- 教材讨论
 - -JH第4章第3节第1、2、3小节

问题1: greedy vs. local search

• greedy和local search有什么异同?

问题2: MIN-VCP

- 这个算法的基本思路是什么?
- 近似比是多少,为什么?
- 时间复杂度是多少?为什么?用什么数据结构来实现?

```
Algorithm 4.3.2.1. Input: A graph G = (V, E). Step 1: C := \emptyset {during the computation C \subseteq V, and at the end C should contain a vertex cover}; A := \emptyset \quad \{ \text{during the computation } A \subseteq E \text{ is a matching, and at the end } A \text{ is a maximal matching} \}; \\ E' := E \quad \{ \text{during the computation } E' \subseteq E, E' \text{ contains exactly the edges that are not covered by the actual } C, \text{ and at the end } E' = \emptyset \}. \\ \text{Step 2: } \text{ while } E' \neq \emptyset \\ \text{ do begin choose an arbitrary edge } \{u,v\} \text{ from } E'; \\ C := C \cup \{u,v\}; \\ A := A \cup \{\{u,v\}\}; \\ E' := E' - \{\text{all edges incident to } u \text{ or } v\} \} end \text{Output: } C.
```

问题2: MIN-VCP (续)

- 你能构造出一些近似比为1的例子吗?
- 你能构造出一些近似比为2的例子吗?

```
Input: A graph G = (V, E).
Algorithm 4.3.2.1.
   Step 1: C := \emptyset {during the computation C \subseteq V, and at the end C should
                        contain a vertex cover);
              A := \emptyset {during the computation A \subseteq E is a matching, and at the
                        end A is a maximal matching};
              E' := E {during the computation E' \subseteq E, E' contains exactly the
                        edges that are not covered by the actual C, and at the end
                        E' = \emptyset.
   Step 2: while E' \neq \emptyset
              do begin choose an arbitrary edge \{u, v\} from E';
                        C := C \cup \{u, v\};
                        A := A \cup \{\{u,v\}\};
                        E' := E' - \{ \text{all edges incident to } u \text{ or } v \}
              end
   Output: C.
```

问题2: MIN-VCP (续)

- 你能构造出一些近似比为1的例子吗?
- 你能构造出一些近似比为2的例子吗?



问题3: SCP

- 这个算法的基本思路是什么?
- 近似比是多少?
- 时间复杂度是多少?为什么?用什么数据结构来实现?

Algorithm 4.3.2.11.

```
Input: (X,\mathcal{F}), where X is a finite set, \mathcal{F}\subseteq \mathcal{P}ot(X) such that X=\bigcup_{Q\in\mathcal{F}}Q. Step 1: C:=\emptyset {during the computation C\subseteq\mathcal{F} and at the end C is a set cover of (X,\mathcal{F})}; U:=X \  \, \{\text{during the computation } U\subseteq X,\, U=X-\bigcup_{Q\in C}Q \text{ for the actual } C, \text{ and at the end } U=\emptyset\}. Step 2: \mathbf{while} \quad U\neq\emptyset \mathbf{do\ begin\ choose\ an\ } S\in\mathcal{F} \text{ such\ that\ } |S\cap U| \text{ is\ maximal;} U:=U-S; C:=C\cup\{S\} end \mathbf{Output:\ } C.
```

问题3: SCP (续)

- 你能构造出一些近似比为1的例子吗?
- 你能构造出一些近似比为Ω(Inn)的例子吗?

Algorithm 4.3.2.11.

```
\begin{array}{ll} \text{Input:} & (X,\mathcal{F}), \quad \text{where } X \text{ is a finite set, } \mathcal{F} \subseteq \mathcal{P}ot(X) \text{ such that } X = \bigcup_{Q \in \mathcal{F}} Q. \\ \text{Step 1:} & C := \emptyset \quad \{ \text{during the computation } C \subseteq \mathcal{F} \text{ and at the end } C \text{ is a set cover of } (X,\mathcal{F}) \}; \\ & U := X \quad \{ \text{during the computation } U \subseteq X, \ U = X - \bigcup_{Q \in C} Q \text{ for the actual } C, \text{ and at the end } U = \emptyset \}. \\ \text{Step 2:} & \text{while} \quad U \neq \emptyset \\ & \text{do begin choose an } S \in \mathcal{F} \text{ such that } |S \cap U| \text{ is maximal; } \\ & U := U - S; \\ & C := C \cup \{S\} \\ & \text{end} \\ \text{Output: } C. \end{array}
```

问题3: SCP (续)

- 你能构造出一些近似比为1的例子吗?
- 你能构造出一些近似比为Ω(Inn)的例子吗?

```
- U = { (x,y) | 1 \le x,y \le n }

- S_1 = \{ (x,y) | x \le n/2 \}

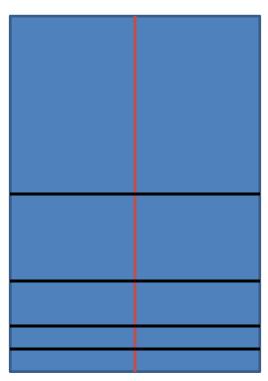
- S_2 = \{ (x,y) | x > n/2 \}

- T_1 = \{ (x,y) | n/2 < y \le n \}

- T_2 = \{ (x,y) | n/4 < y \le n/2 \}

- T_3 = \{ (x,y) | n/8 < y \le n/4 \}

- .....
```



问题3: SCP (续)

• 和MIN-VCP相比,SCP到底难在什么地方?

问题4: greedy

- 任选一个问题,给出一种greedy算法,并尽力给出近似比 (或者给出一些坏例子)
 - longest simple path
 - MAX-SAT
 - MAX-CL
 - MAX-CUT

问题5: MAX-CUT

- 这个算法的基本思路是什么?
- 近似比是多少,为什么?
- 时间复杂度是多少?为什么?用什么数据结构来实现?

Algorithm 4.3.3.1.

```
\begin{array}{ll} \text{Input:} & \mathsf{A} \ \mathsf{graph} \ G = (V, E). \\ \mathsf{Step 1:} & S = \emptyset \\ & \{\mathsf{the cut is considered to be } (S, V - S); \ \mathsf{in fact} \ S \ \mathsf{can be chosen} \\ & \mathsf{arbitrarily in this step}\}; \\ \mathsf{Step 2:} & \ \mathsf{while} & \mathsf{there exists such a vertex} \ v \in V \ \mathsf{that the movement of} \ v \\ & \mathsf{from one side of the cut} \ (S, V - S) \ \mathsf{to the other side of} \ (S, V - S) \ \mathsf{increases the cost of the cut.} \\ & \mathsf{do begin take a} \ u \in V \ \mathsf{whose movement from one side of} \ (S, V - S) \\ & \mathsf{to the other side of} \ (S, V - S) \ \mathsf{increases the cost of the cut,} \\ & \mathsf{and move this} \ u \ \mathsf{to the other side.} \\ & \ \mathsf{end} \\ \mathsf{Output:} \ (S, V - S). \\ \end{array}
```

问题5: MAX-CUT (续)

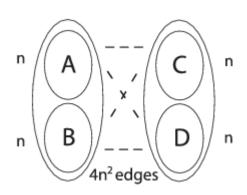
- 你能构造出一些近似比为1的例子吗?
- 你能构造出一些近似比为2的例子吗?

Algorithm 4.3.3.1.

```
Input: A graph G=(V,E). Step 1: S=\emptyset {the cut is considered to be (S,V-S); in fact S can be chosen arbitrarily in this step}; Step 2: while there exists such a vertex v\in V that the movement of v from one side of the cut (S,V-S) to the other side of (S,V-S) increases the cost of the cut. do begin take a u\in V whose movement from one side of (S,V-S) to the other side of (S,V-S) increases the cost of the cut, and move this u to the other side. end Output: (S,V-S).
```

问题5: MAX-CUT (续)

- 你能构造出一些近似比为1的例子吗?
- 你能构造出一些近似比为2的例子吗?



问题6: local search

- 任选一个问题,给出一种local search算法,并尽力给出近似比(或者给出一些坏例子)
 - longest simple path
 - MAX-SAT
 - MAX-CL
 - MIN-VCP
 - SCP