# Signals

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#### Introduction

- · Signals are Software Interrupts to Handle Asynchronous Events
- signal(7)
- Signal Concepts
  - numerous conditions
    - the terminal-generated signals
    - hardware exceptions
    - software conditions
    - the kill(2) function
    - the kill(1) command
  - handle signal
    - ignore the signal
      - but we cannot ignore SIGKILL and SIGSTOP
    - catch the signal, the register customized handler
      - but we cannot register SIGKILL and SIGSTOP
    - let the default action apply
      - every signal has a default action
- List of Linux Signals (e.g. Ubuntu 18.04 x86 64)
  - o POSIX.1-1990 Standard

Signal	Value	Action	Comment
SIGHUP	1	Term	Hangup detected on controlling terminal or death of controlling process
SIGINT	2	Term	Interrupt from keyboard
SIGQUIT	3	Core	Quit from keyboard
SIGILL	4	Core	Illegal Instruction
SIGABRT	6	Core	Abort signal from abort (3)

Signal	Value	Action	Comment	
SIGFPE	8	Core	Floating-point exception	
SIGKILL	9	Term	Kill signal	
SIGSEGV	11	Core	Invalid memory reference	
SIGPIPE	13	Term	Broken pipe: write to pipe with no readers; see pipe(7)	
SIGALRM	14	Term	Timer signal from alarm(2)	
SIGTERM	15	Term	Termination signal	
SIGUSR1	30,10,16	Term	User-defined signal 1	
SIGUSR2	31,12,17	Term	User-defined signal 2	
SIGCHLD	20,17,18	Ign	Child stopped or terminated	
SIGCONT	19,18,25	Cont	Continue if stopped	
SIGSTOP	17,19,23	Stop	Stop process	
SIGTSTP	18,20,24	Stop	Stop typed at terminal	
SIGTTIN	21,21,26	Stop	Terminal input for background process	
SIGTTOU	22,22,27	Stop	Terminal output for background process	
SUSv2 and I	SUSv2 and POSIX.1-2001			
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Signal	Value	Action	Comment
SIGBUS	10,7,10	Core	Bus error (bad memory access)
SIGPOLL		Term	Pollable event (Sys V). Synonym for SIGIO
SIGPROF	27,27,29	Term	Profiling timer expired
SIGSYS	12,31,12	Core	Bad system call (SVr4); see also seccomp(2)
SIGTRAP	5	Core	Trace/breakpoint trap
SIGURG	16,23,21	Ign	Urgent condition on socket (4.2BSD)
SIGVTALRM	26,26,28	Term	Virtual alarm clock (4.2BSD)
SIGXCPU	24,24,30	Core	CPU time limit exceeded (4.2BSD); see setrlimit(2)
SIGXFSZ	25,25,31	Core	File size limit exceeded (4.2BSD); see setrlimit(2)

### • Linux 2.4

Signal	Value	Action	Comment
SIGIOT	6	Core	IOT trap. A synonym for SIGABRT
SIGEMT	7,-,7	Term	Emulator trap
SIGSTKFLT	-,16,-	Term	Stack fault on coprocessor (unused)

Signal	Value	Action	Comment
SIGIO	23,29,22	Term	I/O now possible (4.2BSD)
SIGCLD	-,-,18	lgn	A synonym for SIGCHLD
SIGPWR	29,30,19	Term	Power failure (System V)
SIGINFO	29,-,-		A synonym for SIGPWR
SIGLOST	-,-,-	Term	File lock lost (unused)
SIGWINCH	28,28,20	lgn	Window resize signal (4.3BSD, Sun)
SIGUNUSED	-,31,-	Core	Synonymous with SIGSYS

# The Signal Function

- signal(2) void (\*signal(int signum, void (\*handler)(int)))(int)
  - typedef void (\*sighandler\_t)(int);
  - sighandler\_t signal(int signum, sighandler\_t handler);
  - return: previous disposition of signal, SIGERR on error
  - signum is the name of the signal
  - handler is the function to be called when the signal occurs
    - can also be SIG\_IGN or SIG\_DFL
    - SIG\_IGN ignore signum
    - SIG\_DFL use default action in above table
  - implementations differs among system
    - better to use sigaction function
- Example

```
void sig_usr(int signo) {
 if (signo == SIGUSR1) {
   printf("received SIGUSR1\n");
  } else if (signo == SIGUSR2) {
    printf("received SIGUSR2\n");
  } else {
    printf("receive signal %d\n", signo);
  }
}
int main() {
  if (signal(SIGUSR1, sig_usr) == SIG_ERR) {
    printf("can't catch SIGUSR1");
  }
  if (signal(SIGUSR2, sig_usr) == SIG_ERR) {
    printf("can't catch SIGUSR2");
  }
  while (1) {
    pause();
```

```
}
}
```

```
$ ./a.out &
[1] 15512
$ kill -USR1 15512
received SIGUSR1
$ kill -USR2 15512
received SIGUSR2
$ kill 15512 # default send SIGTERM
[1]+ Terminated ./a.out
$ kill -l # check available signals
1) SIGHUP 2) SIGINT ...
```

- Signal Setup
  - signal is either default or ignore when program is executed
  - fork(2): child inherits parent's signal dispositions
  - exec(3): child changes the signal dispositions to their default action
  - the shell sets SIGINT and SIGQUIT in the background process to be ignored

# Interrupted System Calls

- · A System Call may be Interrupted
  - if process catch a signal while in slow system call
  - then the system call return an error and errno was set to EINTR
    - EINTR: 4, interrupted system call
- · Handle Interrupted System Calls

```
again:
  if ((n = read(fd, buf, BUFFSIZE)) < 0) {
    if (errno == EINTR)
      goto again;
}</pre>
```

- Prevent Interrupted System Calls
  - 4.2BSD introduced automatic restarting of certain interrupted system calls
  - o ioctl, read, readv, write, writev, wait, and waitpid
  - the feature can be disabled if you don't like it

#### Reentrant Functions

· When a Signal is Being Caught

- instruction is interrupted
- · execute signal handler instructions
- if the handler returns, then the process continues the interrupted instructions
- If Handler Uses the Same System Call as the Interrupted Instructions
  - if works well, they are reentrant functions
  - some functions can not work with reentrant, because
    - they use static data structures
    - they are part of the standard I/O library (buffering)
    - they call malloc or free
  - reentrant functions example

```
void handler(int signo) {
  tmpnam(NULL);
}
int main() {
  char *s = tmpnam(NULL);
  signal(SIGALRM, handler);
  alarm(3);
  for (int i = 0; i < 6; ++i) {
    printf("tmpnam = %s\n", s);
    sleep(1);
  }
}</pre>
```

- SIGCLD and SIGCHLD
  - the signal disposition is SIG\_DFL: default ignore
  - zombie avoidance
    - explicitly set disposition to SIG\_IGN
    - no zombie will be created if children of the calling process terminate

# Reliable Signal Model

- · Delivery of a Signal
  - delivered: a process received the signal and action is taken
  - pending: a signal is generated, but not delivered
- Blocking the Delivery of a Signal
  - if signal is blocked and handler is SIG\_DFL or a handler
  - · then the signal remains pending until
    - unblock the signal or
    - change handler to SIG\_IGN
- · Block Signal is Generated More Than Once
  - POSIX.1 allows deliver signal once or more
  - most UNIX systems do not queue signals

- Deliver Signals More Than Once
  - POSIX.1 does not specify the order
  - it suggests that signals related to the current state be delivered before other signals
- Signal Mask
  - defines the blocked signals in a process
  - · each signal has a corresponding bit

### **Functions for Handling Signals**

```
kill(2):int kill(pid_t pid, int sig);
raise(3):int raise(int sig);
      return: 0 OK, -1 error

    send signal to

            pid > 0: pid process
            pid == 0: every process in the process group as same as the caller
            pid == -1: every process the caller has permission to send, except process 1
            pid < -1: every process in the process group -pid</li>
     raise(sig) equals kill(getpid(), sig)

    NULL signal

            may be used to check the existence of a process
            no such process: kill returns -1 and errno is set to ESRCH
                  ESRCH: 3, no such process

    alarm(2): unsigned int alarm(unsigned int seconds);

    return: 0 or number of seconds remaining until previously scheduled alarm

    SIGALRM is generated when the timer expires, default terminate the process

    only one alarm clocks per process

pause(2):int pause(void);
      return: -1 with errno set to EINTR

    suspend a process until it received a signal

abort(3): void abort(void);

    send SIGABRT to the caller

      • the signal can be caught by a handler, but
            the handler will not return if it calls exit(3), _exit(2) or _Exit(2)
            • if the handler returns, then abort terminates the process
            • if the handler calls siglongjmp, the program may continue to execute
```

operation

```
• sigemptyset(3):int sigemptyset(sigset_t *set);
     initializes set to empty
```

```
sigfillset(3):int sigfillset(sigset_t *set);
```

```
initializes set to full
           sigaddset(3):int sigaddset(sigset_t *set, int signum);
                add signum to set
           sigdelset(3):int sigdelset(sigset_t *set, int signum);
                delete signum from set
           return: 0 OK, -1 error

    membership test

           sigismember(3):int sigismember(const sigset_t *set, int signum);
           return: 1 if true. 0 if false. -1 error

    sigprocmask(2):int sigprocmask(int how, const sigset_t *set, sigset_t

  *oldset);
     return: 0 OK, -1 error

    block or unblock signals

     how
           SIG_BLOCK: $newblkset = curblkset \cup set$
           SIG_UNBLOCK: $newblkset = curblkset - set$
           SIG_SETMASK: $newblkset = set$

    if oldset is not NULL, then previous value of the signal mask is stored oldset

    if set is NULL, then the mask remains unchanged

    sigpending(2):int sigpending(sigset_t *set);
```

Example

return: 0 OK, -1 error

get currently pending signals

```
int main() {
 sigset_t newmask, oldmask, pendmask;
 /* sig_quit print caught SIGQUIT and then reset SIGQUIT handler */
 if (signal(SIGQUIT, sig_quit) == SIG_ERR) {
   printf("err signal register\n");
 }
 sigemptyset(&newmask);
 sigaddset(&newmask, SIGQUIT);
 if (sigprocmask(SIG_BLOCK, &newmask, &oldmask) < 0) {
   printf("err sigprocmask\n");
 }
 printf("SIGQUIT blocked\n");
 sleep(5);
 if (sigpending(&pendmask) < 0) {
  printf("err sigpending\n");
 }
 if (sigismember(&pendmask, SIGQUIT)) {
   printf("SIGQUIT in pending\n");
 }
 if (sigprocmask(SIG_SETMASK, &oldmask, NULL) < 0) {
   printf("err sigprocmask\n");
```

```
printf("SIGQUIT unblocked\n");
sleep(5);
}
```

```
$ ./a.out
^\^\^\
SIGQUIT in pending
caught SIGQUIT
SIGQUIT unblocked
```

- sigaction(2):int sigaction(int signum, const struct sigaction \*act, struct sigaction \*oldact);
  - return: 0 OK, -1 error
  - examine or modify (or both) the signal action
  - sigaction

```
struct sigaction {
  void (*sa_handler)(int); // default handler
  void (*sa_sigaction)(int, siginfo_t *, void *); // alternative
  handler when SA_SIGINFO is enabled
  sigset_t sa_mask; // blocked signals when executing the
  handler
  int sa_flags; // various options
  void (*sa_restorer)(void);
};
```

• sa\_flags

sa_flags	Description
SA_INTERRUPT	obsoleted after Linux 2.6.19
SA_NOCLDSTOP	when signum is SIGCHLD, do not receive notification when child processes stop
SA_NOCLDWAIT	when signum is SIGCHLD, do not transform children into zombies
SA_RESETHAND	when handler is established, restore the action to the default upon entry to the handler
SA_RESTART	when handler is established, restart some system calls instead of generating EINTR
SA_NODEFER	when handler is established, do not prevent the signal from being received from within its own handler

sa_flags	Description
SA_SIGINFO	when handler is established, use sa_sigaction handler
SA_ONSTACK	when handler is established, call the handler on an alternate signal stack provided by sigaltstack(2)
SA_RESTORER	used in C libraries, sa_restorer field contains the address of a "signal trampoline", see sigreturn(2)

- sigsuspend(2):int sigsuspend(const sigset\_t \*mask);
  - return: -1 with errno set to EINTR
  - replace signal mask and then suspends the process until
    - a signal whose action is to invoke a handler or
    - a signal to terminate a process, sigsuspend(2) does not return in this condition
  - it is not possible to block SIGKILL or SIGSTOP
  - catching signal between sigprocmask(2) and pause(2)
    - after handler is done, the program is suspended
    - use sigsuspend(2) to prevent this condition
- Example: Wait for a Global Variable to be Set

```
volatile sig_atomic_t quitflag;
void sig_int(int signo) {
  if (signo == SIGINT)
   printf("\ninterrupt\n");
  else if (signo == SIGQUIT)
    quitflag = 1;
}
int main(void) {
  sigset_t newmask, oldmask, zeromask;
  if (signal(SIGINT, sig_int) == SIG_ERR) printf("signal(SIGINT)
error");
  if (signal(SIGQUIT, sig_int) == SIG_ERR) printf("signal(SIGQUIT)
error");
  sigemptyset(&zeromask);
  sigemptyset(&newmask);
  sigaddset(&newmask, SIGQUIT);
  if (sigprocmask(SIG_BLOCK, &newmask, &oldmask) < 0)</pre>
printf("SIG_BLOCK error");
  while (quitflag == 0) sigsuspend(&zeromask);
  quitflag = 0;
  if (sigprocmask(SIG_SETMASK, &oldmask, NULL) < 0)</pre>
printf("SIG_SETMASK error");
}
```

Example: Process Synchronization

```
// library
volatile sig_atomic_t sigflag;
sigset_t newmask, oldmask, zeromask;
void sig_usr(int signo) { sigflag = 1; }
void TELL_WAIT(void) {
  if (signal(SIGUSR1, sig_usr) == SIG_ERR) printf("signal(SIGUSR1)
error");
 if (signal(SIGUSR2, sig_usr) == SIG_ERR) printf("signal(SIGUSR2)
error");
 sigemptyset(&zeromask);
 sigemptyset(&newmask);
 sigaddset(&newmask, SIGUSR1);
 sigaddset(&newmask, SIGUSR2);
 if (sigprocmask(SIG_BLOCK, &newmask, &oldmask) < 0)</pre>
printf("SIG_BLOCK error");
}
void TELL_PARENT(pid_t pid) { kill(pid, SIGUSR2); }
void TELL_CHILD(pid_t pid) { kill(pid, SIGUSR1); }
void WAIT_PARENT(void) {
 while (sigflag == 0) sigsuspend(&zeromask);
  sigflag = 0;
  if (sigprocmask(SIG_SETMASK, &oldmask, NULL) < 0)</pre>
printf("SIG_SETMASK error");
}
void WAIT_CHILD(void) {
 while (sigflag == 0) sigsuspend(&zeromask);
  sigflag = 0;
  if (sigprocmask(SIG_SETMASK, &oldmask, NULL) < 0)</pre>
printf("SIG_SETMASK error");
}
// application
int main(void) {
 pid_t pid;
 TELL_WAIT();
  if ((pid = fork()) < 0) {
  printf("fork error");
  } else if (pid == 0) {
    WAIT_PARENT(); /* parent goes first */
    charatatime("output from child\n");
  } else {
    charatatime("output from parent\n");
    TELL_CHILD(pid);
  }
}
```

### setjmp(3) and longjmp(3)

- setjmp(3):int setjmp(jmp\_buf env);
- longjmp(3):void longjmp(jmp\_buf env, int val);
- setjmp(3) and longjmp(3) can Jump Accross Functions
  - handler perform longjmp(3) to main loop
  - but mask of a signal is set automatically when that signal is caught
  - longjmp(3) exits the handler with the mask, and it will cause problems
- Solutions
  - setjmp(3) should save signal mask
  - the mask is restored when longjmp(3) is called
- POSIX.1 does not define how do setjmp(3) and longjmp(3) handle masks
  - FreeBSD and MAC OS X save and restore the mask automatically
  - however, Linux must work with sigsetjmp(3) and siglongjmp(3)
- sigsetjmp(3):int sigsetjmp(sigjmp\_buf env, int savesigs);
  - return: 0 if called directly, nonzero if returning from a call to siglongjmp(3)
  - if savesigs is nonzero
    - usually set savesig to nonzero
    - the process's current signal mask is saved in env
    - the mask will be restored if a siglong jmp (3) is later performed with this env
- siglongjmp(3): void siglongjmp(sigjmp\_buf env, int val);
- Example

```
sigjmp_buf jmpbuf;
volatile sig_atomic_t canjump;

void sig_usr1(int signo) {
   time_t starttime;
   if (canjump == 0) return;
   pr_mask("starting sig_usr1: ");
   alarm(3);
   starttime = time(NULL);
   for (;;)
    if (time(NULL) > starttime + 5) break;
   pr_mask("finishing sig_usr1: ");
   canjump = 0;
   siglongjmp(jmpbuf, 1);
}

void sig_alrm(int signo) { pr_mask("in sig_alrm: "); }
```

```
int main(void) {
   if (signal(SIGUSR1, sig_usr1) == SIG_ERR) printf("signal(SIGUSR1)
error");
   if (signal(SIGALRM, sig_alrm) == SIG_ERR) printf("signal(SIGALRM)
error");
   pr_mask("starting main: ");
   if (sigsetjmp(jmpbuf, 1)) {
      pr_mask("ending main: ");
      exit(0);
   }
   canjump = 1;
   for (;;) pause();
}
```

```
$ ./a.out &
[1] 8874
starting main:
$ kill -SIGUSR1 8874
starting sig_usr1: SIGUSR1
in sig_alrm: SIGUSR1 SIGALRM
finishing sig_usr1: SIGUSR1
ending main: # SIGUSR1 if savesigs is 0 (i.e. mask in the sig_usr1
function)
```

### The System Function

- In Ch06
  - · a simple system implementation is introduced
  - but it does not handle signals
- POSIX.1 requires the system function
  - ignore SIGINT/SIGQUIT
    - these signals should be sent only to child process
  - block SIGCHLD
    - parent will not confuse between the termination of child and other child
- System Implementation
  - preserve original signal action for SIGINT and SIGQUIT
  - block SIGCHLD
  - child
    - restore handler
    - unblock SIGCHLD
    - execute command
  - parent
    - wait for child
  - · restore handler and unblock SIGCHLD

```
int system(const char *cmdstring) {
 pid_t pid;
 int status;
 struct sigaction ignore, saveintr, savequit;
 sigset_t chldmask, savemask;
 if (cmdstring == NULL) return 1;
 ignore.sa_handler = SIG_IGN;
 sigemptyset(&ignore.sa_mask);
 ignore.sa_flags = 0;
 if (sigaction(SIGINT, &ignore, &saveintr) < 0) return -1;
 if (sigaction(SIGQUIT, &ignore, &savequit) < 0) return -1;
 sigemptyset(&chldmask);
 sigaddset(&chldmask, SIGCHLD);
 if (sigprocmask(SIG_BLOCK, &chldmask, &savemask) < 0) return -1;
 if ((pid = fork()) < 0) {
   status = -1;
 } else if (pid == 0) {
    sigaction(SIGINT, &saveintr, NULL);
    sigaction(SIGQUIT, &savequit, NULL);
    sigprocmask(SIG_SETMASK, &savemask, NULL);
   execl("/bin/sh", "sh", "-c", cmdstring, (char *)0);
   _exit(127);
 } else {
   while (waitpid(pid, &status, 0) < 0) {
     if (errno != EINTR) {
        status = -1;
        break;
      }
 }
 if (sigaction(SIGINT, &saveintr, NULL) < 0) return -1;
 if (sigaction(SIGQUIT, &savequit, NULL) < 0) return -1;
 if (sigprocmask(SIG_SETMASK, &savemask, NULL) < 0) return -1;
 return status;
}
```

#### · Job Control Signals

- except SIGCHLD, most applications don't handle these signals
- vi editor will save/restore terminal state when the process is stopped/continued

#### Signal Description

	<u> </u>
SIGCHLD	Child stopped or terminated
SIGCONT	Continue if stopped
SIGSTOP	Stop signal
SIGTSTP	Stop typed at terminal
SIGTTIN	Terminal input for background process

Signal	Description
SIGTTOU	Terminal output for background process