

# Basic Algebra II: Chapter 6

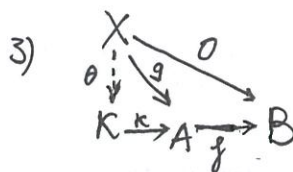
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Motivation: Extend the theory of modules to a bigger category.

- Initial object  $I$ :  $\forall X \in \text{obl}$ ,  $\exists!$  morphism  $I \rightarrow X$ .
- Terminal object  $T$ :  $\forall X \in \text{obl}$ ,  $\exists!$  morphism  $X \rightarrow T$ .
- Zero object  $0$ :  $0$  is an initial object and terminal object.
- If  $\mathcal{C}$  is a category with zero object, given  $f: A \rightarrow B$

$k: K \rightarrow A$  is a kernel of  $f$  if

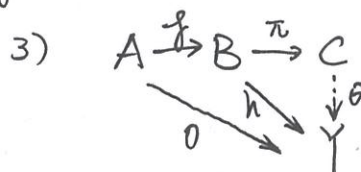
- 1)  $k$  is monic 2)  $f \circ k = 0$



compatible with the usual case!

$\pi: B \rightarrow C$  is a cokernel of  $f$  if

- 1)  $\pi$  is epic 2)  $\pi \circ f = 0$



Def. A category  $\mathcal{C}$  is additive if

- $\mathcal{C}$  has a zero object.
- $\text{Hom}(A, B)$  is an abelian gp  $\forall A, B \in \text{obl}(\mathcal{C})$
- Distribution Laws:  $X \xrightarrow{a} A \xrightarrow[f]{g} B \xrightarrow{b} Y$ , then  
 $(f+g) \circ a = fa + ga$  &  $b \circ (f+g) = bf + bg$
- $\mathcal{C}$  has finite product.

Def. If  $\mathcal{C}$  and  $\mathcal{D}$  are additive categories, a functor  $T: \mathcal{C} \rightarrow \mathcal{D}$  is additive if  $\forall f, g \in \text{Hom}_{\mathcal{C}}(A, B)$ ,  $T(f+g) = Tf + Tg$   
 thus is,  $T: f \mapsto Tf$  is a group hom.  $\text{Hom}_{\mathcal{C}}(A, B) \rightarrow \text{Hom}_{\mathcal{D}}(TA, TB)$

Def. A category  $\mathcal{C}$  is an abelian category if it is an additive cat and

1) every mor has a kernel and cokernel

2) every monic is a kernel and every epic is a cokernel.

$$A \xrightarrow{f} B \xrightarrow{g} C \Rightarrow f=g$$

$$C \xrightarrow{f} A \xrightarrow{g} B \Rightarrow f=g$$

Ex. 1)  ${}_R\text{Mod}$  and  $\text{Mod}_R$  are abelian cats.

2) cat of torsion-free ab gps. is an additive cat.

but not an abelian cat. Consider  $\bigvee \mathbb{Z} \rightarrow 2\mathbb{Z}$ ,  
inclusion

$\Pi_2$  is its cokernel which is not torsion-free.

Rmk. Any Abelian cat has finite  $\bigvee$  limit and  $\bigvee$  direct limit.  
inverse finite

pf. kernel exists  $\Rightarrow$  equalizer exists  $\Rightarrow$  inverse limit exists.  
product exists

Def. A complex in an ab cat  $\mathcal{C}$  is a seq of objs and mors

$$C = (C_\bullet, d_\bullet) = \cdots \rightarrow C_{n+1} \xrightarrow{d_{n+1}} C_n \xrightarrow{d_n} C_{n-1} \rightarrow \cdots$$

$$\text{st. } d_n \circ d_{n+1} = 0, \forall n \in \mathbb{Z}.$$

A chain map  $f = f_\bullet: (C_\bullet, d_\bullet) \rightarrow (C'_\bullet, d'_\bullet)$  is a seq of mors  $f_n: C_n \rightarrow C'_n$

$$\text{st. } d'_{n+1} \circ f_{n+1} = f_n \circ d_{n+1} \quad \forall n \in \mathbb{Z}.$$

Rmk. The cat of all complexes in  $\mathcal{C}$  is denoted by  $\text{Comp}(\mathcal{C})$ , which is an ab cat.

Consider  $\text{Comp}({}_R\text{Mod}) = {}_R\text{Comp} \ni (C_\bullet, d_\bullet)$

$$\bullet \quad Z_i(C) = \ker d_i, \quad B_i(C) = \text{Im } d_{i+1}, \quad H_i(C) = Z_i(C) / B_i(C)$$

$i$ -cycles

$i$ -boundaries

$i$ -homology

- $H_n: {}_R\text{Comp} \rightarrow {}_R\text{Mod}$  is an additive functor.

$$H_n: C \mapsto H_n(C) = Z_n / B_n.$$

$$f \mapsto H_n(f) = f_* \quad \text{"induced map"}$$

Rmt.  $C$  is exact  $\Leftrightarrow H_n(C) = 0, \forall n \in \mathbb{Z}$ .

- We can rewrite a complex

$$\cdots \rightarrow C_{i+1} \xrightarrow{d_{i+1}} C_i \xrightarrow{d_i} C_{i-1} \rightarrow \cdots \quad \text{"decreasing"}$$

$$\text{to } \cdots \rightarrow C^{(-i-1)} \xrightarrow{d^{(-i-1)}} C^i \xrightarrow{d^i} C^{(-i-1)} \rightarrow \cdots \quad \text{cochain complex}$$

$$\text{this is } \cdots \rightarrow C^{i-1} \xrightarrow{d^{i-1}} C^i \xrightarrow{d^i} C^{i+1} \rightarrow \cdots \quad \text{"increasing"}$$

- A cochain complex  $C$ .

$$Z^i(C) = \ker d^i, \quad B^i(C) = \text{Im } d^{i-1}, \quad H^i(C) = Z^i(C) / B^i(C)$$

*i-cocycles*                      *i-coboundaries*                      *i-cohomology*

Def.  $0 \rightarrow C' \xrightarrow{\alpha} C \xrightarrow{\beta} C'' \rightarrow 0$  is SES of complexes  
 $\Leftrightarrow 0 \rightarrow C'_n \xrightarrow{\alpha_n} C_n \xrightarrow{\beta_n} C''_n \rightarrow 0$  is SES of  $\mathcal{C}$ ,  $\forall n \in \mathbb{Z}$

Thm.  $0 \rightarrow C' \xrightarrow{\alpha} C \xrightarrow{\beta} C'' \rightarrow 0$  SES in  ${}^R\text{Comp}$ . then there exists

a morphism  $\partial_n: H_n(C'') \rightarrow H_n(C): \text{cls}(z'') \mapsto \text{cls}(\alpha_n^{-1} \alpha_n^{-1} \beta_n^{-1}(z''))$  "Connecting hom"  
 such that  $\cdots \rightarrow H_n(C') \xrightarrow{\alpha_n*} H_n(C) \xrightarrow{\beta_n*} H_n(C'') \xrightarrow{\partial_n} H_{n-1}(C') \xrightarrow{\alpha_{n-1}*} H_{n-1}(C) \rightarrow \cdots$   
 exact. "Long Exact Seq"

Cor. (Snake Lemma) A commutative diagram with exact rows:

$$\begin{array}{ccccccc} & & \Sigma^f & \Sigma^g & \Sigma^h & & \\ & & \downarrow & \downarrow & \downarrow & & \\ 0 & \rightarrow & A' & \rightarrow & A & \rightarrow & A'' \rightarrow 0 \\ & & \downarrow f & \downarrow g & \downarrow h & & \\ 0 & \rightarrow & B' & \rightarrow & B & \rightarrow & B'' \rightarrow 0 \end{array}$$

$$\text{implies } 0 \rightarrow \ker f \rightarrow \ker g \rightarrow \ker h \rightarrow \text{coker } f \rightarrow \text{coker } g \rightarrow \text{coker } h \rightarrow 0$$

$$0 \rightarrow H_1(\Sigma^f) \rightarrow H_1(\Sigma^g) \rightarrow H_1(\Sigma^h) \rightarrow H_0(\Sigma^f) \rightarrow H_0(\Sigma^g) \rightarrow H_0(\Sigma^h) \rightarrow 0$$



Thm.  $0 \rightarrow C' \xrightarrow{i} C \xrightarrow{p} C'' \rightarrow 0$  commutative and exact in  $\text{Comp}$

$$\begin{array}{ccccccc} & & i & & p & & \\ & & \downarrow & & \downarrow & & \\ 0 & \rightarrow & C' & \xrightarrow{j} & C & \xrightarrow{p} & C'' \rightarrow 0 \\ & & f & & g & & \\ & & \downarrow & & \downarrow & & \\ 0 & \rightarrow & D' & \xrightarrow{j} & D & \xrightarrow{p} & D'' \rightarrow 0 \\ & & & & i & & \end{array}$$

then

$$\begin{array}{ccccccc} \rightarrow H_n(C') & \xrightarrow{i_n^*} & H_n(C) & \xrightarrow{p_n^*} & H_n(C'') & \xrightarrow{\partial_n} & H_{n-1}(C) \rightarrow \dots \\ \downarrow j_n^* & & \downarrow j_n^* & & \downarrow h_n^* & & \downarrow j_{n-1}^* \\ \rightarrow H_n(D') & \xrightarrow{j_n^*} & H_n(D) & \xrightarrow{j_n^*} & H_n(D'') & \xrightarrow{\partial_n'} & H_{n-1}(D') \rightarrow \dots \end{array}$$

Commutative and exact.

Def.  $(C, \varepsilon)$  is a resolution of  $M$  if  $\varepsilon$  is an *augmentation*

$$\rightarrow C_n \xrightarrow{d_n} C_{n-1} \rightarrow \dots \rightarrow C_1 \xrightarrow{d_1} C_0 \xrightarrow{d_0} M \rightarrow 0$$

is exact.

- $H_0(C) = C_0/d_0C_1 = C_0/\ker d_0 = M$
- $C_0$  is projective  $\forall i \in \mathbb{N}$ , then  $(C, d)$  is a projective resolution.
- Every  $R$ -module  $M$  has a  $\text{free}^0$  resolution.

$M$  has a free resolution.

$$\begin{array}{ccccccc} & & & \nearrow K_2 & \searrow i_2 & & \\ \cdots & F_2 & \dashrightarrow & F_1 & \xrightarrow{i_0 \circ \varepsilon_1} & F_0 & \xrightarrow{\varepsilon_0} M \rightarrow 0 \\ & \nwarrow K_3 & & & \swarrow \varepsilon_1 & \nearrow i_1 & \\ & 0 & & & 0 & K_1 & \searrow i_1 \\ & & & & & 0 & \end{array}$$

Thm. (Comparison Theorem)  $f: A \rightarrow A'$  in an ab cat  $\mathcal{C}$ ,  $P$  projective resolution of  $A$   
 $P'$  resolution of  $A'$ , then  $\exists \tilde{f}$  its following diagram commutes

$$\begin{array}{ccccccc} \hookrightarrow P_1 & \rightarrow & P_0 & \rightarrow & A & \rightarrow & 0 \\ \downarrow f_1 & & \downarrow f_0 & & \downarrow f & & \\ \hookrightarrow P'_1 & \rightarrow & P'_0 & \rightarrow & A' & \rightarrow & 0 \end{array}$$

Pf. By def of projective.

Rmk.  $\check{f}$  is not unique, but they are homotopic.

Also proved by def of projectivity.

Def.  $(C, d) \xrightarrow[\beta]{\alpha} (C', d')$ , then  $\alpha$  is homotopic to  $\beta$  if there exists  $s = \{s_i: C_i \rightarrow C'_{i+1}\}$  s.t.  $\alpha_i - \beta_i = d'_{i+1} s_i + s_{i-1} d_i$ . ( $\alpha \simeq \beta$ )

$$\begin{array}{ccccccc} \rightarrow & C_{i+1} & \xrightarrow{d'_{i+1}} & C_i & \xrightarrow{d_i} & C_{i-1} & \rightarrow \\ & \downarrow s_{i+1} & & \downarrow s_i & & \downarrow s_{i-1} & \\ \rightarrow & C'_{i+1} & \rightarrow & C'_i & \rightarrow & C'_{i-1} & \rightarrow \end{array}$$

Rmk. There is also a "injective" version of "Comparison Theorem".

Thm. If  $\alpha \simeq \beta: (C, d) \rightarrow (C', d')$ , then  $\alpha_{n*} = \beta_{n*}: H_n(C) \rightarrow H_n(C')$ .  $\forall n \in \mathbb{Z}$ .  
Homotopic chain maps induce the same morphism in homology.

Pf.  $\forall z \in Z_n(C)$ ,  $(\alpha_n - \beta_n)z = (d'_{n+1} s_n + s_{n-1} d_n)z = d'_{n+1}(s_n z) \in B_{n+1}(C')$

$$\alpha_{n*}(cls(z)) = cls(\alpha_n z) = cls((\alpha_n - \beta_n + \beta_n)z) = cls(\beta_n z) = \beta_{n*}$$

## • Left Derived Functors

Consider an additive covariant functor  $F$  between two ab cats,  $\mathcal{C}$  and  $\mathcal{D}$ .

$\dots \rightarrow P_2 \rightarrow P_1 \rightarrow P_0 \rightarrow A \rightarrow 0$  Take  $f: A \rightarrow A'$  in  $\mathcal{C}$ .

$$\begin{array}{ccccccc} & & \downarrow \check{f} & & \downarrow f & & \\ \dots & \rightarrow & P'_2 & \rightarrow & P'_1 & \rightarrow & P'_0 \rightarrow A' \rightarrow 0 \end{array}$$

$\Downarrow F$  additive covariant

$$\dots \rightarrow FP_2 \rightarrow FP_1 \rightarrow FP_0 \rightarrow FA \rightarrow 0$$

$$\downarrow F\check{f} \qquad \downarrow Ff$$

$$\dots \rightarrow FP'_2 \rightarrow FP'_1 \rightarrow FP'_0 \rightarrow FA' \rightarrow 0$$

$\Downarrow F\check{f}: FP_A \rightarrow FP_{A'}$  chain map

$$(F\check{f})_{n*} = H_n(F\check{f}): H_n(FP_A) \rightarrow H_n(FP_{A'})$$

According to the process,  $H_n(F\check{f})$  seems to depend on the choice of  $\check{f}$  &  $P_A$ . But it is not true.

• Let  $\check{f}$  and  $\check{g}$  be the chain map over  $f$

$$\check{f} \simeq \check{g} \Rightarrow F\check{f} \simeq F\check{g} \Rightarrow H_n(F\check{f}) = H_n(F\check{g})$$

• Introduction to Homological Algebra, Rotman. proposition 6.1

So we def  $L_n F: \mathcal{C} \rightarrow \mathcal{D}$ .

$$L_n F(A) = H_n(FP_A), \quad L_n F(f) = H_n(F\check{f})$$

Thm.  $\mathcal{C}, \mathcal{D}$  ab cats with enough projectives.  $F: \mathcal{C} \rightarrow \mathcal{D}$  additive covariant functor.

$$0 \rightarrow A' \xrightarrow{i} A \xrightarrow{p} A'' \rightarrow 0 \text{ SES} \Rightarrow \cdots \rightarrow (L_n F)A' \xrightarrow{(L_n F)i} (L_n F)A \xrightarrow{(L_n F)p} (L_n F)A'' \xrightarrow{\partial_n} (L_{n-1} F)A' \xrightarrow{(L_{n-1} F)i} (L_{n-1} F)A \rightarrow \cdots \text{ LES}$$

pf.

$$\text{SES } 0 \rightarrow A' \rightarrow A \rightarrow A'' \rightarrow 0$$

Horshoe Lemma

$$\text{SES } 0 \rightarrow P_{A'} \rightarrow P_A \rightarrow P_{A''} \rightarrow 0$$

projective: split SES

$$\text{SES } 0 \rightarrow FP_{A'} \rightarrow FP_A \rightarrow FP_{A''} \rightarrow 0$$

additive functor: preserve split SES.

Long exact seq

$$\cdots \rightarrow (L_n F)A' \rightarrow (L_n F)A \rightarrow (L_n F)A'' \xrightarrow{\partial_n} (L_{n-1} F)A' \rightarrow (L_{n-1} F)A \rightarrow \cdots$$