- 作业讲解
 - JH第3章练习3.3.2.7、3.3.2.8、3.3.2.9

JH第3章练习3.3.2.9

Exercise 3.3.2.9. Let $((X,\mathcal{F}),k)$, $\mathcal{F} \subseteq Pot(X)$, be an instance of the decision problem $Lang_{SC}$. Let, for every $x \in X$, $num_{\mathcal{F}}(x)$ be the number of sets in \mathcal{F} that contain x. Define

$$Pat((X, \mathcal{F}), k) = \max\{k, \max\{num_{\mathcal{F}}(x) \mid x \in X\}\}$$

that is a parameterization of $Lang_{SC}$. Find a Pat-parameterized polynomial-time algorithm for $Lang_{SC}$.

- 分治法
 - 每个x所属的num_E(x)个集合中必选至少一个
 - 最多选k个集合
 - 因此,总共不超过PatPat种尝试

- 教材讨论
 - JH第3章第6节

问题1: local search的基本概念

- 你能解释这些术语的含义吗? 基于此,你能解释local search的基本思想吗?
 - feasible solution
 - transformation
 - neighborhood
 - local optimum

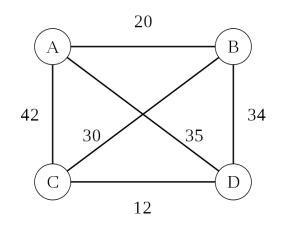
```
{\it LSS(Neigh)}	ext{-}{\it Local Search Scheme} according to a neighborhood Neigh
```

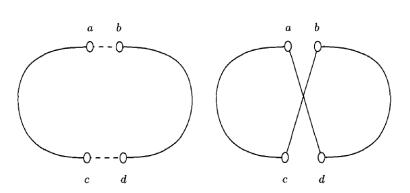
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Input: An input instance x of an optimization problem U. Step 1: Find a feasible solution \alpha \in \mathcal{M}(x). Step 2: while \alpha \notin LocOPT_U(x, Neigh_x) do begin find a \beta \in Neigh_x(\alpha) such that cost(\beta) < cost(\alpha) if U is a minimization problem and cost(\beta) > cost(\alpha) if U is a maximization problem; \alpha := \beta end Output: output(\alpha).
```

- 你能证明LSS的total correctness吗?
- · 决定LSS能否并尽快找到全局最优解的因素有哪些?
 - α
 - Neigh
 - β

问题2: hill climbing

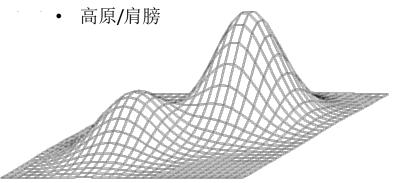
- In computer science, hill climbing is a mathematical optimization technique which belongs to the family of local search. It is an iterative algorithm that starts with an arbitrary solution to a problem, then attempts to find a better solution by incrementally changing a single element of the solution. If the change produces a better solution, an incremental change is made to the new solution, repeating until no further improvements can be found.
 - 这里的 α 、Neigh、β分别是怎么取的?
 - 你能以TSP为例,给出一个具体的算法吗?

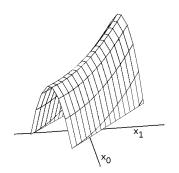




问题2: hill climbing (续)

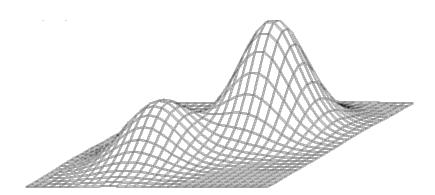
- In computer science, hill climbing is a mathematical optimization technique which belongs to the family of local search. It is an iterative algorithm that starts with an arbitrary solution to a problem, then attempts to find a better solution by incrementally changing a single element of the solution. If the change produces a better solution, an incremental change is made to the new solution, repeating until no further improvements can be found.
 - 你认为hill climbing存在哪些问题?
 - 局部最优
 - 缓升(如:之字形爬升非轴向的山脊)





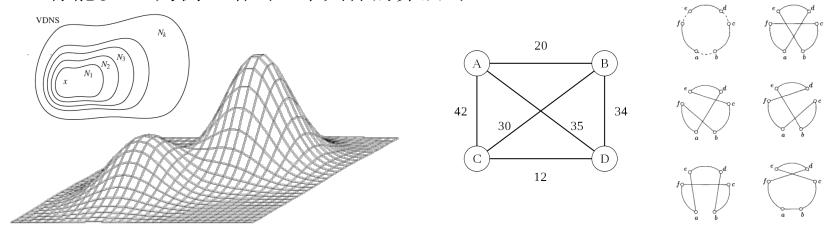


- 你能想到哪些策略来缓解这些问题?
 - $-\alpha$
 - Neigh
 - β



问题3: very large-scale neighborhood search

- A very large-scale neighborhood search is a local search algorithm which makes use of a neighborhood definition, which is large and possibly exponentially sized.
 - 和hill climbing相比,它改变了α、Neigh、β中的哪一个? 这样做有什么好处?
 - 你能以TSP为例,给出一个具体的算法吗?



- 你认为朴素的very large-scale neighborhood search存在什么问题?
- 你能想出折中的策略来应对这个问题吗?

问题3: very large-scale neighborhood search (续)

Input:

- Variable-depth search
 methods are techniques that
 search the k-exchange
 neighborhood partially, hence
 reducing the time used to
 search the neighborhood.
 - KL(Neigh)是如何实现 "partially"的? 你能简述它的思想吗?
 - 你能以TSP为例,给出一个 具体的算法吗?
 - 你还能想到其它方法来实现 "partially"吗?

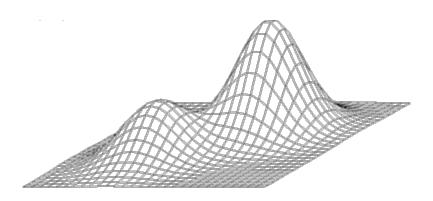
KL(Neigh) Kernighan-Lin Variable-Depth Search Algorithm with respect to the neighborhood Neigh

An input instance I of an optimization problem U.

```
Step 1: Generate a feasible solution \alpha = (p_1, p_2, \dots, p_n) \in \mathcal{M}(I) where (p_1, p_2, \dots, p_n) \in \mathcal{M}(I)
           p_2, \ldots, p_n) is such a parametric representation of \alpha that the local
           transformation defining Neigh can be viewed as an exchange of a few
           of these parameters.
Step 2: IMPROVEMENT := TRUE;
           EXCHANGE := \{1, 2, ..., n\}; J := 0; \alpha_J := \alpha;
           while IMPROVEMENT = TRUE do begin
              while EXCHANGE \neq \emptyset do
                 begin J := J + 1;
                   \alpha_J := a solution from Neigh(\alpha_{J-1}) such that gain(\alpha_{J-1}, \alpha_J)
                   is the maximum of
                   \{gain(\alpha_{J-1}, \delta) | \delta \in Neigh(\alpha_{J-1}) - \{\alpha_{J-1}\} \text{ and } \delta \text{ differs } \}
                   from \alpha_{J-1} in the parameters of EXCHANGE only};
                   EXCHANGE := EXCHANGE - \{ \text{the parameters in which } \}
                   \alpha_J and \alpha_{J-1} differ
              end:
              Compute gain(\alpha, \alpha_i) for i = 1, ..., J;
              Compute l \in \{1, \ldots, J\} such that
                    gain(\alpha, \alpha_i) = \max\{gain(\alpha, \alpha_i) | i \in \{1, 2, ..., J\}\};
              if gain(\alpha, \alpha_l) > 0 then
                 begin \alpha := \alpha_l;
                    EXCHANGE := \{1, 2, \dots, n\}
                 end
              else IMPROVEMENT := FALSE
           end
Step 3: output(\alpha).
```

问题4: Multi-start methods

- Re-start the procedure from a new solution once a region has been explored.
 - 和hill climbing相比,它改变了α、Neigh、β中的哪一个? 这样做有什么好处?
 - 你认为应该如何选择new solution?
 - 你能以TSP为例,给出一个具体的算法吗?



问题4: Multi-start methods (续)

- Greedy Randomized Adaptive Search Procedure (GRASP)
 - The GRASP metaheuristic is a multi-start or iterative process, in which each iteration consists of two phases: construction and local search. The construction phase builds a feasible solution, whose neighborhood is investigated until a local minimum is found during the local search phase. The best overall solution is kept as the result.

```
procedure GRASP(Max_Iterations, Seed)

1 Read_Input();

2 for k = 1,..., Max_Iterations do

3 Solution ← Greedy_Randomized_Construction(Seed);

4 Solution ← Local_Search(Solution);

5 Update_Solution(Solution, Best_Solution);

6 end;

7 return Best_Solution;
end GRASP.
```

Figure 1. Pseudo-code of the GRASP metaheuristic.

Greedy、Randomized、Adaptive 分别是如何体现的?

```
      procedure Greedy Randomized_Construction(Seed)

      1
      Solution \leftarrow \emptyset;

      2
      Evaluate the incremental costs of the candidate elements;

      3
      while Solution is not a complete solution do

      4
      Build the restricted candidate list (RCL);

      5
      Select an element s from the RCL at random;

      6
      Solution \leftarrow Solution \cup {s};

      7
      Reevaluate the incremental costs;

      8
      end;

      9
      return Solution;

      end Greedy_Randomized_Construction.
```

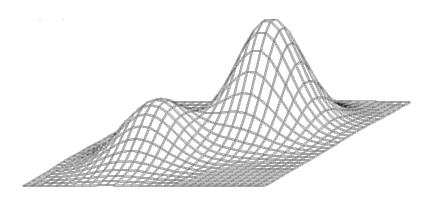
Figure 2. Pseudo-code of the construction phase.

```
 \begin{array}{lll} \textbf{procedure Local\_Search}(\texttt{Solution}) \\ \textbf{1} & \textbf{while Solution is not locally optimal do} \\ \textbf{2} & \textbf{Find } s' \in N(\texttt{Solution}) \ \text{with } f(s') < f(\texttt{Solution}); \\ \textbf{3} & \texttt{Solution} \leftarrow s'; \\ \textbf{4} & \textbf{end}; \\ \textbf{5} & \textbf{return Solution}; \\ \textbf{end Local\_Search}. \end{array}
```

Figure 3. Pseudo-code of the local search phase.

问题5: Stochastic hill climbing

- Stochastic hill climbing chooses at random from among the uphill moves.
 - 和hill climbing相比,它改变了 α 、Neigh、 β 中的哪一个? 这样做有什么好处?
 - 你认为应该如何计算每一种移动的概率?
 - The probability of selection can vary with the steepness of the uphill move.
 - 你能以TSP为例,给出一个具体的算法吗?



- 我们似乎已经讨论了对hill climbing的所有可能的改进策略
 - $-\alpha$: multi-start methods (GRASP)
 - Neigh: very large-scale neighborhood search (variable-depth search)
 - β: stochastic hill climbing
- 你还能想到别的方法吗?
 - 提示: 计算机除了"算"以外,还能做什么?

问题6: Tabu search

- Tabu search enhances the performance of local searches by using memory structures that describe the visited solutions or user-provided sets of rules.
 - Short-term: The list of solutions recently considered. If a potential solution appears on the tabu list, it cannot be revisited until it reaches an expiration point.
 - Intermediate-term: Intensification rules intended to bias the search towards promising areas of the search space.
 - Long-term: Diversification rules that drive the search into new regions (i.e. regarding resets when the search becomes stuck in a plateau or a suboptimal dead-end).
- 你能以TSP为例,给出一个具体的算法吗?

问题7: local search的性能

LSS(Neigh)-Local Search Scheme according to a neighborhood Neigh

```
Input: An input instance x of an optimization problem U. Step 1: Find a feasible solution \alpha \in \mathcal{M}(x). Step 2: while \alpha \notin LocOPT_U(x, Neigh_x) do begin find a \beta \in Neigh_x(\alpha) such that cost(\beta) < cost(\alpha) if U is a minimization problem and cost(\beta) > cost(\alpha) if U is a maximization problem; \alpha := \beta end Output: output(\alpha).
```

- local search的运算时间受哪些因素的影响?
- 什么是exact polynomial-time searchable neighborhood? 为什么cost-bounded integer-valued TSP没有这一性质?

问题8:应用

- 你能综合运用我们讨论的这些策略,分别为下列问题设计一种local search算法吗?
 - longest simple path
 - MAX-SAT
 - MAX-CL
 - MIN-VCP