

SEMINARUL NR 2

- criptosistem cu chei publice

- cheia publică - p, g nr. prime și $g \mid (p-1)$
 $\sim 1024 \text{ bits}$ 160 bits

$\alpha \in \mathbb{Z}_p^*$ el de ordin g , $\alpha^g \equiv 1 \pmod p \wedge \forall 1 \leq g' < g$,
 $\alpha^{g'} \not\equiv 1 \pmod p$

$\beta = \alpha^a \pmod p$

- cheia privată - $a \in \mathbb{Z}_g^*$

$x \in \mathbb{Z}_p^*$ $\text{enc}(x) = (\gamma, \delta)$, $\begin{cases} \gamma = \alpha^k \pmod p \\ \delta = x \cdot \beta^k \pmod p \end{cases}$ $k \in \mathbb{Z}_g^*$ aleator

$\text{dec}(\gamma, \delta) = \delta \cdot (\gamma^{-1})^a \pmod p$

Exerciții

- (5p) 1. Dem. că $\text{dec}(\text{enc}(x)) = x$, $\forall x \in \mathbb{Z}_p^*$
- (3p) 2. Dem. că dacă la criptarea a 2 mesaje dif x_1 și x_2 se folosesc același k , atunci având aceleași (γ, δ_1) și (γ, δ_2) și unul din plaintexte (x_1) , se poate determina ușor al doilea plaintext (x_2) . (fără a avea cheia privată)
- (2p) 3. Dem. că criptarea El Gamal e maleabilă.
 Având un criptotext $(\gamma, \delta) \xleftarrow{\text{enc}} x$

\downarrow ? pot să modific ușor

$(\gamma', \delta') \xrightarrow{\text{dec}} \epsilon x$ ϵ arbitrar ales

Rezepte

$$1. \quad \text{dec}(\text{enc}(x)) = \text{dec}((\gamma, \delta)), \text{ unde } \begin{aligned} \gamma &= \alpha^k \bmod p \\ \delta &= x \cdot \beta^k \bmod p \end{aligned}$$

$$\begin{aligned} \Rightarrow \text{dec}(\text{enc}(x)) &= \text{dec}((\gamma, \delta)) = \delta \cdot (\gamma^{-1})^a \bmod p = \\ &= (x \cdot \beta^k \bmod p) \cdot (\alpha^k \bmod p)^{-1}^a \bmod p = x \cdot \beta^k \cdot (\alpha^k \bmod p)^{-1}^a \bmod p = \\ &= x \cdot (\alpha^a \bmod p)^k \cdot (\alpha^k \bmod p)^{-1}^a \bmod p = \\ &= x \cdot (\alpha^k \bmod p)^a \cdot (\alpha^k \bmod p)^{-1}^a \bmod p = \\ &= x \cdot \left[(\alpha^k \bmod p) (\alpha^k \bmod p)^{-1} \right]^a \bmod p = \\ &= x \cdot 1^a \bmod p = x \quad (\text{da } x \in \mathbb{Z}_p^*) \end{aligned}$$

$$2. \quad \text{enc}(x_1) = (\gamma_1, \delta_1) \text{ unde } \begin{aligned} \gamma_1 &= \alpha^k \bmod p \\ \delta_1 &= x_1 \cdot \beta^k \bmod p \end{aligned}$$

$$\text{enc}(x_2) = (\gamma_2, \delta_2) \text{ unde } \begin{aligned} \gamma_2 &= \alpha^k \bmod p = \gamma_1 \\ \delta_2 &= x_2 \cdot \beta^k \bmod p \end{aligned}$$

$$(\gamma_1, \delta_1)$$

$$(\gamma_2, \delta_2)$$

$$x_1$$

$$x_2 = ?$$

$$x_1 = \text{dec}((\gamma_1, \delta_1)) = \delta_1 \cdot (\gamma_1^{-1})^a \bmod p = (x_1 \cdot \beta^k \bmod p) \cdot (\gamma_1^{-1})^a \bmod p = x_1$$

$$\Rightarrow x_1 \cdot \beta^k \cdot (\gamma_1^{-1})^a \bmod p = x_1 = x_1 \bmod p$$

$$\Rightarrow \beta^k \cdot (\gamma_1^{-1})^a \bmod p = 1 \quad (*)$$

$$\text{dec}((\gamma_2, \delta_2)) = \delta_2 \cdot (\gamma_2^{-1})^a \bmod p \stackrel{(*)}{=} \delta_2 \cdot (\gamma_1^{-1})^a \bmod p =$$

$$= (x_2 \cdot \beta^k \bmod p) \cdot (\gamma_1^{-1})^a \bmod p = x_2 \cdot \underbrace{\beta^k (\gamma_1^{-1})^a}_{1 \bmod p} \bmod p \stackrel{(*)}{=} x_2 \bmod p = \underline{x_2} \quad (2)$$

$$\delta_2 = x_2 \cdot \beta^k \mod p$$

$$\delta_1 = x_1 \cdot \beta^k \mod p \Leftrightarrow \beta^k \mod p = \delta_1 \cdot x_1^{-1} \mod p$$

$$\delta_2 = x_2 \cdot \beta^k \mod p \Leftrightarrow \delta_2 = x_2 \cdot \delta_1 \cdot x_1^{-1} \mod p$$

$$\underline{(-) \quad x_2 = \delta_2 \cdot x_1 \cdot \delta_1^{-1} \mod p}$$

$$3. \quad \delta \cdot (\gamma^{-1})^a = x$$

$$\delta' = \delta \cdot \beta \mod p$$

$$\gamma' = \gamma \cdot \alpha \mod p$$

$$\delta'^a \cdot (\gamma'^{-1})^a \mod p = \delta \cdot \beta \mod p \cdot ((\gamma \cdot \alpha \mod p)^{-1})^a \mod p$$

$$= \delta \cdot \beta \cdot (\gamma^{-1} \mod p \cdot \alpha^{-1} \mod p)^a \mod p =$$

$$= \delta \cdot \beta \cdot (\gamma^{-1})^a \cdot (\alpha^{-1})^a \mod p =$$

$$= \delta \cdot (\gamma^{-1})^a \cdot \alpha^a \cdot \alpha^{-a} \mod p = \delta \cdot (\gamma^{-1})^a \mod p = x$$

$$\left(\begin{array}{l} \delta' = \epsilon \delta \mod p \\ \gamma' = \epsilon \gamma \mod p \\ \delta' \cdot (\gamma'^{-1})^a \mod p = \epsilon \delta \cdot (\epsilon \gamma)^{-1})^a \mod p = \\ = \epsilon^{a+1} \delta \cdot \epsilon^{-a} \cdot (\gamma^{-1})^a \mod p = \delta \cdot (\gamma^{-1})^a \cdot \epsilon^a \cdot \epsilon^{-a} \cdot \epsilon \mod p \\ = \epsilon \cdot \delta \cdot (\gamma^{-1})^a \mod p = \epsilon x \end{array} \right)$$

$$\delta' = \epsilon \delta \mod p$$

$$\gamma' = \epsilon \gamma \mod p$$

$$\delta' (\gamma'^{-1})^a \mod p = \epsilon \delta (\gamma^{-1})^a \mod p$$

$$= \epsilon x$$