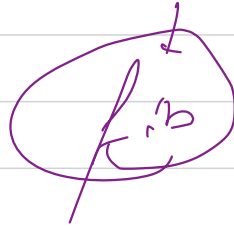


$$f(i) = f(i-1) + f(i-2)$$



$\rightarrow O(1)$  per  
state

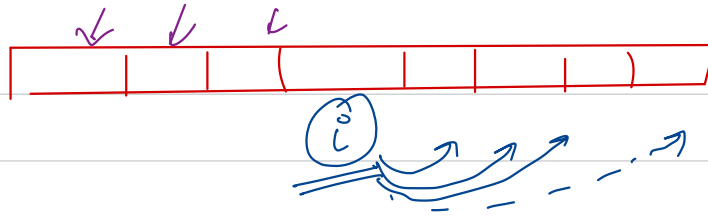
$$O(1) \times (n+1)$$

$$O(n+1)$$

$\downarrow$

$$\underline{O(n)}$$

$\phi$

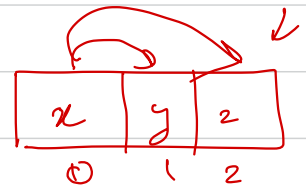


① from any  $i^{\text{th}}$  index I can jump to  
 $(i+1)$   $(i+2)$   $(i+3)$  .....  $(i+k)$

② Min cost to reach from  $(1 \rightarrow N)$

$k=2$

③ Similarly to reach at any  $i^{\text{th}}$  index  
we can come from  
 $(i-1)$   $(i-2)$   $(i-3)$  .....  $(i-k)$



$$\underline{f(i)} = \min \left( \underbrace{|h_i - h_{i-1}| + f(i-1)}_{\downarrow}, \underbrace{|h_i - h_{i-2}| + f(i-2)}_{\downarrow}, \underbrace{|h_i - h_{i-3}| + f(i-3)}_{\downarrow}, \dots, \underbrace{|h_i - h_{i-k}| + f(i-k)}_{\downarrow} \right)$$

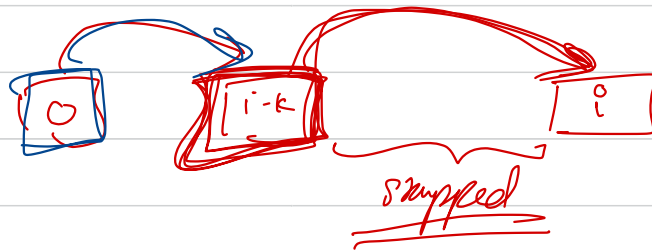
$k$  terms

Calc, min cost  
to reach  $i^{\text{th}}$  index

min cost to  
go from  $0 \rightarrow i$   
via  $i-1$

min cost to reach  
 $0 \rightarrow i$  via  
 $i-2$

min cost to  
reach  $0 \rightarrow i$   
via  $i-k$

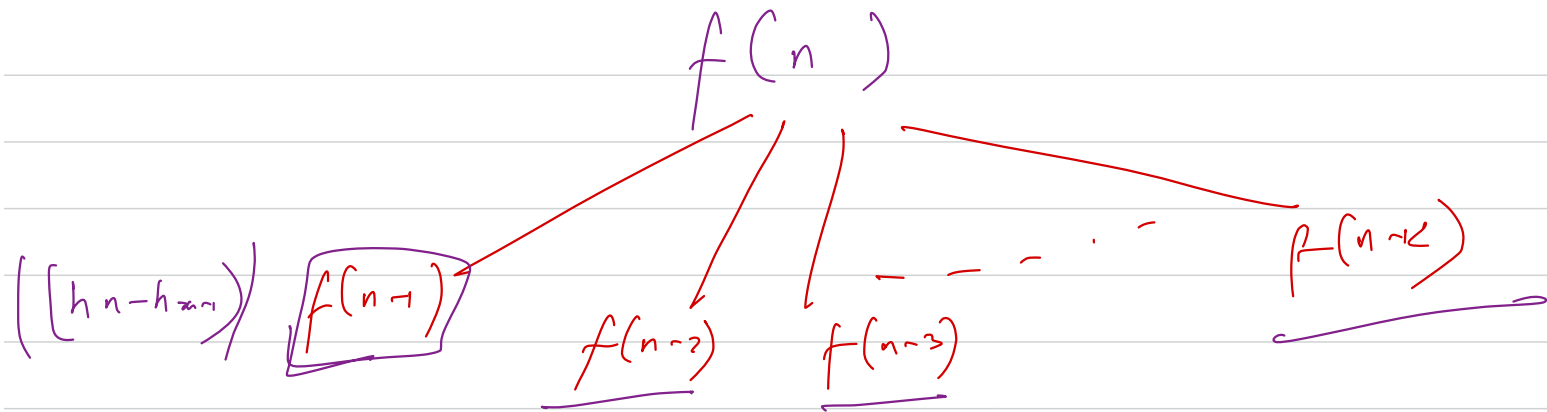


final ans  $\rightarrow \underline{\underline{f(n-1)}}$

$$f(i-k) + |h_i - h_{i-k}|$$

n subproblems

$$\underline{f(i) = \min (|h_i - h_j| + f(i-j)) \quad \forall j \in [1, k]}$$



for  $i^{\text{th}}$  state  $\rightarrow O(k)$

$\rightarrow \underline{\underline{O(nk)}}$

TC  $\rightarrow$  Time reqd to solve ans state  $\times$  no of ans state

$$0 + |h_0 - h_1|$$

4, 9, 2, 5, 6

$n=5$

$$k=3$$

dp

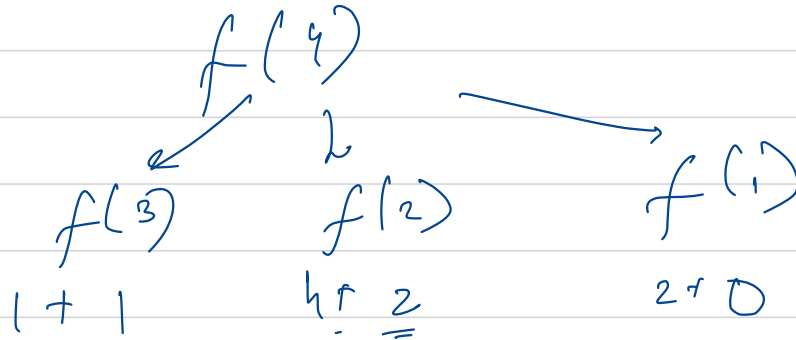


0	0	2	1	2
0	1	2	3	4

dp(i)

final ans

dp[i] → f(i) → min cost to reach i from 0



Q.2

You are given integer  $n$ . On each step, you may subtract from it any one digit that appears on it.

Calc, min steps to make number  $n$  equal to 0.

$$27 \xrightarrow{-7} 20 \xrightarrow{-2} 18 \xrightarrow{-8} 10 \xrightarrow{-1} 9 \xrightarrow{-9} 0$$

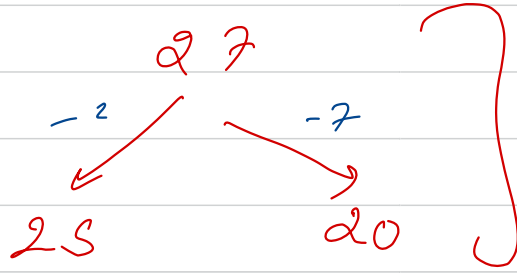
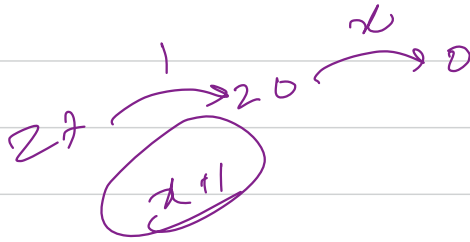
$$n \leq 10^6$$

greedy  
↓  
proof

5 ans

$n = 27 \rightarrow$  min steps to reach 0

We don't know how calc min steps of  $27 \rightarrow 0$



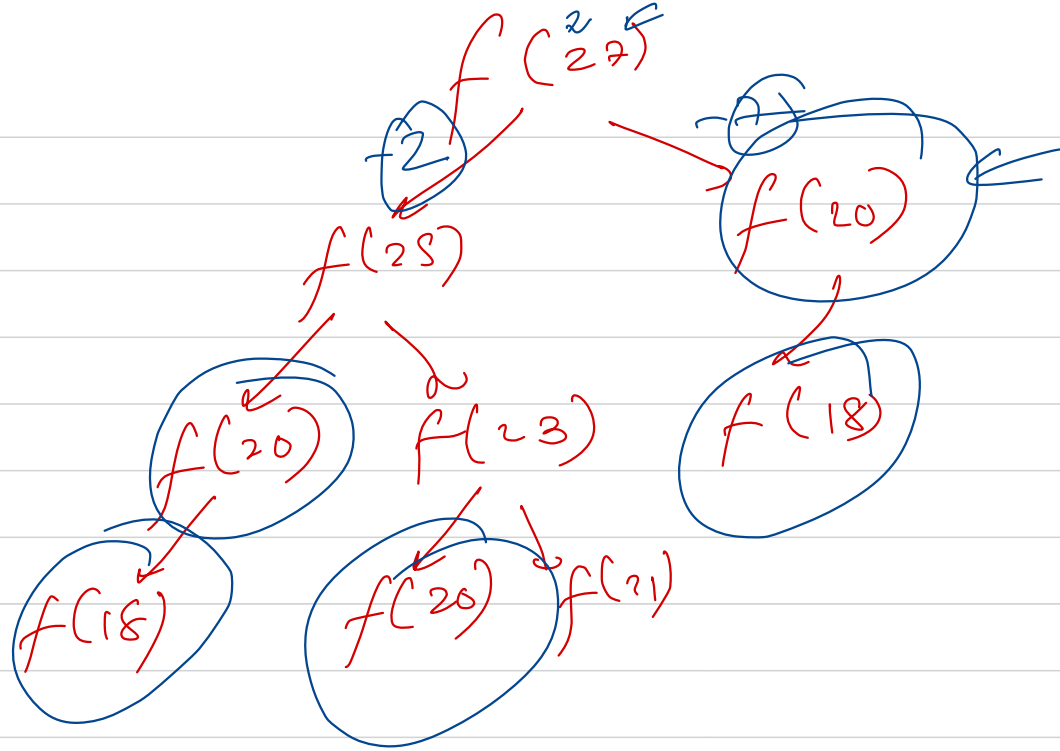
we know this

$$f(27) = 1 + \min(f(25), f(20))$$

min steps to reach 0 from 27

min steps for  $25 \rightarrow 0$

min steps for  $20 \rightarrow 0$





$$f(n) = 1 + \min(dp[n - j])$$

$$\forall j \in \text{digits of } n$$

and

$$j \neq 0$$

$$27 \rightarrow \text{digits of } [2, 7]$$

$$n \in [1, 9]$$

$$f(n) \rightarrow 1$$

$$\begin{array}{l} 27 \rightarrow 25 \\ 27 \rightarrow 20 \end{array}$$

$$\begin{array}{l} 1 \rightarrow 0 \\ 2 \rightarrow 1 \\ 3 \rightarrow 2 \end{array}$$

Q<sup>n</sup> You have a system of coins, of  $n$  coins, Each coin has a +ve value. Calc the no. of <sup>uniq</sup> ways you can produce a money sum  $x$  using the coins.

$[2, 3, 5]_n$  infinite supply

permutation

$$x = 9$$

$$\downarrow$$

$$\underline{\underline{8}}$$

$$x \leq 10^6$$

$$\left. \begin{array}{l} n \leq 10^2 \\ C_i \leq 10^6 \end{array} \right\}$$

2, 2, 5  
2, 5, 2  
5, 2, 2  
3, 3, 3  
2, 2, 2, 3  
2, 2, 5, 2  
2, 3, 2, 2

3, 2, 2, 2

8

[illegible]

8 zeros

We don't know the total ways to reduce  $9 \rightarrow 0$   
 using  $[2, 3, 5]$

Count permutation

$$f(9) = f(7) + f(6) + f(4)$$

no. of ways to  
 reduce  $9 \rightarrow 0$

permutation

no. of  
 words  
 with first element  
 as 2 and reduce  
 $9 \rightarrow 0$

$\left[ \begin{array}{l} 2, 5, 2 \\ 2, 3, 2, 2 \\ 2, 2, 3, 2 \\ 2, 2, 2, 3 \\ 2, 2, 5 \end{array} \right]$

$\rightarrow$  way to go for  
 $7 \rightarrow 0$

$$[c_1, c_2, c_3, \dots, c_n]$$

$$f(x) = \sum_{i=1}^n f(x - c_i)$$

$$dp(x) = \sum_{i=1}^n dp(x - c_i)$$



$$\begin{array}{l} \rightarrow 3 \\ S \rightarrow \textcircled{3} \\ \rightarrow 2, 3 \\ \rightarrow 2, 3 \end{array}$$

$$x=0$$

1 ways

always for  
count

$$\{2, 3\}$$

$\rightarrow$

$$\{2\}$$

$$\{3\}$$

$$\{2, 3\}$$

$$\{0\}$$

$$f(2)$$

$$f(3)$$

$$f(5)$$

$$f(2) = f(0)$$

$$f(3) = f(0) + f(1)$$

$$f(5) = f(0) + f(2) + f(3)$$

Q<sub>2</sub> Consider the same prev question, this time  
calculate the no. of distinct ordered ways  
to produce money

$[2, 3, 5]$

$n = 9$

$\rightarrow (3)$

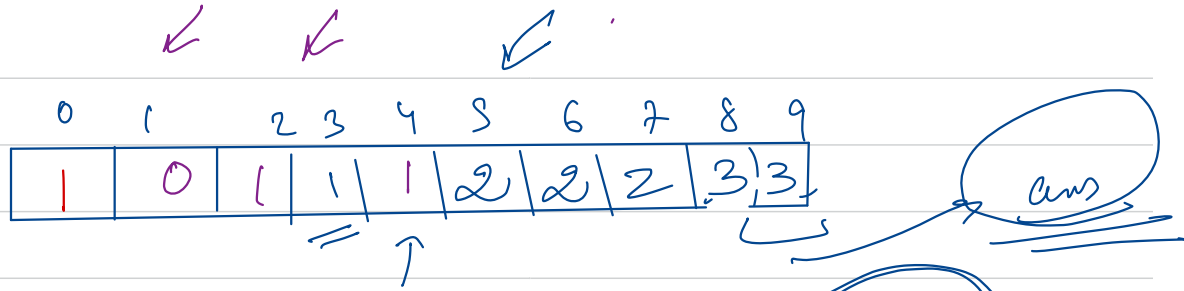
$[2, 2, 2, 3]$

$[2, 2, 5]$

$[3, 3, 3]$

}

$\rightarrow$  <sup>ways</sup> all combinations



{ 2, 3, 6 }

3 2 2 2  
 2 3 3  
 2 2 2 2  
 3, 2, 3

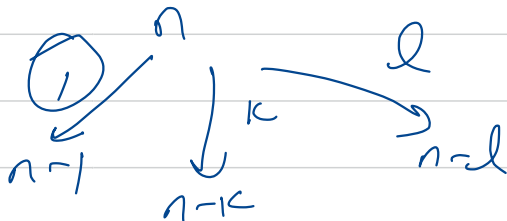
- ① with only 2 what all combination you can make
- ② with only 2, 3 what all combination you can make

$$\sum_{t=1}^{\infty} \text{money}(t) = S, \quad \text{money}(t) = \frac{1}{t} - \frac{1}{t+1}$$

Q is m coins



turn



A, B, A, B, A - - - -

winner  $\rightarrow$  who takes last  
1 coin or  $k$  coins  
 $\infty$  1 coins



$(n) \rightarrow$  no. of coins

$\left. \begin{matrix} k \\ d \end{matrix} \right\} \rightarrow$  On each turn you can have  
 $1, k, d$  coins picked

$$\boxed{f(n, k, l)} = f(n-1, k, l) \quad f(n-k, k, l)$$

if we have  $n$   
coins when

$$f(n-l, k, l)$$

A will win or  
B will win

sub problem

0	0	0	→	1
---	---	---	---	---

0 → loose  
1 → win  
 $k =$

0	1	0	→	1
0	1	1	→	1
1	1	1	→	0
1	0	1	→	1

$$K=3$$

$$l=5$$

$$n=7$$

$$n=2 \rightarrow A$$

$$n=7 \rightarrow A$$



$$n=2 \rightarrow B$$



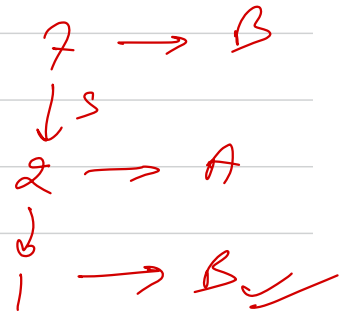
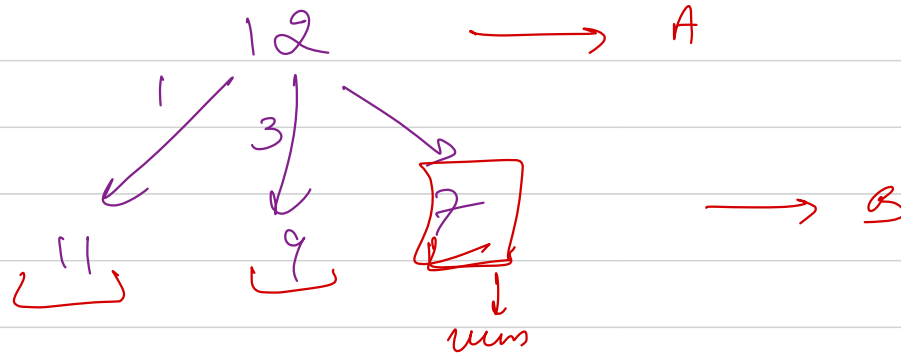
$$n=1$$

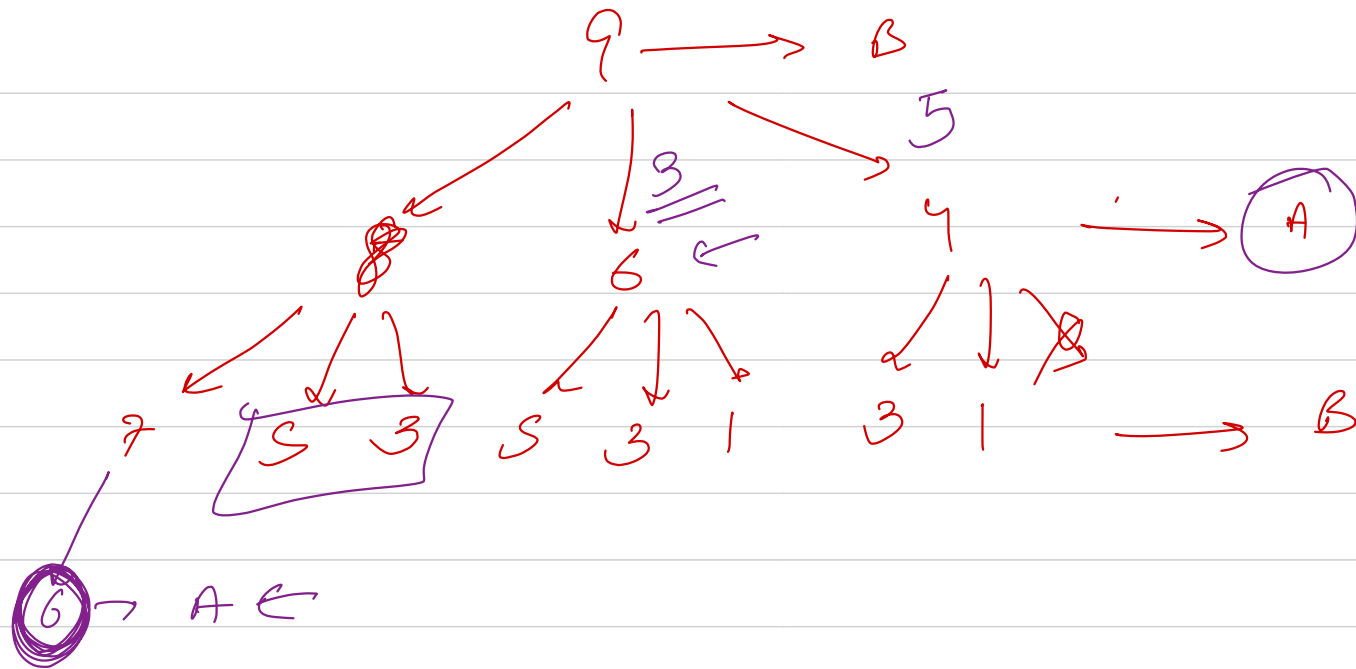


$$n=0$$

$$\rightarrow A \checkmark$$

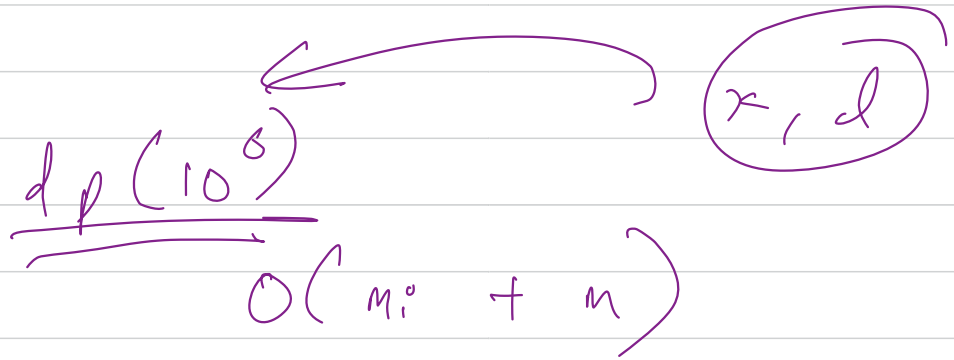
It doesn't only depend on what A is picky  
it also depends on what B gets





$f(n) \rightarrow$  whether  $n$  is a sum state  
or loopy state

$m \rightarrow$



$$dp(n) = dp(n-k) \& dp(n-2) \& dp(n-1)$$

$$dp(n) = \left\{ \begin{array}{ll} dp[n-k] == 0 & \rightarrow 1 \\ \text{or} & \\ dp[n-2] == 0 & \rightarrow 1 \\ \text{or} & \\ dp[n-1] == 0 & \rightarrow 1 \end{array} \right.$$

or any

$$dp[1] \rightarrow 1$$

$$dp[k] \rightarrow 1$$

$$dp[l] \rightarrow 1 \quad \xrightarrow{\quad} \underline{\underline{0}}$$

<https://www.spoj.com/problems/MCOINS/>