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A Lower Bound for the Mixed μ Problem

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Abstract— The mixed μ problem has been shown to be NP hard so that exact analysis appears intractable. Our goal then is to exploit the problem structure so as to develop a polynomial time algorithm that approximates μ and usually gives good answers. To this end it is shown that μ is equivalent to a real eigenvalue maximization problem, and a power algorithm is developed to tackle this problem. The algorithm not only provides a lower bound for μ but has the property that μ is (almost) always an equilibrium point of the algorithm.

Index Terms—Computational methods, control system analysis, robust control, stability analysis, structured singular value.

I. INTRODUCTION

Computation schemes for the complex μ problem, based on upper and lower bounds [1], [2], are now well developed, and software is commercially available as part of the μ-Tools toolbox [3]. The mixed case, however, is a fundamentally more difficult problem and is much less understood. An upper bound for the general mixed μ problem was presented by Fan *et al.* [4] involving a minimization problem on the eigenvalues of a Hermitian matrix, and a practical scheme to compute this upper bound has recently been developed [5]. This paper addresses the problem of computing a lower bound for μ in the mixed case.

We begin with some preliminaries in Section II. It is known that the mixed μ problem is nonconvex and NP hard [6] so that, except for small problems or special cases, one cannot expect to compute exact solutions without an entirely unacceptable amount of computation. Nevertheless, we would like to quickly find approximate solutions to the problem. This motivates the power iteration approach taken in the paper. Previous work on complex μ problems has shown that power iterations are fast, seem to have some nice global properties, and give good answers most of the time. Of course we will not be able to provide guarantees about the global properties of our solution since the problem we are trying to solve is NP hard.

It is shown in Section III that mixed μ can be obtained as the result of a (nonconcave) real eigenvalue maximization. Sections IV–VI present several important theoretical characterizations of the mixed μ problem, including the generalization of the μ decomposition to the mixed case in Section V. This leads to the development of a power algorithm to compute a lower bound for the mixed μ problem which is presented in Section VI. The algorithm performance is very encouraging, both in terms of accuracy of the resulting bound and computational efficiency, and this is briefly discussed in Section VII.

II. NOTATION AND PRELIMINARIES

For any square complex matrix M, we denote the complex conjugate transpose by M^*. The largest singular value and the spectral radius are denoted by $\bar{\sigma}(M)$ and $\rho(M)$, respectively. The real spectral radius is defined as $\rho_R(M) = \max\{|\lambda|: \lambda \text{ is a real eigenvalue of } M\}$, with $\rho_R(M) = 0$ if M has no real eigenvalues. For any complex vector x, then x^T denotes the transpose, x^* the complex conjugate transpose, $|x|$ the Euclidean norm, and $|x|_\infty$ the infinity norm.

The definition of the structured singular value, μ, is dependent upon the underlying block structure of the uncertainties, which is defined as follows. Given a matrix $M \in \mathcal{C}^{n \times n}$ and three nonnegative integers m_r, m_c, and m_C with $m \doteq m_r + m_c + m_C \leq n$, the block structure $\mathcal{K}(m_r, m_c, m_C)$ is an m-tuple of positive integers

$$\mathcal{K} = (k_1, \cdots, k_{m_r}, k_{m_r+1}, \cdots, k_{m_r+m_c}, k_{m_r+m_c+1}, \cdots, k_m) \quad (1)$$

where we require $\sum_{i=1}^m k_i = n$ so these dimensions are compatible with M. This now determines the set of allowable perturbations, namely define

$$X_\mathcal{K} = \Big\{ \Delta = \text{block diag}\big(\delta_1^r I_{k_1}, \cdots, \delta_{m_r}^r I_{k_{m_r}}, \delta_1^c I_{k_{m_r+1}}, \cdots \\ \delta_{m_c}^c I_{k_{m_r+m_c}}, \Delta_1^C, \cdots, \Delta_{m_C}^C \big): \\ \delta_i^r \in \mathcal{R}, \delta_i^c \in \mathcal{C}, \Delta_i^C \in \mathcal{C}^{k_{m_r+m_c+i} \times k_{m_r+m_c+i}} \Big\}. \quad (2)$$

Note that $X_\mathcal{K} \subset \mathcal{C}^{n \times n}$ and that this block structure is sufficiently general to allow for repeated real scalars, repeated complex scalars, and full complex blocks. Note also that the full complex blocks need not be square, but we restrict them as such for notational convenience. The purely complex case corresponds to $m_r = 0$ and the purely real case to $m_c = m_C = 0$.

Definition 1 [1]: The structured singular value, $\mu_\mathcal{K}(M)$, of a matrix $M \in \mathcal{C}^{n \times n}$ with respect to a block structure $\mathcal{K}(m_r, m_c, m_C)$ is defined as

$$\mu_\mathcal{K}(M) = \Big(\min_{\Delta \in X_\mathcal{K}} \{\bar{\sigma}(\Delta): \det(I - \Delta M) = 0\} \Big)^{-1} \quad (3)$$

with $\mu_\mathcal{K}(M) = 0$ if no $\Delta \in X_\mathcal{K}$ solves $\det(I - \Delta M) = 0$.

In this paper we will be concerned directly with the computation of μ rather than how to use μ as a robustness analysis tool. For the reader unfamiliar with μ-based techniques, a fairly comprehensive review is given in [7].

Whilst it is not at all obvious how to compute μ from (3), it is easy to obtain the following crude upper and lower bounds:

$$\rho_R(M) \leq \mu_\mathcal{K}(M) \leq \bar{\sigma}(M). \quad (4)$$

To refine these bounds further we define the following sets of block diagonal matrices (also dependent on the underlying block structure):

$$\mathcal{Q}_\mathcal{K} = \big\{ \Delta \in X_\mathcal{K}: \delta_i^r \in [-1\ 1], \delta_i^{c*}\delta_i^c = 1, \Delta_i^{C*}\Delta_i^C = I_{k_{m_r+m_c+i}} \big\} \quad (5)$$

$$\mathcal{D}_{\mathcal{K}} = \Big\{ \text{block diag}\big(e^{j\theta_1} D_1, \cdots, e^{j\theta_{m_r}} D_{m_r}, D_{m_r+1}, \cdots,$$

$$D_{m_r+m_c}, d_1 I_{k_{m_r+m_c+1}}, \cdots, d_{m_C} I_{k_m} \big) : \theta_i \in \left[-\frac{\pi}{2} \frac{\pi}{2} \right],$$

$$0 < D_i = D_i^* \in \mathcal{C}^{k_i \times k_i}, 0 < d_i \in \mathcal{R} \Big\}. \tag{6}$$

Note that for any $\Delta \in X_{\mathcal{K}}$ and any $D \in \mathcal{D}_{\mathcal{K}}$, $D\Delta = \Delta D$, and consequently we obtain the following lemma.

Lemma 1: For any matrix $M \in \mathcal{C}^{n \times n}$ and any compatible block structure $\mathcal{K}$, for all $D \in \mathcal{D}_{\mathcal{K}}$

$$\mu_{\mathcal{K}}(M) = \mu_{\mathcal{K}}(DMD^{-1}). \tag{7}$$

Now to refine the lower bound we define the unit ball in the perturbation set as

$$\mathbf{B}X_{\mathcal{K}} = \{ \Delta \in X_{\mathcal{K}} : \bar{\sigma}(\Delta) \leq 1 \}. \tag{8}$$

The following lemma results almost immediately from the definition of μ.

Lemma 2: For any matrix $M \in \mathcal{C}^{n \times n}$ and any compatible block structure $\mathcal{K}$

$$\mu_{\mathcal{K}}(M) = \max_{\Delta \in \mathbf{B}X_{\mathcal{K}}} \rho_R(\Delta M). \tag{9}$$

In light of (7) and (9), noting that $\mathcal{Q}_{\mathcal{K}} \subset \mathbf{B}X_{\mathcal{K}}$, we can refine the bounds in (4) to obtain the following lemma.

Lemma 3: For any matrix $M \in \mathcal{C}^{n \times n}$ and any compatible block structure $\mathcal{K}$

$$\max_{Q \in \mathcal{Q}_{\mathcal{K}}} \rho_R(QM) \leq \mu_{\mathcal{K}}(M) \leq \inf_{D \in \mathcal{D}_{\mathcal{K}}} \bar{\sigma}(DMD^{-1}). \tag{10}$$

We introduce one further piece of notation. For any two vectors $x, y \in \mathcal{C}^n$, partition x and y according to the block structure as

$$x = \begin{bmatrix} x_{r_1}^T & \cdots & x_{r_{m_r}}^T & x_{c_1}^T & \cdots & x_{c_{m_c}}^T & x_{C_1}^T & \cdots & x_{C_{m_C}}^T \end{bmatrix}^T$$
$$y = \begin{bmatrix} y_{r_1}^T & \cdots & y_{r_{m_r}}^T & y_{c_1}^T & \cdots & y_{c_{m_c}}^T & y_{C_1}^T & \cdots & y_{C_{m_C}}^T \end{bmatrix}^T \tag{11}$$

where $x_{r_i}, y_{r_i} \in \mathcal{C}^{k_i}$, $x_{c_i}, y_{c_i} \in \mathcal{C}^{k_{m_r+i}}$, $x_{C_i}, y_{C_i} \in \mathcal{C}^{k_{m_r+m_c+i}}$. These will be referred to as the "block components" of x and y, and we define the "nondegeneracy" assumption to be that for every i (in the appropriate set), $|y_{r_i}^* x_{r_i}| \neq 0, |y_{c_i}^* x_{c_i}| \neq 0, |y_{C_i}||x_{C_i}| \neq 0$.

III. LOWER BOUND AS A MAXIMIZATION

In this section we show that the lower bound for the mixed case (10) holds with equality, and hence it is sufficient to consider the complex uncertainties on their boundary. Note, however, that the definition of $\mathcal{Q}_{\mathcal{K}}$ requires us to search over the full range of the real perturbations. The following lemma is taken from [1].

Lemma 4 [1]: Let $p : \mathcal{C}^k \to \mathcal{C}$ be a (multivariable) polynomial and define $\beta = \min\{|z|_\infty : p(z) = 0\}$, then there exists a $z \in \mathcal{C}^k$ such that $p(z) = 0$ and for every i, $|z_i| = \beta$.

This is now used to prove the main result of the section.

Theorem 1: For any matrix $M \in \mathcal{C}^{n \times n}$ and any compatible block structure $\mathcal{K}$

$$\max_{Q \in \mathcal{Q}_{\mathcal{K}}} \rho_R(QM) = \mu_{\mathcal{K}}(M). \tag{12}$$

Proof: It is trivial from (10) if $\mu_{\mathcal{K}}(M) = 0$. So assume $\mu_{\mathcal{K}}(M) = \beta > 0$, and this value is achieved for some perturbation $\hat{\Delta}$, i.e., $\det(I - \hat{\Delta} M) = 0$ and $\bar{\sigma}(\hat{\Delta}) \leq \frac{1}{\beta}$. Now *fix* the real perturbations at these "optimal" values ($\delta_i^r = \hat{\delta}_i^r, i = 1, \cdots, m_r$ with $|\hat{\delta}_i^r| \leq \frac{1}{\beta}$). Then allow the complex part of Δ to vary, and consider minimizing $\bar{\sigma}(\Delta)$ subject to $\det(I - \Delta M) = 0$. Performing a singular value

decomposition (SVD) on Δ, we obtain $\det(I - U\Sigma V M) = 0$ where U and V are (block diagonal) unitary matrices and

$$\Sigma = \text{diag}\big(\hat{\delta}_1^r I_{k_1}, \cdots \hat{\delta}_{m_r}^r I_{k_{m_r}}, \delta_1^c I_{k_{m_r+1}}, \cdots$$
$$\delta_{m_c}^c I_{k_{m_r+m_c}}, \gamma_1^c, \cdots, \gamma_k^c \big)$$

with $k = \sum_{i=m_r+m_c+1}^{m} k_i$. This is a polynomial in $\delta_1^c, \cdots, \delta_{m_c}^c, \gamma_1^c, \cdots, \gamma_k^c$ and so applying Lemma 4 we have a solution with $|\hat{\delta}_1^c| = \cdots = |\hat{\delta}_{m_c}^c| = |\hat{\gamma}_1^c| = \cdots = |\hat{\gamma}_k^c| = \frac{1}{\beta}$ and $\hat{\beta} \geq \beta$. Now suppose $\hat{\beta} > \beta$, say $\hat{\beta} = \beta + \epsilon$ for some $\epsilon > 0$, then since the roots of a polynomial are continuous functions of the coefficients, we can find a $\delta > 0$ so that

$$\left| \delta_i^r - \hat{\delta}_i^r \right| < \delta, i = 1, \cdots, m_r \Rightarrow \left| \delta_i^c - \hat{\delta}_i^c \right| < \frac{\epsilon}{2}, i = 1, \cdots, m_c$$
$$\left| \gamma_i^c - \hat{\gamma}_i^c \right| < \frac{\epsilon}{2}, i = 1, \cdots, k.$$

Then, move each $|\delta_i^r|$ down by $\frac{\delta}{2}$, and we can find a Δ solving $\det(I - \Delta M) = 0$ with $\bar{\sigma}(\Delta) < \frac{1}{\beta}$ contradicting the definition of μ. Thus $\hat{\beta} = \beta$, and it is now easy to check that for this solution $\beta \hat{\Delta} = \hat{Q} \in \mathcal{Q}_{\mathcal{K}}$ with $\rho_R(\hat{Q}M) = \beta = \mu_{\mathcal{K}}(M)$. $\square$

IV. CHARACTERIZATION OF A MAXIMUM POINT

We are interested in computing $\mu_{\mathcal{K}}(M)$, which by (9) and (12) is given by

$$\mu_{\mathcal{K}}(M) = \max_{\Delta \in \mathbf{B}X_{\mathcal{K}}} \rho_R(\Delta M) = \max_{Q \in \mathcal{Q}_{\mathcal{K}}} \rho_R(QM).$$

For reasons of tractability we choose to consider the problem $\max_{Q \in \mathcal{Q}_{\mathcal{K}}} \rho_R(QM)$. However, since this is a nonconcave problem we will in general only be able to find local maxima, and hence we will obtain a lower bound for $\mu_{\mathcal{K}}(M)$ (which is the global maximum). We would like this lower bound be "tight" (i.e., close to μ) and so wish to rule out the maxima of $\rho_R(QM)$ which we know are only local. Thus we only consider $Q \in \mathcal{Q}_{\mathcal{K}}$ which are local maxima of $\rho_R(QM)$ with respect not only to $Q \in \mathcal{Q}_{\mathcal{K}}$ but also to $Q \in \mathbf{B}X_{\mathcal{K}}$. In this section we will develop a characterization of such local maxima.

Note that for any $Q \in \mathcal{Q}_{\mathcal{K}}$ and any $\Delta \in \mathbf{B}X_{\mathcal{K}}$, $Q\Delta \in \mathbf{B}X_{\mathcal{K}}$ and $\Delta Q \in \mathbf{B}X_{\mathcal{K}}$. Now suppose some matrix $Q \in \mathcal{Q}_{\mathcal{K}}$ achieves a local maximum of $\rho_R(QM)$ over $Q \in \mathbf{B}X_{\mathcal{K}}$. Then it is easy to show that the matrix $\hat{M} := QM$ has a local maximum of $\rho_R(\hat{Q}\hat{M})$ over $\hat{Q} \in \mathbf{B}X_{\mathcal{K}}$ at $\hat{Q} = I$. However, since the real elements of Q are not restricted to be on their boundary, we can say more than this. For any matrix $Q \in \mathcal{Q}_{\mathcal{K}}$ [see (5)] define the index sets

$$\mathcal{J}(Q) = \left\{ i \leq m_r : \left| \delta_i^r \right| = 1 \right\} \tag{13}$$
$$\hat{\mathcal{J}}(Q) = \left\{ i \leq m_r : \left| \delta_i^r \right| < 1 \right\} \tag{14}$$

and define the allowable perturbation set

$$\hat{\mathbf{B}}\Delta_\epsilon(\mathcal{J}, \hat{\mathcal{J}}) = \Big\{ \Delta \in X_{\mathcal{K}} : \left| \delta_i^r \right| \leq 1, i \in \mathcal{J}, \left| \delta_i^r \right| < 1 + \epsilon,$$
$$i \in \hat{\mathcal{J}}, \left| \delta_i^c \right| \leq 1, i = 1, \cdots, m_c,$$
$$\bar{\sigma}\big(\Delta_i^C \big) \leq 1, i = 1, \cdots, m_C \Big\}. \tag{15}$$

We see that for sufficiently small $\epsilon > 0$, for any $Q \in \mathcal{Q}_{\mathcal{K}}$ and any $\Delta \in \hat{\mathbf{B}}\Delta_\epsilon(\mathcal{J}(Q), \hat{\mathcal{J}}(Q))$, $Q\Delta \in \mathbf{B}X_{\mathcal{K}}$ and $\Delta Q \in \mathbf{B}X_{\mathcal{K}}$. The point of all this is that if some matrix $Q \in \mathcal{Q}_{\mathcal{K}}$ achieves a local maximum of $\rho_R(QM)$ over $Q \in \mathbf{B}X_{\mathcal{K}}$, then the matrix $\hat{M} := QM$ has a local maximum of $\rho_R(\hat{Q}\hat{M})$ over $\hat{Q} \in \hat{\mathbf{B}}\Delta_\epsilon(\mathcal{J}(Q), \hat{\mathcal{J}}(Q))$ (for some $\epsilon > 0$) at $\hat{Q} = I$ (and in fact the converse is true, provided we assume that for every i, $\delta_i^r \neq 0$).

Before proving the main result we need some preliminary lemmas. The following two linear algebra lemmas are due to Packard [2].

Lemma 5 [2]: Let $y \in \mathcal{C}^n$ and $x \in \mathcal{C}^n$ be nonzero vectors. Then there exists a $d \in \mathcal{R}, d > 0$ such that $y = dx$ iff $\mathrm{Re}(y^* G x) \leq 0$ for every $G \in \mathcal{C}^{n \times n}$ satisfying $G + G^* \leq 0$.

Lemma 6 [2]: Let $y \in \mathcal{C}^n$ and $x \in \mathcal{C}^n$ be nonzero vectors. Then there exists a Hermitian, positive definite $D \in \mathcal{C}^{n \times n}$ such that $y = Dx$ iff $y^* x \in \mathcal{R}$ and $y^* x > 0$.

Now define the closed half-space in the complex plane as, for some scalar $\psi \in \mathcal{R}$

$$H^\psi = \{z : \mathrm{Re}(e^{j\psi} z) \leq 0\}. \tag{16}$$

Then we have the following elementary linear algebra lemmas.

Lemma 7: Given any set of complex scalars $\mathcal{Z} = \{z_i : i = 1, \cdots, m\}$ and any real scalar ψ, then $\mathcal{Z} \subset H^\psi$ iff $\sum_{i=1}^m \alpha_i z_i \in H^\psi$ for all real nonnegative scalars $\alpha_i, i = 1, \cdots, m$.

Proof ($\Leftarrow$): For each z_k choose $\alpha_k = 1$ and $\alpha_i = 0$ for $i \neq k$

$$(\Rightarrow) \qquad \mathrm{Re}\left(e^{j\psi} \sum_{i=1}^m \alpha_i z_i\right) = \sum_{i=1}^m \alpha_i \mathrm{Re}(e^{j\psi} z_i) \leq 0. \quad \square$$

Lemma 8: Given any set of complex scalars $\mathcal{Z} = \{z_i : i = 1, \cdots, m\}$, define $\lambda := \sum_{i=1}^m \alpha_i z_i$ where $\alpha_i, i = 1, \cdots, m$ are real nonnegative scalars. Then λ is not real and positive for any choice of the above α_i's iff $\mathcal{Z} \subset H^\psi$ for some $\psi \in (-\frac{\pi}{2}, \frac{\pi}{2})$.

Proof ($\Leftarrow$): By Lemma 7 $\mathcal{Z} \subset H^\psi$ implies $\lambda \in H^\psi$, and hence $\mathrm{Re}(e^{j\psi} \lambda) \leq 0$. Suppose λ is real and positive. Then this implies $\mathrm{Re}(e^{j\psi}) \leq 0$ which means $\psi \notin (-\frac{\pi}{2}, \frac{\pi}{2})$, which is a contradiction.

($\Rightarrow$): Assume λ is never real and positive. Now suppose $\mathcal{Z} \not\subset H^\psi$ for any $\psi \in (-\frac{\pi}{2}, \frac{\pi}{2})$. First choose $\psi = 0$. Then, we must have at least one $z \in \mathcal{Z}$ with $\mathrm{Re}(z) > 0$. Now we choose $\hat{z}_1$ as the element with $\mathrm{Re}(z) > 0$ having minimum $|\arg(z)|$ (which must be nonzero). Now choose $\psi = \arg(\hat{z}_1)$ and define $\hat{\psi} = \frac{\pi}{2} - \psi$. Then, since $\hat{z}_1 \in H^{\hat{\psi}}$, we must have a (nonzero) $\hat{z}_2 \in \mathcal{Z}$ with $\hat{z}_2 \notin H^{\hat{\psi}}$. Suppose

$$\hat{z}_1 = r_1(\cos\psi + j\sin\psi), \quad \hat{z}_2 = r_2(\cos\phi + j\sin\phi).$$

Then by our choice of $\hat{z}_1$ and $\hat{z}_2$, straightforward trigonometry yields the following facts: $|\sin\phi| \geq |\sin\psi|$, $\mathrm{sgn}(\sin\phi) = -\mathrm{sgn}(\sin\psi)$, $|\cos\phi| \leq |\cos\psi|$, and if $|\cos\phi| = |\cos\psi|$, then $\cos\phi = \cos\psi$. Now choose $\hat{\alpha}_1 = \frac{1}{r_1|\sin\psi|}$ and $\hat{\alpha}_2 = \frac{1}{r_2|\sin\phi|}$. Then we have

$$\hat{\lambda} = \hat{\alpha}_1 \hat{z}_1 + \hat{\alpha}_2 \hat{z}_2 = \frac{\cos\psi}{|\sin\psi|} + \frac{\cos\phi}{|\sin\phi|}.$$

Thus $\hat{\lambda}$ is real and positive, which is a contradiction. $\square$

Putting all this together we obtain the following alignment condition.

Theorem 2: Suppose the matrix $M \in \mathcal{C}^{n \times n}$ has a distinct real eigenvalue $\lambda_0 > 0$ with right and left eigenvectors, x and y, respectively, satisfying the nondegeneracy assumption. Further, suppose that $\rho_R(M) = \lambda_0$. Then if the function $\rho_R(QM)$ attains a local maximum over the set $Q \in \hat{\mathbf{B}}\Delta_\epsilon(\mathcal{J}, \hat{\mathcal{J}})$ (for some $\epsilon > 0$) at $Q = I$, then there exists a matrix $D \in \mathcal{D}_\mathcal{K}$, with $\theta_i = \pm\frac{\pi}{2}$ for every $i \in \hat{\mathcal{J}}$, and a real scalar $\psi \in (-\frac{\pi}{2}, \frac{\pi}{2})$ such that $y = e^{j\psi} Dx$.

Proof: First we parameterize the perturbation set. Consider $G \in X_\mathcal{K}$ with

$$G = \text{block diag}\big(g_1^r I_{k_1}, \cdots, g_{m_r}^r I_{k_{m_r}}, g_1^c I_{k_{m_r+1}}, \cdots,$$
$$g_{m_c}^c I_{k_{m_r+m_c}}, G_1^C, \cdots, G_{m_C}^C\big) \tag{17}$$

and the added restrictions

$$\begin{aligned} g_i^r &\leq 0, \quad i \in \mathcal{J} \\ \mathrm{Re}(g_i^c) &\leq 0, \quad i = 1, \cdots, m_c \\ G_i^C + G_i^{C*} &\leq 0, \quad i = 1, \cdots, m_C. \end{aligned} \tag{18}$$

Now it can be shown that for some $\delta > 0$, the set of all matrices $E(t) := (I + Gt)(I - Gt)^{-1}$ for G as above and t such that

$t\bar{\sigma}(G) \in [0 \ \delta)$ is an open neighborhood of $\hat{\mathbf{B}}\Delta_\epsilon(\mathcal{J}, \hat{\mathcal{J}})$ about $E(0) = I$. So now define the matrix $R(t) := E(t)M$. Then it is clear that $\rho_R(QM)$ has attained a local maximum over the set $Q \in \hat{\mathbf{B}}\Delta_\epsilon(\mathcal{J}, \hat{\mathcal{J}})$ at $Q = I$ iff $\rho_R(R(t))$ has attained a local maximum over $t \geq 0$ at $t = 0$ *for arbitrary G as above.*

Since $R(0) = M$ has a distinct real eigenvalue λ_0, we have (for some nonempty interval about the origin) an analytic function $\lambda(t)$, with $\lambda(0) = \lambda_0$, and $\lambda(t)$ an eigenvalue of $R(t)$. Thus we can differentiate to obtain

$$\dot{\lambda}(0) = y^* \dot{R}(0) x = 2y^* GMx = 2\lambda_0 y^* Gx. \tag{19}$$

In block notation this becomes

$$\dot{\lambda}(0) = 2\lambda_0 \left(\sum_{i=1}^{m_r} g_i^r y_{r_i}^* x_{r_i} + \sum_{i=1}^{m_c} g_i^c y_{c_i}^* x_{c_i} + \sum_{i=1}^{m_C} y_{C_i}^* G_i^C x_{C_i}\right). \tag{20}$$

Define the set of points

$$\begin{aligned} \mathcal{Z} = \{z_i : i = 1, \cdots, m\} &= \big\{g_i^r y_{r_i}^* x_{r_i} : i = 1, \cdots, m_r\big\} \\ &\cup \big\{g_i^c y_{c_i}^* x_{c_i} : i = 1, \cdots, m_c\big\} \\ &\cup \big\{y_{C_i}^* G_i^C x_{C_i} : i = 1, \cdots, m_C\big\} \end{aligned} \tag{21}$$

with the obvious identification for the elements z_i. Now since we are at a maximum point we have that $\dot{\lambda}(0)$ is never real and positive. Thus, noting that we may independently scale g_i^r, g_i^c, G_i^C by arbitrary nonnegative scalars and still satisfy (18), applying Lemma 8 to (20) and (21) gives that this is true iff $\mathcal{Z} \subset H^\psi$ for some $\psi \in (-\frac{\pi}{2}, \frac{\pi}{2})$ for each $G \in X_\mathcal{K}$ satisfying (18). Furthermore, since any summation of G's satisfying (18) also satisfies (18), Lemma 8 gives that this is true iff there is one H^ψ which works for *every* G, i.e., there exists $\psi \in (-\frac{\pi}{2}, \frac{\pi}{2})$ such that $\mathcal{Z} \subset H^\psi$ for *all* $G \in X_\mathcal{K}$ satisfying (18). From the definition of H^ψ in (16), and G in (17), (18), this is equivalent to

$$\begin{aligned} &\mathrm{Re}\big(e^{j\psi} g_i^r y_{r_i}^* x_{r_i}\big) \leq 0, \\ &\qquad \text{for all } g_i^r \in \mathcal{R} \text{ with } g_i^r \leq 0, \ i = 1, \cdots, m_r \\ &\mathrm{Re}\big(e^{j\psi} g_i^r y_{r_i}^* x_{r_i}\big) \leq 0 \\ &\qquad \text{for all } g_i^r \in \mathcal{R}, \ i \in \hat{\mathcal{J}} \\ &\mathrm{Re}\big(e^{j\psi} g_i^c y_{c_i}^* x_{c_i}\big) \leq 0, \\ &\qquad \text{for all } g_i^c \in \mathcal{C} \text{ with } \mathrm{Re}(g_i^c) \leq 0, \ i = 1, \cdots, m_c \\ &\mathrm{Re}\big(e^{j\psi} y_{C_i}^* G_i^C x_{C_i}\big) \leq 0, \\ &\qquad \text{for all } G_i^C \text{ with } G_i^C + G_i^{C*} \leq 0, \ i = 1, \cdots, m_C \end{aligned} \tag{22}$$

for some $\psi \in (-\frac{\pi}{2}, \frac{\pi}{2})$. It is now easy to check that the above conditions may be equivalently expressed as

$$\begin{aligned} &\mathrm{Re}\big(e^{j\psi} y_{r_i}^* x_{r_i}\big) \geq 0, \quad i = 1, \cdots, m_r \\ &\mathrm{Re}\big(e^{j\psi} y_{r_i}^* x_{r_i}\big) = 0, \quad i \in \hat{\mathcal{J}} \\ &e^{j\psi} y_{c_i}^* x_{c_i} \in (0 \ \infty), \quad i = 1, \cdots, m_c \\ &\mathrm{Re}\big(e^{j\psi} y_{C_i}^* G_i^C x_{C_i}\big) \leq 0, \quad \text{for all } G_i^C \text{ with } G_i^C + G_i^{C*} \leq 0, \\ &\qquad\qquad\qquad\qquad\qquad\qquad i = 1, \cdots, m_C. \end{aligned} \tag{23}$$

Since the scalar $e^{j\psi}$ terms may simply be absorbed into one of the vectors, we can apply Lemmas 5 and 6 to each block component of x and y to obtain the equivalent conditions

$$\begin{aligned} y_{r_i} &= e^{j\psi} e^{j\theta_i} D_i x_{r_i}, \quad 0 < D_i = D_i^*, \ \theta_i \in \left[-\tfrac{\pi}{2}, \tfrac{\pi}{2}\right], \\ &\qquad\qquad\qquad\qquad\qquad\qquad i = 1, \cdots, m_r \\ y_{r_i} &= e^{j\psi} e^{j\theta_i} D_i x_{r_i}, \quad 0 < D_i = D_i^*, \ \theta_i = \pm\tfrac{\pi}{2}, \ i \in \hat{\mathcal{J}} \\ y_{c_i} &= e^{j\psi} D_i x_{c_i}, \qquad 0 < D_i = D_i^*, \ i = 1, \cdots, m_c \\ y_{C_i} &= e^{j\psi} d_i x_{C_i}, \qquad 0 < d_i \in \mathcal{R}, \ i = 1, \cdots, m_C. \end{aligned} \tag{24}$$

Stacking these relations in matrix form yields $y = e^{j\psi} D x$ with D of the required form. □

Remarks: Note from the proof that we immediately have a partial converse to Theorem 2, namely that if $y = e^{j\psi} D x$ under the above assumptions, then no directional derivative (in the above sense) of the eigenvalue achieving $\rho_R(QM)$ over the set $Q \in \hat{\mathbf{B}}\Delta_\epsilon(\mathcal{J}, \hat{\mathcal{J}})$ is real and positive at $Q = I$.

V. A Decomposition at μ

Theorem 2 gives us a characterization of a maximum point of $\rho_R(QM)$ in terms of an alignment of the right and left eigenvectors of QM. This leads directly to the following decomposition.

Theorem 3: Suppose $Q \in \mathcal{Q}_\mathcal{K}$ achieves a local maximum of $\rho_R(QM)$ over $Q \in \mathbf{B}X_\mathcal{K}$ and that the eigenvalue achieving $\rho_R(QM)$, denoted β, is distinct and positive. Then, if the right and left eigenvectors of QM, denoted x and y, respectively, satisfy the nondegeneracy assumption, there exists a matrix $D \in \mathcal{D}_\mathcal{K}$ with $D^2 \in \mathcal{D}_\mathcal{K}$ and $\theta_i = \pm\frac{\pi}{4}$ for $i \in \hat{\mathcal{J}}(Q)$ such that

$$QDMD^{-1}(Dx) = \beta Dx$$
$$(x^* D^*)QD^* M(D^*)^{-1} = \beta x^* D^* \qquad (25)$$

with $\beta \le \mu_\mathcal{K}(M)$. Furthermore, if the above maximum is global, then $\beta = \mu_\mathcal{K}(M)$.

Proof: Since $Q \in \mathcal{Q}_\mathcal{K}$ is a local maximum of $\rho_R(QM)$ over $Q \in \mathbf{B}X_\mathcal{K}$, the matrix $\hat{M} := QM$ achieves a local maximum of $\rho_R(\hat{Q}\hat{M})$ over $\hat{Q} \in \hat{\mathbf{B}}\Delta_\epsilon(\mathcal{J}(Q), \hat{\mathcal{J}}(Q))$ (for some $\epsilon > 0$) at $\hat{Q} = I$. Now apply Theorem 2 to conclude $y = e^{j\psi}\hat{D}x$ with $\hat{D} \in \mathcal{D}_\mathcal{K}$ and $\hat{\theta}_i = \pm\frac{\pi}{2}$ for $i \in \hat{\mathcal{J}}(Q)$, then define D as the unique matrix such that $D \in \mathcal{D}_\mathcal{K}$ and $D^2 = \hat{D}$. Substitution of this into the right and left eigenvalue equations of QM and simple manipulations (note that for any $Q \in \mathcal{Q}_\mathcal{K}$ and any $D \in \mathcal{D}_\mathcal{K}$, Q and D commute) yield the results in (25). Finally, note from Theorem 1 that we have $\beta \le \mu_\mathcal{K}(M)$, and if the above maximum is global then $\beta = \mu_\mathcal{K}(M)$. □

Remarks: Employing simple manipulations of (25) yields a partial converse of this theorem. If we have a decomposition as in (25) with β real and positive and x nonzero, then we have that β is an eigenvalue of QM with right and left eigenvectors, x and y, respectively [thus β is a lower bound for $\mu_\mathcal{K}(M)$], where $y = re^{j\psi}D^2x$ with D as above, r a positive real scalar (which we could thus absorb into D), and $\psi \in [-\frac{\pi}{2}, \frac{\pi}{2}]$. Thus defining $\hat{D} = rD^2$ we have $y = e^{j\psi}\hat{D}x$ with $\hat{D}$ as in Theorem 2 and $\psi \in [-\frac{\pi}{2}, \frac{\pi}{2}]$. If we add the further technical assumption that we are not in the special case of $\theta_i = \pm\frac{\pi}{4}$ for all $i = 1, \cdots, m_r$ and $m_c = 0$, $m_C = 0$, then we have $\psi \in (-\frac{\pi}{2}, \frac{\pi}{2})$.

Thus, we (almost) always have a decomposition at μ of the form (25), and any such decomposition gives us a lower bound for μ. Now we reformulate this condition into a set of vector equations.

Lemma 9: Suppose we have matrices $Q \in \mathcal{Q}_\mathcal{K}$ with $\delta_i^r \ne 0$ for $i = 1, \cdots, m_r$ and $\hat{D} \in \mathcal{D}_\mathcal{K}$ with $\hat{D}^2 \in \mathcal{D}_\mathcal{K}$ and $\hat{\theta}_i = \pm\frac{\pi}{4}$ for $i \in \hat{\mathcal{J}}(Q)$. Then we have a nonzero vector $\hat{x}$ and a real positive scalar β such that

$$Q\hat{D}M\hat{D}^{-1}(\hat{D}\hat{x}) = \beta\hat{D}\hat{x}$$
$$(\hat{x}^*\hat{D}^*)Q\hat{D}^* M(\hat{D}^*)^{-1} = \beta\hat{x}^*\hat{D}^* \qquad (26)$$

iff there exists a matrix $D \in \mathcal{D}_\mathcal{K}$ with $\theta_i = \pm\frac{\pi}{2}$ for $i \in \hat{\mathcal{J}}(Q)$ and nonzero vectors b, a, z, w such that

$$Mb = \beta a \qquad M^*z = \beta w$$
$$b = Qa \qquad b = D^{-1}w$$
$$z = Q^*QDa \qquad z = Q^*w. \qquad (27)$$

Proof ($\Rightarrow$): Define $x = \hat{D}\hat{x}$ and b, a, z, w as $b = \hat{D}^{-1}x, a = \hat{D}^{-1}Q^{-1}x, z = \hat{D}Q^*x, w = \hat{D}x$. Finally, define $D = \hat{D}^2$; the result follows.

($\Leftarrow$): Define $\hat{D}$ as the unique matrix $\hat{D} \in \mathcal{D}_\mathcal{K}$ such that $\hat{D}^2 = D$, and $\hat{x} = b$, the result follows directly. □

VI. A Power Algorithm for the Lower Bound

In light of Lemma 9, the problem of computing a lower bound for $\mu_\mathcal{K}(M)$ is reduced to one of finding a solution to the set of equations in (27) which gives us a decomposition as in (25). We would like to develop an algorithm for computing such a solution. First note that if we partition b, a, z, w compatibly with the block structure as in (11), then the set of constraint equations

$$b = Qa \qquad b = D^{-1}w$$
$$z = Q^*QDa \qquad z = Q^*w$$

can be broken down into a series of m similar independent constraint equations on the block components (since Q and D are block diagonal). These equations are of three types corresponding to a repeated real scalar block, a repeated complex scalar block, or a full complex block. We now consider a generic constraint of each type. The following two lemmas are due to Packard [2].

Lemma 10 (Repeated Complex Scalar Block [2]): Let $b, a, z, w \in \mathcal{C}^k$ be nonzero vectors with $a^*w \ne 0$. Then there exists a complex scalar q with $|q| = 1$ and a complex matrix $D \in \mathcal{C}^{k \times k}$ with $0 < D = D^*$ such that

$$b = qa \qquad b = D^{-1}w$$
$$z = q^*qDa \qquad z = q^*w$$

if and only if

$$z = \frac{w^*a}{|w^*a|}w \qquad b = \frac{a^*w}{|a^*w|}a. \qquad (28)$$

Lemma 11 (Full Complex Block [2]): Let $b, a, z, w \in \mathcal{C}^k$ be nonzero vectors. There exists a complex matrix $Q \in \mathcal{C}^{k \times k}$ with $Q^*Q = I_k$ and a real positive scalar d such that

$$b = Qa \qquad b = d^{-1}w$$
$$z = Q^*Qda \qquad z = Q^*w$$

if and only if

$$z = \frac{|w|}{|a|}a \qquad b = \frac{|a|}{|w|}w. \qquad (29)$$

Now we consider a repeated real scalar block, bearing in mind that we have additional constraints if the real perturbation is not on the boundary (i.e., for $i \in \hat{\mathcal{J}}(Q)$).

Lemma 12 (Repeated Real Scalar Block): Let $b, a, z, w \in \mathcal{C}^k$ be nonzero vectors with $a^*w \ne 0$. We have a real scalar q with $|q| \le 1$, a real scalar $\theta \in [-\frac{\pi}{2}, \frac{\pi}{2}]$, and a complex matrix $D \in \mathcal{C}^{k \times k}$ with $0 < D = D^*$ such that

$$b = qa \qquad b = e^{-j\theta}D^{-1}w$$
$$z = q^*qe^{j\theta}Da \qquad z = q^*w$$

with $\theta = \pm\frac{\pi}{2}$ for $|q| < 1$ iff

$$z = qw \qquad b = qa \qquad (30)$$

with

$$\mathrm{Re}(a^*w) \ge 0 \quad \text{for } q = 1$$
$$\mathrm{Re}(a^*w) \le 0 \quad \text{for } q = -1$$
$$\mathrm{Re}(a^*w) = 0 \quad \text{for } |q| < 1. \qquad (31)$$

Proof ($\Rightarrow$)*:* Immediately we have $z = qw$ and $b = qa$. Thus $a^*w = \frac{1}{q}b^*w = \frac{1}{q}e^{j\theta}w^*(D^*)^{-1}w$. Now $q = 1$ implies $\arg(a^*w) = \theta$, and hence $\text{Re}(a^*w) \geq 0$. Similarly, $q = -1$ implies $\arg(a^*w) = \theta + \pi$ and hence $\text{Re}(a^*w) \leq 0$. Finally, $|q| < 1$ implies $\arg(a^*w) = \theta$ or $\theta + \pi$ with $\theta = \pm\frac{\pi}{2}$. Thus $\arg(a^*w) = \pm\frac{\pi}{2}$, and so $\text{Re}(a^*w) = 0$.

($\Leftarrow$)*:* Immediately we have $b = qa$ and $z = q^*w$, and so $b^*w = qa^*w$. Denoting $\theta = \arg(b^*w)$, we see that for $q = 1$ $\text{Re}(a^*w) \geq 0$, which implies $\text{Re}(b^*w) \geq 0$, and so $\theta \in [-\frac{\pi}{2}, \frac{\pi}{2}]$. Similarly for $q = -1$ $\text{Re}(a^*w) \leq 0$, which implies $\text{Re}(b^*w) \geq 0$, and so $\theta \in [-\frac{\pi}{2}, \frac{\pi}{2}]$. Finally for $|q| < 1$, $\text{Re}(a^*w) = 0$ which implies $\text{Re}(b^*w) = 0$, and so $\theta = \pm\frac{\pi}{2}$. Now $b^*(e^{-j\theta}w)$ is real and positive, and so applying Lemma 6 we have a matrix $\hat{D}$ with $0 < \hat{D} = \hat{D}^*$ such that $b = e^{-j\theta}\hat{D}w$. Define $D = \hat{D}^{-1}$, and we have $b = e^{-j\theta}D^{-1}w$ and $z = q^*w = q^*e^{j\theta}Db = q^*qe^{j\theta}Da$. $\square$

These lemmas now allow us (with a few technical assumptions) to eliminate matrices Q and D from (27). To avoid the notation becoming excessive, we consider a simple block structure with $m_r = m_c = m_C = 1$ for the remainder of this section. We stress that this is purely for notational convenience, and the general formulas for an arbitrary block structure, as defined in Section II, are simply obtained by duplicating the appropriate formulas for each block. So given $\mathcal{K} = (k_1, k_2, k_3)$, the appropriate scaling sets become

$$\mathcal{Q}_{\text{sub}} = \big\{\text{block diag}\big(q^r I_{k_1}, q^c I_{k_2}, Q^C\big) : q^r \in [-1\ 1],$$
$$q^{c*}q^c = 1, Q^{C*}Q^C = I_{k_3}\big\} \tag{32}$$

$$\mathcal{D}_{\text{sub}} = \big\{\text{block diag}\big(e^{j\theta}D_1, D_2, dI_{k_3}\big) : \theta \in \Big[-\frac{\pi}{2}, \frac{\pi}{2}\Big],$$
$$0 < D_i = D_i^* \in \mathcal{C}^{k_i \times k_i}, 0 < d \in \mathcal{R}\big\} \tag{33}$$

and we partition b, a, z, w compatibly with this block structure as

$$b = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix}, \quad a = \begin{bmatrix} a_1 \\ a_2 \\ a_3 \end{bmatrix}, \quad z = \begin{bmatrix} z_1 \\ z_2 \\ z_3 \end{bmatrix}, \quad w = \begin{bmatrix} w_1 \\ w_2 \\ w_3 \end{bmatrix} \tag{34}$$

where $b_i, a_i, z_i, w_i \in \mathcal{C}^{k_i}$. Then we obtain our final form of (27) as in the following theorem, which will form the basis of a power iteration to compute a lower bound for $\mu_\mathcal{K}(M)$.

Theorem 4: Suppose we have vectors $b, a, z, w \in \mathcal{C}^n$ partitioned as in (34) with $b_i, a_i, z_i, w_i \neq 0$ and $a_1^*w_1, a_2^*w_2 \neq 0$. Then there exist matrices $Q \in \mathcal{Q}_{\text{sub}}$ and $D \in \mathcal{D}_{\text{sub}}$ and a positive real scalar β such that

$$Mb = \beta a \qquad M^*z = \beta w$$
$$b = Qa \qquad b = D^{-1}w$$
$$z = Q^*QDa \qquad z = Q^*w$$

with $\theta \in [-\frac{\pi}{2}, \frac{\pi}{2}]$, and $\theta = \pm\frac{\pi}{2}$ for $|q^r| < 1$ iff

$$Mb = \beta a$$
$$z_1 = qw_1 \quad z_2 = \frac{w_2^*a_2}{|w_2^*a_2|}w_2 \quad z_3 = \frac{|w_3|}{|a_3|}a_3$$
$$M^*z = \beta w$$
$$b_1 = qa_1 \quad b_2 = \frac{a_2^*w_2}{|a_2^*w_2|}a_2 \quad b_3 = \frac{|a_3|}{|w_3|}w_3$$

$$\tag{35}$$

for some real scalar $q \in [-1\ 1]$ with

$$\text{Re}\big(a_1^*w_1\big) \geq 0 \quad \text{for } q = 1$$
$$\text{Re}\big(a_1^*w_1\big) \leq 0 \quad \text{for } q = -1 \tag{36}$$
$$\text{Re}\big(a_1^*w_1\big) = 0 \quad \text{for } |q| < 1.$$

Proof: Apply Lemmas 10–12 to the appropriate block components. $\square$

Remarks: Since (35) and (36) are unaffected if we multiply b and a by an arbitrary positive real scalar α, and z and w by an arbitrary positive real scalar γ, then in searching for solutions to these equations we may impose the additional restriction $|a| = |w| = 1$.

Any solution to (35) and (36) immediately gives us a decomposition as in (25), and hence β is a lower bound for $\mu_\mathcal{K}(M)$. We also note that under certain technical assumptions (as given), there always exists a solution to these equations with $\beta = \mu_\mathcal{K}(M)$. Since we would like to find the largest β we can that solves (35) and (36), we now propose finding a solution to this system of equations via the following power iteration:

$$\tilde{\beta}_{k+1}a_{k+1} = Mb_k$$
$$z_{1_{k+1}} = \tilde{q}_{k+1}w_{1_k} \quad z_{2_{k+1}} = \frac{w_{2_k}^*a_{2_{k+1}}}{|w_{2_k}^*a_{2_{k+1}}|}w_{2_k}$$
$$z_{3_{k+1}} = \frac{|w_{3_k}|}{|a_{3_{k+1}}|}a_{3_{k+1}}$$
$$\hat{\beta}_{k+1}w_{k+1} = M^*z_{k+1} \tag{37}$$
$$b_{1_{k+1}} = \hat{q}_{k+1}a_{1_{k+1}} \quad b_{2_{k+1}} = \frac{a_{2_{k+1}}^*w_{2_{k+1}}}{|a_{2_{k+1}}^*w_{2_{k+1}}|}a_{2_{k+1}}$$
$$b_{3_{k+1}} = \frac{|a_{3_{k+1}}|}{|w_{3_{k+1}}|}w_{3_{k+1}}$$

where $\tilde{q}_{k+1}$ and $\hat{q}_{k+1}$ evolve as

$$\tilde{\alpha}_{k+1} = \text{sgn}(\hat{q}_k)\frac{|b_{1_k}|}{|a_{1_{k+1}}|} + \text{Re}\big(a_{1_{k+1}}^*w_{1_k}\big)$$

If $|\tilde{\alpha}_{k+1}| \geq 1$ Then $\tilde{q}_{k+1} = \frac{\tilde{\alpha}_{k+1}}{|\tilde{\alpha}_{k+1}|}$ Else $\tilde{q}_{k+1} = \tilde{\alpha}_{k+1}$

$$\hat{\alpha}_{k+1} = \text{sgn}(\tilde{q}_{k+1})\frac{|b_{1_k}|}{|a_{1_{k+1}}|} + \text{Re}\big(a_{1_{k+1}}^*w_{1_{k+1}}\big)$$

If $|\hat{\alpha}_{k+1}| \geq 1$ Then $\hat{q}_{k+1} = \frac{\hat{\alpha}_{k+1}}{|\hat{\alpha}_{k+1}|}$ Else $\hat{q}_{k+1} = \hat{\alpha}_{k+1}$ $\tag{38}$

and $\tilde{\beta}_{k+1}, \hat{\beta}_{k+1}$ are chosen positive real so that $|a_{k+1}| = |w_{k+1}| = 1$.

It is now straightforward to verify that if the algorithm converges to some equilibrium point, then we satisfy the appropriate constraints on each block component; hence by Lemmas 10–12 we have nonzero vectors $b, a, z, w \in \mathcal{C}^n$, matrices $Q \in \mathcal{Q}_{\text{sub}}, D \in \mathcal{D}_{\text{sub}}$, and positive real scalars $\tilde{\beta}, \hat{\beta}$ such that

$$Mb = \tilde{\beta}a \qquad M^*z = \hat{\beta}w$$
$$b = Qa \qquad b = D^{-1}w \tag{39}$$
$$z = Q^*QDa \qquad z = Q^*w.$$

Thus if $\tilde{\beta} = \hat{\beta}$, then we satisfy (27) and so have a decomposition as in (25); hence $\tilde{\beta}$ is a lower bound for $\mu_\mathcal{K}(M)$ [associated with a local maximum of $\rho_R(QM)$].

We note that if $\tilde{\beta} \neq \hat{\beta}$, then we have not found a decomposition as in (25). However, from (39) we find that $QMb = \tilde{\beta}b$ and $w^*QM = \hat{\beta}w^*$. Thus we have that both $\tilde{\beta}$ and $\hat{\beta}$ are real positive eigenvalues of QM, and so by Lemma 3, $\max(\tilde{\beta}, \hat{\beta})$ still gives us a lower bound for $\mu_\mathcal{K}(M)$.

VII. CONCLUDING REMARKS

The algorithm described here has been implemented in software and is commercially available as part of the μ-Tools toolbox [3]. We now have a good deal of numerical experience with the algorithm on benchmark problems, and in addition the code has been used for a number of real engineering applications which are detailed elsewhere

in the literature (e.g., see [8]). We have found that the algorithm typically performs very well in terms of convergence, accuracy of the resulting bound, and required computation. Space constraints preclude our including this material here, but we refer the interested reader to [9] and the references therein for a detailed numerical study of the algorithm performance as well as [10] for recent efforts at further enhancing the performance.

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Correction to "Stability Conditions for Multiclass Fluid Queueing Networks"

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In the above-mentioned paper,[1] the recommending Associate Editor line in the first paragraph of the footnote on p. 1618 should have read: "Recommended by Associate Editor, W.-B. Gong."

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