

Price of Anarchy

Lecture 7.4

$$\Gamma = \langle N, (S_i)_{i \in N}, (u_i)_{i \in N} \rangle$$

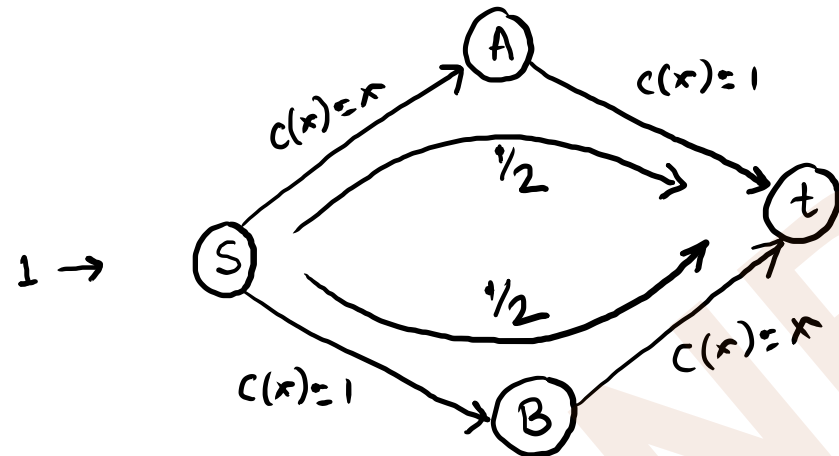
Social welfare function:

$$w: \prod_{i \in N} S_i \rightarrow \mathbb{R}_{\geq 0}$$

Definition (PoA):

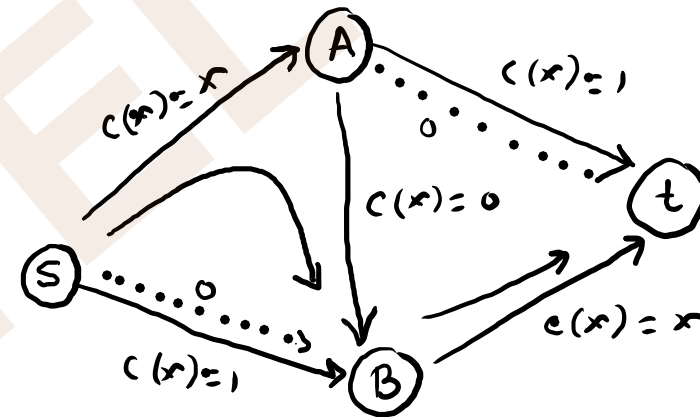
$$\text{PoA} = \frac{\max_{s \in \prod_{i \in N} S_i} w(s)}{\max_{s^{eq} \in \text{PSNE}(\Gamma)} w(s^{eq})}$$

Braess's Paradox



$(S \rightarrow A \rightarrow t : \frac{1}{2}, S \rightarrow B \rightarrow t : \frac{1}{2})$ is a PSNE

$$\text{Total delay} = \frac{3}{2}$$



$S \rightarrow A \rightarrow B$ always take a cost of $x (\leq 1)$ which is at most the cost of $S \rightarrow B$

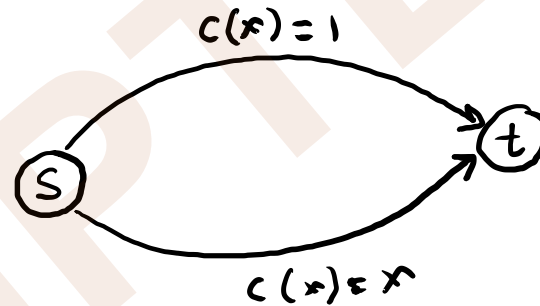
$S \rightarrow A \rightarrow B \rightarrow t$ is a weakly dominant strategy. Thus $(S \rightarrow A \rightarrow B \rightarrow t : 1)$ is a WDSE.

$$\text{Total delay} = 2$$

If the social welfare function is $\frac{1}{\text{total delay}}$, then

$$\text{PoA} \geq \frac{2}{3/2} = \frac{4}{3}$$

Pigou's Network:



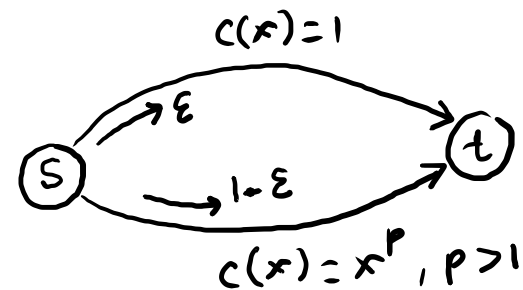
The bottom edge weakly dominates the top edge. So all traffic using the bottom edge is a WDSE.

Total cost = 1

$\frac{1}{2}$ traffic using the top edge and the other $\frac{1}{2}$ traffic

using the bottom edge incurs a cost of $\frac{1}{2} \cdot 1 + \frac{1}{2} \cdot \frac{1}{2} = \frac{3}{4}$.

$$\text{PoA} \geq \frac{4}{3}$$



All traffic following the bottom path is a WDSE. Total cost is 1.


Total cost is $\epsilon \cdot 1 + (1-\epsilon) \cdot (1-\epsilon)^p, \epsilon > 0$
 $= \epsilon + (1-\epsilon)^{p+1}$

$$\text{PoA} \geq \frac{1}{\epsilon + (1-\epsilon)^{p+1}} \rightarrow \infty \text{ as } p \rightarrow \infty$$

As non-linearity in the cost function increases, the PoA also increases.

PoA for Selfish Networks

Pigou's Network:

- Define:
- It has two vertices: source s and sink t
 -  edges from s to t .
 - Traffic rate is λ (> 0)
 - The cost of one edge is $c(x)$; $c(\cdot)$ is a non-negative, non-decreasing, continuous function

- The cost of the other edge is $c(r)$.

