



Problem Solving

Course on Game Theory and Greedy Algorithms

Greedy solution May not work

$\tau_i, \underline{d_i}$

$$\text{penalty} = \underline{d - d_i} \quad \text{if } d > d_i$$

τ_i

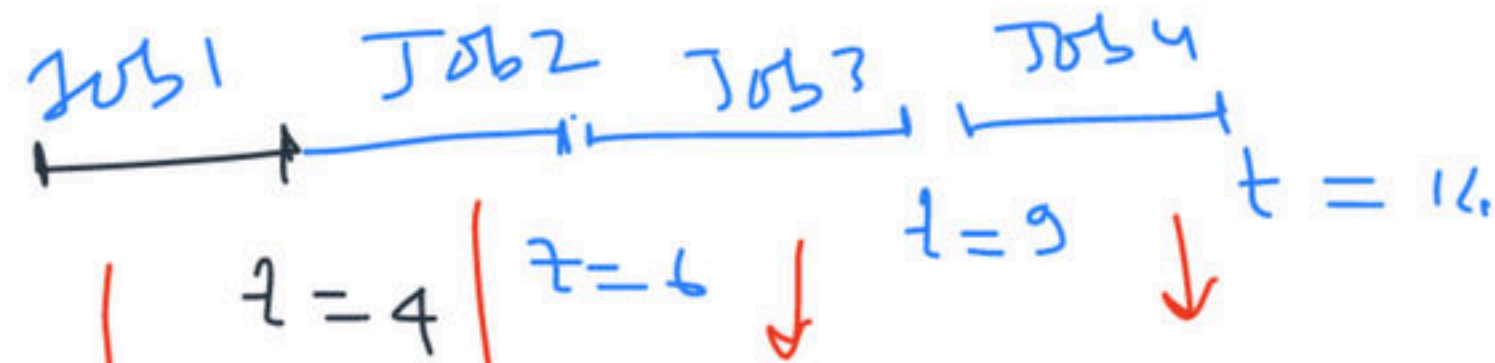
* perform them such that the penalty is min.

τ_i

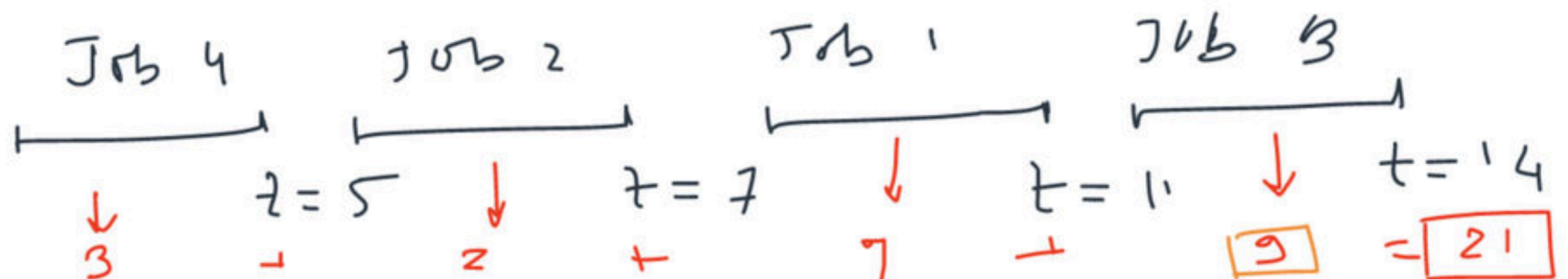
The order in which these tasks should be performed?? Day starts at $t=0$.

τ_i

Job 1	$\tau_i \downarrow$ 4	$d_i = \underline{4}$
Job 2	2	$= \underline{3}$
Job 3	3	$= \underline{5}$
Job 4	5	$= \underline{2}$



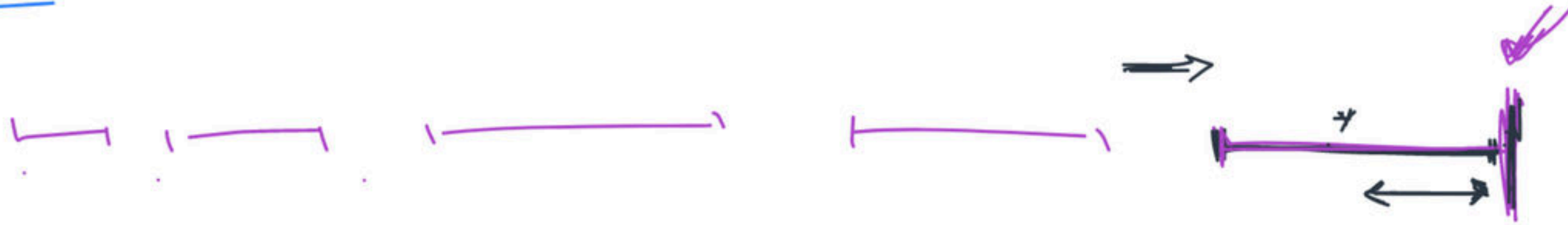
$$0 + 3 + 4 + \boxed{12} = \boxed{19} \text{ penalty.}$$



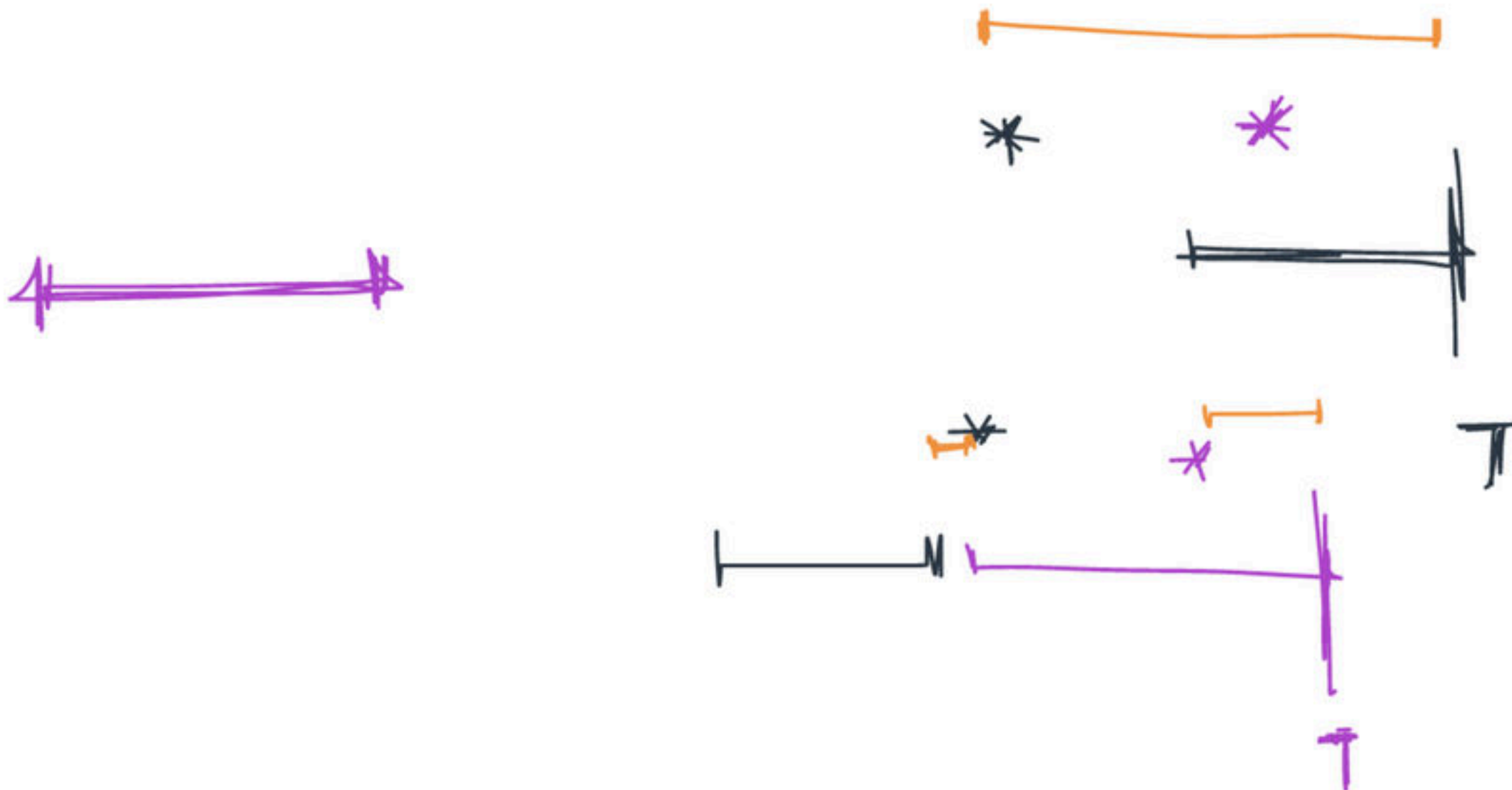
$$\boxed{9} = \boxed{21}$$



Modification: what if we want to minimize the maximum penalty.

Solution??



$$T = \sum t_i$$




 $\rightarrow t$
 job whose
 deadline
 is late
 $d \neq \frac{T}{2}$



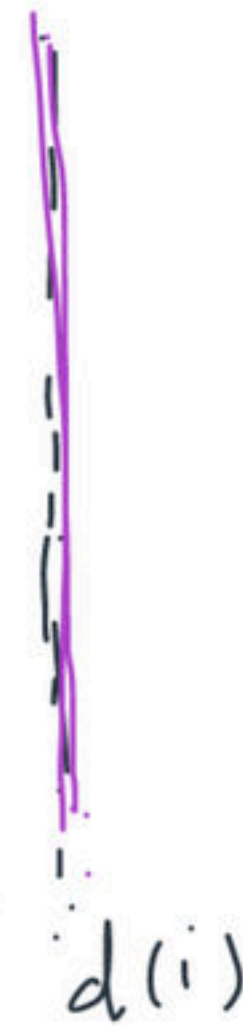
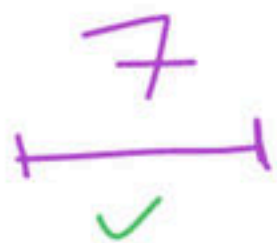


$\underline{d_i}$ $\boxed{p_i}$

You want to earn maximum possible profit.

NOTE: All of them takes same time to finish. $t_i = 1$

- 1) Sort in dec order of profit.
- 2) Now pick a job and do it as late possible but within the timeline.



'S' → 'S1'
↑
you removed the jobs which don't earn profit.
S*

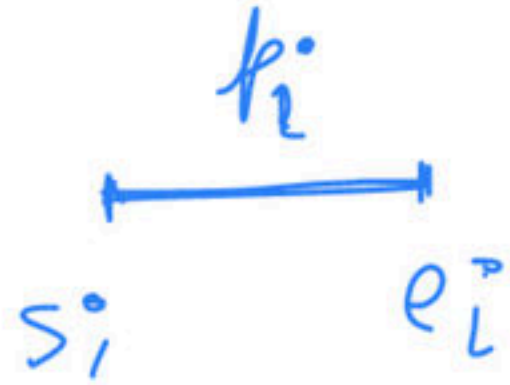
$\boxed{d(j) > d(i)}$
 $A(j) < d(i)$

$d(j)$ $d(i)$

→ You put the guy at the end without losing any profit

Therefore from the exchange argument, it always makes sense to put a job 'i' as late as possible provided that still earns you profit.

modification



now do you schedule to maximize the profit ??

↓
* This cannot be solved greedily.

↓
"Dynamic Programming"

Huffman Coding → lossless compression

"abc" → 3 × 8 = 24 bits
↓ ↓ ↓
'A byte long'

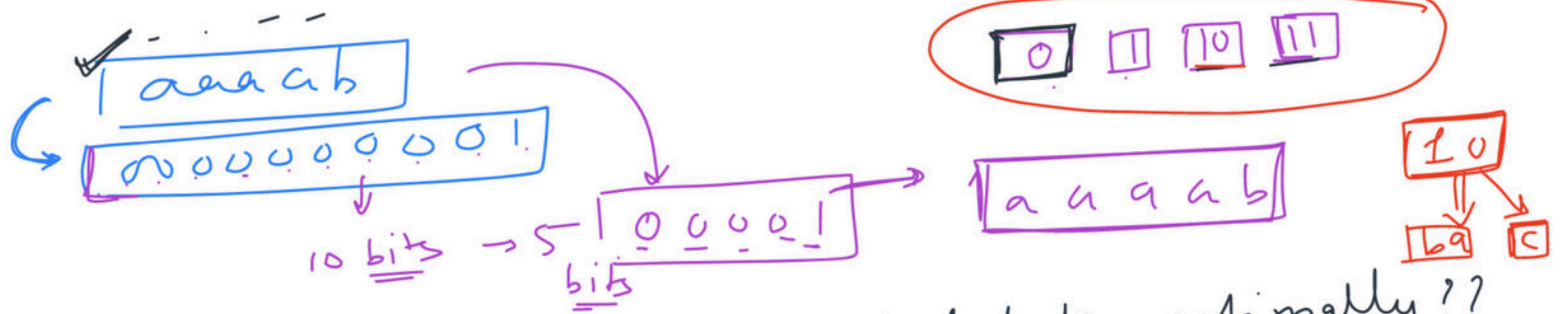
aaaaa → 8 + 8 = 16 bits
↑
a7
8 × 7 = 56 bits

'i' → 8 bits

a → 2
b → 3
⋮
z → 12
z → 8

⇒ on an average the total length is minimized.

$\Sigma = \{a, b, c, d\}$
00 01 10 11



How to assign bit sequences to the alphabets optimally??

To ensure no ambiguities, we need to make sure that bit seq of any character is not a prefix of another bit seq.



\Rightarrow prefix codings

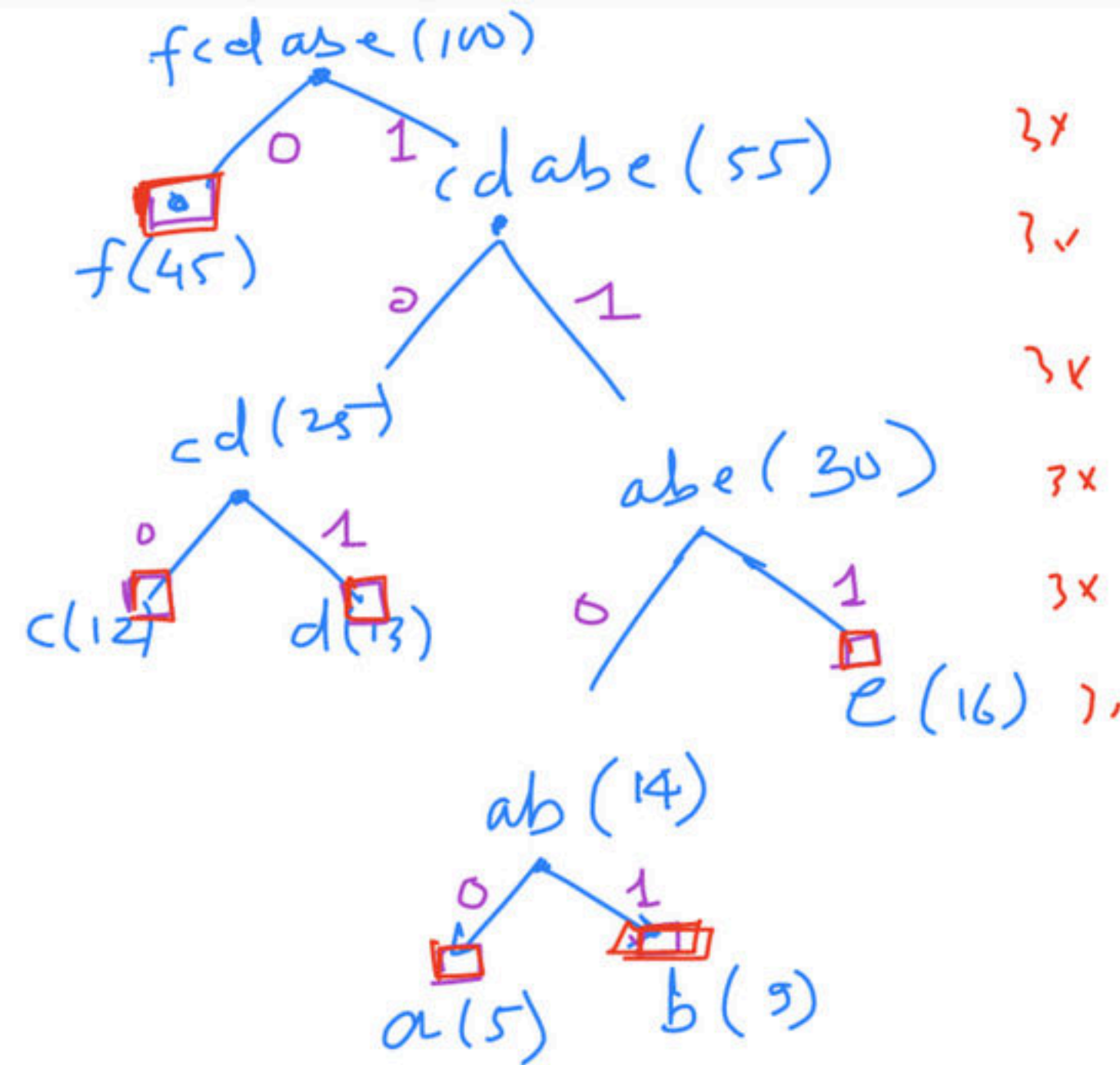
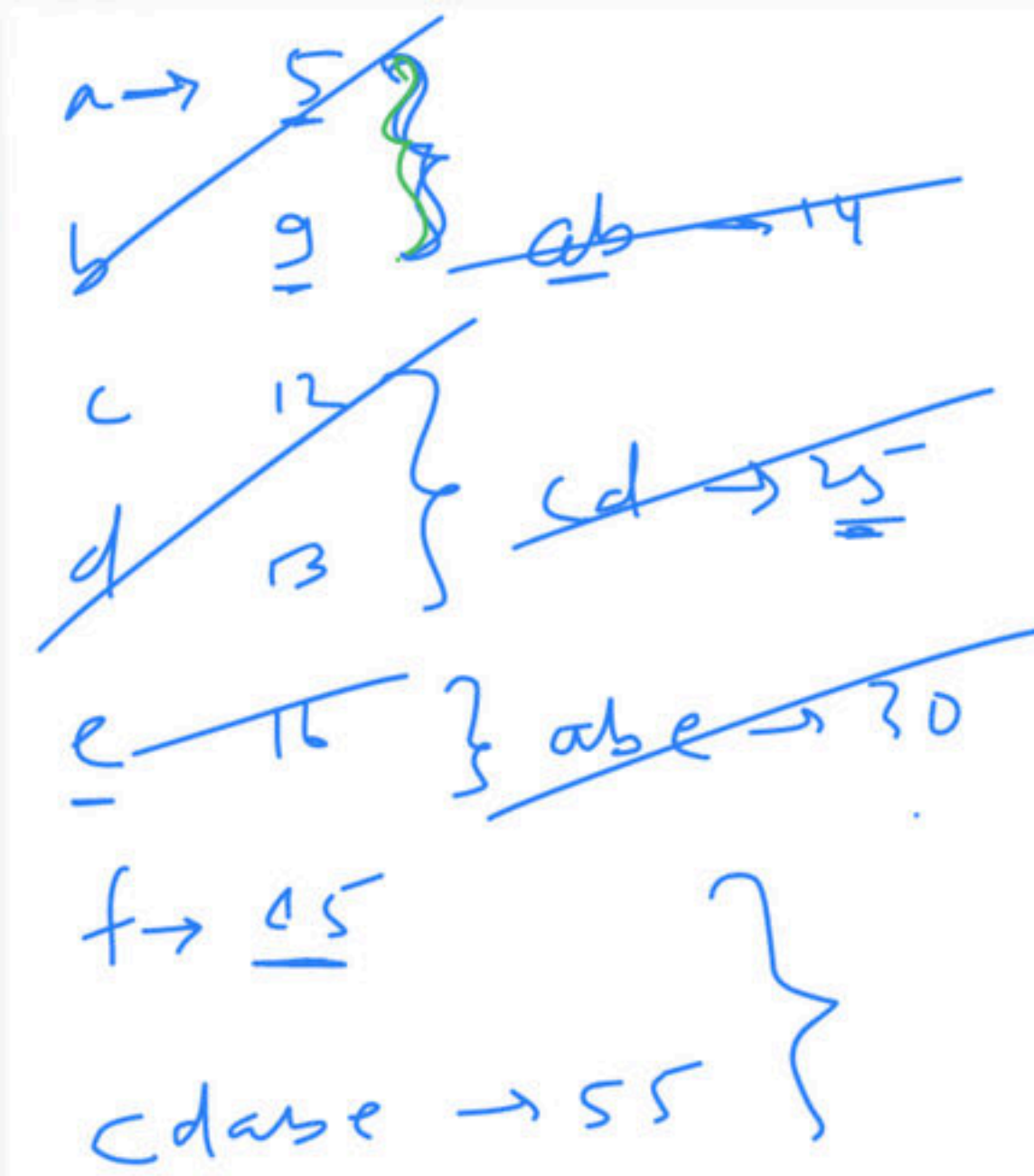
ascii characters are prefix encoding or not!! \rightarrow Yes

How to ensure optimality??

Σ	a	b	c	d	...
f	f_1	f_2	f_3	f_n	...

8 bit length

How to determine the bit seq for each of the characters??



3x	1100	$a \rightarrow 5 \times 4 = 20$
3x	1101	$b \rightarrow 9 \times 4 = 36$
3x	100	$c \rightarrow 12 \times 3 = 36$
3x	101	$d \rightarrow 13 \times 3 = 39$
3x	111	$e \rightarrow 16 \times 3 = 48$
3x	10	$f \rightarrow 45 \times 1 = 45$
		224

1300 bits for normal bit rep.

$$\frac{76}{300} \times 100 = 25.3\%$$

You're a robber and broke into a grocery store which contains 'n' kind of grains (rice, wheat etc). Amount of each grain is given and also its monetary value. Your bag has a size 'W'. What is the maximum monetary value can you steal from the store??

	rice	wheat-	x	y	z
weight	5	4	5	3	4
price	5 u	12 3	13 2	17 1	2 5

$$W = 11$$

$$3 + 5 + 3 = 11$$

$$17 + 13 + 12 \times \frac{3}{4} = \boxed{37}$$

$$3y + 4 \text{ wheat-} + 4 \text{ unit of 'x'} = 17 + 12 + \frac{13}{5} \times 4 = \boxed{39.4}$$

$$W = 39$$

What is the max monetary value that you can steal??

Ans: Sort according to price/weight (price per unit wt)
then pick them one by one.

If splitting / breaking is not allowed then what??
→ Greedy fails || 0-1 knapsack problem → dp

□ $\begin{cases} \text{you take it} \\ \text{don't take it} \end{cases}$

b_i time to destroy that letter

'S' \leftarrow burn as little ^{letters} as possible, but the burning should go on to at least 'S' unit of time.



$$T = \sum b_i > S$$

$$T = S$$

$$T < S$$

$$T > S$$

Minimize $\|T - S\|$. T.c. = $O(n^2 \log n)$

$$T = S$$

