

# Normalization of Database Tables

# In this chapter, you will learn:

- What normalization is and what role it plays in the database design process
- About the normal forms 1NF, 2NF, 3NF, BCNF, and 4NF
- How normal forms can be transformed from lower normal forms to higher normal forms
- How normalization and ER modeling are used concurrently to produce a good database design
- How some situations require denormalization to generate information efficiently

# Database Tables and Normalization

- Normalization
  - Process for evaluating and correcting table structures to minimize data redundancies
    - Reduces data anomalies
  - Works through a series of stages called normal forms:
    - First normal form (1NF)
    - Second normal form (2NF)
    - Third normal form (3NF)

# Database Tables and Normalization (continued)

- Normalization (continued)
  - 2NF is better than 1NF; 3NF is better than 2NF
  - For most business database design purposes, 3NF is as high as we need to go in normalization process
  - Highest level of normalization is not always most desirable

# The Need for Normalization

- Example: Company that manages building projects
  - Charges its clients by billing hours spent on each contract
  - Hourly billing rate is dependent on employee's position
  - Periodically, report is generated that contains information displayed in Table 5.1

# The Need for Normalization (continued)

TABLE  
5.1

A Sample Report Layout

PROJ. NUM.	PROJECT NAME	EMPLOYEE NUMBER	EMPLOYEE NAME	JOB CLASS.	CHG/ HOUR	HOURS BILLED	TOTAL CHARGE
15	Evergreen	103	June E. Arbough	Elec. Engineer	\$ 85.50	23.8	\$ 2,011.10
		101	John G. News	Database Designer	\$105.00	19.4	\$ 2,037.00
		105	Alice K. Johnson*	Database Designer	\$105.00	35.7	\$ 3,748.50
		106	William Smithfield	Programmer	\$ 35.75	12.6	\$ 450.45
		102	David H. Senior	Systems Analyst	\$ 96.75	23.8	\$ 2,302.65
				<b>Subtotal</b>			<b>\$10,549.70</b>
18	Amber Wave	114	Annelise Jones	Applications Designer	\$ 48.10	25.6	\$ 1,183.26
		118	James J. Frommer	General Support	\$ 18.36	45.3	\$ 831.71
		104	Anne K. Ramoras*	Systems Analyst	\$ 96.75	32.4	\$ 3,135.70
		112	Darlene M. Smithson	DSS Analyst	\$ 45.95	45.0	\$ 2,021.80
				<b>Subtotal</b>			<b>\$ 7,172.47</b>
22	Rolling Tide	105	Alice K. Johnson	Database Designer	\$105.00	65.7	\$ 6,793.50
		104	Anne K. Ramoras	Systems Analyst	\$ 96.75	48.4	\$ 4,682.70
		113	Delbert K. Joenbrood*	Applications Designer	\$ 48.10	23.6	\$ 1,135.16
		111	Geoff B. Wabash	Clerical Support	\$ 26.87	22.0	\$ 591.14
		106	William Smithfield	Programmer	\$ 35.75	12.8	\$ 457.60
				<b>Subtotal</b>			<b>\$13,660.10</b>
25	Starflight	107	Maria D. Alonso	Programmer	\$ 35.75	25.6	\$ 879.45
		115	Travis B. Bawangi	Systems Analyst	\$ 96.75	45.8	\$ 4,431.15
		101	John G. News*	Database Designer	\$105.00	56.3	\$ 5,911.50
		114	Annelise Jones	Applications Designer	\$ 48.10	33.1	\$ 1,592.11
		108	Ralph B. Washington	Systems Analyst	\$ 96.75	23.6	\$ 2,283.30
		118	James J. Frommer	General Support	\$ 18.36	30.5	\$ 559.98
		112	Darlene M. Smithson	DSS Analyst	\$ 45.95	41.4	\$ 1,902.33
				<b>Subtotal</b>			<b>\$17,559.82</b>
				<b>Total</b>			<b>\$48,942.09</b>

Note: \* indicates project leader.

# The Need for Normalization (continued)

FIGURE  
5.1

Tabular representation of the report format

Table name: RPT\_FORMAT

Database name: Ch05\_ConstructCo

	PROJ_NUM	PROJ_NAME	EMP_NUM	EMP_NAME	JOB_CLASS	CHG_HOUR	HOURS
► 15	Evergreen	103	June E. Arbough	Elect. Engineer	\$84.50	23.8	
			101	John G. News	Database Designer	\$105.00	19.4
			105	Alice K. Johnson *	Database Designer	\$105.00	35.7
			106	William Smithfield	Programmer	\$35.75	12.6
			102	David H. Senior	Systems Analyst	\$96.75	23.8
18	Amber Wave	114	Annelise Jones	Applications Designer	\$48.10	24.6	
			118	James J. Frommer	General Support	\$18.36	45.3
			104	Anne K. Ramoras *	Systems Analyst	\$96.75	32.4
			112	Darlene M. Smithson	DSS Analyst	\$45.95	44.0
			22	Alice K. Johnson	Database Designer	\$105.00	64.7
22	Rolling Tide	105	104	Anne K. Ramoras	Systems Analyst	\$96.75	48.4
			113	Delbert K. Joenbrood *	Applications Designer	\$48.10	23.6
			111	Geoff B. Wabash	Clerical Support	\$26.87	22.0
			106	William Smithfield	Programmer	\$35.75	12.8
			25	Maria D. Alonso	Programmer	\$35.75	24.6
25	Starflight	107	115	Travis B. Bawangi	Systems Analyst	\$96.75	45.8
			101	John G. News *	Database Designer	\$105.00	56.3
			114	Annelise Jones	Applications Designer	\$48.10	33.1
			108	Ralph B. Washington	Systems Analyst	\$96.75	23.6
			118	James J. Frommer	General Support	\$18.36	30.5
			112	Darlene M. Smithson	DSS Analyst	\$45.95	41.4

# The Need for Normalization (continued)

- Structure of data set in Figure 5.1 does not handle data very well
- The table structure appears to work; report generated with ease
- Unfortunately, report may yield different results depending on what data anomaly has occurred

# The Normalization Process

- Each table represents a single subject
- No data item will be unnecessarily stored in more than one table
- All attributes in a table are dependent on the primary key

# The Normalization Process (continued)

TABLE  
5.2

Normal Forms

NORMAL FORM	CHARACTERISTIC	SECTION
First normal form (1NF)	Table format; no repeating groups and PK identified	5.3.1
Second normal form (2NF)	1NF and no partial dependencies	5.3.2
Third normal form (3NF)	2NF and no transitive dependencies	5.3.3
Boyce-Codd normal form (BCNF)	Every determinant is a candidate key (special case of 3NF)	5.6.1
Fourth normal form (4NF)	3NF and no independent multivalued dependencies	5.6.2

# Conversion to First Normal Form

- Repeating group
  - Derives its name from the fact that a group of multiple entries of same type can exist for any single key attribute occurrence
- Relational table must not contain repeating groups
- Normalizing table structure will reduce data redundancies
- Normalization is three-step procedure

# Conversion to First Normal Form (continued)

- Step 1: Eliminate the Repeating Groups
  - Present data in tabular format, where each cell has single value and there are no repeating groups
  - Eliminate repeating groups, eliminate nulls by making sure that each repeating group attribute contains an appropriate data value

# Conversion to First Normal Form (continued)

FIGURE  
5.2

A table in first normal form

Table name: DATA_ORG_1NF					Database name: Ch05_ConstructCo		
	PROJ_NUM	PROJ_NAME	EMP_NUM	EMP_NAME	JOB_CLASS	CHG_HOUR	HOURS
▶	15	Evergreen	103	June E. Arbough	Elect. Engineer	\$84.50	23.8
	15	Evergreen	101	John G. News	Database Designer	\$105.00	19.4
	15	Evergreen	105	Alice K. Johnson *	Database Designer	\$105.00	35.7
	15	Evergreen	106	William Smithfield	Programmer	\$35.75	12.6
	15	Evergreen	102	David H. Senior	Systems Analyst	\$96.75	23.8
	18	Amber Wave	114	Annelise Jones	Applications Designer	\$48.10	24.6
	18	Amber Wave	118	James J. Frommer	General Support	\$18.36	45.3
	18	Amber Wave	104	Anne K. Ramoras *	Systems Analyst	\$96.75	32.4
	18	Amber Wave	112	Darlene M. Smithson	DSS Analyst	\$45.95	44.0
	22	Rolling Tide	105	Alice K. Johnson	Database Designer	\$105.00	64.7
	22	Rolling Tide	104	Anne K. Ramoras	Systems Analyst	\$96.75	48.4
	22	Rolling Tide	113	Delbert K. Joenbrood *	Applications Designer	\$48.10	23.6
	22	Rolling Tide	111	Geoff B. Wabash	Clerical Support	\$26.87	22.0
	22	Rolling Tide	106	William Smithfield	Programmer	\$35.75	12.8
	25	Starflight	107	Maria D. Alonzo	Programmer	\$35.75	24.6
	25	Starflight	115	Travis B. Bawangi	Systems Analyst	\$96.75	45.8
	25	Starflight	101	John G. News *	Database Designer	\$105.00	56.3
	25	Starflight	114	Annelise Jones	Applications Designer	\$48.10	33.1
	25	Starflight	108	Ralph B. Washington	Systems Analyst	\$96.75	23.6
	25	Starflight	118	James J. Frommer	General Support	\$18.36	30.5
	25	Starflight	112	Darlene M. Smithson	DSS Analyst	\$45.95	41.4

# Conversion to First Normal Form (continued)

- Step 2: Identify the Primary Key
  - Primary key must uniquely identify attribute value
  - New key must be composed

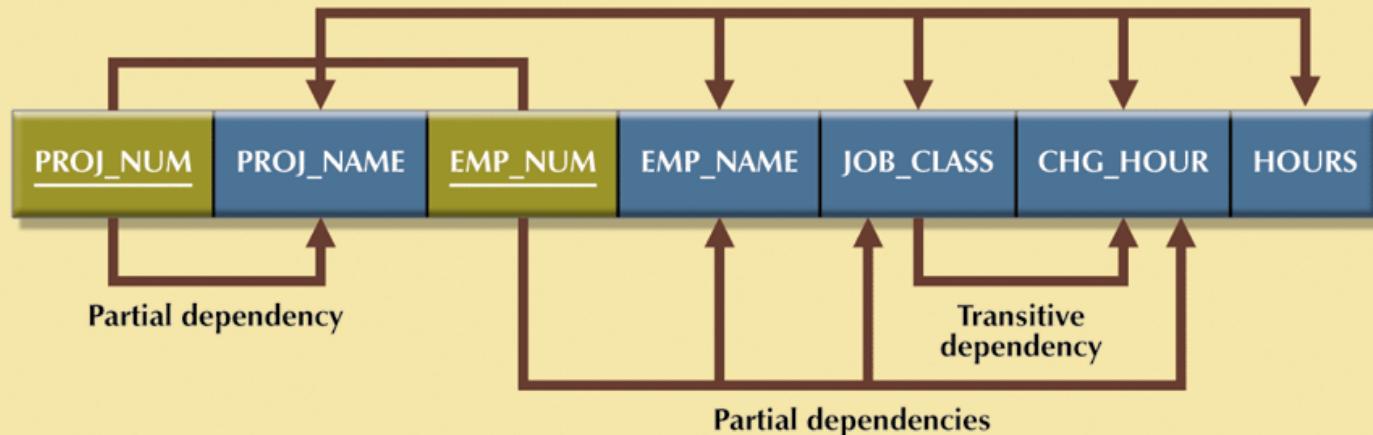
# Conversion to First Normal Form (continued)

- Step 3: Identify All Dependencies
  - Dependencies can be depicted with help of a diagram
  - Dependency diagram:
    - Depicts all dependencies found within given table structure
    - Helpful in getting bird's-eye view of all relationships among table's attributes
    - Makes it less likely that will overlook an important dependency

# Conversion to First Normal Form (continued)

FIGURE  
5.3

First normal form (1NF) dependency diagram



1NF (PROJ\_NUM, EMP\_NUM, PROJ\_NAME, EMP\_NAME, JOB\_CLASS, CHG\_HOURS, HOURS)

PARTIAL DEPENDENCIES:

(PROJ\_NUM → PROJ\_NAME)

(EMP\_NUM → EMP\_NAME, JOB\_CLASS, CHG\_HOUR)

TRANSITIVE DEPENDENCY:

(JOB CLASS → CHG\_HOUR)

# Conversion to First Normal Form (continued)

- First normal form describes tabular format in which:
  - All key attributes are defined
  - There are no repeating groups in the table
  - All attributes are dependent on primary key
- All relational tables satisfy 1NF requirements
- Some tables contain partial dependencies
  - Dependencies based on only part of the primary key
  - Sometimes used for performance reasons, but should be used with caution
  - Still subject to data redundancies

# Conversion to Second Normal Form

- Relational database design can be improved by converting the database into second normal form (2NF)
- Two steps

# Conversion to Second Normal Form (continued)

- Step 1: Write Each Key Component on a Separate Line
  - Write each key component on separate line, then write original (composite) key on last line
  - Each component will become key in new table

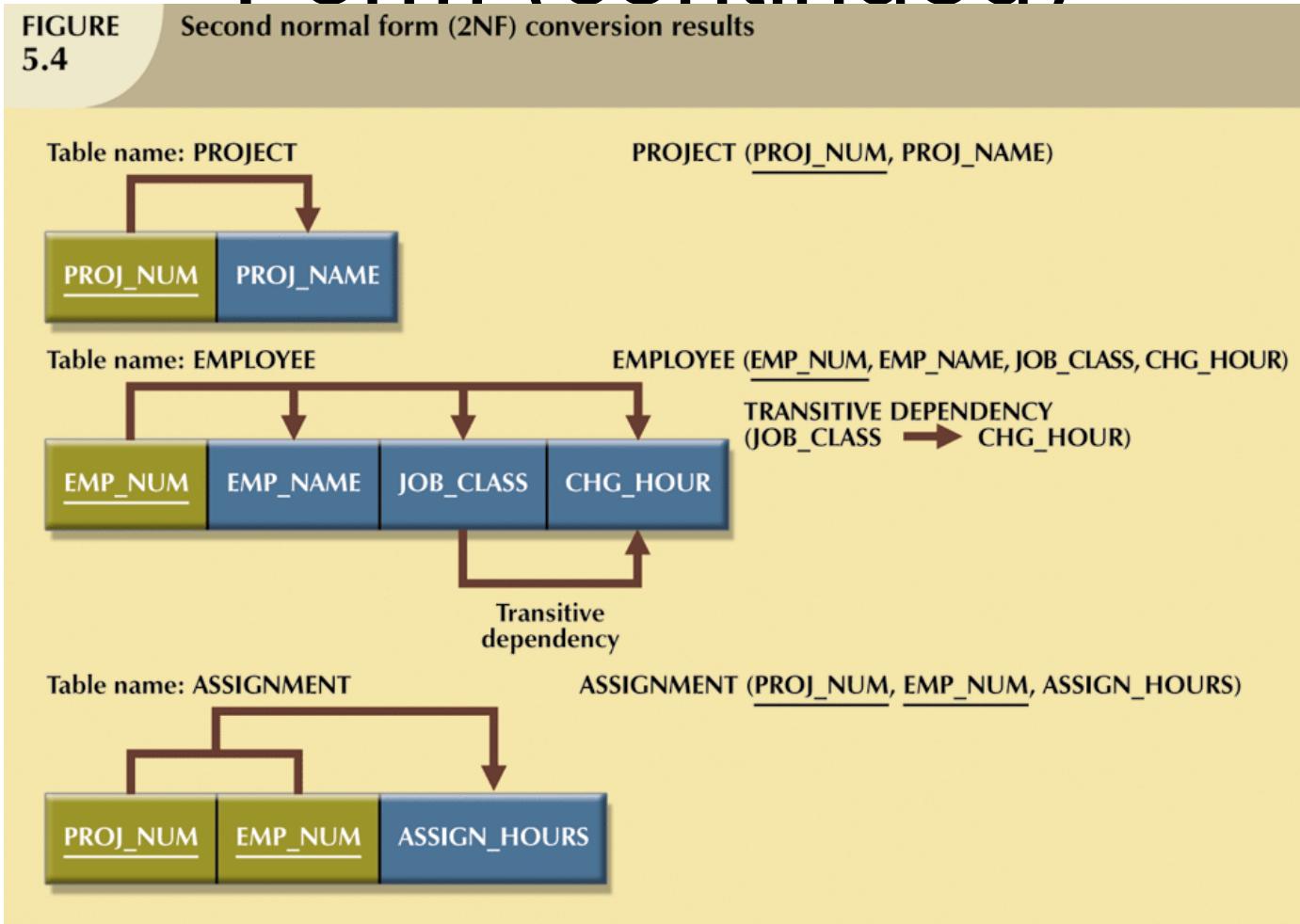
# Conversion to Second Normal Form (continued)

- Step 2: Assign Corresponding Dependent Attributes
  - Determine those attributes that are dependent on other attributes
  - At this point, most anomalies have been eliminated

# Conversion to Second Normal Form (continued)

FIGURE  
5.4

Second normal form (2NF) conversion results



# Conversion to Second Normal Form (continued)

- Table is in second normal form (2NF) when:
  - It is in 1NF and
  - It includes no partial dependencies:
    - No attribute is dependent on only portion of primary key

# Conversion to Third Normal Form

- Data anomalies created are easily eliminated by completing three steps
- Step 1: Identify Each New Determinant
  - For every transitive dependency, write its determinant as PK for new table
    - Determinant
      - Any attribute whose value determines other values within a row

# Conversion to Third Normal Form (continued)

- Step 2: Identify the Dependent Attributes
  - Identify attributes dependent on each determinant identified in Step 1 and identify dependency
  - Name table to reflect its contents and function

# Conversion to Third Normal Form (continued)

- Step 3: Remove the Dependent Attributes from Transitive Dependencies
  - Eliminate all dependent attributes in transitive relationship(s) from each of the tables that have such a transitive relationship
  - Draw new dependency diagram to show all tables defined in Steps 1–3
  - Check new tables as well as tables modified in Step 3 to make sure that each table has determinant and that no table contains inappropriate dependencies

# Conversion to Third Normal Form (continued)

FIGURE  
5.5

Third normal form (3NF) conversion results

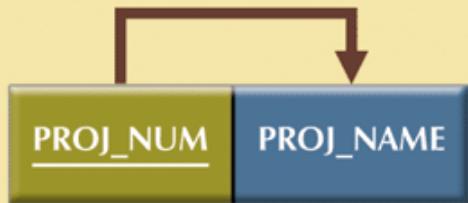


Table name: PROJECT

PROJECT (PROJ\_NUM, PROJ\_NAME)



Table name: JOB

JOB (JOB\_CLASS, CHG\_HOUR)

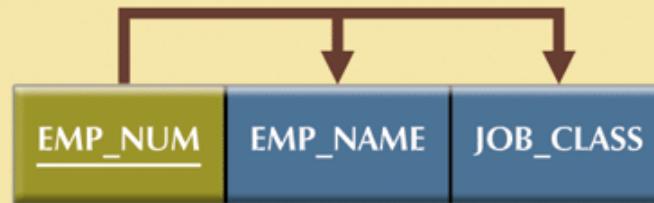


Table name: EMPLOYEE

EMPLOYEE (EMP\_NUM, EMP\_NAME, JOB\_CLASS)

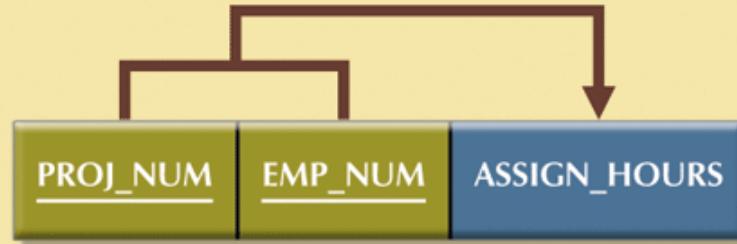


Table name: ASSIGNMENT

ASSIGNMENT (PROJ\_NUM, EMP\_NUM, ASSIGN\_HOURS)

# Conversion to Third Normal Form (continued)

- A table is in third normal form (3NF) when both of the following are true:
  - It is in 2NF
  - It contains no transitive dependencies

# Improving the Design

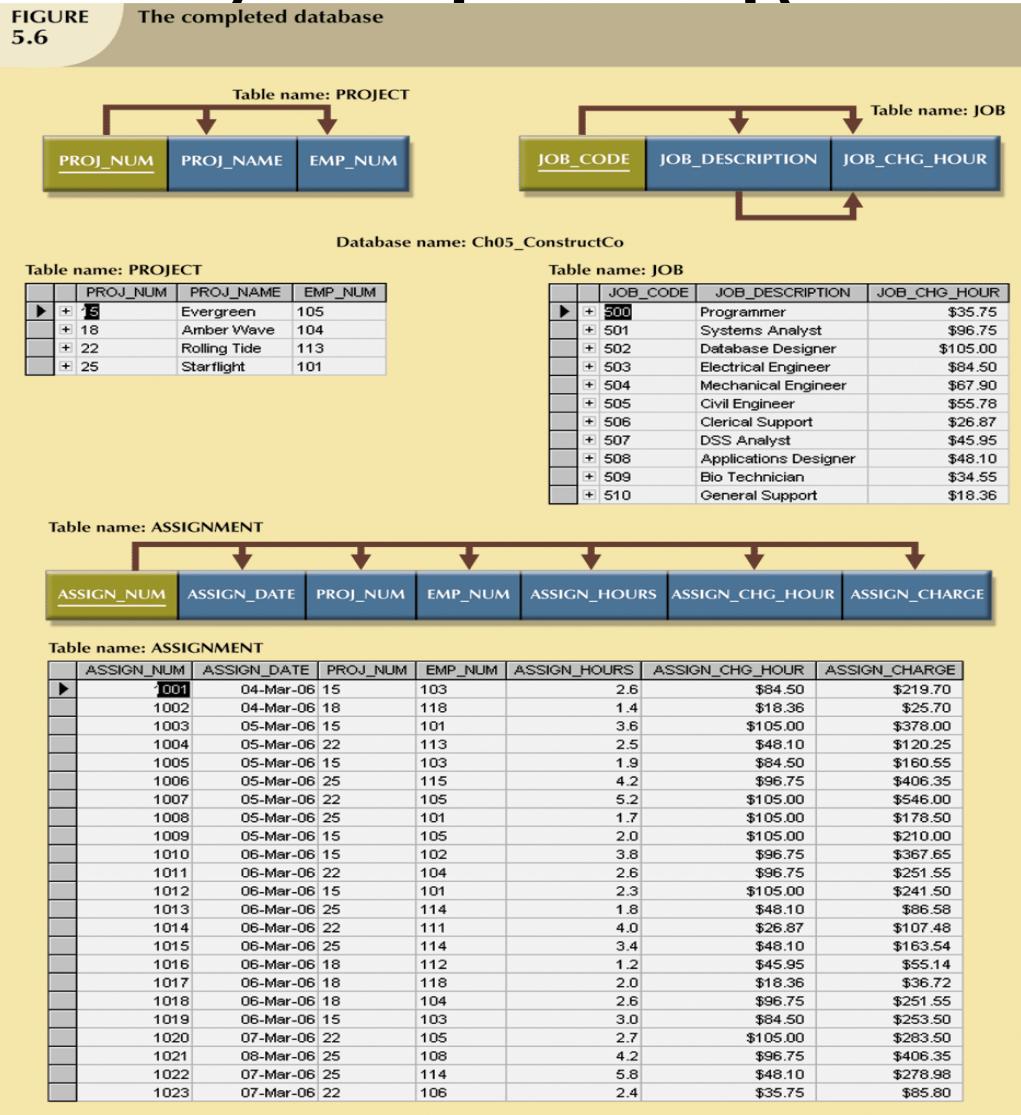
- Table structures are cleaned up to eliminate troublesome initial partial and transitive dependencies
- Normalization cannot, by itself, be relied on to make good designs
- It is valuable because its use helps eliminate data redundancies

# Improving the Design (continued)

- Issues to address in order to produce a good normalized set of tables:
  - Evaluate PK Assignments
  - Evaluate Naming Conventions
  - Refine Attribute Atomicity
  - Identify New Attributes
  - Identify New Relationships
  - Refine Primary Keys as Required for Data Granularity
  - Maintain Historical Accuracy
  - Evaluate Using Derived Attributes

# Improving the Design

**FIGURE  
5.6**



# Improving the Design (continued)

FIGURE  
5.6

The completed database (continued)

Table name: EMPLOYEE



Table name: EMPLOYEE

	EMP_NUM	EMP_LNAME	EMP_FNAME	EMP_INITIAL	EMP_HIREDATE	JOB_CODE
►	+ 101	News	John	G	08-Nov-00	502
	+ 102	Senior	David	H	12-Jul-89	501
	+ 103	Arbough	June	E	01-Dec-97	503
	+ 104	Ramoras	Anne	K	15-Nov-88	501
	+ 105	Johnson	Alice	K	01-Feb-94	502
	+ 106	Smithfield	William		22-Jun-05	500
	+ 107	Alonzo	Maria	D	10-Oct-94	500
	+ 108	Washington	Ralph	B	22-Aug-89	501
	+ 109	Smith	Larry	W	18-Jul-99	501
	+ 110	Olenko	Gerald	A	11-Dec-96	505
	+ 111	Wabash	Geoff	B	04-Apr-89	506
	+ 112	Smithson	Darlene	M	23-Oct-95	507
	+ 113	Joenbrood	Delbert	K	15-Nov-94	508
	+ 114	Jones	Annelise		20-Aug-91	508
	+ 115	Bawangi	Travis	B	25-Jan-90	501
	+ 116	Pratt	Gerald	L	05-Mar-95	510
	+ 117	Williamson	Angie	H	19-Jun-94	509
	+ 118	Frommer	James	J	04-Jan-06	510

# Surrogate Key Considerations

- When primary key is considered to be unsuitable, designers use surrogate keys
- Data entries in Table 5.3 are inappropriate because they duplicate existing records
  - Yet there has been no violation of either entity integrity or referential integrity

# Surrogate Key Considerations (continued)

TABLE  
5.3

Duplicate Entries in the Job Table

JOB_CODE	JOB_DESCRIPTION	JOB_CHG_HOUR
511	Programmer	\$35.75
512	Programmer	\$35.75

# The Boyce-Codd Normal Form (BCNF)

- Every determinant in table is a candidate key
  - Has same characteristics as primary key, but for some reason, not chosen to be primary key
- When table contains only one candidate key, the 3NF and the BCNF are equivalent
- BCNF can be violated only when table contains more than one candidate key

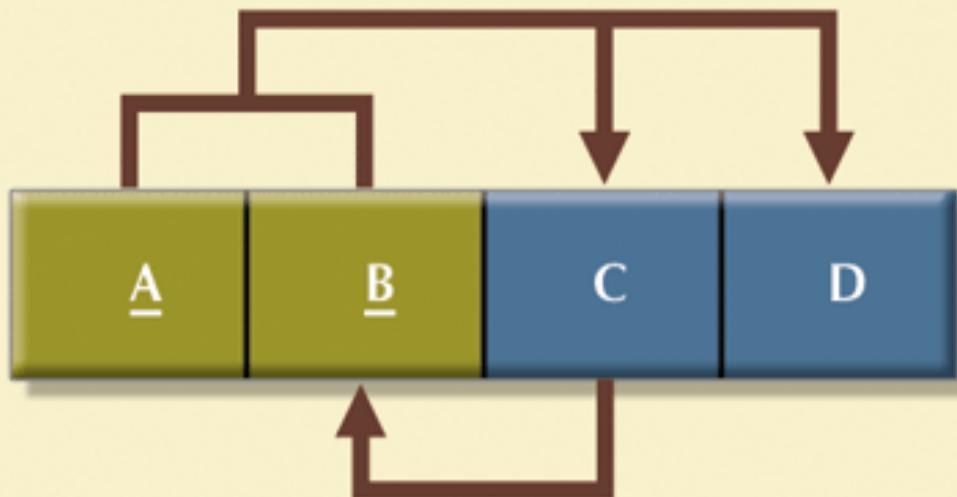
# The Boyce-Codd Normal Form (BCNF) (continued)

- Most designers consider the BCNF as special case of 3NF
- Table is in 3NF when it is in 2NF and there are no transitive dependencies
- Table can be in 3NF and fails to meet BCNF
  - No partial dependencies, nor does it contain transitive dependencies
  - A nonkey attribute is the determinant of a key attribute

# The Boyce-Codd Normal Form (BCNF) (continued)

**FIGURE  
5.7**

A table that is in 3NF but not in BCNF



# The Boyce-Codd Normal Form (BCNF) (continued)

TABLE  
5.4

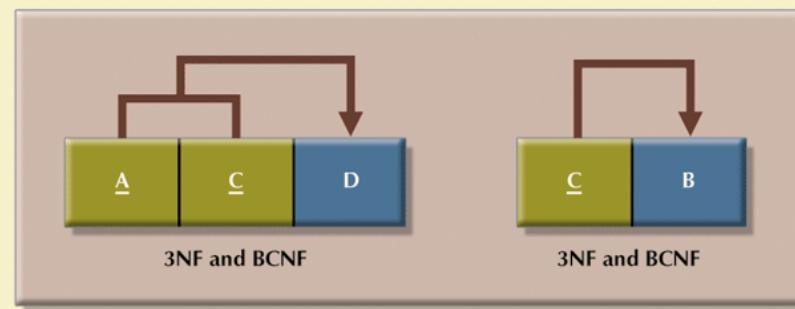
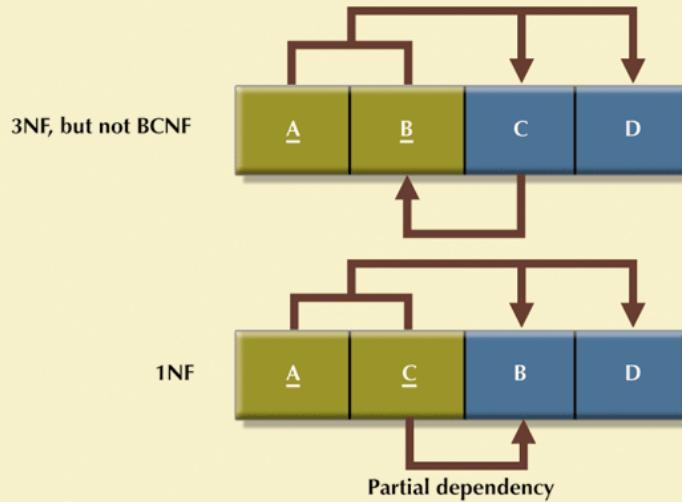
Sample Data for a BCNF Conversion

STU_ID	STAFF_ID	CLASS_CODE	ENROLL_GRADE
125	25	21334	A
125	20	32456	C
135	20	28458	B
144	25	27563	C
144	20	32456	B

# The Boyce-Codd Normal Form (BCNF) (continued)

FIGURE  
5.8

Decomposition to BCNF



# Fourth Normal Form (4NF)

- Table is in fourth normal form (4NF) when both of the following are true:
  - It is in 3NF
  - Has no multiple sets of multivalued dependencies
- 4NF is largely academic if tables conform to following two rules:
  - All attributes must be dependent on primary key, but independent of each other
  - No row contains two or more multivalued facts about an entity

# Fourth Normal Form (4NF) (continued)

**FIGURE  
5.10**

Tables with multivalued dependencies

Database name: Ch05\_Service

Table name: VOLUNTEER\_V1

	EMP_NUM	ORG_CODE	ASSIGN_NUM
▶	10123	RC	1
	10123	LW	3
	10123		4

Table name: VOLUNTEER\_V2

	EMP_NUM	ORG_CODE	ASSIGN_NUM
▶	10123	RC	
	10123	LW	
	10123		1
	10123		3
	10223		4

Table name: VOLUNTEER\_V3

	EMP_NUM	ORG_CODE	ASSIGN_NUM
▶	10123	RC	1
	10123	RC	3
	10123	LW	4

# Fourth Normal Form (4NF)

## (continued)

FIGURE  
5.11

A set of tables in 4NF

Table name: EMPLOYEE

	EMP_NUM	EMP_LNAME
▶	10121	Rogers
+	10122	O'Leary
+	10123	Panera
+	10124	Johnson

Database name: Ch05\_Service

Table name: PROJECT

	PROJ_CODE	PROJ_NAME	PROJ_BUDGET
▶	1	BeThere	\$1,023,245.00
+	2	BlueMoon	\$20,198,608.00
+	3	GreenThumb	\$3,234,456.00
+	4	Gofast	\$5,674,000.00
+	5	GoSlow	\$1,002,500.00

Table name: ORGANIZATION

	ORG_CODE	ORG_NAME
▶	RC	Red Cross
+	UW	United Way
+	WF	Wildlife Fund

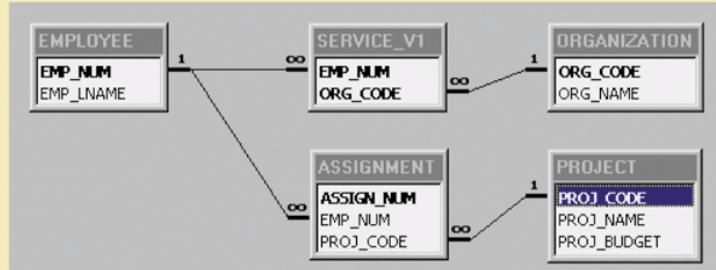
Table name: ASSIGNMENT

	ASSIGN_NUM	EMP_NUM	PROJ_CODE
▶	1	10123	1
2	10121	2	
3	10123	3	
4	10123	4	
5	10121	1	
6	10124	2	
7	10124	3	
8	10124	5	

Table name: SERVICE\_V1

	EMP_NUM	ORG_CODE
▶	10123	RC
10123	UW	
10123	WF	

The relational diagram



# Normalization and Database Design

- Normalization should be part of design process
- Make sure that proposed entities meet required normal form before table structures are created
- Many real-world databases have been improperly designed or burdened with anomalies if improperly modified during course of time
- You may be asked to redesign and modify existing databases

# Normalization and Database Design (continued)

- ER diagram
  - Provides big picture, or macro view, of an organization's data requirements and operations
  - Created through an iterative process
    - Identifying relevant entities, their attributes and their relationship
    - Use results to identify additional entities and attributes

# Normalization and Database Design (continued)

- Normalization procedures
  - Focus on characteristics of specific entities
  - Represents micro view of entities within ER diagram
- Difficult to separate normalization process from ER modeling process
- Two techniques should be used concurrently

# Normalization and Database Design (continued)

**FIGURE  
5.12**

Initial contracting company  
ERD

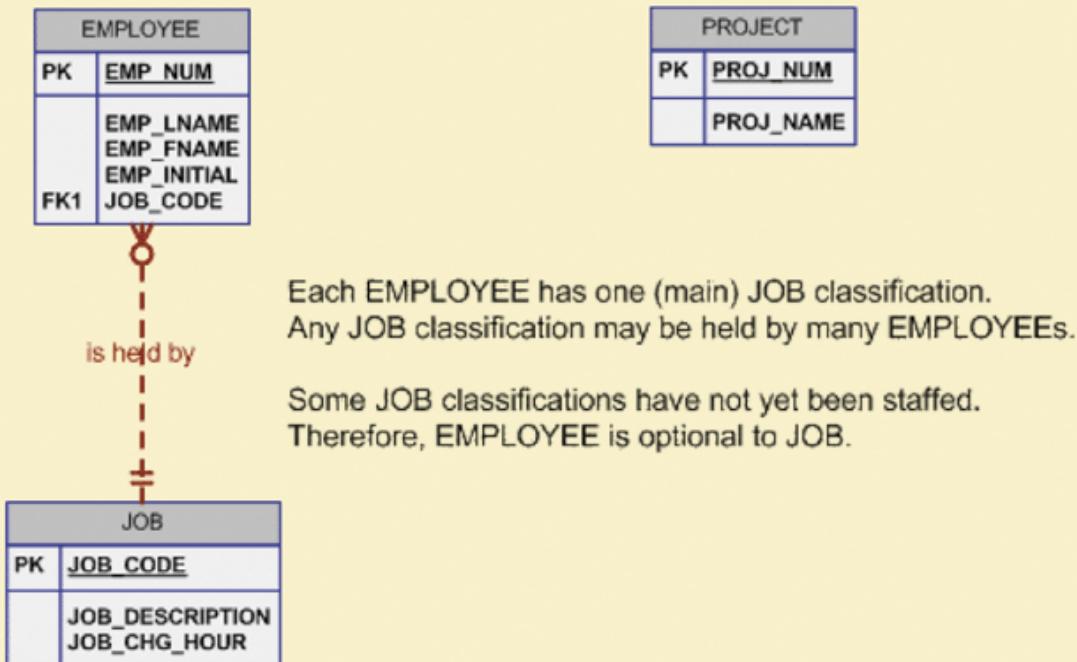
EMPLOYEE	
PK	<u>EMP_NUM</u>
	EMP_LNAME
	EMP_FNAME
	EMP_INITIAL
	JOB_DESCRIPTION
	JOB_CHG_HOUR

PROJECT	
PK	<u>PROJ_NUM</u>
	PROJ_NAME

# Normalization and Database Design (continued)

FIGURE  
5.13

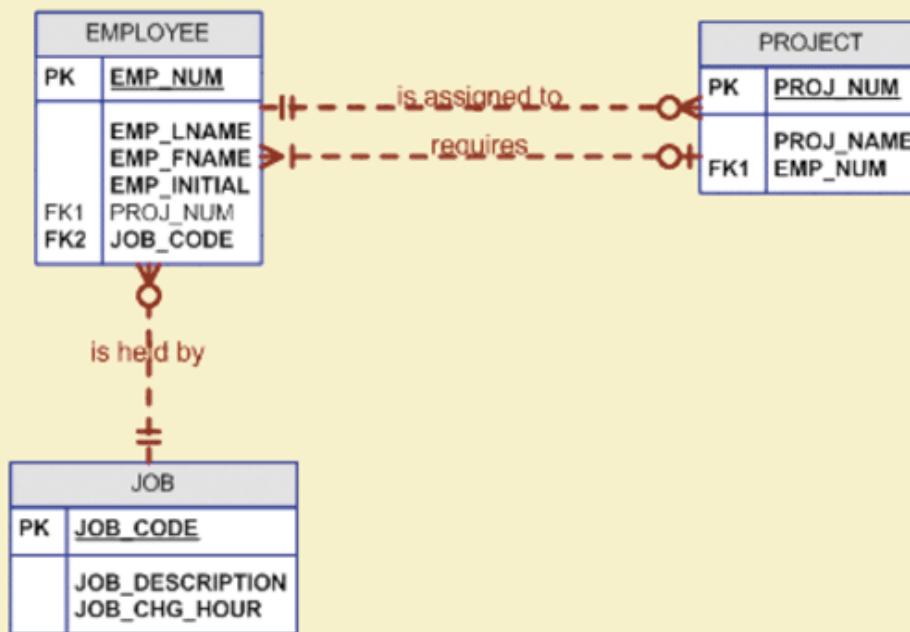
Modified contracting company ERD



# Normalization and Database Design (continued)

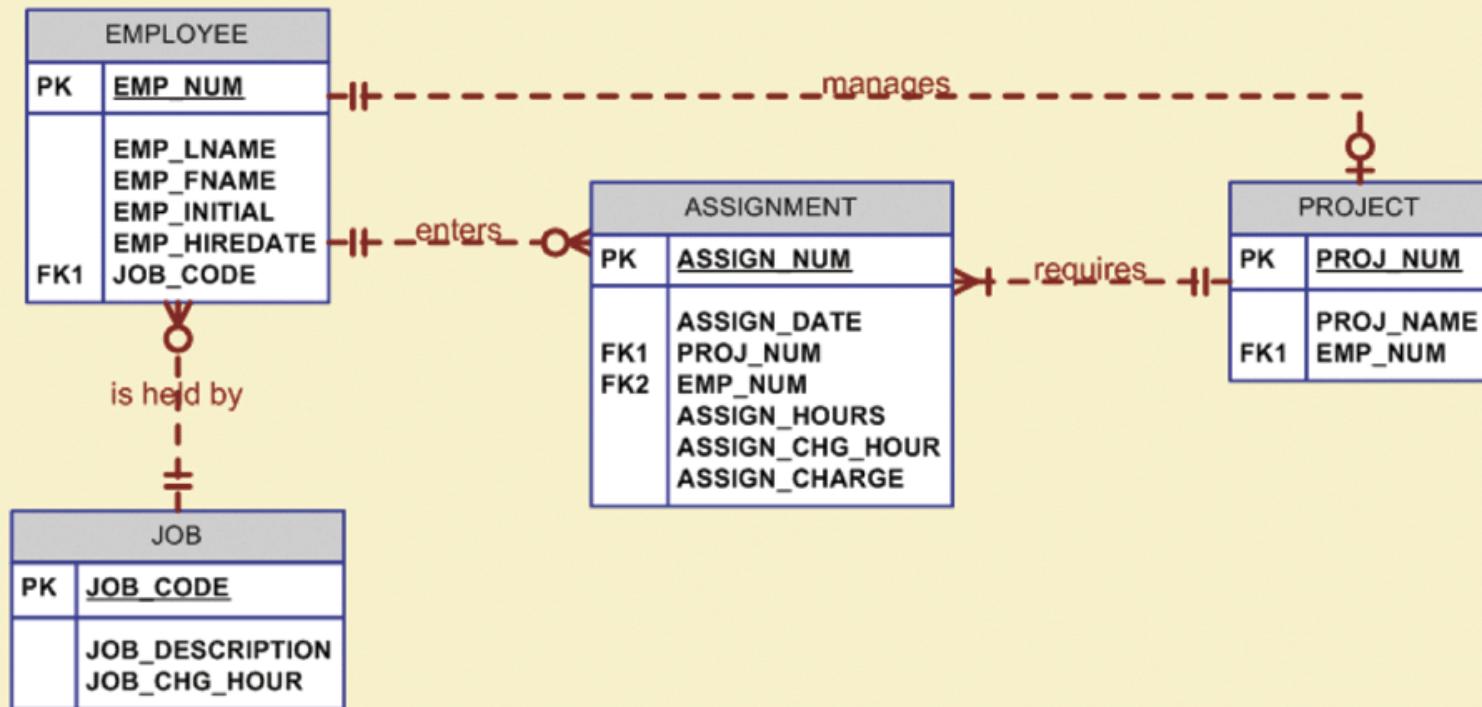
FIGURE  
5.14

Incorrect M:N relationship representation



# Normalization and Database Design (continued)

FIGURE 5.15 Final contracting company ERD



# Normalization and Database Design (continued)

FIGURE  
5.16

The implemented database

Table name: EMPLOYEE						Database name: Ch05_ConstructCo			
	EMP_NUM	EMP_LNAME	EMP_FNAME	EMP_INITIAL	EMP_HIREDATE	JOB_CODE	JOB_CODE	JOB_DESCRIPTION	JOB_CHG_HOUR
► +	101	News	John	G	08-Nov-00	502	► +	500	\$35.75
► +	102	Senior	David	H	12-Jul-89	501	► +	501	\$96.75
► +	103	Arbrough	June	E	01-Dec-97	503	► +	502	\$105.00
► +	104	Ramoras	Anne	K	15-Nov-88	501	► +	503	\$84.50
► +	105	Johnson	Alice	K	01-Feb-94	502	► +	504	\$67.90
► +	106	Smithfield	William		22-Jun-05	500	► +	505	\$55.78
► +	107	Alonzo	Maria	D	10-Oct-94	500	► +	506	\$26.87
► +	108	Washington	Ralph	B	22-Aug-89	501	► +	507	\$45.95
► +	109	Smith	Larry	V	18-Jul-99	501	► +	508	\$48.10
► +	110	Olenko	Gerald	A	11-Dec-98	505	► +	509	\$34.55
► +	111	Wabash	Geoff	B	04-Apr-89	506	► +	510	\$18.36
► +	112	Smithson	Darlene	M	23-Oct-95	507			
► +	113	Joinbrood	Delbert	K	15-Nov-94	508			
► +	114	Jones	Annelise		20-Aug-91	508			
► +	115	Bawangi	Travis	B	25-Jan-90	501			
► +	116	Prat	Gerald	L	05-Mar-95	510			
► +	117	Williamson	Angie	H	19-Jun-94	509			
► +	118	Frommer	James	J	04-Jan-06	510			

Table name: PROJECT			
	PROJ_NUM	PROJ_NAME	EMP_NUM
► +	5	Evergreen	105
► +	18	Amber Wave	104
► +	22	Rolling Tide	113
► +	25	Starflight	101

Table name: ASSIGNMENT						
	ASSIGN_NUM	ASSIGN_DATE	PROJ_NUM	EMP_NUM	ASSIGN_HOURS	ASSIGN_CHG_HOUR
► +	1001	04-Mar-06 15	103		2.6	\$84.50
	1002	04-Mar-06 18	118		1.4	\$18.36
	1003	05-Mar-06 15	101		3.6	\$105.00
	1004	05-Mar-06 22	113		2.5	\$48.10
	1005	05-Mar-06 15	103		1.9	\$84.50
	1006	05-Mar-06 25	115		4.2	\$96.75
	1007	05-Mar-06 22	105		5.2	\$105.00
	1008	05-Mar-06 25	101		1.7	\$105.00
	1009	05-Mar-06 15	105		2.0	\$105.00
	1010	06-Mar-06 15	102		3.8	\$96.75
	1011	06-Mar-06 22	104		2.6	\$96.75
	1012	06-Mar-06 15	101		2.3	\$105.00
	1013	06-Mar-06 25	114		1.8	\$48.10
	1014	06-Mar-06 22	111		4.0	\$26.87
	1015	06-Mar-06 25	114		3.4	\$48.10
	1016	06-Mar-06 18	112		1.2	\$45.95
	1017	06-Mar-06 18	118		2.0	\$18.36
	1018	06-Mar-06 18	104		2.6	\$96.75
	1019	06-Mar-06 15	103		3.0	\$84.50
	1020	07-Mar-06 22	105		2.7	\$105.00
	1021	08-Mar-06 25	108		4.2	\$96.75
	1022	07-Mar-06 25	114		5.8	\$48.10
	1023	07-Mar-06 22	106		2.4	\$35.75

# Denormalization

- Creation of normalized relations is important database design goal
- Processing requirements should also be a goal
- If tables decomposed to conform to normalization requirements:
  - Number of database tables expands

# Denormalization (continued)

- Joining the larger number of tables takes additional input/output (I/O) operations and processing logic, thereby reducing system speed
- Conflicts between design efficiency, information requirements, and processing speed are often resolved through compromises that may include denormalization

# Denormalization (continued)

- Unnormalized tables in production database tend to suffer from these defects:
  - Data updates are less efficient because programs that read and update tables must deal with larger tables
  - Indexing is more cumbersome
  - Unnormalized tables yield no simple strategies for creating virtual tables known as views

# Denormalization (continued)

- Use denormalization cautiously
- Understand why—under some circumstances—unnormalized tables are better choice

# Summary

- Normalization is technique used to design tables in which data redundancies are minimized
- First three normal forms (1NF, 2NF, and 3NF) are most commonly encountered
- Table is in 1NF when all key attributes are defined and when all remaining attributes are dependent on primary key

# Summary (continued)

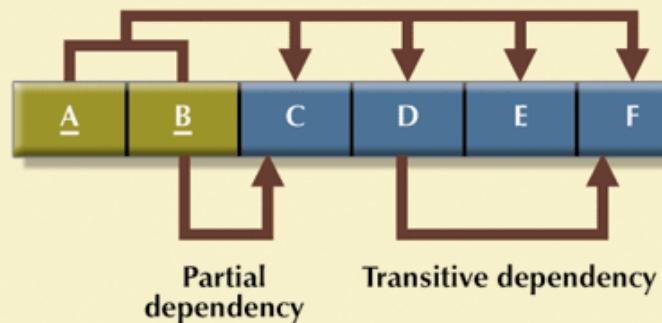
- Table is in 2NF when it is in 1NF and contains no partial dependencies
- Table is in 3NF when it is in 2NF and contains no transitive dependencies
- Table that is not in 3NF may be split into new tables until all of the tables meet 3NF requirements
- Normalization is important part—but only part —of design process

# Summary (continued)

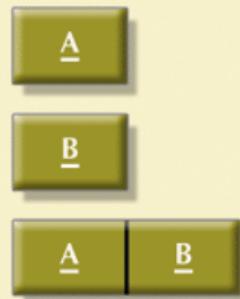
FIGURE  
5.17

The initial 1NF structure

The Initial 1NF Structure



Step 1: Write each PK component on a separate line; then write the original (composite) PK on the last line.

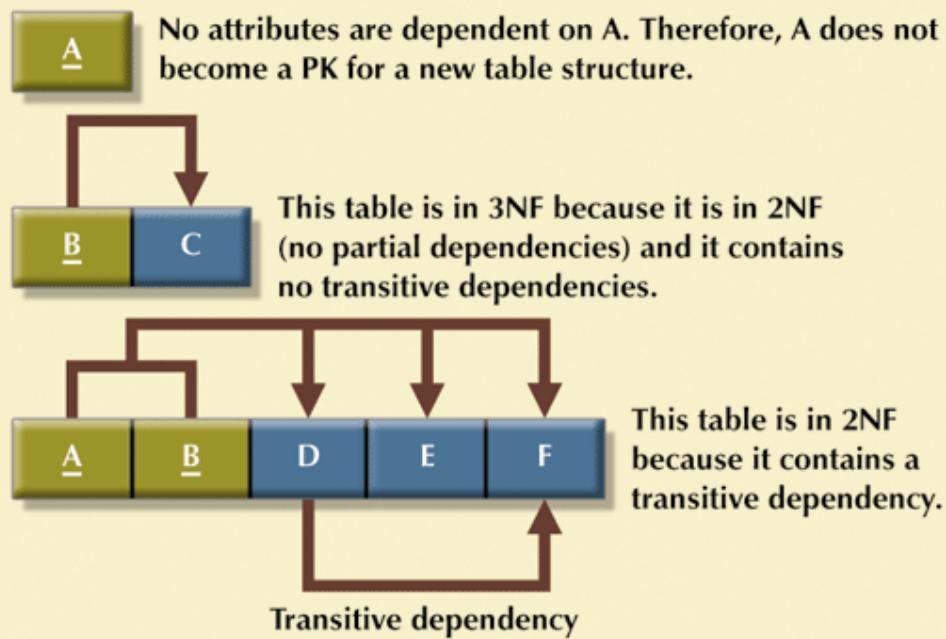


# Summary (continued)

FIGURE  
5.18

## Identifying possible PK attributes

Step 2: Place all dependent attributes with the PK attributes identified in Step 1.



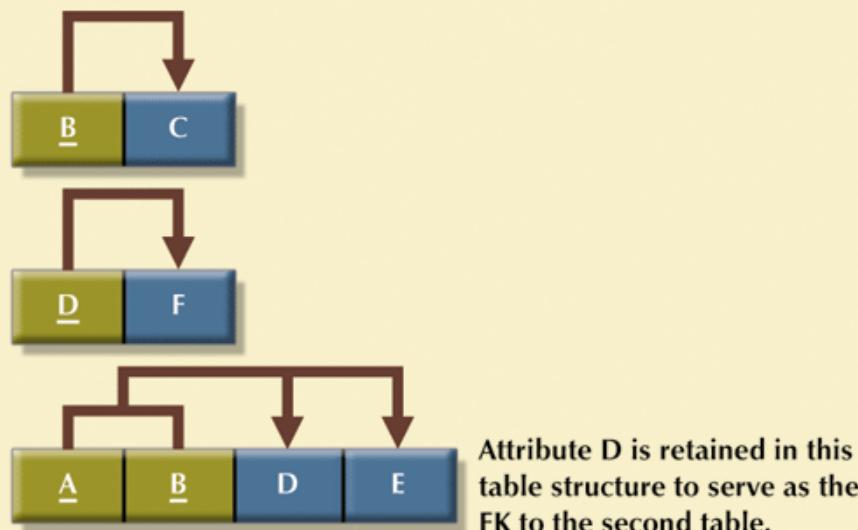
# Summary (continued)

FIGURE  
5.19

Table structures based on the selected PKs

Step 3: Remove all transitive dependencies identified in Step 2 and retain all 3NF structures.

All tables are in 3NF because they are in 2NF (no partial dependencies) and they do not contain transitive dependencies.



# Summary (continued)

- Table in 3NF may contain multivalued dependencies that produce either numerous null values or redundant data
- It may be necessary to convert 3NF table to fourth normal form (4NF) by
  - Splitting table to remove multivalued dependencies
- Tables are sometimes denormalized to yield less I/O which increases processing speed