       iter = (A[r][k] < 0);
     nosol = !iter:
     if (iter) {
       1 = k-1;
       for (j = k; j < n; j++) {
         if (A[r][j] < 0) {
           for (i = 0; !(s = A[i][j] - A[i][l]); i++)
           if (s < 0) {
              1 = \dot{j};
       for (s = 0; !A[s][1]; s++)
       num = -A[r][1];
       denom = 1;
       for (j = 0; j < n; j++) {
         if (A[r][j] < 0 && j != 1) {
           for (i = s-1, b = 1; b && i >= 0; i--) {
              b = (A[i][j] == 0);
           if (b) {
              i = A[s][j];
              r1 = A[s][1];
              t = euclid(i, r1);
              if (t*r1 == i \&\& t > 1) {
                for (i = s+1; !(r1 = t*A[i][1] - A[i][j]); i++)
                if (r1 > 0) {
                  t--;
```

```
int_prog.c
                                                                           Page 3/3
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              c = -A[r][j];
              if (c*denom > t*num) {
                num = c;
                denom = t;
        for (j = 0; j <= n; j++) {
          if (j != 1) {
            c = euclid(A[r][j]*denom, num);
            if (c) {
              for (i = 0; i <= n+m; i++) {
               A[i][j] += c*A[i][1];
 } while (iter && !nosol);
  *value = -A[0][n];
 for (j = 0; j < n; j++) {
   x[j] = A[m+1+j][n];
 return !nosol;
int main(void)
 int A[MAX_ROWS][MAX_COLS];
 int x[MAX_VARS];
 int val, t;
 int m, n, i, j;
 while (scanf("%d%d", &n, &m) == 2 && n > 0 && m > 0) {
    /* read cost function */
    printf("Input cost function to minimize:\n");
    for (i = \hat{0}; i < n; i++) {
      scanf("%d", &A[0][i]);
    /* read constraints */
    for (i = 1; i <= m; i++) {</pre>
      printf("Input constraint #%d:\n", i);
      for (j = \bar{0}; j < n+1; j++) {
        scanf("%d", &A[i][j]);
    t = int_prog(A, n, m, \&val, x);
    if (t) {
      printf("Minimum cost = %d\n", val);
      for (i = 0; i < n; i++)
        printf("x[\%2d] = \%2d\n", i, x[i]);
    } else {
      printf("No solution exists.\n");
 return 0;
```

```
josephus.cc
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                                                                            Page 1/1
// Josephus Problem
// Author: Darcy Best
// Date : September 4, 2010
// The Josephus problem:
    A group of n people are in a circle, and you start by killing
       person f. Then, you kill every kth person until only one person
// Two implementations are given here (Note that neither depend on k):
    1. Determine the survivor -- O(n)
     2. Determine the full killing order -- O(n^2)
// If there are 17 people, with every 5th person killed (killing the // lst person first), the kill order is:
       1, 6, 11, 16, 5, 12, 2, 9, 17, 10, 4, 15, 14, 3, 8, 13, 7 (survivor = 7)
// NOTE: This is 1-based, not 0-based.
#include <iostream>
using namespace std;
const int MAX_N = 100;
int survivor(int n, int f, int k) {
 return (n==1 ? 1 : (survivor(n-1,k,k) + (f-1)) % n + 1);
void killOrder(int n, int f, int k, int A[]){
 if(n == 0) return;
 A[0] = 0;
 killOrder(n-1,k,k,A+1);
 for(int i=0;i<n;i++)
    A[i] = (A[i] + (f-1)) % n + 1;
int main(){
 int n,f,k,kOrder[MAX_N];
 while(cin >> n >> f >> k && (n || f || k)){
    killOrder(n,f,k,kOrder);
    for(int i=0;i<n;i++)
      cout << kOrder[i] << endl;</pre>
    cout << "Survivor: " << survivor(n,f,k) << endl;</pre>
 return 0;
```

```
kmp.cc
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                                                                           Page 1/2
 * KMP String Matching
 * Author: Howard Cheng
 ^{\star} The prepare_pattern routine takes in the pattern you wish to search
 * for, and perform some processing to give a "failure array" to be used
 * by the actual search. The complexity is linear in the length of the
 * pattern.
 * The find_pattern routine takes in a string s, a pattern pat, and a
* vector T computed by prepare pattern. It returns the index of the * first occurrence of pat in s, or -1 if it does not occur in s.
 * The complexity is linear in the length of the string s.
 */
#include <iostream>
#include <string>
#include <vector>
#include <algorithm>
using namespace std;
void prepare_pattern(const string &pat, vector<int> &T)
 int n = pat.length();
 T.resize(n+1);
 fill(T.begin(), T.end(), -1);
 for (int i = 1; i <= n; i++) {
    int pos = T[i-1];
    while (pos != -1 && pat[pos] != pat[i-1]) {
      pos = T[pos];
    T[i] = pos + 1;
int find_pattern(const string &s, const string &pat, const vector<int> &T)
 int sp = 0, kp = 0;
 int slen = s.length(), plen = pat.length();
 while (sp < slen) {
    while (kp != -1 \&\& (kp == plen || pat[kp] != s[sp])) {
      kp = T[kp];
    kp++; sp++;
    if (kp == plen)
     // a match is found
      return sp - plen;
      // if you want more than one match (i.e. all matches), do not return
      // in the above but rather record the location of the match. Continue
      // the loop with:
      // kp = T[kp];
 return -1;
int main (void)
 string str, pat;
 while (cin >> str >> pat) {
   vector<int> T;
    prepare_pattern(pat, T);
    cout << "index = " << find_pattern(str, pat, T) << endl;</pre>
```

```
Printed by Howard Cheng
                                       kmp.cc
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                                                                       Page 2/2
return 0;
```

```
linsolve.cc
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                                                                       Page 1/3
 * Solution of systems of linear equations
 * Author: Howard Cheng
 * Reference:
    K.E. Atkinson. "An Introduction to Numerical Analysis." 2nd Ed., John
    Wiley & Sons, 1988, pages 520-521. ISBN 0-471-62489-6
 * To solve the system Ax = b where A is an n \times n matrix, first call
 * LU decomp on A to obtain its LU decomposition. Once the LU
 * decomposition is obtained, it can be used to solve linear systems with
 * the same coefficient matrix A but different vectors of b using the
 * LU_solve routine. This routine is numerically stable (provided that
 * the original coefficient matrix has a small condition number).
 * The inputs to LU_decomp are the matrix A, the dimension n, an
 * output array pivot of n-1 elements such that pivot[i] = j means
 * that rows i and j were swapped during the i-th step, and an output
 * parameter to return the determinant of the matrix. The function
 * returns 1 if successful, and 0 if the matrix is singular. The
 ^{\star} matrix A is overwritten by its LU decomposition on return. If the
 * matrix is singular, the content of A should not be used (it represents
 * intermediate results during the decomposition).
 * The inputs to LU_solve are the LU decomposition of A, the dimension
 * n, the pivot array from LU_decomp, and n-dimensional vectors b and
 * x. This function should be called only if the original matrix A
 * has a small condition number. You can check this by checking that
 * the determinant returned by LU_decomp is not too close to 0. This is
 * only a crude check: you should really be computing the condition number
 * of the matrix.
 */
#include <iostream>
#include <cmath>
using namespace std;
const int MAX_N = 10;
int LU_decomp(double A[MAX_N][MAX_N], int n, int pivot[], double &det)
 double s[MAX_N];
                             /* factors used in implicit scaling */
 double c, t;
 int i, j, k;
 det = 1.0;
  /* compute s[i] */
 for (i = 0; i < n; i++) {
   s[i] = 0.0;
    for (j = 0; j < n; j++) {
     if ((t = fabs(A[i][j])) > s[i]) {
       s[i] = t;
    if (s[i] == 0.0) {
      /* a row of zeroes: singular */
     det = 0.0;
     return 0;
  /* do the row reductions */
 for (k = 0; k < n-1; k++) {
   c = fabs(A[k][k]/s[k]);
    pivot[k] = k;
    for (i = k+1; i < n; i++) {
```

```
linsolve.cc
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                                                                           Page 2/3
      t = fabs(A[i][k]/s[i]);
      if (t > c) {
       c = t;
        pivot[k] = i;
    if (c == 0) {
      /* pivot == 0: singular */
      det = 0.0;
      return 0;
    /* do row exchange */
    if (k != pivot[k]) {
      det *= -1.0;
      for (j = k; j < n; j++)  {
       t = A[k][j];
        A[k][j] = A[pivot[k]][j];
        A[pivot[k]][j] = t;
        t = s[k];
       s[k] = s[pivot[k]];
        s[pivot[k]] = t;
    /* do the row reduction */
    for (i = k+1; i < n; i++) {</pre>
     A[i][k] /= A[k][k];
     for (j = k+1; j < n; j++) {
   A[i][j] -= A[i][k] * A[k][j];</pre>
    det *= A[k][k];
  /* note that the algorithm as state in the book is incorrect. The */
  /* following is need to ensure that the last row is not all 0's.
  /* (maybe the book is correct, depending on what you think it's
  /* supposed to do.)
 if (\bar{A[n-1]}[n-1] == 0.0) {
   det = 0.0;
   return 0;
 } else {
   det *= A[n-1][n-1];
    return 1;
void LU_solve(double A[MAX_N][MAX_N], int n, int pivot[], double b[],
              double x[])
 double t;
 int i, j, k;
 for (i = 0; i < n; i++) {
   x[i] = b[i];
 for (k = 0; k < n-1; k++) {
    /* swap if necessary */
    if (k != pivot[k]) {
      t = x[k];
     x[k] = x[pivot[k]];
     x[pivot[k]] = t;
    for (i = k+1; i < n; i++) {
      x[i] -= A[i][k] * x[k];
```

```
linsolve.cc
                                                                            Page 3/3
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 x[n-1] /= A[n-1][n-1];
 for (i = n-2; i >= 0; i--) {
   for (j = i+1; j< n; j++) {
     x[i] -= A[i][j] * x[j];
   x[i] /= A[i][i];
int main(void)
 double A[MAX_N][MAX_N], x[MAX_N], b[MAX_N];
                         /* only n-1 is needed, but what the heck */
 int pivot[MAX_N];
 int n, i, j;
 double det;
 while (cin >> n && 0 < n && n <= MAX_N) {
    cout << "Enter A:" << endl;</pre>
    for (i = 0; i < n; i++) {
     for (j = 0; j < n; j++) {
       cin >> A[i][j];
     }
    cout << "Enter b:";
    for (i = 0; i < n; i++) {</pre>
     cin >> b[i];
    if (LU_decomp(A, n, pivot, det)) {
      LU_solve(A, n, pivot, b, x);
      cout << "LU decomposition of A:" << endl;</pre>
      for (i = 0; i < n; i++) {
        for (j = 0; j < n; j++) {
  cout << A[i][j] << "";</pre>
        cout << endl;
      cout << "det = " << det << endl;
      cout << "x = ";
      for (i = 0; i < n; i++) {
       cout << x[i] << "";
      cout << endl;
    } else {
      cout << "A is singular" << endl;</pre>
 return 0;
```

```
matching.c
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                                                                         Page 1/4
/* unweighted matching in a bipartite graph.
 * author: Matthew McNaughton, Jan 16, 1999.
 * mcnaught@cs.ualberta.ca
 * The bipartite graph {\it G} is split into two sets, {\it U} and {\it V},
 * of user-defined maximum size MAXU and MAXV.
 * the input graph is in bipgraph[MAXU][MAXV].
 * there is an edge between node u \in U and node v \in V
 * iff bipgraph[u][v] != 0.
 * The output is in matching[MAXU].
 * node u \setminus in U and node v \setminus in V are matched iff matching[u] == v.
 * parameters match(int u, int v) mean: u is the number of vertices
 * in U, v the number in V. They are assumed to be numbered 0 .. u-1
 * and 0 .. v-1, respectively.
 * Technique: given a non-maximum matching M on G, find an "alternating path"
 * u_1 v_1 ... u_n v_n so that u_1 and v_n are not matched in M, but
 * v_k u_k+1 are matched with each other. Then "flip" the edges so
 * that edges on this path which were not in the matching are, and edges
 * which were are not. This increases the size of the matching by one.
 ^{\star} It is a fact that if no such path exists, then M is maximum.
 * This algorithm finds several alternating paths at once by performing
 * bfs starting at all unmatched nodes u \in U. Paths which do not
 * have intersecting nodes can be alternated in the same bfs run.
 * bfs is performed repeated until the matching cannot be expanded.
#include <stdio.h>
#include <string.h>
#include <assert.h>
FILE *in, *out;
/* change these as necessary */
#define MAXU 100
#define MAXV 100
#define U(i) (i)
#define V(i) ((i) + MAXU)
#define isU(i) ((i) < MAXU)
#define isV(i) ((i) >= MAXU)
#define isMatched(i) (isU(i) ? flagUmatched[(i)] : flagVmatched[(i)-MAXU])
#define isUsed(i) (isU(i) ? flagUused[(i)] : flagVused[(i)-MAXU])
#define isVisited(i) (isU(i) ? flagUvisited[(i)] : flagVvisited[(i)-MAXU])
#define setMatched(i) (isU(i)?(flagUmatched[(i)]=1):(flagVmatched[(i)-MAXU]=1))
#define setUsed(i) (isU(i)?(flagUused[(i)]=1):(flagVused[(i)-MAXU]=1))
#define setVisited(i) (isU(i)?(flaqUvisited[(i)]=1):(flaqVvisited[(i)-MAXU]=1))
char bipgraph [MAXU] [MAXV];
int matching[MAXU]; /* matching[u] == v, _not_ plus MAXU */
char flagUmatched[MAXU], flagVmatched[MAXV];
char flagUvisited[MAXU], flagVvisited[MAXV];
char flagUused[MAXU], flagVused[MAXV];
int predecessor[MAXU+MAXV], queue[MAXU+MAXV];
/* u and v are the number of vertices in sets U, and V, respectively,
* filling up bipgraph[0..u-1][0..v-1].
* result:
 * matching[u0] == v0 iff u0 and v0 are in the matching,
* otherwise matching[u0] = -1 */
void
match(int u, int v) {
 int i, j, head, tail, bad, last, increased;
```

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for( i = 0; i < u; i++ ) {
  matching[i] = -1;
  flagUmatched[i] = 0;
for( i = 0; i < v; i++ ) flagVmatched[i] = 0;</pre>
do { /* find alternating paths by repeated bfs. */
  for( i = 0; i < MAXU+MAXV; i++ ) predecessor[i] = -1;</pre>
  for( i = 0; i < MAXU; i++ ) flagUused[i] = flagUvisited[i] = 0;</pre>
  for( i = 0; i < MAXV; i++ ) flagVused[i] = flagVvisited[i] = 0;</pre>
  head = tail = 0;
  /* put all the unmatched u's on the queue. They start the
    * alternating path. */
  for( i = 0; i < u; i++ )
    if( ! isMatched(U(i))) {
      queue[tail++] = U(i);
      predecessor[i] = -1; /* redundant statement */
      setVisited(U(i));
  /* flag that at least one path was found by the bfs.
   * when the bfs does not find an alternating path we are done. */
  increased = 0;
  while( head != tail ) {
    i = queue[head++];
    ^{\prime *} this node appeared on some previously found alternating path. ^{*\prime}
    if( isUsed(i) ) continue;
    if( isV(i) && !isMatched(i) ) {
      /* we got to the end of an alternating path. see if
       * it is disjoint with other paths found so far. only
        * then can we mess it up a bit. */
      for( j = i; j != -1; j = predecessor[j]) {
        if( isUsed(i)) {
          bad = 1;
          break;
      if(! bad) {
        /* this path is pristine. switch "polarity" of edges
         * in the matching on this path. */
        /* flag and instrumention - whether (not) to guit,
         * and how many paths we found this bfs. */
        increased++;
        for( j = i; j != -1; last = j, j = predecessor[j] ) {
          if( isV(j) && !isMatched(j)) {
            /* the only unmatched v - actually this means we
              * are on the first iteration of this loop. */
            setMatched(i);
          } else if( isU(j) ) {
            if( isMatched(j) ) {
              /* the node we saw in the previous iteration of
                * this loop must be a V. We will match with it
                * instead of the one we used to match with, which
                * must be the next node visited in this loop. */
              assert(isV(last));
              matching[j] = last - MAXU;
             } else {
               /* we are the very first u, one of the ones the
                * bfs queue was "seeded" with. We should have ...*/
```

matching.c Jan 16, 14 15:11 Page 3/4 assert (predecessor[i] == -1); setMatched(j); assert(isV(last)); matching[j] = last - MAXU; setUsed(j); /* this node cannot be used for other * paths we might run across in the future * on this bfs. */ } /* for */ } /* if ! bad */ } /* isV and !isMatched */ else if(isV(i)) { /* this must be a matched V - find the matching U and put it on * the queue if it is not visited or used. */ bad = 1;for (j = 0; j < u; j++) { if(isMatched(U(j)) && matching[j] == i - MAXU){ /* this is the one. */ if(! isVisited(U(j)) && !isUsed(U(j))) { setVisited(U(j)); queue[tail++] = U(j); predecessor[U(j)] = i; bad = 0;break; assert (!bad); } /* isV */ else if(isU(i)) /* we are at U. whether it is unmatched (a "seed"), * or matched, we do the same thing - put on the queue * all V's which it is connected to in the graph but * which it is _not_ paired to in the current matching. */ $for(j = 0; j < v; j++) {$ if (bipgraph[i][j] && !isVisited(V(j)) && !isUsed(V(j)) && matching[i] != j) { /* we can put this one on the queue. */ queue[tail++] = V(j); predecessor[V(j)] = i;setVisited(V(j)); } else { assert(0); /* should be no other cases. */ /st this is the end of the bfs. st/} while (increased); return; main() int i,j,u,v,setnum; in = stdin; out = stdout; setnum = 0; while (fscanf (in, "%d %d", &u, &v) == 2) {

mincostmaxflowdense.cc Jan 16, 14 15:11 Page 1/3 * Min Cost Max Flow for Dense graphs * Authors: Frank Chu, Igor Naverniouk * http://shygypsy.com/tools/mcmf3.cpp * Min cost max flow * (Edmonds-Karp relabelling + Dijkstra) * This implementation takes a directed graph where each edge has a * capacity ('cap') and a cost per unit of flow ('cost') and returns a * maximum flow network of minimal cost ('fcost') from s to t. - cap (global): adjacency matrix where cap[u][v] is the capacity of the edge $u \rightarrow v$. cap[u][v] is 0 for non-existent edges. - cost (global): a matrix where cost[u][v] is the cost per unit of flow along the edge $u \rightarrow v$. If cap[u][v] == 0, cost[u][v] is ignored. ALL COSTS MUST BE NON-NEGATIVE! - n: the number of vertices ([0, n-1] are considered as vertices). - s: source vertex. - t: sink. * RETURNS: - the flow - the total cost through 'fcost' - fnet contains the flow network. Careful: both fnet[u][v] and fnet[v][u] could be positive. Take the difference. * COMPLEXITY: - Worst case: O(n^2*flow <? n^3*fcost) * REFERENCE: Edmonds, J., Karp, R. "Theoretical Improvements in Algorithmic Efficieincy for Network Flow Problems". This is a slight improvement of Frank Chu's implementation. **/ #include <iostream> #include <algorithm> #include <climits> using namespace std; // the maximum number of vertices + 1 const int NN = 1024; // adjacency matrix (fill this up) int cap[NN][NN]; // cost per unit of flow matrix (fill this up) int cost[NN][NN]; // flow network and adjacency list int fnet[NN][NN], adj[NN][NN], deg[NN]; // Dijkstra's successor and depth int par[NN], d[NN]; // par[source] = source; // Labelling function int pi[NN]; const int Inf = INT_MAX/2; // Dijkstra's using non-negative edge weights (cost + potential) #define Pot(u,v) (d[u] + pi[u] - pi[v])bool dijkstra(int n, int s, int t) for (int i = 0; i < n; i++) {</pre> d[i] = Inf;par[i] = -1;

```
mincostmaxflowdense.cc
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                                                                        Page 2/3
  d[s] = 0;
 par[s] = -n - 1;
  while (1)
   // find u with smallest d[u]
    int u = -1, bestD = Inf;
    for (int i = 0; i < n; i++) {
     if (par[i] < 0 && d[i] < bestD) {</pre>
       bestD = d[u = i];
    if (bestD == Inf) break;
    // relax edge (u,i) or (i,u) for all i;
    par[u] = -par[u] - 1;
    for (int i = 0; i < deg[u]; i++) {
     // try undoing edge v->u
     int v = adj[u][i];
     if (par[v] >= 0) continue;
      if (fnet[v][u] && d[v] > Pot(u,v) - cost[v][u]) {
       d[v] = Pot(u, v) - cost[v][u];
       par[v] = -u-1;
      // try edge u->v
      if (fnet[u][v] < cap[u][v] && d[v] > Pot(u,v) + cost[u][v]) {
       d[v] = Pot(u,v) + cost[u][v];
       par[v] = -u - 1;
  for (int i = 0; i < n; i++) {
   if (pi[i] < Inf) {
     pi[i] += d[i];
 return par[t] >= 0;
#undef Pot
int mcmf (int n, int s, int t, int &fcost )
  // build the adjacency list
  fill(deg, deg+NN, 0);
 for (int i = 0; i < n; i++) {
   for (int j = 0; j < n; j++) {
     if (cap[i][j] || cap[j][i]) {
        adj[i][deg[i]++] = j;
  for (int i = 0; i < NN; i++) {</pre>
   fill(fnet[i], fnet[i]+NN, 0);
  fill(pi, pi+NN, 0);
  int flow = fcost = 0;
  // repeatedly, find a cheapest path from s to t
  while (dijkstra(n, s, t)) {
   // get the bottleneck capacity
    int bot = INT_MAX;
    for (int v = t, u = par[v]; v != s; u = par[v = u]) {
     bot = min(bot, fnet[v][u] ? fnet[v][u] : (cap[u][v] - fnet[u][v]));
```

```
mincostmaxflowdense.cc
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    // update the flow network
    for (int v = t, u = par[v]; v != s; u = par[v = u]) {
     if (fnet[v][u]) {
       fnet[v][u] -= bot;
        fcost -= bot * cost[v][u];
     } else {
       fnet[u][v] += bot;
       fcost += bot * cost[u][v];
    flow += bot;
 return flow;
#include <iostream>
using namespace std;
int main()
 int numV;
 cin >> numV;
 for (int i = 0; i < NN; i++) {</pre>
   fill(cap[i], cap[i]+NN, 0);
 int m, a, b, c, cp;
 int s, t;
 cin >> m;
 cin >> s >> t;
 // fill up cap with existing capacities.
 // if the edge u \rightarrow v has capacity 6, set cap[u][v] = 6.
 // for each cap[u][v] > 0, set cost[u][v] to the
  // cost per unit of flow along the edge i \rightarrow v
 for (int i=0; i<m; i++) {
   cin >> a >> b >> cp >> c;
   cost[a][b] = c; // cost[b][a] = c;
    cap[a][b] = cp; // cap[b][a] = cp;
 int fcost;
 int flow = mcmf( numV, s, t, fcost );
 cout << "flow: " << flow << endl;
 cout << "cost: " << fcost << endl;
 return 0;
```

mincostmaxflowsparse.cc Jan 16, 14 15:11 Page 1/3 // MIN COST MAX FLOW // Authors: Frank Chu, Igor Naverniouk /******** * Min cost max flow * (Edmonds-Karp relabelling + fast heap Dijkstra) * Takes a directed graph where each edge has a capacity ('cap') and a * cost per unit of flow ('cost') and returns a maximum flow network * of minimal cost ('fcost') from s to t. USE mcmf3.cpp FOR DENSE GRAPHS! - cap (global): adjacency matrix where cap[u][v] is the capacity of the edge $u \rightarrow v$. cap[u][v] is 0 for non-existent edges. - cost (global): a matrix where cost[u][v] is the cost per unit of flow along the edge $u \rightarrow v$. If cap[u][v] == 0, cost[u][v] is ignored. ALL COSTS MUST BE NON-NEGATIVE! - n: the number of vertices ([0, n-1] are considered as vertices). - s: source vertex. - t: sink. * RETURNS: - the flow - the total cost through 'fcost' - fnet contains the flow network. Careful: both fnet[u][v] and fnet[v][u] could be positive. Take the difference. COMPLEXITY: - Worst case: O(m*log(m)*flow <? n*m*log(m)*fcost) * FIELD TESTING: - Valladolid 10594: Data Flow * REFERENCE: Edmonds, J., Karp, R. "Theoretical Improvements in Algorithmic Efficieincy for Network Flow Problems". This is a slight improvement of Frank Chu's implementation. **/ #include <iostream> #include <algorithm> #include <climits> using namespace std; // the maximum number of vertices + 1 #define NN 1024 // adjacency matrix (fill this up) int cap[NN][NN]; // cost per unit of flow matrix (fill this up) int cost[NN][NN]; // flow network and adjacency list int fnet[NN][NN], adj[NN][NN], deg[NN]; // Dijkstra's predecessor, depth and priority queue int par[NN], d[NN], q[NN], inq[NN], qs; // Labelling function int pi[NN]; #define Inf (INT_MAX/2) #define BUBL { \ $t = q[i]; q[i] = q[j]; q[j] = t; \$ $t = inq[q[i]]; inq[q[i]] = inq[q[j]]; inq[q[j]] = t; }$ // Dijkstra's using non-negative edge weights (cost + potential) #define Pot(u,v) (d[u] + pi[u] - pi[v])

```
mincostmaxflowsparse.cc
 Jan 16, 14 15:11
                                                                             Page 2/3
bool dijkstra ( int n, int s, int t )
  fill(d, d+NN, Inf);
  fill(par, par+NN, -1);
  fill(ing, ing+NN, -1);
  d[s] = qs = 0;
  inq[q[qs++] = s] = 0;
 par[s] = n;
    // get the minimum from q and bubble down
    int u = q[0];
    inq[u] = -1;
    q[0] = q[--qs];
    if( qs ) inq[q[0]] = 0;
    for (int i = 0, j = 2*i + 1, t; j < qs; i = j, j = 2*i + 1) {
   if (j + 1 < qs && d[q[j + 1]] < d[q[j]]) j++;</pre>
      if (d[q[j]] >= d[q[i]]) break;
      BUBL;
    // relax edge (u,i) or (i,u) for all i;
    for (int k = 0, v = adj[u][k]; k < deg[u]; v = adj[u][++k]) {
      // try undoing edge v->u
      if (fnet[v][u] && d[v] > Pot(u,v) - cost[v][u])
        d[v] = Pot(u,v) - cost[v][par[v] = u];
      // try using edge u->v
      if (fnet[u][v] < cap[u][v] && d[v] > Pot(u,v) + cost[u][v])
        d[v] = Pot(u,v) + cost[par[v] = u][v];
      if (par[v] == u) {
        // bubble up or decrease key
        if( inq[v] < 0 ) { inq[q[qs] = v] = qs; qs++; }
        for( int i = inq[v], j = ( i - 1 )/2, t;
    d[q[i]] < d[q[j]]; i = j, j = ( i - 1 )/2 )</pre>
 for( int i = 0; i < n; i++ ) if( pi[i] < Inf ) pi[i] += d[i];</pre>
 return par[t] >= 0;
int mcmf( int n, int s, int t, int &fcost )
  // build the adjacency list
  fill(deg, deg+NN, 0);
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++)
      if (cap[i][j] || cap[j][i]) adj[i][deg[i]++] = j;
  for (int i = 0; i < NN; i++) {</pre>
    fill(fnet[i], fnet[i]+NN, 0);
  fill(pi, pi+NN, 0);
  int flow = fcost = 0;
  // repeatedly, find a cheapest path from s to t
 while (dijkstra(n, s, t)) {
    // get the bottleneck capacity
    int bot = INT_MAX;
    for (int v = t, u = par[v]; v != s; u = par[v = u]) {
      bot = min(bot, fnet[v][u] ? fnet[v][u] : ( cap[u][v] - fnet[u][v] ));
```

mincostmaxflowsparse.cc Jan 16, 14 15:11 Page 3/3 // update the flow network for (int v = t, u = par[v]; v != s; u = par[v = u]) if (fnet[v][u]) { fnet[v][u] -= bot; fcost -= bot * cost[v][u]; } else { fnet[u][v] += bot; fcost += bot * cost[u][v]; } flow += bot; return flow; int main() int numV; int m, a, b, c, cp; int s, t; cin >> numV; cin >> m; cin >> s >> t; // fill up cap with existing capacities. // if the edge $u \rightarrow v$ has capacity 6, set cap[u][v] = 6. // for each cap[u][v] > 0, set cost[u][v] to the // cost per unit of flow along the edge u->v **for** (int i=0; i<m; i++) { cin >> a >> b >> cp >> c; cost[a][b] = c; // cost[b][a] = c; cap[a][b] = cp; // cap[b][a] = cp; int fcost; int flow = mcmf(numV, s, t, fcost); cout << "flow: " << flow << endl; cout << "cost: " << fcost << endl; return 0;

```
mst.cc
 Jan 16, 14 15:08
                                                                        Page 1/3
 * Implementation of Kruskal's Minimum Spanning Tree Algorithm
 * Author: Howard Cheng
 * This is a routine to find the minimum spanning tree. It takes as
       n: number of vertices
       m: number of edges
    elist: an array of edges (if (u,v) is in the list, there is no need
           for (v,u) to be in, but it wouldn't hurt, as long as the weights
           are the same).
  The following are returned:
    index: an array of indices that shows which edges from elist are in
           the minimum spanning tree. It is assumed that its size is at
     size: the number of edges selected in "index". If this is not
           n-1, the graph is not connected and we have a "minimum
           spanning forest."
  The weight of the MST is returned as the function return value.
 * The run time of the algorithm is O(m \log m).
 * Note: the elements of elist may be reordered.
 * Modified by Rex Forsyth using C++ Aug 28, 2003
 * This version defines the unionfind and edge as classes and provides
 * constructors. The edge class overloads the < operator. So the sort does
 * not use a * cmp function. It uses dynamic arrays.
#include <cmath>
#include <iostream>
#include <iomanip>
#include <cstdlib>
#include <cassert>
#include <algorithm>
using namespace std;
class UnionFind
     struct UF { int p; int rank; };
  public:
      UnionFind(int n) {
                                 // constructor
        howMany = n;
        uf = new UF[howMany];
        for (int i = 0; i < howMany; i++) {</pre>
           uf[i].p = i;
           uf[i].rank = 0;
      ~UnionFind() {
        delete[] uf;
      int find(int x) { return find(uf,x); }
                                                 // for client use
     bool merge(int x, int y) {
        int res1, res2;
        res1 = find(uf, x);
        res2 = find(uf, y);
         if (res1 != res2) {
            if (uf[res1].rank > uf[res2].rank) {
```

```
mst.cc
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                                                                          Page 2/3
               uf[res2].p = res1;
            else
               uf[res1].p = res2;
if (uf[res1].rank == uf[res2].rank) {
                  uf[res2].rank++;
            return true;
         return false;
  private:
      int howMany;
      UF* uf;
      int find(UF uf[], int x) {
                                              // for internal use
         if (uf[x].p != x) {
            uf[x].p = find(uf, uf[x].p);
         return uf[x].p;
};
class Edge {
  public:
      Edge (int i=-1, int j=-1, double weight=0) {
        v1 = i;
         v2 = j;
         w = weight;
      bool operator<(const Edge& e) const { return w < e.w; }</pre>
      int v1, v2;
                            /* two endpoints of edge
      double w;
                            /* weight, can be double instead of int */
};
double mst(int n, int m, Edge elist[], int index[], int& size)
 UnionFind uf(n):
 sort (elist, elist+m);
 double w = 0.0;
 size = 0;
 for (int i = 0; i < m && size < n-1; i++) {
   int c1 = uf.find(elist[i].v1);
    int c2 = uf.find(elist[i].v2);
    if (c1 != c2) {
     index[size++] = i;
     w += elist[i].w;
     uf.merge(c1, c2);
 }
 return w;
int main (void)
  cout << fixed << setprecision(2);</pre>
  cin >> n;
   double* x = new double[n];
   double* y = new double[n];
```

```
mult.cc
 Jan 16, 14 15:08
                                                                         Page 1/1
 ^{\star} Multiplication/division without overflow
 * Author: Howard Cheng
 * Given a list of factors in the numerator (num, size n) and a list \,
 st of factors in the denominator (dem, size m), it returns the product
 ^{\star} of the numerator divided by the denominator, while reducing the
 * result as soon as it is larger than some BOUND.
#include <iostream>
#include <cassert>
using namespace std;
const int BOUND = (1 << 16);</pre>
double mult(double num[], int n, double dem[], int m)
   int i, j;
  double prod = 1.0;
   i = j = 0;
   while (i < n || j < m) {
      if (prod >= BOUND && j < m) {
         prod /= dem[j++];
      } else if (i < n) {
         prod *= num[i++];
      } else {
         assert(j < m);
         prod /= dem[j++];
   return prod;
int main(void)
 double A[1000], B[1000];
 int n, m, i;
 while (cin >> n >> m && n > 0 && m > 0) {
    for (i = 0; i < n; i++) {
      cin >> A[i];
    for (i = 0; i < m; i++) {
      cin >> B[i];
    cout << "prod = " << mult(A, n, B, m) << endl;
 return 0;
```

```
networkflow2.cc
 Jan 16, 14 15:08
                                                                        Page 1/3
 * Network Flow (Relabel-to-front)
 * Author: Howard Cheng
 * The routine network\_flow() finds the maximum flow that can be
 * pushed from the source (s) to the sink (t) in a flow network
 ^{\star} (i.e. directed graph with capacities on the edges). The maximum
 * flow is returned. The flow is given in the flow array (look for
 * positive flow).
 * The complexity of this algorithm is O(n^3), which is good if the
 * graph is small but the maximum flow can be large. Since the
 * algorithm is O(n^3) we are going to use the adjacency matrix
 ^{\star} representation.
#include <iostream>
#include <list>
#include <cassert>
using namespace std;
const int MAX_NODE = 102;
void clear_graph(int graph[MAX_NODE][MAX_NODE], int n)
 for (int i = 0; i < n; i++)
   for (int j = 0; j < n; j++) {
     graph[i][j] = 0;
 }
void push(int graph[MAX_NODE][MAX_NODE], int flow[MAX_NODE][MAX_NODE],
          int e[], int u, int v)
 int cf = graph[u][v] - flow[u][v];
 int d = (e[u] < cf) ? e[u] : cf;
 flow[u][v] += d;
 flow[v][u] = -flow[u][v];
 e[u] -= d;
 e[v] += d;
void relabel(int graph[MAX_NODE][MAX_NODE], int flow[MAX_NODE][MAX_NODE],
             int n, int h[], int u)
 h[u] = -1;
 for (int v = 0; v < n; v++) {
   if (graph[u][v] - flow[u][v] > 0 &&
        (h[u] == -1 \mid | 1 + h[v] < h[u])) {
     h[u] = 1 + h[v];
 assert (h[u] >= 0);
void discharge(int graph[MAX_NODE][MAX_NODE], int flow[MAX_NODE][MAX_NODE],
               int n, int e[], int h[], list<int>& NU,
               list<int>::iterator &current, int u)
 while (e[u] > 0) {
   if (current == NU.end()) {
     relabel(graph, flow, n, h, u);
     current = NU.begin();
    } else {
      int v = *current;
```

```
networkflow2.cc
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                                                                        Page 2/3
      if (qraph[u][v] - flow[u][v] > 0 && h[u] == h[v] + 1) {
        push (graph, flow, e, u, v);
      } else
        ++current;
int network_flow(int graph[MAX_NODE][MAX_NODE], int flow[MAX_NODE][MAX_NODE],
                 int n, int s, int t)
 int e[MAX_NODE], h[MAX_NODE];
 int u, v, oh;
 list<int> N[MAX_NODE], L;
 list<int>::iterator current[MAX_NODE], p;
 for (u = 0; u < n; u++) {
   h[u] = e[u] = 0;
 for (u = 0; u < n; u++)
   for (v = 0; v < n; v++) {
     flow[u][v] = 0;
     if (graph[u][v] > 0 || graph[v][u] > 0) {
       N[u].push_front(v);
 h[s] = n;
 for (u = 0; u < n; u++) {
   if (graph[s][u] > 0) {
     e[u] = flow[s][u] = graph[s][u];
     e[s] += flow[u][s] = -graph[s][u];
   if (u != s && u != t) {
     L.push_front(u);
   current[u] = N[u].begin();
 p = L.begin();
 while (p != L.end()) {
   u = *p;
   oh = h[u];
   discharge (graph, flow, n, e, h, N[u], current[u], u);
   if (h[u] > oh) {
     L.erase(p);
     L.push_front(u);
     p = L.begin();
   ++p;
  int maxflow = 0;
 for (u = 0; u < n; u++) {
   if (flow[s][u] > 0) {
     maxflow += flow[s][u];
 return maxflow;
void print_flow(int flow[MAX_NODE][MAX_NODE], int n)
 for (int i = 0; i < n; i++) {
   for (int j = 0; j < n; j++) {
     if (flow[i][j] > 0) {
        cout << i << "->" << j << ":" << flow[i][j] << endl;
```

```
networkflow2.cc
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                                                                           Page 3/3
int main(void)
 int graph[MAX_NODE][MAX_NODE];
 int s, t;
 int n, m, u, v, c;
int flow[MAX_NODE][MAX_NODE];
 int maxflow;
 while (cin >> n && n > 0) {
    clear_graph(graph, n);
    cin >> m >> s >> t;
    while (m-- > 0) {
      cin >> u >> v >> c;
      graph[u][v] = c;
   maxflow = network_flow(graph, flow, n, s, t);
    cout << "flow = " << maxflow << endl;</pre>
   print_flow(flow, n);
 return 0;
```

```
networkflow.cc
 Jan 16, 14 15:08
                                                                         Page 1/3
 * Network Flow
 * Author: Howard Cheng
 * The routine network_flow() finds the maximum flow that can be
 * pushed from the source (s) to the sink (t) in a flow network
 * (i.e. directed graph with capacities on the edges). The maximum
 * flow is returned. Note that the graph is modified. If you wish to
 * recover the flow on an edge, it is in the "flow" field, as long as
 * is_real is set to true.
 * Note: if you have an undirected network. simply call add_edge twice
 ^{\star} with an edge in both directions (same capacity). Note that 4 edges
 * will be added (2 real edges and 2 residual edges). To discover the
 * actual flow between two vertices u and v, add up the flow of all
 ^{\star} real edges from u to v and subtract all the flow of real edges from
 ^{st} v to u. (In fact, for a residual edge the flow is always 0 in this
 * implementation.)
 * This code can also be used for bipartite matching by setting up an
 * appropriate flow network.
 ^{\star} The code here assumes an adjacency list representation since most
 * problems requiring network flow have sparse graphs.
 ^{\star} This is the basic augmenting path algorithm and it is not the most
 * efficient. But it should be good enough for most programming contest
 * problems. The complexity is O(f m) where f is the size of the flow
 * and m is the number of edges. This is good if you know that f
 * is small, but can be exponential if f is large.
 */
#include <iostream>
#include <algorithm>
#include <vector>
#include <list>
#include <cassert>
using namespace std;
struct Edge:
typedef list<Edge>::iterator EdgeIter;
struct Edge {
 int to:
 int cap;
 int flow:
 bool is_real;
 EdgeIter partner;
 Edge(int t, int c, bool real = true)
   : to(t), cap(c), flow(0), is_real(real)
 { };
 int residual() const
    return cap - flow;
};
struct Graph
 list < Edge > *nbr;
 int num_nodes;
 Graph (int n)
   : num_nodes(n)
    nbr = new list<Edge>[num_nodes];
```

```
networkflow.cc
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                                                                        Page 2/3
  ~Graph()
    delete[] nbr;
  // note: this routine adds an edge to the graph with the specified capacity,
  // as well as a residual edge. There is no check on duplicate edge, so it
  // is possible to add multiple edges (and residual edges) between two
  // vertices
  void add_edge(int u, int v, int cap)
    nbr[u].push_front(Edge(v, cap));
    nbr[v].push_front(Edge(u, 0, false));
    nbr[v].begin()->partner = nbr[u].begin();
    nbr[u].begin()->partner = nbr[v].begin();
};
void push_path(Graph &G, int s, int t, const vector<EdgeIter> &path, int flow)
  for (int i = 0; s != t; i++) {
    if (path[i]->is_real) {
      path[i]->flow += flow;
      path[i]->partner->cap += flow;
    } else
      path[i]->cap -= flow;
      path[i]->partner->flow -= flow;
    s = path[i] -> to;
// the path is stored in a peculiar way for efficiency: path[i] is the
// i-th edge taken in the path.
int augmenting_path(const Graph &G, int s, int t, vector<EdgeIter> &path,
                    vector<br/>bool> &visited, int step = 0)
  if (s == t) {
    return -1;
  for (EdgeIter it = G.nbr[s].begin(); it != G.nbr[s].end(); ++it) {
    int v = it -> to;
    if (it->residual() > 0 && !visited[v]) {
      path[step] = it;
      visited[v] = true;
      int flow = augmenting_path(G, v, t, path, visited, step+1);
      if (flow == -1) {
        return it->residual();
      } else if (flow > 0) {
        return min(flow, it->residual());
 return 0;
// note that the graph is modified
int network_flow(Graph &G, int s, int t)
  vector<bool> visited(G.num_nodes);
  vector<EdgeIter> path(G.num_nodes);
  int flow = 0, f;
    fill(visited.begin(), visited.end(), false);
    if ((f = augmenting_path(G, s, t, path, visited)) > 0) {
      push_path(G, s, t, path, f);
```

```
networkflow.cc
                                                                               Page 3/3
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      flow += f;
  while (f > 0);
  return flow;
int main(void)
 Graph G(100);
 int's, t, u, v, cap, flow;
  cin >> s >> t;
 while (cin >> u >> v >> cap) {
   G.add_edge(u, v, cap);
 flow = network_flow(G, s, t);
cout << "maximum flow = " << flow << endl;</pre>
  return 0;
```

```
pointpoly.cc
 Jan 16, 14 15:11
                                                                          Page 1/2
 * Point-in-polygon test
 * Author: Howard Cheng
 * Reference:
    http://www.exaflop.org/docs/cgafaq/cga2.html
 * Given a polygon as a list of n vertices, and a point, it returns
 ^{\star} whether the point is in the polygon or not.
 * One has the option to define the behavior on the boundary.
 */
#include <iostream>
#include <cmath>
#include <cassert>
using namespace std;
/* how close to call equal */
const double EPSILON = 1E-8;
/* what should be returned on the boundary? */
const bool BOUNDARY = true;
struct Point {
 double x, y;
};
/* counterclockwise, clockwise, or undefined */
enum Orientation {CCW, CW, CNEITHER};
Orientation ccw(Point a, Point b, Point c)
 double dx1 = b.x - a.x;
 double dx2 = c.x - b.x;
 double dy1 = b.y - a.y;
 double dy2 = c.y - b.y;
 double t\bar{1} = dy2^* dx1;
 double t2 = dy1 * dx2;
 if (fabs(t1 - t2) < EPSILON) {</pre>
    if (dx1 * dx2 < 0 | | dy1 * dy2 < 0) {
      if (dx1*dx1 + dy1*dy1 >= dx2*dx2 + dy2*dy2 - EPSILON) {
        return CNEITHER;
      } else {
        return CW;
    } else {
      return CCW;
 } else if (t1 > t2) {
    return CCW;
 } else {
    return CW;
bool point_in_poly(Point poly[], int n, Point p)
 int i, j, c = 0;
  /* first check to see if point is one of the vertices */
 for (i = 0; i < n; i++) {
    if (fabs(p.x - poly[i].x) < EPSILON && fabs(p.y - poly[i].y) < EPSILON) {</pre>
      return BOUNDARY;
```

```
pointpoly.cc
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                                                                        Page 2/2
 /* now check if it's on the boundary */
 for (i = 0; i < n-1; i++)
   if (ccw(poly[i], poly[i+1], p) == CNEITHER) {
     return BOUNDARY;
 if (ccw(poly[n-1], poly[0], p) == CNEITHER) {
   return BOUNDARY;
 /* finally check if it's inside */
 for (i = 0, j = n-1; i < n; j = i++)
   if (((poly[i].y <= p.y && p.y < poly[j].y) ||</pre>
       (poly[j].y <= p.y && p.y < poly[i].y)) &&
       (p.x < (poly[j].x - poly[i].x) * (p.y - poly[i].y)
         / (poly[j].y - poly[i].y) + poly[i].x))
     c = !c;
 return c;
int main(void)
 Point *polygon, p;
 int n;
 int i;
 while (cin >> n && n > 0)
   polygon = new Point[n];
   assert (polygon);
   for (i = 0; i < n; i++) {
     cin >> polygon[i].x >> polygon[i].y;
   while (cin >> p.x >> p.y) {
     if (point_in_poly(polygon, n, p)) {
       cout << "yes";
     } else {
       cout << "no";
     cout << endl;
   delete[] polygon;
 return 0;
```

```
polygon inter.cc
 Jan 16, 14 15:11
                                                                           Page 1/5
 * Convex Polygon Intersection
 * Author: Howard Cheng
 * This routine takes two convex polygon, and returns the intersection
 * which is also convex. If the intersection contains less than
 ^{\star} 3 points, it is considered empty.
#include <iostream>
#include <iomanip>
#include <cmath>
#include <algorithm>
#include <cassert>
using namespace std;
/* how close to call equal */
const double EPSILON = 1E-8;
struct Point {
 double x, y;
};
const bool BOUNDARY = true;
/* counterclockwise, clockwise, or undefined */
enum Orientation {CCW, CW, CNEITHER};
/* Global point for computing convex hull */
Point start_p;
Orientation ccw(Point a, Point b, Point c)
 double dx1 = b.x - a.x;
 double dx2 = c.x - b.x;
 double dy1 = b.y - a.y;
 double dy^2 = c.y - b.y;
double t1 = dy^2 * dx^1;
 double t2 = dy1 * dx2;
 if (fabs(t1 - t2) < EPSILON) {</pre>
    if (dx1 * dx2 < 0 | | dy1 * dy2 < 0) {
      if (dx1*dx1 + dy1*dy1) >= dx2*dx2 + dy2*dy2 - EPSILON) {
        return CNEITHER:
      } else {
        return CW;
    } else {
     return CCW;
 } else if (t1 > t2) {
    return CCW;
 } else {
    return CW;
bool ccw_cmp (const Point &a, const Point &b)
 return ccw(start_p, a, b) == CCW;
int convex_hull(Point polygon[], int n, Point hull[]) {
 int count, best_i, i;
 if (n == 1) {
```

```
polygon inter.cc
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                                                                         Page 2/5
    hull[0] = polygon[0];
    return 1;
  /* find the first point: min y, and then min x */
  start_p = polygon[0];
  best i = 0;
  for (i = 1; i < n; i++) {
   if ((polygon[i].y < start_p.y) ||</pre>
        (polygon[i].y == start_p.y && polygon[i].x < start_p.x)) {</pre>
      start_p = polygon[i];
     best_i = i;
 polygon[best_i] = polygon[0];
 polygon[0] = start_p;
  /* get simple closed polygon */
  sort (polygon+1, polygon+n, ccw_cmp);
  /* do convex hull */
 count = 0;
  hull[count] = polygon[count]; count++;
 hull[count] = polygon[count]; count++;
  for (i = 2; i < n; i++) {
    while (count > 1 &&
           ccw(hull[count-2], hull[count-1], polygon[i]) == CW) {
      /* pop point */
      count --;
   hull[count++] = polygon[i];
 return count;
bool point_in_poly(Point poly[], int n, Point p)
 int i, j, c = 0;
  /* first check to see if point is one of the vertices */
  for (i = 0; i < n; i++) {
   if (fabs(p.x - poly[i].x) < EPSILON && fabs(p.y - poly[i].y) < EPSILON) {</pre>
      return BOUNDARY:
  /* now check if it's on the boundary */
  for (i = 0; i < n-1; i++) {
   if (ccw(poly[i], poly[i+1], p) == CNEITHER) {
      return BOUNDARY;
  if (ccw(poly[n-1], poly[0], p) == CNEITHER) {
    return BOUNDARY;
  /* finally check if it's inside */
  for (i = 0, j = n-1; i < n; j = i++) {
   if (((poly[i].y <= p.y && p.y < poly[j].y) ||</pre>
        (poly[j].y \le p.y \&\& p.y < poly[i].y)) \&\&
        (p.x < (poly[j].x - poly[i].x) * (p.y - poly[i].y)
         / (poly[j].y - poly[i].y) + poly[i].x))
      c = !c;
 return c;
/* returns 1 if intersect at a point, 0 if not, -1 if the lines coincide */
int intersect_line (Point a, Point b, Point c, Point d, Point &p)
```

```
polygon inter.cc
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                                                                        Page 3/5
 double r, s;
 double denom, num1, num2;
 num1 = (a.y - c.y)*(d.x - c.x) - (a.x - c.x)*(d.y - c.y);
 num2 = (a.y - c.y) * (b.x - a.x) - (a.x - c.x) * (b.y - a.y);
 denom = (b.x - a.x)*(d.y - c.y) - (b.y - a.y)*(d.x - c.x);
 if (fabs(denom) >= EPSILON) {
   r = num1 / denom;
    s = num2 / denom;
   if (-EPSILON <= r && r <= 1+EPSILON && -EPSILON <= s && s <= 1+EPSILON) {
     p.x = a.x + r*(b.x - a.x);
     p.y = a.y + r*(b.y - a.y);
     return 1;
    } else
     return 0;
 } else {
    if (fabs(num1) >= EPSILON) {
     return 0;
   } else {
     return -1;
int intersect_polygon(Point poly1[], int n1, Point poly2[], int n2,
                      Point *&out)
 Point *newpoly, p;
 char *used;
 int new_n = n1 + n2 + n1*n2;
 int count, i, i2, j, j2, new_count;
 newpoly = new Point[new_n];
 out = new Point[new n];
 used = new char[new_n];
 assert (newpoly && out && used);
 count = 0;
 fill(used, used+new_n, 0);
 for (i = 0; i < n1; i++) {
   if (point_in_poly(poly2, n2, poly1[i])) {
     newpoly[count++] = poly1[i];
 for (i = 0; i < n2; i++) {
   if (point_in_poly(poly1, n1, poly2[i])) {
     newpoly[count++] = poly2[i];
 for (i = 0; i < n1; i++) {
    i2 = (i+1 == n1) ? 0 : i+1;
    for (j = 0; j < n2; j++) {
     i2 = (j+1 == n2) ? 0 : j+1;
     if (intersect_line(poly1[i], poly1[i2], poly2[j], poly2[j2], p) == 1) {
       newpoly[count++] = p;
 if (count >= 3)
   n = convex_hull(newpoly, count, out);
   if (n < 3) {
     delete[] out;
      n = 0;
```

```
polygon inter.cc
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                                                                         Page 4/5
 } else {
   delete[] out;
   n = 0;
  /* eliminate duplicates */
  for (i = 0; i < n-1; i++) {
    for (j = i+1; j < n; j++) {
     if (out[i].x == out[j].x && out[i].y == out[j].y) {
        used[i] = 1;
 j = 0;
 new\_count = 0;
 for (i = 0; i < n; i++) {
   if (!used[i])
     out[new_count++] = out[i];
 n = new_count;
 delete[] newpoly;
 delete[] used;
 return n;
int read_poly(Point *&poly)
 int n, i;
 cin >> n;
 if (n == 0) {
   return 0;
 poly = new Point[n];
 assert (polv);
 for (i = 0; i < n; i++)
   cin >> poly[i].x >> poly[i].y;
 return n;
int main (void)
 Point *poly1, *poly2, *intersection;
 int n1, n2, n3, i;
 while ((n1 = read_poly(poly1))) {
   n2 = read_poly(poly2);
   n3 = intersect_polygon(poly1, n1, poly2, n2, intersection);
   delete[] poly1;
   delete[] poly2;
    if (n3 >= 3) {
     for (i = 0; i < n3; i++) {
        cout << fixed << setprecision(2);</pre>
        cout << "(" << intersection[i].x << "," << intersection[i].y</pre>
             << ") ";
     cout << endl;
     delete[] intersection;
    } else ⊦
     cout << "Empty Intersection" << endl;
 return 0;
```

Jan 16, 14 15:11	polygon_inter.cc	Page 5/5
}		

```
ratlinsolve.cc
 Jan 16, 14 15:08
// Performs guassian elimination over the rationals.
// Author: Darcy Best
// Date : September 22, 2010
// pair<int, int> means first = numerator, second = denominator
#include <iostream>
#include <iomanip>
#include <cstdlib>
using namespace std;
#define pii pair<int, int>
const int MAX N = 100;
pii *r_m, m_m;
void print(pii x) {
 if(x.second == 1)
    cout << x.first;
 else
    cout << x.first << "/" << x.second;
void print(pii A[MAX_N][MAX_N], int m, int n) {
 for(int i=0;i<m;i++){
    for(int j=0; j<n; j++) {</pre>
     cout << setw(5);
     print(A[i][j]);
    cout << endl;
 cout << endl;
void read(pii& x) {
 cin >> x.first;
 char ch;
 if(cin.peek() == '/')
   cin >> ch >> x.second;
 else
    x.second = 1;
int gcd(int a, int b) {
 while (b) {
    int r = a % b;
   a = b;
   b = r;
 return a;
pii reduce(pii a) {
 if(a.first == 0){
    a.second = 1;
 } else {
    if(a.second < 0) {</pre>
     a.first *=-1;
      a.second *=-1;
    int g = gcd(abs(a.first),a.second);
    a.first /= g;
    a.second /= q;
 return a;
pii operator* (pii a, pii b) {
```

```
ratlinsolve.cc
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                                                                             Page 2/3
  return reduce(pii(a.first*b.first,a.second*b.second));
pii operator+(pii a,pii b) {
 return reduce(pii(a.first*b.second+b.first*a.second,a.second*b.second));
void multRow(pii& x) {
x = x * m_m;
void addMultRow(pii& x) {
x = x + (m_m * (*r_m++));
int rowReduction(pii A[MAX_N][MAX_N], int rows, int cols) {
  int rank = 0;
  for(int c=0; c<cols; c++) {
    for(int r=rank; r<rows; r++) {</pre>
      if(A[r][c].first){
        if(r != rank) // Swap rows
          swap_ranges(A[rank],A[rank]+cols,A[r]);
        if(c == cols-1) // Inconsistent
          return -1;
        // Make first entry 1
        m_m = pii(A[rank][c].second, A[rank][c].first);
        for_each(A[rank]+c+1,A[rank]+cols,multRow);
        A[rank][c] = pii(1,0);
        for (int i=(arb?rank+1:0);i<rows;i++)</pre>
          if(A[i][c].first && i != rank) {
             // Make the other rows 0
            m_m = pii(-A[i][c].first,A[i][c].second);
            r_m = A[rank] + c + 1;
             for_each(A[i]+c+1,A[i]+cols,addMultRow);
            A[i][c] = pii(0,1);
        rank++;
        break:
  return rank;
int main() {
 int C=0;
 int T, m, n, rank;
  pii A[MAX_N][MAX_N];
  while (cin >> T && T) {
    if(C++)
      cout << endl;
    cout << "Solution for Matrix System # " << T << endl;
    cin >> n >> m;
    for(int i=0; i<m; i++)
      for(int j=0; j<=n; j++)
        read(A[i][j]);
    if((rank = rowReduction(A, m, n+1)) < 0){
      cout << "No Solution." << endl;
      if(rank != n) {
        cout << "Infinitely many solutions containing " << n-rank << " arbitrary constants." << endl
        for(int i=0;i<n;i++){</pre>
          cout << "x[" << i+1 << "] = "; print(A[i][n]); cout << endl;</pre>
```

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```
ratlinsolve.cc
                                                                        Page 3/3
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return 0;
                                                                           src/ratlinsolve.cc
```

```
roman numerals.cc
 Jan 16, 14 15:11
                                                                         Page 1/1
// Converts Roman Numerals to Arabic Numbers (and vice versa)
// Author: Darcy Best
// Date : September 5, 2010
// If you are given a valid integer (0 < x < 4000), then it will give
     the standard roman numeral representation of it. Note that if you give
      it a number such that x \ge 4000, then it will just append as many "M"s
      as needed.
// If you are given a valid roman numberal, then it will give you the answer
      as a base 10 number.
#include <iostream>
#include <string>
#include <map>
using namespace std;
const string Roman[13] = {"M", "CM", "D", "CD", "C", "XC", "L", "XL", "X", "IX", "V", "I
V", "I"};
const int Arabic[13] = {1000,900,500,400,100,90,50,40,10,9,5,4,1};
string toRoman(int x) {
 string roman;
 for (int i=0; i<13; i++)
    while(x >= Arabic[i]){
     x -= Arabic[i];
     roman += Roman[i];
 return roman;
int toInt(string s){
 int L1, L2, ind=0, ans=0;
 while(ind < 13) {
   L1 = s.length();
    L2 = Roman[ind].length();
    if(s.substr(0,min(L1,L2)) == Roman[ind]) {
      ans += Arabic[ind];
      s.erase(0,min(L1,L2));
    } else {
     ind++;
 return ans;
int main() {
 char c;
 int x;
 string s;
  // Checks to see if the line is Roman Numerals or Arabic Numbers,
  // then converts to the opposite.
 while(cin >> c) {
    cin.putback(c);
   if(c >= '0' && c <= '9') {
     cin >> x;
      cout << toRoman(x) << endl;</pre>
    } else {
     cin >> s;
      cout << toInt(s) << endl;</pre>
 return 0;
```

```
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                                        SCC.CC
                                                                         Page 1/2
// Compresses a directed graph into its strongly connected components
// Author: Darcy Best
// Date : October 1, 2010
// A set of nodes is "strongly connected" if for any pair of nodes in
// the set, there is a path from u to v AND from v to u.
// Compressing a graph into its strongly connected components means
// converting each strongly connected component into a super-node.
// We then build a "compressed" graph made with the super-nodes. We
// add an edge in the compressed graph between U and V if there is a
// vertex u in U and v in V such that there was an edge from u to v in
// the original graph. The compressed graph will be a Directed Acyclic
// Graph (DAG), and the list of components will be in REVERSE
// topological order.
// If you are only concerned with the number of strongly connected
// components, you do not need to build the graph. See comments below
// on how to remove the SCC graph.
// The complexity of this algorithm is O(|V| + |E|).
#include <iostream>
#include <algorithm>
#include <stack>
#include <cassert>
#include <vector>
using namespace std;
const int MAX_NODES = 100005;
struct Graph{
 int numNodes;
 vector<int> adj[MAX_NODES];
 void clear() {
    numNodes = 0;
    for(int i=0;i<MAX_NODES;i++)</pre>
      adj[i].clear();
 void add_edge(int u,int v) {
    if(find(adj[u].begin(),adj[u].end(),v) == adj[u].end())
      adj[u].push_back(v);
};
int po[MAX_NODES], comp[MAX_NODES];
void DFS (int v, const Graph& G, Graph& G_scc, int& C,
         stack<int>& P, stack<int>& S) {
 po[v] = C++;
 S.push(v); P.push(v);
 for (unsigned int i=0; i < G.adj[v].size(); i++) {</pre>
    int w = G.adj[v][i];
    if (po[w] == -1) {
      DFS(w,G,G_scc,C,P,S);
    } else if(comp[w] == -1){
      while(!P.empty() && (po[P.top()] > po[w]))
       P.pop();
 if(!P.empty() && P.top() == v){
    while(!S.empty()){
      int t = S.top();
      S.pop();
      comp[t] = G_scc.numNodes;
```

```
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                                           SCC.CC
                                                                              Page 2/2
      if(t == v)
         break;
    G scc.numNodes++;
    P.pop();
int SCC(const Graph& G, Graph& G_scc) {
  G_scc.clear();
  int C=1;
  stack<int> P,S;
  fill(po,po+G.numNodes,-1);
  fill(comp,comp+G.numNodes,-1);
  for (int i=0; i < G.numNodes; i++)</pre>
    if (po[i] == -1)
      DFS(i,G,G_scc,C,P,S);
  // You do not need this if you are only interested in the number of
  // strongly connected components.
  for (int i=0; i < G.numNodes; i++) {</pre>
    for(unsigned int j=0; j<G.adj[i].size(); j++) {</pre>
      int w = G.adj[i][j];
      if(comp[i] != comp[w])
         G_scc.add_edge(comp[i],comp[w]);
  return G_scc.numNodes;
// Declare these as a global variable if MAX_NODES is large to
// avoid Runtime Error.
Graph G, G_scc;
int main(){
  int u, v, m, n;
  int n scc;
  while(cin >> n >> m && (n || m)){
    G.clear();
    G.numNodes = n;
    for (int i=0; i < m; i++) {</pre>
      cin >> u >> v;
      G.add_edge(u,v);
    n\_scc = SCC(G, G\_scc);
    cout << "# of Strongly Connected Components: " << n_scc << endl;</pre>
  return 0;
```

```
simplex.cc
 Jan 16, 14 15:08
                                                                         Page 1/2
#include <algorithm>
using namespace std;
const int MAX_CONSTRAINTS = 100;
const int MAX_VARS = 100;
const int MAXM = MAX CONSTRAINTS + 1;
const int MAXN = MAX_VARS + 1;
const double EPS = 1e-9;
const double INF = 1.0/0.0;
double A[MAXM][MAXN];
int basis[MAXM], out[MAXN];
void pivot(int m, int n, int a, int b)
 int i, j;
 for (i = 0; i <= m; i++)
   if (i != a)
      for (j = 0; j <= n; j++)
        if (j != b)
          A[i][j] -= A[a][j] * A[i][b] / A[a][b];
 for (j = 0; j \le n; j++)
   if (j != b) A[a][j] /= A[a][b];
 for (i = 0; i <= m; i++)
   if (i != a) A[i][b] = -A[i][b] / A[a][b];
 A[a][b] = 1 / A[a][b];
 swap(basis[a], out[b]);
double simplex(int m, int n, double C[][MAXN], double X[])
 int i, j, ii, jj;
 for (i = 1; i <= m; i++)
    copy(C[i], C[i]+n+1, A[i]);
 for (j = 0; j <= n; j++)
   A[0][j] = -C[0][j];
 for (i = 0; i <= m; i++)
   basis[i] = -i;
 for (j = 0; j <= n; j++)
   out[j] = j;
 for (;;) {
    for (i = ii = 1; i <= m; i++)
      if (A[i][n] < A[ii][n] || (A[i][n] == A[ii][n] && basis[i] < basis[ii]))</pre>
       ii = i;
    if (A[ii][n] >= -EPS) break;
    for (j = jj = 0; j < n; j++)
      if (A[ii][j] < A[ii][jj] - EPS ||
          (A[ii][j] < A[ii][jj] - EPS && out[i] < out[j]))
        jj = j;
    if (A[ii][jj] >= -EPS) return -INF;
   pivot(m, n, ii, jj);
 for (;;) {
    for (j = jj = 0; j < n; j++)
      if (A[0][j] < A[0][jj] || (A[0][j] == A[0][jj] && out[j] < out[jj]))</pre>
    if (A[0][jj] > -EPS) break;
    for (i=1, ii=0; i <= m; i++)
      if ((A[i][jj]>EPS) &&
          (!ii || (A[i][n]/A[i][jj] < A[ii][n]/A[ii][jj]-EPS) ||
           ((A[i][n]/A[i][jj] < A[ii][n]/A[ii][jj]+EPS) &&
            (basis[i] < basis[ii]))))
        ii = i;
    if (A[ii][jj] <= EPS) return INF;</pre>
    pivot(m, n, ii, jj);
 fill(X, X+n, 0);
```

```
Printed by Howard Cheng
                                      simplex.cc
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                                                                        Page 2/2
  for (i = 1; i <= m; i++)
    if (basis[i] >= 0)
     X[basis[i]] = A[i][n];
  return A[0][n];
#include <iostream>
#include <iomanip>
int main (void)
  double C[MAXM] [MAXN], X[MAX_VARS];
                C[0][1] = -3; C[0][2] = 0; C[0][3] = 0;
  C[0][0] = -1;
  C[1][0] = -2; C[1][1] = -3; C[1][2] = -6; C[1][3] = 250;
                 C[2][1] = -5; C[2][2] = -5; C[2][3] = 400;
  C[2][0] = -1;
  C[3][0] = -59; C[3][1] = -35; C[3][2] = -160; C[3][3] = 30;
  double val = simplex(2, 2, C, X);
  cout << fixed << setprecision(3);</pre>
  cout << "val = " << val << endl;
  cout << "X[0] = " << X[0] << endl;
  cout << "X[1] = " << X[1] << endl;
  // cout << "X[2] = " << X[2] << endl;
  return 0;
```

```
str rotation period.cc
 Jan 16, 14 15:08
                                                                         Page 1/1
 \mbox{*} Finding the lexicographically least rotation of a string, and finding
 * the smallest period of a string.
  Author: Sumudu Fernando
 * Given a string, the algorithm can be used to compute two things:
     a) the position at which the lexicographically least rotation starts.
       If there are ties, give the first position.
    b) the length of the shortest substring such that the original string
       is a concatenation of copies of that substring
 * Complexity: O(n) where n = length of the string
 * Tested on: 719
                             Glass Beads
              10298
                              Power Strings
              ACPC 2011 H Let's call a SPaDE a SPaDE
#include <iostream>
#include <string>
#include <algorithm>
using namespace std;
// pos = position of the start of the lexicographically least rotation
// period = the period
void compute (string s, int &pos, int &period)
 s += s;
 int len = s.length();
 int i = 0, j = 1;
 for (int k = 0; i+k < len && j+k < len; k++) {
   if (s[i+k] > s[j+k]) {
      i = \max(i+k+1, j+1);
      k = -1:
   } else if (s[i+k] < s[j+k]) {</pre>
      j = \max(j+k+1, i+1);
      \bar{k} = -1;
 pos = min(i, j);
 period = (i > j) ? i - j : j - i;
int main (void)
 string s;
 while (cin >> s) {
   int pos, period;
    compute(s, pos, period);
   int n = s.length();
   s += s;
   cout << "least rotation = " << s.substr(pos, n) << endl;</pre>
    cout << "period = " << s.substr(0, period) << endl;</pre>
 return 0;
```

```
suffixarray.cc
 Jan 16, 14 15:08
                                                                       Page 1/4
 * Suffix array
 * Author: Howard Cheng
 * References:
    Manber, U. and Myers, G. "Suffix Arrays: a New Method for On-line
    String Searches." SIAM Journal on Computing. 22(5) p. 935-948, 1993.
     T. Kasai, G. Lee, H. Arimura, S. Arikawa, and K. Park. "Linear-time
     Longest-common-prefix Computation in Suffix Arrays and Its
     Applications." Proc. 12th Annual Conference on Combinatorial
    Pattern Matching, LNCS 2089, p. 181-192, 2001
    J. Kärkkäinen and P. Sanders. Simple linear work suffix array
    construction. In Proc. 13th International Conference on Automata,
    Languages and Programming, Springer, 2003
 * The build_sarray routine takes in a string str of n characters (null-
 * terminated), and construct an array sarray. Optionally, you can also
 * construct an lcp array from the sarray computed. The properties
 * - If p = sarray[i], then the suffix of str starting at p (i.e.
    str[p..n-1] is the i-th suffix when all the suffixes are sorted in
    lexicographical order
    NOTE: the empty suffix is not included in this list, so sarray[0] != n.
 * - lcp[i] contains the length of the longest common prefix of the suffixes
    pointed to by sarray[i-1] and sarray[i]. lcp[0] is defined to be 0.
 \star - To see whether a pattern P occurs in str, you can look for it as
    the prefix of a suffix. This can be done with a binary search in
    O(|P| \log n) time. Call find() to return a pair \langle L, R \rangle such that
    all occurrences of the pattern are at positions sarray[i] with
    L \le i \le R. If L == R then there is no match.
 * The construction of the suffix array takes O(n) time.
 */
#include <iostream>
#include <iomanip>
#include <string>
#include <algorithm>
#include <climits>
using namespace std;
bool leg(int al, int a2, int b1, int b2)
 return(a1 < b1 || a1 == b1 && a2 <= b2);
bool leg(int al, int a2, int a3, int b1, int b2, int b3)
 return(a1 < b1 || a1 == b1 && leg(a2,a3, b2,b3));
void radixPass(int* a, int* b, int* r, int n, int K)
 int* c = new int[K + 1];
 fill(c, c+K+1, 0);
 for (int i = 0; i < n; i++) c[r[a[i]]]++;</pre>
 for (int i = 0, sum = 0; i <= K; i++) {
     int t = c[i]; c[i] = sum; sum += t;
 for (int i = 0; i < n; i++) b[c[r[a[i]]]++] = a[i];
 delete [] c;
```

```
suffixarray.cc
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                                                                       Page 2/4
#define GetI() (SA12[t] < n0 ? SA12[t] * 3 + 1 : (SA12[t] - n0) * 3 + 2)
void sarray_int(int* s, int* SA, int n, int K) {
 int n0=(n+2)/3, n1=(n+1)/3, n2=n/3, n02=n0+n2;
 int* s12 = new int[n02 + 3]; s12[n02] = s12[n02+1] = s12[n02+2] = 0;
 int* SA12 = new int[n02 + 3]; SA12[n02]=SA12[n02+1]=SA12[n02+2]=0;
 int* s0 = new int[n0];
 int* SA0 = new int[n0];
 for (int i=0, j=0; i < n+(n0-n1); i++) if (i%3 != 0) s12[j++] = i;
 radixPass(s12 , SA12, s+2, n02, K);
 radixPass(SA12, s12, s+1, n02, K);
 radixPass(s12 , SA12, s , n02, K);
 int name = 0, c0 = -1, c1 = -1, c2 = -1;
 for (int i = 0; i < n02; i++) {
   if (s[SA12[i]] != c0 || s[SA12[i]+1] != c1 || s[SA12[i]+2] != c2) {
     name++; c0 = s[SA12[i]]; c1 = s[SA12[i]+1]; c2 = s[SA12[i]+2];
   if (SA12[i] % 3 == 1) { s12[SA12[i]/3]
                          \{ s12[SA12[i]/3 + n0] = name; \}
   else
 if (name < n02)
   sarray_int(s12, SA12, n02, name);
   for (int i = 0; i < n02; i++) s12[SA12[i]] = i + 1;</pre>
 } else
   for (int i = 0; i < n02; i++) SA12[s12[i] - 1] = i;
 for (int i=0, j=0; i < n02; i++) if (SA12[i] < n0) s0[j++] = 3*SA12[i];
 radixPass(s0, SAO, s, n0, K);
 for (int p=0, t=n0-n1, k=0; k < n; k++) {
   int i = GetI();
   int j = SA0[p];
   if (SA12[t] < n0 ?
       leq(s[i],
                       s12[SA12[t] + n0], s[j],
                                                      s12[j/3]:
       leg(s[i], s[i+1], s12[SA12[t]-n0+1], s[j], s[j+1], s12[j/3+n0]))
     SA[k] = i; t++;
     if (t == n02) {
       for (k++; p < n0; p++, k++) SA[k] = SA0[p];
   } else {
     SA[k] = j; p++;
     if (p == n0) {
       for (k++; t < n02; t++, k++) SA[k] = GetI();
 delete [] s12; delete [] SA12; delete [] SA0; delete [] s0;
void build_sarray(string str, int sarray[])
 int n = str.length();
 if (n <= 1) {
   for (int i = 0; i < n; i++) {
     sarray[i] = i;
   return;
 int *s = new int[n+3];
 int *SA = new int[n+3];
```

```
suffixarray.cc
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                                                                         Page 3/4
 for (int i = 0; i < n; i++)
    s[i] = (int)str[i] - CHAR_MIN + 1;
 s[n] = s[n+1] = s[n+2] = 0;
 sarray_int(s, SA, n, 256);
 copy(SA, SA+n, sarray);
 delete[] s;
 delete[] SA;
void compute_lcp(string str, int sarray[], int lcp[])
 int n = str.length();
 int *rank = new int[n];
 for (int i = 0; i < n; i++) {</pre>
   rank[sarray[i]] = i;
 int h = 0;
 for (int i = 0; i < n; i++) {
    int k = rank[i];
    if (k == 0)
     lcp[k] = -1;
    } else {
      int j = sarray[k-1];
      while (i + h < n \&\& j + h < n \&\& str[i+h] == str[j+h]) {
       h++;
     lcp[k] = h;
    if (h > 0) {
     h--;
 lcp[0] = 0;
 delete[] rank;
pair<int, int> find(const string &str, const int sarray[],
                   const string &pattern)
 int n = str.length(), p = pattern.length();
 int L, R;
 if (pattern <= str.substr(sarray[0], p)) {</pre>
   L = 0;
 } else if (pattern > str.substr(sarray[n-1], p)) {
   L = n;
 } else {
    int lo = 0, hi = n-1;
   while (hi - lo > 1) {
     int mid = lo + (hi - lo)/2;
     if (pattern <= str.substr(sarray[mid], p)) {</pre>
       hi = mid;
     } else {
       lo = mid;
    L = hi;
 if (pattern < str.substr(sarray[0], p)) {</pre>
   R = 0;
 } else if (pattern >= str.substr(sarray[n-1], p)) {
    R = n;
 } else {
    int lo = 0, hi = n-1;
    while (hi - lo > 1) {
      int mid = lo + (hi - lo)/2;
```

```
suffixarray.cc
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                                                                                Page 4/4
      if (pattern < str.substr(sarray[mid], p)) {</pre>
        hi = mid;
      } else {
         lo = mid;
    R = hi;
  if (L > R) R = L;
  return make_pair(L, R);
int main (void)
 string str;
  int sarray[100], lcp[100];
  unsigned int i;
  while (cin >> str) {
   build_sarray(str, sarray);
    compute_lcp(str, sarray, lcp);
    for (i = 0; i < str.length(); i++) {
  cout << setw(3) << i << ":" << setw(2) << lcp[i] << ","</pre>
            << str.substr(sarray[i], str.length()-sarray[i]) << endl;
 return 0;
```

```
top sort.cc
 Jan 16, 14 15:08
                                                                         Page 1/2
 * Topological sort
 * Author: Howard Cheng
 * Given a directed acyclic graph, the topological_sort routine
 * returns a vector of integers that gives the vertex number (0 to n-1)
 * such that if there is a path from v1 to v2, then v1 occurs earlier
 ^{\star} than v2 in the order. Note that the topological sort result is not
 * necessarily unique.
  topological_sort returns true if there is no cycle. Otherwise it
 * returns false and the sorting is unsuccessful.
 * The complexity is O(n + m).
#include <iostream>
#include <algorithm>
#include <vector>
#include <queue>
using namespace std;
typedef int Edge;
typedef vector<Edge>::iterator EdgeIter;
struct Graph {
 vector<Edge> *nbr;
 int num_nodes;
 Graph (int n)
    : num_nodes(n)
    nbr = new vector<Edge>[num_nodes];
  ~Graph()
    delete[] nbr;
 // note: There is no check on duplicate edge, so it is possible to
  // add multiple edges between two vertices
 void add_edge(int u, int v)
    nbr[u].push_back(Edge(v));
};
bool topological_sort(const Graph &G, vector<int> &order)
 vector<int> indeg(G.num_nodes);
 fill(indeg.begin(), indeg.end(), 0);
 for (int i = 0; i < G.num_nodes; i++) {</pre>
    for (int j = 0; j < G.nbr[i].size(); j++) {</pre>
      indeg[G.nbr[i][j]]++;
 // use a priority queue if you want to get a topological sort order
  // with ties broken by lexicographical ordering
  queue<int> q;
 for (int i = 0; i < G.num_nodes; i++) {</pre>
    if (indeg[i] == 0) {
      q.push(i);
```

```
top sort.cc
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                                                                           Page 2/2
  order.clear();
  while (!q.empty())
    int v = q.front();
    q.pop();
    order.push_back(v);
    for (int i = 0; i < G.nbr[v].size(); i++) {</pre>
      if (--indeg[G.nbr[v][i]] == 0) {
        q.push(G.nbr[v][i]);
  return order.size() == G.num_nodes;
int main (void)
 int n, m;
 while (cin >> n >> m && (n || m)) {
    Graph G(n);
    for (int i = 0; i < m; i++) {
      int u, v;
      cin >> u >> v;
      G.add_edge(u, v);
    vector<int> order;
    if (topological_sort(G, order)) {
      for (int i = 0; i < n; i++) {
        if (i) cout << '';
        cout << order[i];</pre>
      cout << endl;
    } else
      cout << "there is a cycle" << endl;</pre>
 return 0:
```

```
unionfind.cc
 Jan 16, 14 15:11
                                                                       Page 1/1
// UnionFind class -- based on Howard Cheng's C code for UnionFind
// Modified to use C++ by Rex Forsyth, Oct 22, 2003
// Constuctor -- builds a UnionFind object of size n and initializes it
// find -- return index of x in the UnionFind
// merge -- updates relationship between x and y in the UnionFind
class UnionFind
      struct UF { int p; int rank; };
   public:
     UnionFind(int n) {
                                 // constructor
        howMany = n;
        uf = new UF[howMany];
         for (int i = 0; i < howMany; i++) {
           uf[i].p = i;
           uf[i].rank = 0;
      ~UnionFind() {
         delete[] uf;
      int find(int x) { return find(uf,x); }
                                                   // for client use
      bool merge(int x, int y) {
         int res1, res2;
        res1 = find(uf, x);
         res2 = find(uf, y);
         if (res1 != res2) {
            if (uf[res1].rank > uf[res2].rank) {
               uf[res2].p = res1;
           else
               uf[res1].p = res2;
               if (uf[res1].rank == uf[res2].rank) {
                 uf[res2].rank++;
           return true;
         return false;
   private:
      int howMany;
     UF* uf;
      int find(UF uf[], int x) { // recursive funcion for internal use
         if (uf[x].p != x) {
           uf[x].p = find(uf, uf[x].p);
         return uf[x].p;
};
```

```
vecsum.cc
 Jan 16, 14 15:11
                                                                          Page 1/2
 * Largest subvector sum
 * Author: Howard Cheng
 * Reference: Programming Pearl, page 74
 * Given an array of integers, we find the continguous subvector that
 * gives the maximum sum. If all entries are negative, it returns
 * an empty vector with sum = 0.
 * If we want the subvector to be nonempty, we should first scan for the
 * largest element in the vector (1-element subvector) and combine the
 * result in this routine.
 ^{\star} The sum is returned, as well as the start and the end position
 * (inclusive). If start > end, then the subvector is empty.
 */
#include <iostream>
#include <cassert>
using namespace std;
int vecsum(int v[], int n, int &start, int &end)
 int maxval = 0;
 int \max end = 0;
 int max_end_start, max_end_end;
 int i;
 start = max_end_start = 0;
 end = max_end_end = -1;
 for (i = 0; i < n; i++)
   if (v[i] + max_end >= 0)
      max\_end = v[i] + max\_end;
      max\_end\_end = i;
    } else {
      max_end_start = i+1;
     max\_end\_end = -1;
      max_end = 0;
    if (maxval < max_end) {</pre>
     start = max_end_start;
      end = max_end_end;
      maxval = max_end;
    } else if (maxval == max_end) {
      /* put whatever preferences we have for a tie */
      /* ^{st} eg. longest ^{st} subvector, and then the one that starts the earliest ^{st}/
      if (max_end_end - max_end_start > end - start | |
          (max_end_end - max_end_start == end - start &&
          max_end_start < start)) {</pre>
        start = max_end_start;
       end = max_end_end;
       maxval = max_end;
 return maxval;
int main (void)
 int n;
 int *v;
 int i;
 int sum, start, end;
```

```
Jan 16, 14 15:08
                                      zero one.c
                                                                         Page 1/4
 * Zero-one programming
 * Author: Howard Cheng
 * Reference:
    http://www.cs.sunysb.edu/~algorith/implement/syslo/distrib/processed/
 * This algorithm is based on BALAS branching testing.
  This algorithm solves the BINARY linear program:
                                [cost function]
        min cx
        s.t.
             Ax \le b
                                [constraints]
             x[i] = 0 \text{ or } 1.
  where A is an m x n matrix,
        c and x are n-dimensional vectors,
         b is an m-dimensional vector.
 * n = number of variables
 * m = number of constraints
 * It returns whether there exists a solution.
 * The optimal value of the cost function is returned in value.
 ^{\star} The assignment giving the optimal cost function value is returned in x.
 * Important Notes:
 ^{\star} 1. The matrices and arrays start their indices at 1!!!!!!
 * 2. If we want to have constraints that are >=, just multiply all the
    coefficients by -1.
 * 3. If we want to \hat{h} ave constraints that are ==, do both >= and <=.
 ^{\star} 4. The content of A, b, and c is preserved after this routine.
 * 5. The coefficients in the cost vector c must be positive. If not,
      make a change of variable x'[i] = 1-x[i] and adjust all constraints
      as well as the returned optimal value. This is especially useful
      if you wish to maximize the cost function.
 */
#include <stdio.h>
#include <limits.h>
#include <assert.h>
#define MAX_VAR 1000
#define MAX CONS 100
#define MAX_ROWS MAX_CONS+1
#define MAX_COLS MAX_VAR+1
int zero_one(int A[MAX_ROWS][MAX_COLS], int *b, int *c, int n, int m,
             int *val, int *x)
 int exist;
 int alpha, beta, gamma, i, j, mnr, nr;
 int p, r, r1, r2, s, t, z;
 int y[MAX_ROWS], w[MAX_ROWS], zr[MAX_ROWS];
 int ii[MAX_COLS], jj[MAX_COLS], xx[MAX_COLS];
 int kk[MAX_COLS+1];
 for (i = 1; i <= m; i++) {
   y[i] = b[i];
 z = 1;
 for (j = 1; j <= n; j++) {
   xx[j] = 0;
   z += c[j];
 *val = z+z;
```

```
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                                      zero one.c
                                                                          Page 2/4
 s = t = z = exist = 0;
 kk[1] = 0;
m10:
 p = mnr = 0;
 for (i = 1; i <= m; i++) {
   if ((r = y[i]) < 0) {
     p++;
     gamma = 0;
      alpha = r;
     beta = -INT MAX;
      for (j = 1; j <= n; j++) {
       if (xx[j] <= 0) {
          if (c[i] + z >= *val) {
            xx[j] = 2;
            kk[s+1]++;
            jj[++t] = j;
            else {
            if ((r1 = A[i][j]) < 0) {
              alpha -= r1;
              gamma += c[j];
              if (beta < r1) {
                beta = r1;
      if (alpha < 0) {
        goto m20;
      if (alpha + beta < 0) {</pre>
        if (gamma + z >= *val) {
          goto m20;
        for (j = 1; j \le n; j++) {
         r1 = A[i][j];
          r2 = xx[j];
          if (r1 < 0) {
            if (!r2) {
              xx[j] = -2;
              for (nr = 1; nr <= mnr; nr++) {</pre>
                zr[nr] -= A[w[nr]][j];
                if (zr[nr] < 0) {
                  goto m20;
          } else {
            if (r2 < 0) {
              alpha -= r1;
              if (alpha < 0) {
                goto m20;
              gamma += c[j];
              if (gamma + z >= *val) {
                goto m20;
        mnr++;
        w[mnr] = i;
        zr[mnr] = alpha;
 if (!p) {
    *val = z;
```

```
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                                     zero one.c
                                                                        Page 3/4
    exist = 1;
    for (j = 1; j <= n; j++) {
     x[j] = (xx[j] == 1) ? 1 : 0;
    goto m20;
 if (!mnr) {
   p = 0;
    gamma = -INT\_MAX;
    for (j = 1; j <= n; j++) {
     if (!xx[j]) {
       beta = 0;
        for (i = 1; i <= m; i++) {
         r = y[i];
          r1 = A[i][j];
          if (r < r1) {
           beta += r - r1;
       r = c[j];
        if ((beta > gamma) ||
            (beta == gamma && r < alpha)) {
          alpha = r;
          gamma = beta;
         p = j;
   if (!p) {
     goto m20;
    s++;
    kk[s+1] = 0;
    jj[++t] = p;
    ii[s] = xx[p] = 1;
    z += c[p];
    for (i = 1; i <= m; i++) {
     y[i] -= A[i][p];
 } else {
    s++;
    ii[s] = kk[s+1] = 0;
    for (j = 1; j \le n; j++) {
     if (xx[j] < 0) {
        jj[++t] = j;
        ii[s]--;
       z += c[j];
       xx[j] = 1;
       for (i = 1; i <= m; i++) {
         y[i] -= A[i][j];
 goto m10;
m20:
 for (j = 1; j <= n; j++) {
   if (xx[j] < 0) {
     xx[j] = 0;
 if (s > 0) {
   do {
     p = t;
     t -= kk[s+1];
     for (j = t+1; j <= p; j++) {
       xx[jj[j]] = 0;
```

```
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                                        zero one.c
                                                                             Page 4/4
      p = (ii[s] >= 0) ? ii[s] : -ii[s];
      kk[s] += p;
      for (j = t-p+1; j \le t; j++)  {
        p = jj[j];
xx[p] = 2;
        z \stackrel{-}{-} c[p];
        for (i = 1; i <= m; i++) {
         y[i] += A[i][p];
      s--:
      if (ii[s+1] >= 0) {
        goto m10;
    } while (s);
 return exist;
int main (void)
  int A[MAX_ROWS][MAX_COLS];
 int c[MAX_COLS], x[MAX_COLS], b[MAX_ROWS];
 int val, t;
 int m, n, i, j;
 while (scanf("%d%d", &n, &m) == 2 && n > 0 && m > 0) {
    /* read cost function */
    printf("Input cost function to minimize:\n");
    for (i = \hat{1}; i \le n; i++) {
      scanf("%d", &c[i]);
    /* read constraints */
    for (i = 1; i <= m; i++) {
      printf("Input constraint #%d:\n", i);
      for (j = 1; j <= n; j++) {
       scanf("%d", &A[i][j]);
      scanf("%d", &b[i]);
    t = zero_one(A, b, c, n, m, \&val, x);
    if (t) {
      printf("Minimum cost = %d\n", val);
      for (i = 1; i <= n; i++) {
        printf("x[\%2d] = \%2d\n", i, x[i]);
    } else {
      printf("No solution exists.\n");
 return 0;
```