

# **Programming Applications and Programming Languages** M30235

TB1&2

University of Portsmouth
BSc Computer Science
2nd Year

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# **Lecture - Introduction**

14:00 22/01/24 Jiacheng Tan

• No content from TB1 is assessed as part of TB2

# Lecture - Introduction to Programming Languages

14:00 22/01/24 Jiacheng Tan

Since there are many different types of application, there are also many types of programming language. The main programming domains are as follows

- Scientific (e.g. ForTran)
- Business (e.g. COBOL)
- AI (e.g. LISP)
- Systems Programming (e.g. C, C++)
- Web Software (e.g. HTML, JavaScript)

### **Language Categories**

There are several ways to categorise programming languages, such as by uses, paradigms, abstraction level, etc

#### **Machine Languages**

- Machine languages directly run on the hardware, using the instruction set of the processor
- Machine code is usually written in hexadecimal as this is a more efficient way of displaying the binary which represent the instructions
- It is very hard for programmers to directly write machine code, as it is not easy to remember instructions and it lacks features such as jump targets, subroutines, etc

### **Assembly Languages**

- A slight abstraction over machine languages
- Each instruction is replaced with an alphanumeric symbol which is easier for programmers to remember and understand
- They also include features such as subroutines, jump targets, etc which make it much easier to create complex programs

#### **System Programming Languages**

- More abstracted from machine languages, but you are still concerned with low-level functions such as memory management
- Used to create operating systems, and for embedded applications where low system requirements do not allow the use of high-level languages

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### **High Level Languages**

- Languages that are machine-independent (are not written directly in machine code, and are therefore portable between CPU architectures)
- Need to be compiled or otherwise translated from text to machine code before they can be run

### **Scripting Languages**

- Used to create programs which perform a single, simple task
- · These are used for system administration
- Usually interpreted languages
- More akin to pseudocode than other programming languages

### **Domain-Specific Languages**

- Some languages are designed to perform a specific task much more efficiently
- The specific purpose could be just about anything, but are specific to that task and either cannot be used otherwise or are not well suited for it

### **Programming Paradigms**

There are several different paradigms which are used in programming

- Procedural
  - Most programming languages are procedural
  - A program is made up of one or more routines which are run in a specific order
- Functional
  - Applies mathematical functions to inputs to get a result
  - Useful for data processing applications such as data analysis or big data
- Logical

There are also two major types of programming languages, which are designed for different purposes

- · Imperative Languages
  - Programs are defined as a sequence of commands for the computer to perform
  - Like a recipe for exactly how to get the desired output
- Declarative Languages
  - Programs describe the desired results without actually specifying how the program should complete the task
  - Functional and logical programming languages are examples of this

# **Lecture - Implementation and Compilation**

14:00 02/02/24 Jiacheng Tan

There are 3 main methods of implementing a language:

- Compilation Programs are translated into machine language, either before (Compilation) or during (JIT) execution
- Pure Interpretation Programs are interpreted by another program, known as an interpreter
- Hybrid Interpretation A compromise between the two, code is compiled into an intermediary language, which is then interpreted with a Language Virtual Machine

### **Compilation**

- High-level code is translated into machine code for a specific platform
- This results in slow translation, but much faster execution
- The compilation process has multiple stages
  - Lexical Analysis Converts characters in the source into lexical units
  - Syntax Analysis Transforms lexical units into parse trees which represent the syntactic structure of the program
  - Semantics Analysis Generate intermediary code
  - Code Generation Intermediary code is translated into platform-specific machine code
- The program which completes this process is known as the Compiler
- During this process, the compiler uses a "Symbol Table", which each stage interacts with

### **Lexical Analysis**

The scanner reads the source code one character at a time and returns a sequence of tokens which are sent to the next phase. Tokens are symbolic names for elements of the source language. An example of a token in C++ is the keyword 'void', which is a type definition, another example is ';', which delimits the end of a statement. Each token is also stored in the symbol table, along with it's attributes.

#### **Symbol Table**

The symbol table stores all of the identifiers of a source program, along with their attributes. These attributes include information such as the type of a variable, the size or length of a string or array, the arguments to be used with a function and the types of each argument, etc.

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### **Syntax Analysis or Parsing**

The parser analysis the structure of the source code. The parser takes the output of the lexical analyser as a sequence of tokens. It attempts to apply the syntactic rules (or grammar) of the language to the sequence of tokens. The parser uses the language's grammar to derive a parse tree for each statement. Parsers usually construct Abstract Syntax Trees (ASTs), which are slightly simpler and easier to represent with a computer, but which still represent the same syntax. If the syntax tree is invalid for the language's grammar, a syntax error is generated and the compilation process stops

### **Semantic Analysis**

The semantic analyser catches any other issues that are still valid syntax. For example, if you attempt to add a string to a float, it could still be syntactically correct, but semantically makes no sense and is not possible to compute. It is also able to find issues with the variable types of function arguments, such as attempting to use a string in the place of an integer or float.

### **Code Generation and Optimisation**

The code optimiser attempts to improve the time and space efficiency of the program. It can do this in several ways, such as simplifying constants (e.g. replacing 10 \* 10 with 100), removing unreachable code, optimising the flow of code, etc.

The final task of the compiler is to generate the final output code. This could be in the form of platform-specific machine code, or intermediary code for use with a virtual machine. This stage also deals with scheduling and assigning registers for use during execution

### **Pure Interpretation**

- High level code is directly executed by another program known as the interpreter
- There is no syntax or semantics analysis, and there is no optimisation
- Only really suitable for small, non-real-time applications
- It also often requires more space as it needs to store the symbol table during execution
- · Very few modern languages use interpreters, other than Python, JavaScript and PHP

## **Hybrid Interpretation**

- A compromise between compilers and pure interpreters
- High-level languages are translated or compiled into an intermediary language, using the same compilation steps as before
- The intermediary code is then run by a platform-specific virtual machine, which interprets the code into machine language

#### **Just-in-Time**

- Programs are initially translated into an intermediary language
- This is then loaded into memory and segments of the program are then translated into machine code just before execution
- The machine code is then kept in case the function is called again somewhere else in the program
- This drastically improves the execution speed as compared to pure interpretation, but is still slower and typically less space and memory efficient than a compiled program