Greedy algorithms Greedy algorithms always make the choice lest taking ente consideration

consequences. op timesed 1. Greedy algorithms are always based on intuition 2 observation.

- 2. Greedy algorithmes can only work when local optimal extition contributes to global answer.
- Two properties: 1) Greedy choice proporty: If by choosing the best choice at each step welthout reconsiduing the previous step chosen.
- (2) optimal substructure If the optimal overall solution to the problem corresponds to the optimal solution of êts subproblems. Q Tosk scheduler (leetcode)

Total CPV ydes = Busy + Fidle ??

- No. of busy slots = no. of terets = len(ar) Max no. of falle slots is defined by the frequency of most frequent tosk. A B C A A
- Mar possible follo elots = (1 - max - 1)= (5-1) * 2 A ->5

(8) - (1)

ans = -1

 $\rightarrow B \rightarrow (2)$

ABCAB - A - A

ABACBACBCBCB > Total → (17)

D. Minimum subset product

Eg: {-1,03

Total time = 5+5 = (13)

 $fg: \{-1,-1,-2,4,3\}$ ans = -24

Eg: [-1,0,2] ans = -2 (1) what if all tre no?? [no regaline no] [2 1 3 5] [213045] ans -> min ele of the array

what if 1 -re no. precent??

[odd no. of

[2 1 3 0 4 5 x-1] negative

all other no.

negative except o what if 2-ve no are present?? [2 1 3 3 4 5 -2 -1] no] -re with all paritire except o. magnitude Defining factor -> no. of negative elements.

1. no regative no: min (ars.)

2. even no. of negative no.

Cases:

3. odd no. of negative no. Je a[i]] ali] to Mice 2 holes / mece = [4, -4, 2]

holes = [4,0,5]

atil #0

II [a[i]] » regative no. of min value

the max distance depends on the max distar any mice has to travel. Sost the mice positions 2) Sost the holes position mice [i] - holes [i]

[-4-0] = 4

i=1, |4-2| = 2

Meeting rooms ! (HW)