# Xilinx Zynq FPGA, TI DSP, MCU 기반의 프로그래밍 및 회로 설계 전문가 과정 #fork

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# 1.배운내용 복습.

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# SYSCALL\_FORK Driving

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SYSCALL\_DEFINEO(fork)

systemcall 로 정의 되어있는 fork를 찾는다.

\_do\_fork 함수에 인자로 SIGCHLD와 0, 0, NULL, NULL, 0 이 넘어가 받아온 리턴값을 리턴한다. 그러므로 \_do\_fork를 확인한다.

## 2. \_do\_fork

```
long do fork(unsigned long clone flags,
              unsigned long stack start,
              unsigned long stack size,
              int __user *parent_tidptr,
              int user *child tidptr.
              unsigned long tls)
       clone flags = SIGCHLD
       stack start = 0
       stack size = 0
        int_user *parent_tidptr = NULL
        int user *child tiptr = NULL
        tls = 0
       struct task struct *p;
        int trace = 0;
        long nr;
        /* trace 는 디버깅 용도
           task struct *p 는 자식 프로세스
nr 은 자식 프로세스의 pid 가 된다. */
         * Determine whether and which event to report to ptracer. When
         * called from kernel thread or CLONE UNTRACED is explicitly
         * requested, no event is reported; otherwise, report if the event
         * for the type of forking is enabled.
         /*#define CLONE_UNTRACED 0x008000000
if문 값이 0일때 통과이기 때문에 통과 */
        if (!(clone flags & CLONE UNTRACED)) {
                /* #define CLONE VFORK 0x00004000
                SIGCHLD = 17 과 &연산하면 0 이프문 통과 못함 */
                if (clone flags & CLONE VFORK)
                        trace = PTRACE EVENT VFORK;
                else if ((clone flags & CSIGNAL) != SIGCHLD)//여기서도 엘즈 통과함
                        trace = PTRACE EVENT CLONE:
                        /*결국 fork 했을시에 PTRACE EVENT FORK통과한다 clone_flags에 들어오는 인자값에 의해 어떤 동작을 실행시켰는지 검사하는 로직 */
                        trace = PTRACE EVENT FORK;
                /* current = 현재 구동중인 task
                   trace = 1
                if (likely(!ptrace_event_enabled(current, trace)))
                        trace = 0:
        /* trace = 0
           clone flags = SIGCHLD
           stack start = 0
           stack size = 0
           child_tidptr = NULL
           NULL
           tls = 0
        /*복사된 프로세스가 페이 저장*/
        p = copy_process(clone_flags, stack_start, stack_size,
                         child_tidptr, NULL, trace, tls);
         * Do this prior waking up the new thread - the thread pointer
         * might get invalid after that point, if the thread exits quickly.
        if (!IS_ERR(p)) {
```

```
if (!IS ERR(p)) {
        struct completion vfork:
        struct pid *pid:
        trace sched process fork(current, p);
        /*pid 값을 얻어옴 */
        pid = get_task_pid(p, PIDTYPE_PID);
/*자식프로세스의 pid를 nr에 저장*/
        nr = pid vnr(pid);
        if (clone flags & CLONE PARENT SETTID)
                put_user(nr, parent_tidptr);
        if (clone flags & CLONE_VFORK) {
                p->vfork done = &vfork;
                init_completion(&vfork);
                get task struct(p);
        /*run q 에 새로운 task 대기 */
        wake up new task(p);
        /* forking complete and child started to run, tell ptracer */
        if (unlikely(trace))
                ptrace event pid(trace, pid);
        if (clone flags & CLONE VFORK) {
                if (!wait for vfork done(p, &vfork))
                        ptrace_event_pid(PTRACE_EVENT_VFORK DONE, pid);
        }
        put pid(pid);
} else {
        nr = PTR ERR(p);
return nr;
/*pid return */
```

long \_do\_fork(unsigned long clone\_flags, unsigned long stack\_start, unsigned long stack\_size, int \_user \*parent\_tidptr, int \_user \*child\_tidptr, unsigned long tls)

do\_fork 항수는 외안같은 외자를 같는다 이 이자로 널어오 갔들을 나열해보

do\_fork 함수는 위와같은 인자를 갖는다. 이 이자로 넘어온 값들을 나열해보면

clone flags = SIGCHLD
stack\_start = 0
stack\_size = 0
parent\_tidptr = NULL
child\_tidptr = NULL
tls = 0

가 되고 여기서 clone\_flags 에 의해 fork 인지 vfork 인지 clone 인지 디버깅을 거쳐 trace 값이 바뀌고 그 trace 값에 의해 그에 해당하는 작업을 수행시켜준다.

p= copy\_process(clone\_flags, stack\_start, stack\_size, child\_tidptr, NULL, trace, tls) 의 값이 넘어간다.

그럼 copy\_process를 확인한다.

copy\_process 는 복사할 프로세스의 데이터를 저장한다.

밑에 pid = get\_task\_pid(p,PIDTYPE\_PID);

복사할 프로세스 즉 부모 프로세스의 pid 값을 얻어온다.

nr = pid\_vnr(pid) 자식 프로세스의 pid 값을 얻어 nr에 저장한다.

#### - copy\_process

```
/* trace = 0
  clone flags = SIGCHLD
  stack start = 0
  stack size = 0
  child tidptr = NULL
  pid *pid = NULL
  tls = 0
static struct task_struct *copy_process(unsigned long clone_flags,
                                        unsigned long stack_start,
                                        unsigned long stack_size, int user *child tidptr.
                                        struct pid *pid.
                                        int trace.
                                        unsigned long tls)
        int retval:
        struct task_struct *p;
        void *cgrp ss priv[CGROUP CANFORK COUNT] = {};
        /* 정상적인 fork를 했을경우 밑에 있는 에러검출문은 다 통과하지 않는다. */
       /*에러일 경우 에러 내용
새로운 mount namespace로 태스크 생성을 요청했지만 파일 시스템 정보를 공유할 수 없으므로 실패로 함수를 빠져나간다.*/
        if ((clone_flags & (CLONE_NEWNS|CLONE_FS)) == (CLONE_NEWNS|CLONE_FS))
        (**Control ERR_PTR(-EINMAL);
/*에러일 경우 에러 내용
새로운 user namespace로 태스크 생성을 요청했지만 파일 시스템 정보를 공유할 수 없으므로 실패로 함수를 빠져나간다.*/
        if ((clone_flags & (CLONE_NEWUSER|CLONE_FS)) == (CLONE_NEWUSER|CLONE_FS))
                return ERR PTR(-EINVAL);
        * Thread groups must share signals as well, and detached threads
         * can only be started up within the thread group.
        /*thread 생성을 요청했지만 시그널 핸들러의 공유가 필요하다. 그렇지 않은 경우 실패로 함수를 빠져나간다.
         (thread 는 signal 공유 */
        if ((clone flags & CLONE THREAD) && !(clone flags & CLONE SIGHAND))
                return ERR PTR(-EINVAL);
       /*
* Shared signal handlers imply shared VM. By way of the above,
        * thread groups also imply shared VM. Blocking this case allows * for various simplifications in other code.
       /*시그널 핸들러를 공유할 때 같은 가상 주소(VM)를 사용해야한다. 그렇지 않은 경우 실패로 함수를 빠져나간다.*/
if ((clone flags & CLONE SIGHAND) && !(clone flags & CLONE VM))
                return ERR PTR(-EINVAL):
       /*
* Siblings of global init remain as zombies on exit since they are
        * not reaped by their parent (swapper). To solve this and to avoid
* multi-rooted process trees, prevent global and container-inits
         * from creating siblings.
         /* 부모 태스크에 SIGNAL UNKILLABLE 시그널 플래그가 있는 경우 부모 태스크 클론을 요청한 경우 실패로 함수를 빠져나간다.*/
          ((clone_flags & CLONE_PARENT) &&
                                current->signal->flags & SIGNAL UNKILLABLE
                return ERR PTR(-EINVAL);
```

```
if (clone flags & CLONE THREAD) {
               if ((clone flags & (CLONE NEWUSER | CLONE NEWPID)) ||
                   (task active pid ns(current) !=
                               current->nsproxy->pid ns for children))
                       return ERR PTR(-EINVAL):
       /* task ops return 지금 만들어져 아무것도 없기 때문에 값은 0 */
       retval = security task create(clone flags);
       /*retval = 0*/
       if (retval)
               goto fork out:
       retval = -ENOMEM:
       /*current = 현재 구동중인 task */
       p = dup_task_struct(current);
       /* p 에 원래의 task가 복사됨 */
       if (!p)
               goto fork out;
       ftrace_graph_init_task(p);
       rt mutex init task(p):
#ifdef CONFIG PROVE LOCKING
       DEBUG LOCKS WARN ON(!p->hardirgs enabled);
       DEBUG LOCKS WARN ON(!p->softirgs enabled);
#endif
       retval = -EAGAIN:
       if (atomic read(&p->real cred->user->processes) >=
                       task rlimit(p, RLIMIT NPROC)) {
               if (p->real cred->user != INIT USER &&
                   !capable(CAP SYS RESOURCE) && !capable(CAP SYS ADMIN))
                       goto bad fork free;
       /*define PF NPROC EXCEEDED 0x00001000 */
       current->flags &= ~PF NPROC EXCEEDED;
       /*새로운 namespace 할당? */
       retval = copy_creds(p, clone_flags);
       if (retval < 0)
               goto bad fork free;
        * If multiple threads are within copy process(), then this check
        * triggers too late. This doesn't hurt, the check is only there
        * to stop root fork bombs.
        */
       retval = -EAGAIN:
       if (nr threads >= max threads)
               goto bad fork cleanup count;
       delayacct tsk init(p); /* Must remain after dup task struct() */
       p->flags &= ~(PF SUPERPRIV | PF WO WORKER);
       p->flags |= PF FORKNOEXEC:
       INIT LIST HEAD(&p->children);
       INIT LIST_HEAD(&p->sibling);
       rcu copy process(p):
       p->vfork done = NULL;
       spin lock init(&p->alloc lock);
       init sigpending(&p->pending);
      //signal pending
```

```
p->utime = p->stime = p->atime = 0:
        p->utimescaled = p->stimescaled = 0:
        prev cputime init(&p->prev cputime);
#ifdef CONFIG VIRT CPU ACCOUNTING GEN
        sealock init(&p->vtime sealock):
        p->vtime snap = 0;
        p->vtime snap whence = VTIME SLEEPING;
#endif
#if defined(SPLIT RSS COUNTING)
        memset(&p->rss stat, 0, sizeof(p->rss stat));
#endif
        p->default timer slack ns = current->timer slack ns:
        task io accounting init(&p->ioac);
        acct clear integrals(p);
        posix cpu timers init(p);
        p->start time = ktime get ns();
        p->real start time = ktime get boot ns();
        p->io context = NULL;
        p->audit context = NULL;
        threadgroup change begin(current):
        cgroup fork(p);
#ifdef CONFIG NUMA
        p->mempolicy = mpol_dup(p->mempolicy);
        if (IS ERR(p->mempolicy)) {
                retval = PTR ERR(p->mempolicy):
                p->mempolicy = NULL;
                goto bad fork cleanup threadgroup lock;
#endif
#ifdef CONFIG CPUSETS
        p->cpuset mem spread rotor = NUMA NO NODE;
        p->cpuset slab spread rotor = NUMA NO NODE;
        segcount init(&p->mems allowed seg);
#endif
#ifdef CONFIG TRACE IROFLAGS
        p->irq events = 0;
        p->hardirgs enabled = 0;
        p->hardirg enable ip = 0;
        p->hardirg enable event = 0;
        p->hardirq disable ip = THIS IP;
        p->hardirq disable event = 0;
        p->softirgs enabled = 1;
        p->softirg enable ip = THIS IP;
        p->softirq_enable_event = 0;
        p->softirq disable ip = 0;
        p->softirg disable event = 0:
        p->hardirg context = 0;
        p->softirg context = 0;
#endif
        p->pagefault disabled = 0;
#ifdef CONFIG LOCKDEP
       p->lockdep depth = 0; /* no locks held yet */
       p->curr chain key = 0:
```

```
p->lockdep depth = 0: /* no locks held vet */
        p->curr chain kev = 0:
        p->lockdep recursion = 0;
#endif
#ifdef CONFIG DEBUG MUTEXES
        p->blocked on = NULL; /* not blocked yet */
#endif
#ifdef CONFIG BCACHE
        p->sequential io
                                = 0:
        p->sequential io avq
#endif
        /* Perform scheduler related setup. Assign this task to a CPU. */
        retval = sched fork(clone flags, p);
        if (retval)
                goto bad fork cleanup policy:
        retval = perf event init task(p);
        if (retval)
                goto bad fork cleanup policy:
        retval = audit alloc(p);
        if (retval)
                goto bad fork cleanup perf;
        /* copy all the process information */
        shm init task(p):
        retval = copy semundo(clone flags, p);
        if (retval)
                goto bad fork cleanup audit;
        retval = copy files(clone flags, p);
        if (retval)
                goto bad fork cleanup semundo;
        retval = copy fs(clone flags, p);
        if (retval)
                goto bad fork cleanup files;
        retval = copy_sighand(clone_flags, p);
        if (retval)
                goto bad fork cleanup fs:
        retval = copy_signal(clone_flags, p);
        if (retval)
                goto bad fork cleanup sighand;
        retval = copy mm(clone flags, p);
        if (retval)
                goto bad fork cleanup signal;
        retval = copy namespaces(clone flags, p);
        if (retval)
                goto bad fork cleanup mm;
        retval = copy_io(clone_flags, p);
        if (retval)
                goto bad fork cleanup namespaces;
        retval = copy thread tls(clone flags, stack start, stack size, p, tls);
        if (retval)
                goto bad fork cleanup io:
        if (pid != &init struct pid) {
                pid = alloc pid(p->nsproxy->pid ns for children);
                if (IS ERR(pid)) {
                        retval = PTR ERR(pid);
                        goto bad fork cleanup io;
```

```
list add tail rcu(&p->thread node,
                                          &p->signal->thread head):
                attach pid(p. PIDTYPE PID):
                nr threads++:
       }
       total forks++;
       spin unlock(&current->sighand->siglock);
       syscall tracepoint update(p);
       write unlock irg(&tasklist lock);
       proc_fork_connector(p);
       cgroup post fork(p, cgrp ss priv);
       threadgroup change end(current):
       perf event fork(p):
       trace task newtask(p, clone flags);
       uprobe copy process(p, clone flags);
       return p;
bad fork_cancel_cgroup:
       cgroup cancel fork(p, cgrp ss priv);
bad fork free pid:
       if (pid != &init struct pid)
               free pid(pid);
bad fork cleanup to:
       if (p->io context)
               exit io context(p);
bad fork cleanup namespaces:
       exit task namespaces(p);
bad fork_cleanup_mm:
       if (p->mm)
                mmput(p->mm);
bad fork cleanup signal:
       if (!(clone flags & CLONE THREAD))
               free signal struct(p->signal);
bad_fork_cleanup_sighand:
__cleanup_sighand(p->sighand);
bad_fork_cleanup_fs:
       exit_fs(p); /* blocking */
bad fork_cleanup_files:
       exit_files(p); /* blocking */
bad_fork_cleanup_semundo:
       exit sem(p):
bad fork cleanup audit:
       audit free(p);
bad fork cleanup_perf:
       perf event free task(p);
bad_fork_cleanup_policy:
#ifdef CONFIG NUMA
       mpol put(p->mempolicy);
bad_fork_cleanup_threadgroup_lock:
#endif
        threadgroup_change_end(current);
       delayacct_tsk_free(p);
bad fork cleanup count:
       atomic dec(&p->cred->user->processes);
       exit creds(p);
ad fork free:
```

### - dup\_task\_struct

```
static struct task struct *dup task struct(struct task struct *orig)
/*부모 프로세스가 인자로 함수 진입 */
       struct task struct *tsk:
       struct thread info *ti:
        /* 접근 가능한 node(메모리)를 할당해준다. */
        int node = tsk_fork_get_node(orig);
       int err;
/*위에서 할당한 노드가 들어감 */
/*실제 할당된 slab메모리 리턴 */
       tsk = alloc_task_struct_node(node);
       if (!tsk)
       return NULL;
/*원래 태스크의 메모리가 복사됨 페이지크기보다 클때는 다른페이지 연결리스트로 관리됨 */
        ti = alloc thread info node(tsk, node);
       if (!ti)
               goto free tsk:
       err = arch dup task struct(tsk, orig);
       if (err)
               goto free ti:
       tsk->stack = ti:
#ifdef CONFIG SECCOMP
        * We must handle setting up seccomp filters once we're under
        * the sighand lock in case orig has changed between now and
        * then. Until then, filter must be NULL to avoid messing up
        * the usage counts on the error path calling free task.
        tsk->seccomp.filter = NULL;
#endif
       setup thread stack(tsk, orig);
       clear user return notifier(tsk);
       clear_tsk_need_resched(tsk);
       set task stack end magic(tsk);
#ifdef CONFIG CC STACKPROTECTOR
        tsk->stack canary = get random int();
#endif
        * One for us, one for whoever does the "release_task()" (usually
        * parent)
       atomic_set(&tsk->usage, 2);
#ifdef CONFIG BLK DEV IO TRACE
       tsk->btrace_seq = 0;
#endif
        tsk->splice pipe = NULL:
       tsk->task_frag.page = NULL;
       tsk->wake q.next = NULL;
       account_kernel_stack(ti, 1);
       /*복사된 task가 리턴됨 */
        return tsk:
free ti:
        free_thread_info(ti);
free_tsk:
```

```
/* called from do fork() to get node information for about to be created task */
int tsk fork get node(struct task struct *tsk)
/*부모 task 들어옴 */
#ifdef CONFIG NUMA
       if (tsk == kthreadd task)
              return tsk->pref node fork:
               /*NUMA 일때 노드가 여러개이므로 task struct 에서 노드를 얻어온다.
               부모프로세스 태스크와 kthreadd_task(task_struct)가 같을때, */
#endif
       return NUMA NO NODE:
       /*그게 아니면 우마이기 때문에 노드는 한개이므로 리턴 *1
     static inline struct task struct *alloc task struct node(int node)
           /* origin이 사용하는 노드
             task가 갖고있는 시피유 캐시와, GFP KERNEL = 메모리 할당 옵션 , 원래 태스크가 쓰던 노드*/
           return kmem cache alloc node(task struct cachep, GFP KERNEL, node);
void *kmem cache_alloc(struct kmem_cache *cachep, gfp_t flags)
        /* task_cpu_cache , GFP KERNEL 메모리 할당 옵션이 넘어감 */
        void *ret = slab alloc(cachep, flags, RET IP );
        trace kmem_cache_alloc(_RET_IP_, ret,
                               cachep->object size, cachep->size, flags):
        return ret:
EXPORT SYMBOL(kmem cache alloc);
```

```
slab alloc node(struct kmem cache *cachep, gfp t flags, int nodeid,
                  unsigned long caller)
       unsigned long save flags:
       void *ptr;
       /* NUMA ID 값 설정 */
       int slab node = numa mem id():
       /*flags = GFP KERNEL */
       flags &= gfp allowed mask:
       Nockdep trace alloc(flags):
       if (slab should failslab(cachep, flags))
               return NULL:
        /* slab을 할당받아 반환해줌 */
       cachep = memcg kmem get cache(cachep, flags);
       cache alloc debugcheck before(cachep, flags);
       /*Interrupt 를 끄고 flags가 리턴됨 */
       local irg save(save flags);
       if (nodeid == NUMA NO NODE)
               nodeid = slab node;
       /*aet node 원래 노드 s에 노드 할당 */
       if (unlikely(!get node(cachep, nodeid))) {
               /* Node not bootstrapped yet */
               ptr = fallback alloc(cachep, flags):
               goto out:
       /* nodeid 와 slab node 는 같지 않다 */
       if (nodeid == slab node) {
                * Use the locally cached objects if possible.
                * However ____cache_alloc does not allow fallback
                * to other nodes. It may fail while we still have
                * objects on other nodes available.
                       __cache_alloc(cachep, flags);
               ptr =
               if (ptr)
                       goto out:
       /* cache alloc node can fall back to other nodes */
       /*obi return 실제 메모리 할당 */
       ptr = cache alloc node(cachep, flags, nodeid);
 out:
       local irg restore(save flags);
       ptr = cache alloc debugcheck after(cachep, flags, ptr, caller);
       kmemleak alloc recursive(ptr, cachep->object size, 1, cachep->flags,
                                flags):
       if (likely(ptr)) {
               kmemcheck slab alloc(cachep, flags, ptr, cachep->object size);
               if (unlikely(flags & __GFP_ZERO))
                       memset(ptr, 0, cachep->object_size);
       memcg kmem put cache(cachep);
       return ptr;
```

```
struct kmem cache * memcg kmem get cache(struct kmem cache *cachep)
        struct mem caroup *memca:
        struct kmem cache *memcg cachep;
        int kmemca id:
        VM BUG ON(!is root cache(cachep));
        /* current = task struct
          CONFIG MEMCG 가 설정되었으면
          memcq kmem skip account:1;
          1비트만 사용한다. fork를 하는경우 설정 안되있다. */
        if (current->memcg kmem skip account)
               return cachep:
        /*retrun memca */
        memcg = get mem cgroup from mm(current->mm);
        /*kernel memory cg id 값 구해옴 */
        kmemca id = READ ONCE(memca->kmemca id):
        if (kmemcg id < 0)
               goto out:
       /*cache index return */
/*여기서 slab을 할당 받음*/
        memca cachep = cache from memca idx(cachep, kmemca id):
        if (likely(memcg cachep))
               return memcg cachep;
         * If we are in a safe context (can wait, and not in interrupt
         * context), we could be be predictable and return right away.
         * This would guarantee that the allocation being performed
         * already belongs in the new cache.
         * However, there are some clashes that can arrive from locking.
         * For instance, because we acquire the slab mutex while doing
         * memca create kmem cache, this means no further allocation
         * could happen with the slab mutex held. So it's better to
         * defer everything.
         */
        memcg schedule kmem cache create(memcg, cachep);
out:
        css put(&memca->css):
        return cachep;
```

```
static struct mem cgroup *get mem cgroup from mm(struct mm struct *mm)
       struct mem cgroup *memcg = NULL;
       rcu read lock():
       do {
                * Page cache insertions can happen withou an
                * actual mm context, e.g. during disk probing
                * on boot, loopback IO, acct() writes etc.
                /*mm 의 유무 확인
                 mm 이 존재하므로 else 로 갑니다.*/
               if (unlikely(!mm))
                      memcq = root mem cgroup;
               else {
                      /*css 있으면 css리턴 없으면 NULL리턴
                        정보 있기때문에 css리턴되서 if문 통과 못함*/
                       memcg = mem cgroup from task(rcu dereference(mm->owner));
                      if (unlikely(!memcq))
                              memcg = root mem cgroup;
       } while (!css tryget online(&memcg->css));
       rcu read unlock();
       return memcg;
/*memcg 값 있어서 리턴 */
struct mem cgroup *mem cgroup from task(struct task struct *p)
         * mm update next owner() may clear mm->owner to NULL
         * if it races with swapoff, page migration, etc.
         * So this can be called with p == NULL.
        if (unlikely(!p))
               return NULL;
        /*css 있으면 css리턴 없으면 NULL리턴 */
        return mem cgroup from css(task css(p, memory cgrp id));
EXPORT SYMBOL(mem cgroup from task);
```

```
cache_from_memcg_idx(struct kmem_cache *s, int idx)
{
    struct kmem_cache *cachep;
    struct memcg_cache_array *arr;

    rcu_read_lock();
    /* s = cachep */
    /* 부노가 쓰던 메모리를 복사했기 때문에 s에 값이 있다 그걸 arr에 넣고*/
    arr = rcu_dereference(s->memcg_params.memcg_caches);

/*
    * Make sure we will access the up-to-date value. The code updating
    * memcg_caches issues a write barrier to match this (see
    * memcg_create_kmem_cache()).
    */
        /*arr 를 cachep로 가져온다*//
        cachep = lockless_dereference(arr->entries[idx]);
    rcu_read_unlock();

return cachep;
}
```

```
static void * cache alloc node(struct kmem cache *cachep, qfp t fla
                                 int nodeid)
        struct list head *entry;
        struct page *page:
        struct kmem_cache_node *n;
        void *obj;
        int x;
        VM BUG ON(nodeid < 0 || nodeid >= MAX NUMNODES);
        n = get_node(cachep, nodeid);
/*n = node 의 번호값을 가져옴 */
        BUG ON(!n):
        /*BUG ON = debug 할때만 사용 */
retry:
        check irg off():
        spin lock(&n->list lock);
        entry = n->slabs partial.next;
        if (entry == &n->slabs partial) {
                n->free touched = 1:
                entry = n->slabs free.next;
                if (entry == &n->slabs_free)
                         goto must grow:
        /*offset 으로 시작 값 가져옴 */
        /*entry - offset 페이지 값 셋팅 */
        page = list entry(entry, struct page, lru);
        check spinlock acquired node(cachep, nodeid);
        STATS INC NODEALLOCS(cachep);
        STATS INC ACTIVE(cachep);
        STATS SET HIGH(cachep);
        BUG ON(page->active == cachep->num);
        /*retrun objp*/
        obj = slab_get_obj(cachep, page, nodeid);
/*할당완료_했기 때문에 프리한 오브젝트 하나 지움 */
        n->free objects--;
        /* move slabp to correct slabp list: */
        list del(&page->lru):
        if (page->active == cachep->num)
                list add(&page->lru, &n->slabs full);
        else
                list add(&page->lru, &n->slabs partial);
        spin_unlock(&n->list_lock);
        goto done;
must grow:
        spin unlock(&n->list lock);
        x = cache grow(cachep, gfp exact node(flags), nodeid, NULL);
        if (x)
                goto retry;
        return fallback alloc(cachep, flags);
done:
```

```
-security_task_create
int security_task_create(unsigned long clone_flags)
{
    /* clone_flags = SIGCHLD
    return security_ops->task_create(clone_flags);*/
    return call_int_hook(task_create, 0, clone_flags);
    /*task ops retrun*/
}
```

결국 do\_fork에서 리턴되는건 자식 프로세스의 pid 값이며, 그 pid 값을 가진 프로세스 정보는 부모프로세스에서 복사된 정보로 차있다.