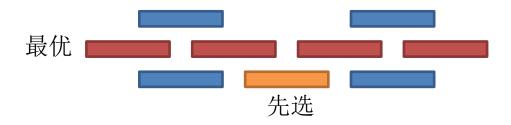
- 作业讲解
 - -TC第16.1节练习2、3
 - -TC第16.2节练习1、2
 - -TC第16.3节练习2、5、8
 - -TC第16章问题1

TC第16.1节练习3

Overlaps the fewest other remaining activities.



TC第16.2节练习1

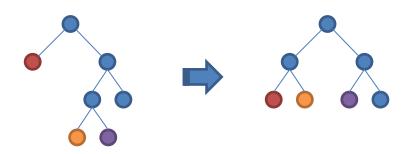
- 什么是greedy-choice property?
 - 一般的证明方法是怎样的?
- 按性价比降序,如果第一个i未被选入最优解,而是选了之后的某个j,那么会怎样?
 - 若j比i的性价比低,替换之后可得更优解,与最优解矛盾
 - 若j和i的性价比相等,替换之后仍是最优解

TC第16.3节练习5

- 这题并不限于Huffman编码,而是讨论任意的编码
- 频率: F₁≥F₂≥...≥F_n
- 如果最优解不满足L₁≤L₂≤...≤L_n,即对于i<j,有L_i>L_j,那么交换i和j,新旧总长度之差为(F_iL_i+F_iL_i)-(F_iL_i+F_iL_i)=(F_i-F_i)(L_i-L_i)≤0
 - 如果<0,与最优解矛盾
 - 如果=0,不断这样交换,可以调到 $L_1 ≤ L_2 ≤ ... ≤ L_n$

TC第16.3节练习8

- 关于binary tree的用词
 - A full binary tree is a tree in which every node other than the leaves has two children.
 - A complete binary tree is a binary tree in which every level, except possibly the last, is completely filled, and all nodes are as far left as possible.
 - A perfect binary tree is a full binary tree in which all leaves are at the same depth or same level, and in which every parent has two children.
- 为什么这里Huffman编码树一定是perfect(深度都为8)?
 - 否则,必有一个深度<8,两个深度>8
 - 改变树形,提2降1,总长度减小,与最优解矛盾



TC第16章问题1

• (b)

- 如果greedy给出的不是最优解,那么最优解可以视作这种形式:
 - c^k, ..., cⁱ⁺¹都用满了(与greedy一致)
 - cⁱ没有用满(比greedy只少不多)
 - cⁱ⁻¹, ..., c⁰先不管
- 最优解需要用剩余的ci-1, ..., c⁰填补缺口: 至少ci
- 在填补中,如果每个cⁱ都只用了少于c个,最多只能填补cⁱ-1,不合要求
- 因此,必有一个cⁱ用了不少于c个,那么可以换成一个cⁱ⁺¹,与最优解矛盾

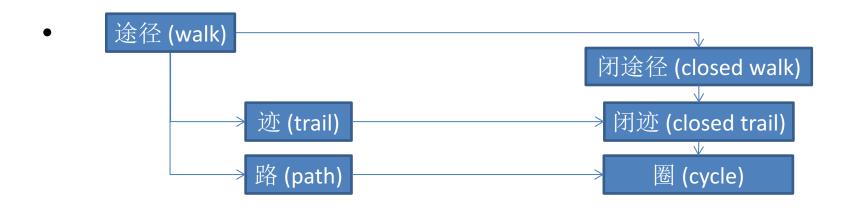
- 教材讨论
 - GC第1章
 - -GC第2章第1、2节
 - GC第3章第1节



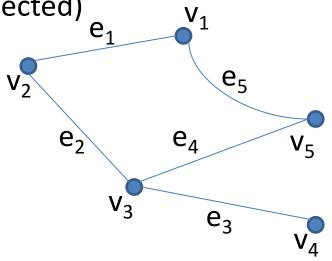
图论的学习, 仿佛逻辑思维训练的"艺术体操"

问题1: 图论中的术语

- 图 (graph)、顶点 (vertex/node)、边 (edge)
- 相邻 (adjacent)、关联 (incident)、邻居 (neighbor)
- 阶 (order)、边数 (size)
 - 边数的上下界是多少? (与阶为单位)
- 简单图 (simple graph)
- 平行边/重边 (parallel/multiple edges)、环边 (loop)
- 有向图
 - 有向边 (arc/directed edge)、有向图 (digraph)
 - 定向图 (oriented graph)、底图 (underlying graph)*



- 连通 (connected)、不连通 (disconnected)
- 连通分支 (component)
 - n个顶点的连通分支至少含几条边?
- 距离(distance)、直径 (diameter)



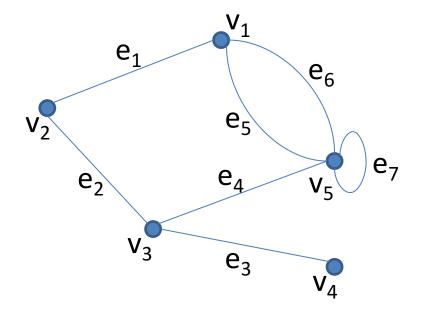
- 度 (degree)
 - 为什么奇度顶点的个数总是偶数?
- 孤立点 (isolated vertex)
- 叶子 (leaf)
- 最小度 (δ)、最大度 (Δ)
- 有向图
 - 入度 (indegree)、出度 (outdegree)

- 零图 (null graph)*
- 空图 (empty graph)
- 平凡图 (trivial graph)
- 完全图 (complete graph)
 - 完全图的边数和阶是什么关系?
- r-正则图 (r-regular graph)
 - 你能构造一个阶为5的3-正则图吗?
- 二部图/二分图 (bipartite graph)
 - 为什么不含奇圈是二部图的充要条件?
- 完全二部图 (complete bipartite graph)
 - 星 (star)
- k-部图 (k-partite graph)

- 子图 (subgraph)、真子图 (proper subgraph)
- 生成子图 (spanning subgraph)
- 导出子图 (induced subgraph)
- 同构 (isomorphism)
 - 为什么同构是一种等价关系?
- 补图 (complement)
- 自补图 (self-complementary)
 - 你能举出例子吗?
- 并 (union)
- 联/连接 (join)
- 笛卡儿积 (Cartesian product)

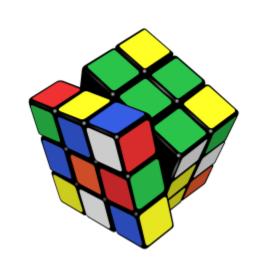
• 图的邻接矩阵表示

	v ₁	V ₂	V ₃	V ₄	v ₅
V ₁		1			2
V ₂	1		1		
V ₃		1		1	1
V ₄			1		
V ₅	2		1		1



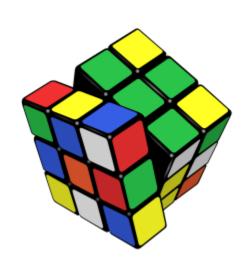
问题2: 用图来建模问题

- 解决给定的魔方,最快需要多少步?
- 你能构造一个需要解决步数最多的魔方吗?



问题2:用图来建模问题(续)

- 魔方状态转换图
 - 阶是多少?
 - 是连通图吗?
 - 是正则图吗?



问题2: 用图来建模问题(续)

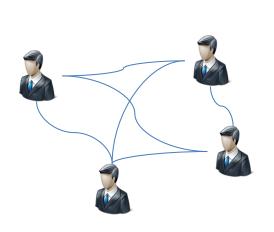
• the suspects

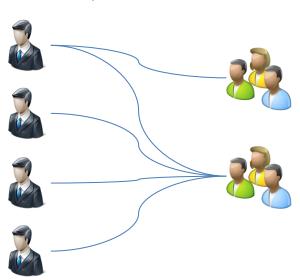
- Severe acute respiratory syndrome (SARS), an atypical pneumonia of unknown aetiology, was recognized as a global threat in mid-March 2003. To minimize transmission to others, the best strategy is to separate the suspects from others.
- In the Not-Spreading-Your-Sickness University (NSYSU), there are many student groups. Students in the same group intercommunicate with each other frequently, and a student may join several groups.
 To prevent the possible transmissions of SARS, the NSYSU collects the member lists of all student groups, and makes the following rule in their standard operation procedure (SOP).
 - Once a member in a group is a suspect, all members in the group are suspects.
- However, they find that it is not easy to identify all the suspects when a student is recognized as a suspect. Your job is to write a program which finds all the suspects.

问题2: 用图来建模问题(续)

the suspects

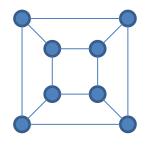
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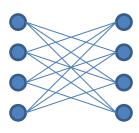




问题3: 图论中的常用证明方法

- 构造法,例如:
 - 证明以下两个图同构。





- 极端情况法(为自己新增一个条件),例如:
 - 证明: 最小度为2的简单图必有圈,最小度为3的简单图必有偶圈。

- 反证法,例如:
 - 任何简单图必有至少两个顶点具有相等的度。
 - 若G是简单图且 δ (G)≥k,则G有长为至少k的路。

- 分情况讨论法,例如:
 - 如果图G不连通,则其补图必连通。

- 数学归纳法,例如:
 - 设A是r阶图G的邻接矩阵,则An的第i行第j列元素 $a_{ij}^{(n)}$ 等于G中从 v_i 到 v_i 的长度为n的途径的数目($1 \le n < r$)。

- 数学归纳法,例如:
 - 设A是r阶图G的邻接矩阵,则An的第i行第j列元素 $a_{ij}^{(n)}$ 等于G中从 v_i 到 v_i 的长度为n的途径的数目($1 \le n < r$)。

$$a_{ij}^{(n)} = \sum_{r=1}^{\nu} a_{ir}^{(n-1)} a_{rj}$$

问题4: 图的集合表示

- 如何用集合的语言来表示一个无向图? 你能想到几种方式?
 - 方法1
 - V={v1, v2, ...}
 - E={e1, e2, ...}
 - endpoints(e1)={v1, v2}, endpoints(e2)={v3, v3}, ...
 - 方法2
 - V={v1, v2, ...}
 - E={{v1, v2}, {v3}, ...}
- 它们各有什么优缺点?



- 方法1
 - V={v1, v2, ...}
 - E={e1, e2, ...}
 - tail(e1)=v1, tail(e2)=v3, ...
 - head(e1)=v2, head(e2)=v3, ...
- 方法2
 - V={v1, v2, ...}
 - E={{{v1}, {v1, v2}}, {{v3}}, ...}







