1. Use the simplex method to describe all the optimal solutions of the following problem.

Take x_3 and s_2 leaves:

Take x_2 and s_1 leaves:

Finally, we can obtain that the max z = 8, $s_1, s_2 = 0$

Thus

$$\begin{cases} & \forall x_i \geq 0 \\ x_1 + 7x_4 - x_2 = -1 \\ x_1 + x_3 + 5x_4 = -1 \end{cases}$$

It's a polygon. We find its vertices assigning 2 coordinates to 0 and get only 3 vertices:

$$(1, 2, 0, 0), (0, 1, 1, 0), (0, 2.4, 0, 0.2)$$

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2. Solve the following LP problem.

Take x_2 enters and s_1 leaves:

Note that $x \le 5$ and obtain that $x_2 = 5 - x_1 + \frac{1}{2}x_3 - \frac{1}{2}S_1$

It is unbound. Then we can make x_3 as large as we wish.

$$x_1 = 0$$
, $x_2 = 5 + \frac{1}{2}t$, $x_3 = t$, $x_1 = 0$, $x_2 = 20$, $x_3 = 25 + \frac{1}{2}t$, $x_3 = 15 + \frac{1}{2}t$

Where t can be any positive number.

Solve the following problems by the two-phase simplex method:

(a)

(b)

(c)

(a)

max

 x_2

-1

 $3x_1$

max

 $\frac{1}{2}S_1$

 $S_1, S_2, S_3, x_0, x_1, x_2$

 $\{-x_0\}$ artificial obj

 x_1

Finally z = 5

$$x_1 = 1$$
, $x_2 = 2$, $S_1 = 0$, $S_2 = 0$, $S_3 = 0$

(b)

$$\Rightarrow \begin{array}{c} \max \quad -3 \ + \ x_1 \ + \ x_2 \ - \ S_2 \\ \text{s.t.} \quad S_1 \ = \ 2 \ - \ 2x_1 \ + \ S_2 \\ x_0 \ = \ 3 \ + \ x_1 \ - \ x_2 \ + \ S_2 \\ S_3 \ = \ 5 \ - \ 3x_1 \ - \ 2x_2 \ + \ S_2 \\ S_1, S_2, S_3, x_1, x_2 \ \geq \ 0 \\ \hline x_1 \ = \ 1 \ + \ \frac{1}{2}S_2 \ - \ \frac{1}{2}S_1 \ - \ \frac{1}{2}S_2 \\ S_3 \ = \ 2 \ - \ 2x_2 \ + \ \frac{1}{2}S_1 \ + \ \frac{1}{2}S_2 \\ S_3 \ = \ 2 \ - \ 2x_2 \ + \ \frac{3}{2}S_1 \ - \ \frac{1}{2}S_2 \\ S_1, S_2, S_3, x_0, x_1, x_2 \ \geq \ 0 \\ \hline x_2 \ = \ 1 \ + \ \frac{3}{4}S_1 \ - \ \frac{1}{4}S_2 \ - \ \frac{1}{2}S_3 \end{array}$$

$$\Rightarrow \begin{array}{c} \max \quad -1 \quad + \quad \frac{1}{4}S_{1} \quad - \quad \frac{3}{4}S_{2} \quad - \quad \frac{1}{2}S_{3} \\ \text{s.t.} \quad x_{1} \quad = \quad 1 \quad - \quad \frac{1}{2}S_{1} \quad + \quad \frac{1}{2}S_{2} \\ x_{0} \quad = \quad 1 \quad - \quad \frac{1}{4}S_{1} \quad + \quad \frac{3}{4}S_{2} \quad + \quad \frac{1}{2}S_{3} \\ x_{2} \quad = \quad 1 \quad + \quad \frac{3}{4}S_{2} \quad - \quad \frac{1}{4}S_{2} \quad - \quad \frac{1}{2}S_{3} \\ x_{3} \quad = \quad 1 \quad + \quad \frac{3}{4}S_{2} \quad - \quad \frac{1}{4}S_{2} \quad - \quad \frac{1}{2}S_{3} \\ x_{4} \quad = \quad 1 \quad + \quad \frac{3}{4}S_{2} \quad - \quad \frac{1}{4}S_{2} \quad - \quad \frac{1}{2}S_{3} \\ x_{5} \quad = \quad 1 \quad + \quad \frac{1}{2}S_{5} \quad + \quad \frac{1}{2}S_{5} \\ x_{7} \quad = \quad 1 \quad + \quad \frac{1}{2}S_{7} \quad + \quad \frac{1}{2}S_{7} \quad + \quad \frac{1}{2}S_{7} \quad + \quad \frac{1}{2}S_{7} \\ x_{7} \quad = \quad 1 \quad + \quad \frac{1}{2}S_{7} \quad + \quad \frac{1}{$$

 \implies Optimal obj < 0. The Linear Problem is infeasible

(c)

$$\begin{array}{c}
\text{max} \quad -3 \quad + \quad x_1 \quad + \quad x_2 \quad - \quad S_2 \\
\text{s.t.} \quad S_1 \quad = \quad 2 \quad - \quad 2x_1 \quad + \quad S_2 \\
x_0 \quad = \quad 3 \quad - \quad x_1 \quad - \quad x_2 \quad + \quad S_2 \\
S_3 \quad = \quad 5 \quad - \quad 3x_1 \quad + \quad S_2 \\
& \quad S_1, S_2, S_3, x_0, x_1, x_2 \\
\hline
x_1 \quad = \quad 1 \quad + \quad \frac{1}{2}S_2 \quad - \quad \frac{1}{2}S_1
\end{array}$$

$$\Rightarrow \begin{array}{c}
\text{max} \quad -2 \quad + \quad x_2 \quad - \quad \frac{1}{2}S_1 \quad - \quad \frac{1}{2}S_2 \\
\text{s.t.} \quad x_1 \quad = \quad 1 \quad + \quad \frac{1}{2}S_2 \quad - \quad \frac{1}{2}S_1 \quad + \quad \frac{1}{2}S_2 \\
S_3 \quad = \quad 2 \quad - \quad x_2 \quad + \quad \frac{1}{2}S_1 \quad - \quad \frac{1}{2}S_2 \quad - \quad$$

$$\Longrightarrow \begin{array}{c} \max \quad 5 \quad -S_1 \quad + \quad 2S_2 \\ \text{s.t.} \quad x_1 \quad = \quad 1 \quad - \quad \frac{1}{2}S_1 \quad + \quad \frac{1}{2}S_2 \\ x_2 \quad = \quad 2 \quad + \quad \frac{1}{2}S_1 \quad - \quad \frac{1}{2}S_2 \\ S_3 \quad = \quad 2 \quad + \quad \frac{3}{2}S_1 \quad - \quad \frac{3}{2}S_2 \\ S_1, S_2, S_3, x_1, x_2 \quad \geq \quad 0 \\ \hline S_2 \quad = \quad \frac{4}{3} \quad + \quad S_1 \quad - \quad \frac{2}{3}S_3 \\ \hline S_2 \quad = \quad \frac{4}{3} \quad + \quad S_1 \quad - \quad \frac{2}{3}S_3 \\ \hline S_1, S_2, S_3, x_1, x_2 \quad \geq \quad 0 \\ \hline \end{array} \right)$$

$$x_1 = \frac{5}{3}, \ x_2 = \frac{8}{3} + t, \ S_1 = t, \ S_2 = \frac{4}{3} + t, \ S_3 = 0$$

$$z = \frac{33}{5} + t$$

Where $\ t$ can be any positive number, unbound

4. Consider the following dictionaries in a cycling example.

The initial dictionary:

After the first iteration:

Dictionary D:

...

...

After the fifth iteration:

Dictionary D*:

Let B be the index set of the basic variables in D, and let B* be the index set of the basic variables in D*. At the first iteration, let x_2 be the entering variable, while x_6 is obviously the leaving variable. Therefore, $x_2 = t$,

 $x_3 = x_4 = x_5 = 0$, $x_1 = 11t$, $x_6 = -4t$, $x_7 = 1 - 11t$, z = 53t. Verify that at Dictionary D^* , z = 53t if x_j in D^* are substituted by x_j in D and $c_j^* = 0$ $\forall j \in B^*$.

Let

$$x_2 = t$$
, $x_3 = x_4 = x_5 = 0$, $x_1 = 11t$, $x_6 = -4t$ and $x_7 = 1 - 11t$

Then

$$z = 24 \times (-4t) + 22 \times (11t) - 93 \times (t) - 21 \times (0) = 53t$$

5. For a maximization problem, let D denote a dictionary in which x_t leaves the basis and x_s will enter the basis. D:

$$x_i = b_i - \sum_{j \notin \mathcal{B}} a_{ij} x_j \qquad i \in \mathcal{B}$$

$$z = v + \sum_{j \notin \mathbf{B}} c_j x_j$$

where B is the index set of basis variables in D, and $s \notin B$ and $t \in B$.

Now let D^* be a dictionary in which x_t enters the basis.

 D^* :

$$x_i = b_i^* - \sum_{j \notin \mathcal{B}^*} a_{ij}^* x_j \qquad i \in \mathcal{B}^*$$

$$z = v^* + \sum_{j \notin \mathbf{B}^*} c_j^* x_j$$

Show $c_t^* a_{ts} > 0$. (Hint: you need to show $c_t^* > 0$ and $a_{ts} > 0$)

Since C_s is always positive, we can choose S to enter the basis

$$x_t = -a_{ts}x_s + \cdots$$

Therefore $a_{ts} > 0$ and we can obtain the minimal ratio $x_s = \frac{x_t}{a_{ts}}$ if choosing x_t to leave.

Then $c_t^* = C_s \times \frac{x_t}{a_{ts}} a_{ts} > 0$ since $C_s > 0$ and $a_{ts} > 0$.

Finally $c_t^* a_{ts} > 0$ is proven.