

- The generalized voltage  $v_k$ 
  - $\Pr[\text{A random walk starting at } k \text{ visits } j \text{ before } i]$
- Calculating  $\mathbf{v} \triangleq (v_1 \ v_2 \ \dots \ v_n)^\top$ 
  - $v_i = 0$  and  $v_j = 1$
  - $\forall k \neq i, j, v_k = \sum_l p_{kl} \cdot v_l$
  - Split  $\mathbf{P} = \begin{pmatrix} \hat{\mathbf{P}} & \mathbf{c}_i & \mathbf{c}_j \\ \mathbf{r}_i^\top & 0 & p(i, j) \\ \mathbf{r}_j^\top & p(j, i) & 0 \end{pmatrix}$  and  $\mathbf{v} = (\hat{\mathbf{v}} \ 0 \ 1)^\top$
  - Then  $\hat{\mathbf{v}} = \hat{\mathbf{P}}\hat{\mathbf{v}} + \mathbf{r}_j \Rightarrow \hat{\mathbf{v}} = (\mathbf{I} - \hat{\mathbf{P}})^{-1}\mathbf{r}_j$
- $\text{ep}(i \rightarrow j) = \sum_k p_{ik} \cdot v_k = \mathbf{r}_i^\top (\mathbf{I} - \hat{\mathbf{P}})^{-1} \mathbf{c}_j + p(i, j)$
- Computing all  $\text{ep}(i \rightarrow j)$  requires  $\Theta(n^2)$  matrix inversions

**Theorem.** Let  $\mathbf{Q} = [q(i, j)]_{n \times n} \triangleq (\mathbf{I} - c\mathbf{P})^{-1}$ .  $\forall i \neq j$ , there is

$$\text{ep}(i \rightarrow j) = \frac{q(i, j)}{q(i, i)q(j, j) - q(i, j)q(j, i)}.$$

**Lemma.** The expected time  $r_i$  for a random walk starting at node  $i$  to return to  $i$  is the reciprocal of the stationary probability of  $i$ . That is

$$r_i = \frac{1}{\pi_i}.$$

- Intuitively
  - A long walk always ends up in stationary distribution  $\pi$
  - Suppose the walk length is  $T$ , then the expected number of times it visits  $i$  is  $\pi_i \cdot T$
  - The average length between two visits is  $\frac{T}{\pi_i \cdot T} = \frac{1}{\pi_i}$
  - A rigorous proof requires the Strong Law of Large Numbers

**Theorem.** The probability that a random walk starting at node  $i$  visits  $j$  before returning to  $i$ , which equals  $\text{ep}(i \rightarrow j)$ , satisfies

$$\text{ep}(i \rightarrow j)c(i, j) = \frac{1}{\pi_i},$$

where  $c(i, j)$  is the commute time between  $i$  and  $j$ .

- Proof
  - Consider a random walk  $w$  starting at  $i$ , and random variables

- \*  $X$  = the first time  $w$  returns to  $i$
- \*  $Y$  = the first time  $w$  returns to  $i$  after visiting  $j$
- By definition  $E(X) = \frac{1}{\pi_i}$  and  $E(Y) = c(i, j)$
- Clearly  $X \leq Y$ , and  $\Pr[X = Y] = p \triangleq \text{ep}(i \rightarrow j)$ 
  - \*  $E(Y - X) = p \cdot 0 + (1 - p) \cdot E(Y) = (1 - p)c(i, j)$
- Also  $E(Y - X) = E(Y) - E(X) = c(i, j) - \frac{1}{\pi_i}$
- $\text{ep}(i \rightarrow j) + \text{ep}(j \rightarrow i) = \frac{1}{c(i, j)} \left( \frac{1}{\pi_i} + \frac{1}{\pi_j} \right)$
- Recall that  $h(i, j)$  is small whenever  $\pi_j$  is large, which is bad for personalization
  - Sarkar et al. (2008) alleviate this by restricting the length of random walk
  - Tong et al. (2007) alleviate this by reducing the dependence on stationary distribution