

(1973) (1977)
Gibbard-Satterthwaite Theorem

Lecture 9.4

For any player $i \in [n]$ and any type $\theta_i \in \Theta_i$, we
get a partial order R_i on the set X of outcomes.
 $x, y \in X$, $x R_i y \stackrel{\text{def}}{=} u_i(x, \theta_i) \geq u_i(y, \theta_i)$

$R_i^{\theta_i}$ is called the rational preference relation
of player i when its type is θ_i .

We call a preference relation to be strict if that rational preference relation is a linear order/complete order. We denote the set of all possible strict rational preference relations on X by $\mathcal{L}(X)$.

Recall: $f: \prod_{i \in [n]} \mathcal{H}_i \longrightarrow X$

Social choice function (G-S Theorem) $f: \mathcal{L}(X)^n \longrightarrow X$

Unanimity: A social choice function $f: L(X)^n \rightarrow X$ is called unanimous if $\forall p_1, \dots, p_n \in L(X)$ such that the best outcomes in all p_1, \dots, p_n are the same, say x . we have $f(p_1, \dots, p_n) = x$

Ex-post efficiency \Rightarrow unanimity.

Theorem: Let $f: L(X)^n \rightarrow X$ be a social choice function
such that

(i) We have at least 3 outcomes. That is $|X| \geq 3$.

(ii) f is unanimous.

(iii) Every player has a strict rational preference relation.

Then f is dominant strategy incentive compatible
if and only if f is a dictatorship.

Way-Outs from GS Impossibility

- (1) Assume "more structure" on outcomes and "more structure" on the utility functions of the players.
 ← Quasilinear setting.
- (2) Be satisfied with Bayesian incentive compatibility.
- (3) "Computational barriers" can the manipulation problem be NP-hard?

