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CPSC 335-11  
Project 1

Algorithm 1:

Pseudocode:

Function seat\_swaps(row<vector>):

```
    count = 0
    for i from 0 to length(row) - 2, i increments by 2
        if row[i] + 1 != row[i + 1] or row[i] - 1 != row[i + 1]: // Check if each pair is a
            couple
            for j from i + 1 to length(row) - 1:
                if row[i] + 1 == row[i + 1] or row[i] - 1 == row[i + 1]: // Find matching person
                    swap(row[i + 1], row[j]) // Swap match
                    count++
            break // Prevents over incrementation
    return count
```

Efficiency Analysis:

I will analyze this algorithm using the step count method:

There are an  $n$  total number of individuals, there are an  $m$  number of couples, the relationship between  $n$  and  $m$  being  $m = n/2$ .

The total amount of steps in this algorithm is  $m * n / 2$ , this is because the outer loop of the algorithm must *at least* run  $m$  number of times, as we must check through every couple pairing, while the inside loop is checking two people at a time reducing the amount of loops to the number of people  $n / 2$ . So for  $m$  loops through the outer loop, two people are compared an  $n / 2$  amount of times, therefore the total amount of steps is represented by  $m * n / 2$  or  $mn/2$ . Because both  $n$  and  $m$  are linear in complexity the resulting algorithm is  $O(n)$  complexity.

Algorithm 2:

Pseudocode:

Function optimal\_start\_city(city\_distances<vector>, city\_fuel<vector>, car\_mpg):

```
    n = total number of cities in vector
    for start from 0 to n, start increments by 1's
        check = 1 // used as a flag to bail if starting city is not a valid start
        car_fuel = 0
        for i from 0 to n
            j = (i + start) % 2 // index of current city
            fuel = city_fuel[j]
```

```
distance = city_distance[j]
car_fuel += fuel * car_mpg - distance //subtract fuel for driving
if car_fuel < 0
    check = 0 // sets bailout condition for invalid city
if check == 1
    return start
```

#### Efficiency Analysis:

I will be using step count analysis to determine the efficiency class of this algorithm. Let  $n$  = the number of cities in our test vector. Both the outside loop and inside loop of my algorithm loops at least  $n$  number of times. Which means the worst case scenario for this algorithm is that the algorithm runs an  $n$ -squared number of times. But because both the inside and outside loop of the algorithm has a complexity of  $O(n)$  the resulting algorithm is also  $O(n)$  complexity