







DSA Data

Data Structures Algorithms

Interview Preparation

Data Science

Topic-wise Practice

C++

Mathematics | Graph Theory Basics - Set 1

Difficulty Level: Easy • Last Updated: 03 Jan, 2023







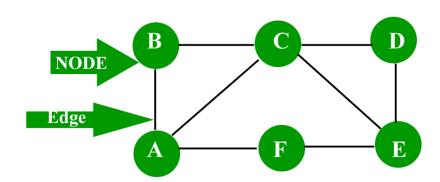
Read

Discuss

A graph is a data structure that is defined by two components:

- 1. A **node** or a vertex.
- 2. An edge E or **ordered pair is a connection between two nodes u,v** that is identified by unique pair(u,v). The pair (u,v) is ordered because (u,v) is not same as (v,u) in case of directed graph. The edge may have a weight or is set to one in case of unweighted graph.

Consider the given below graph,



To know about "Graph

representation" click here

Applications: Graph is a data structure which ______ed extensively in our real-life.

 Social Network: Each friend list are represer

2. <u>Google Maps:</u> Various represented as edges

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Not now

ctivities, suggestion and

des and the roads are th between two nodes.

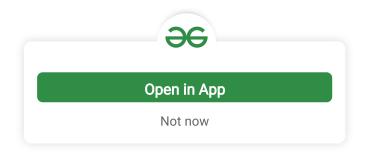
- 3. <u>Recommendations on e-commerce websites:</u> The "Recommendations for you" section on various e-commerce websites uses graph theory to recommend items of similar type to user's choice.
- 4. Graph theory is also used to study molecules in chemistry and physics.

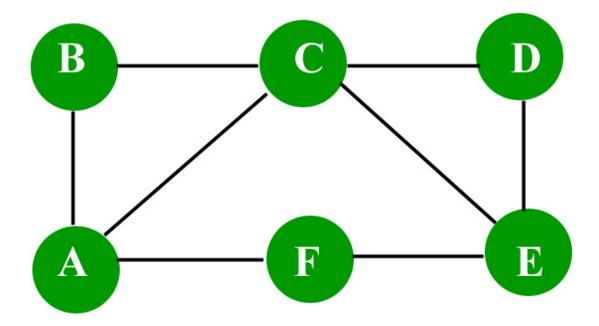
More on graphs: Characteristics of graphs:

- 1. <u>Adjacent node:</u> A node 'v' is said to be adjacent node of node 'u' if and only if there exists an edge between 'u' and 'v'.
- 2. <u>Degree of a node:</u> In an undirected graph the number of nodes incident on a node is the degree of the node. In case of directed graph ,**Indegree** of the node is the **number of arriving edges** to a node. **Outdegree** of the node is the **number of departing edges to a node**.

Note: 1 a self-loop is counted twice

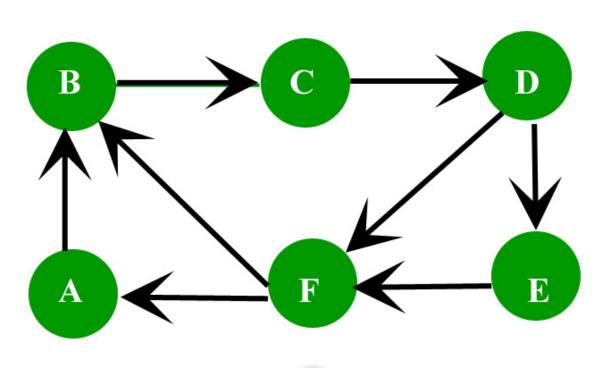
2 the sum of degree of all the vertices in a graph G is even.





Un-directed graph

1.



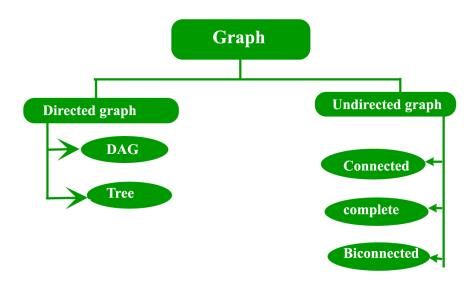


2. Path: A path of length

juence of n+1 nodes.

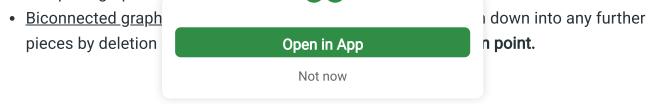
$$P(u,v)=(v0,v1,v2,v3.....vn)$$

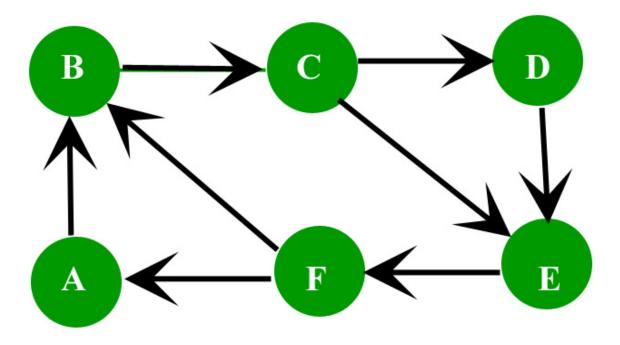
- 1. A path is simple if all the nodes are distinct, exception is source and destination are same.
- 2. <u>Isolated node:</u> A node with degree 0 is known as isolated node. Isolated node can be found by Breadth first search (BFS). It finds its application in **LAN network** in finding whether a **system is connected or not.**



Types of graphs:

- 1. <u>Directed graph:</u> A graph in which the direction of the edge is defined to a particular node is a directed graph.
 - <u>Directed Acyclic graph:</u> It is a directed graph with no cycle. For a vertex 'v' in DAG there is no directed edge starting and ending with vertex 'v'. a) Application: Critical game analysis, expression tree evaluation, game evaluation.
 - <u>Tree:</u> A tree is just a restricted form of graph. That is, it is a **DAG with a restriction that a** child can have only one parent.
- 2. <u>Undirected graph:</u> A graph in which the direction of the edge is not defined. So if an edge exists between node 'u' and 'v', then there is a path from node 'u' to 'v' and vice versa.
 - Connected graph: A graph is connected when there is a path between every pair of vertices. In a connected graph there is no unreachable node.
 - <u>Complete graph:</u> A graph in which each pair of graph vertices is connected by an edge. In other words, every node 'u' is adjacent to every other node 'v' in graph 'G'. A complete graph would have n(n-1)/2 ee below for proof.

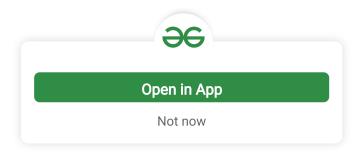


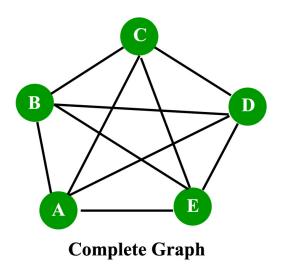


Connected Graph

Proof for complete graph:

- 1. Consider a complete graph with n nodes. Each node is connected to other n-1 nodes. Thus it becomes n * (n-1) edges. But this counts each edge twice because this is a undirected graph so divide it by 2.
- 2. Thus it becomes n(n-1)/2.



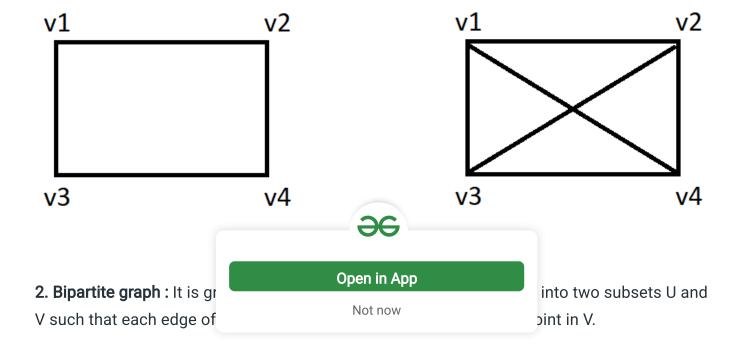


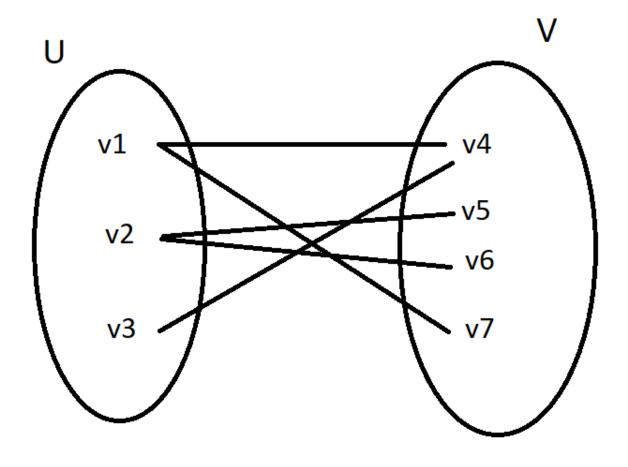
Consider the given graph, //Omit the repetitive edges Edges on node A = (A,B),(A,C),(A,E),(A,C). Edges on node B = (B,C),(B,D),(B,E). Edges on node C = (C,D),(C,E). Edges on node D = (D,E). Edges on node E = EMPTY.https://en.wikipedia.org/wiki/Graph_theory Total edges = 4+3+2+1+0=10 edges. Number of node = 5. Thus n(n-1)/2=10 edges. Thus proven. Read next set - <u>Graph Theory Basics</u>

Some more graphs:

1. Regular graph: A graph in which every vertex x has same/equal degree.k-regular graph means every vertex has k degree.

Every complete graph K_n will have (n-1)-regular graph which means degree is n-1.





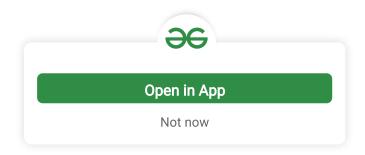
Bipartite graph

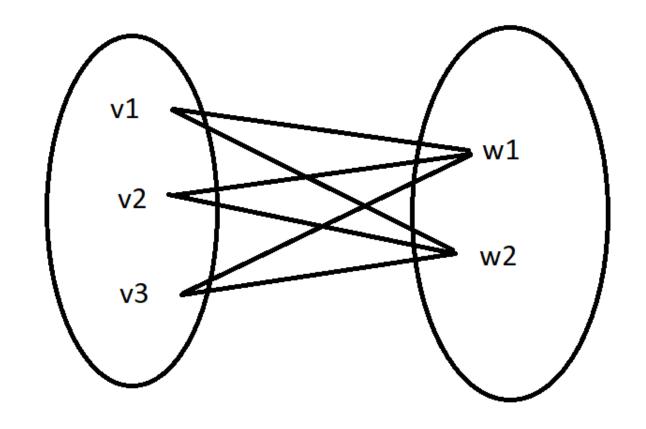
3. Complete Bipartite graph: it is a simple graph with vertex set partitioned into two subsets:

$$U=\{v_1, v_2, ..., v_m\}$$
 and $W=$

$$\{w_1, w_2, \dots, w_n\}$$

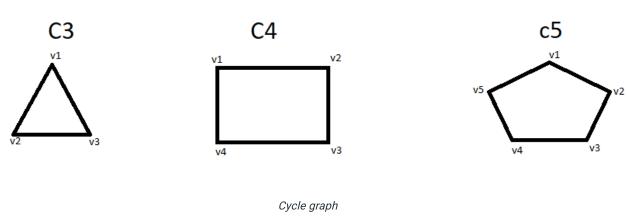
- i. There is an edge from each \boldsymbol{v}_i to each $\boldsymbol{w}_j.$
- ii. there is not an selp loop.





Complete Bipartite graph

4. Cycle graph : A graph of n vertices $(n \ge 3)$. v_1, v_2, \dots, v_n with edges $(v_1, v_2), (v_2, v_3), \dots, (v_{n-1}, v_n), (v_n, v_1)$.



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