ECE 385

Fall 2021

Experiment # 6

Simple Computer SLC-3.2 in SystemVerilog

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I. Introduction

In this experiment, we overall design a simple microprocessor using SystemVerilog. It consists of the LC-3 ISA, a 16-bit processor with 16-bit Program Counter, 16-bit instructions, and 16-bit registers.

More specifically, in week 1,we implement the FETCH instruction both in simulation and on the FPGA, which includes loading PC into MAR and loading MDR into IR. And in week 2, we implement the rest of the SLC3, including the DECODE instruction, which will check for the BEN bit and move to other parts of the state machine based on the opcode. Besides, we also implement the EXECUTE instruction: ADD/ADDi, AND/ANDi, NOT, BR, JMP, JSR, LDR, STR or PAUSE.

II. Written Description and Diagrams of SLC-3

a. Summary of Operation

LC-3 is an assembly language created by the dynamic duo Yale N Patel and Sanjay J Patel. The SLC-3 that we create in this experiment differs from the original LC-3 for simplicity but also uses FETCH to get instructions, decodes them, and executes. The difference is that the SLC-3 lacks some instructions such as TRAP, STI, and there is no use of R signal. We deal with the lack of our ready signal R by adding more states to wait until the correct time based on the worst possible case. The reason for the lack of the R signal is due to the fact that the SRAM in use does not support it.

ADD Adds the contents of SR1 and SR2 and stores the result to DR. Sets the status register.

ADDi Add Immediate. Adds the contents of SR to the sign-extended value imm5 and stores the result to DR. Sets the status register.

AND ANDs the contents of SR1 with SR2 and stores the result to DR. Sets the status register.

ANDi And Immediate. ANDs the contents of SR with the sign-extended value imm5 and stores the result to DR. Sets the status register.

NOT Negates SR and stores the result to DR. Sets the status register.

BR Branch. If any of the condition codes match the condition stored in the status register, takes the branch; otherwise, continues execution. (An unconditional jump can be specified by setting NZP to 111.) Branch location is determined by adding the sign-extended PCoffset9 to the PC.

JMP Jump. Copies memory address from BaseR to PC.

JSR Jump to Subroutine. Stores current PC to R(7), adds sign-extended PCoffset11 to PC.

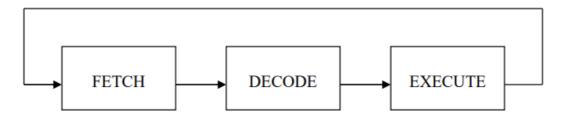
LDR Load using Register offset addressing. Loads DR with memory contents pointed to by (BaseR + SEXT(offset6)). Sets the status register.

STR Store using Register offset addressing. Stores the contents of SR at the memory location pointed to by (BaseR + SEXT(offset6)).

PAUSE Pauses execution until Continue is asserted by the user. Execution should only unpause if Continue is asserted during the current pause instruction; that is, when multiple

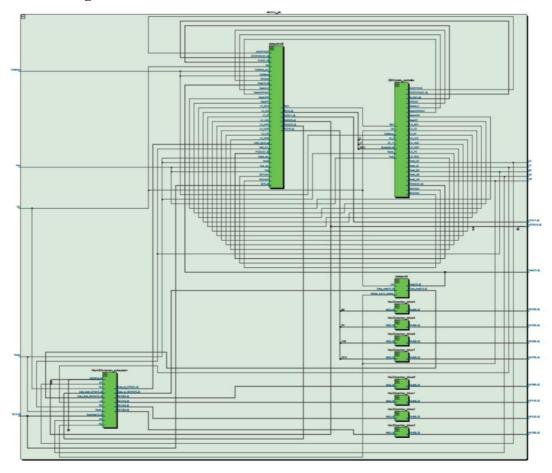
pause instructions are encountered, only one should be "cleared" per press of Continue. While paused, ledVect12 is displayed on the board LEDs. See I/O Specification section for usage notes.

b. Performance of Function



In our design, instructions are FETCHed from memory and stored in the IR, from which we DECODE the opcode and thus go to the specified next state from the OPCODE. The OPCODE is responsible for choosing between the following operations: ADD/ADDi, AND/ANDi, NOT, BR, JMP, JSR, LDR, STR or PAUSE. The EXECUTE cycle is responsible for actually performing the operation, then looping back to initial state. In each state the FSM outputs control signals will get inputted into the various components such as muxes, registers, etc.

c. Block Diagram of slc3.sv



d. Written Description of all .sv modules

1) slc3.sv

```
Immodule slc3(
    input logic [15:0] S,
    input logic Clk, Reset, Run, Continue,
    output logic [11:0] LED,
    output logic [6:0] HEXO, HEX1, HEX2, HEX3, HEX4, HEX5, HEX6, HEX7,
    output logic CE, UB, LB, OE, WE,
    output logic [19:0] ADDR,
    inout wire [15:0] Data //tristate buffers need to be of type wire
);

// Declaration of push button active high signals
logic Reset_ah, Continue_ah, Run_ah;

//assign Reset_ah = ~Reset;
//assign Continue_ah = ~Continue;
//assign Run_ah = ~Run;
sync reset_sync (Clk, ~Reset, Reset_ah);
sync continue_sync (Clk, ~Continue, Continue_ah);
sync run_sync (Clk, ~Run, Run_ah);

// Internal connections
logic BEN;
logic BEN;
logic DEMAR, LD_MDR, LD_IR, LD_BEN, LD_CC, LD_REG, LD_PC, LD_LED;
logic GatePC, GateMDR, GateALU, GateMARMUX;
logic [1:0] PCMUX, ADDR2MUX, ALUK;
logic GatePC, GateMDR, GateALU, GateMARMUX;
logic [15:0] DRML, SRIMUX, SR2MUX, ADDRIMUX;
logic [15:0] MDR_In;
logic [15:0] MDR_In;
logic [15:0] MDR_In;
logic [15:0] Data_from_SRAM, Data_to_SRAM;

// Signals being displayed on hex display
logic [3:0][3:0] hex_4;
```

Input: [15:0] S, Clk, Reset, Run, Continue

Output: [11:0] LED, [6:0] HEX0, [6:0] HEX1, [6:0] HEX2, [6:0] HEX3, [6:0] HEX4, [6:0] HEX5, [6:0] HEX6, [6:0] HEX7, CE, UB, LB, OE, WE, [19:0] ADDR

Inouts: [15:0] Data

Description: This module is the main part of SLC-3 processor, in which contain all the sub modules that included in the processor, including datapath, Mem2IO, ISDU, etc. It will receive the input keys and switches to control the running of processor, and the output will be connected to the outship memory with data transmission.

Purpose: Connect the modules together and route the data and signals.

2) AlU.sv

```
Impodule ALU(
    input logic [15:0] SR1OUT, SR2OUT,IMME,
    input logic sr2mux,
    input logic [1:0] ALUK,
    output logic [15:0] OUT
);

logic [15:0] B, out_add,out_and,out_not;
    ripple_adder adder(.A(SR1OUT),.B(B),.Sum(out_add),.Co());
    assign B = sr2mux ? IMME : SR2OUT;
    assign out_and = B & SR1OUT;
    assign out_not = ~SR1OUT;
    always_comb
begin

unique case (ALUK)
    3'b00: OUT = out_add;
    3'b01: OUT = out_and;
    3'b10: OUT = sR1OUT;
    endcase
end
endmodule
```

Inputs: [15:0] SR1OUT, SR2OUT, IMME [1:0] ALUK Sr2mux

Outputs: [15:0] OUT

Description: This module receives the outputs of register file and immediate number, ALUK control signal as input and the output is the result of the ALU calculation module.

Purpose: This module generates the arithmetic unit of SLC-3 processor.

3) ADDRMUX.sv

Inputs: [15:0] IR, data_from_SR1OUT, data_from_PC [1:0] ADDR2MUX ADDR1MUX Outputs: [15:0] Data Out, SEXT [3:0] Data to controller [2:0] Data to BEN

Description: This module receives inputs from ALU, IR, PC, and control signals from state machine. And then it will decode the IR to generate the immediate value with signal extension and data to control unit and BEN module.

Purpose: This module decodes the IR instruction to get values that as the input to Ben module and controller.

4) Ben.sv

```
#module ben_module(
    input logic Clk, reset,
    input logic [15:0] BUS,
    input logic [15:0] BUS,
    input logic [2:0] IR11_9,
    output logic [2:0] NZP;
    logic n_out,z_out,p_out,ben_input;
    uni_reg #(.N(1)) N(.Clk(Clk),.Load(LD_CC),.reset(reset),.D(NZP[1]),.Data_Out(n_out));
    uni_reg #(.N(1)) Z(.Clk(Clk),.Load(LD_CC),.reset(reset),.D(NZP[1]),.Data_Out(z_out));
    uni_reg #(.N(1)) P(.Clk(Clk),.Load(LD_CC),.reset(reset),.D(NZP[0])),.Data_Out(z_out));
    assign ben_input = IR11_9[2]&n_out | IR11_9[1]&z_out | IR11_9[0]&p_out;
    uni_reg #(.N(1)) Ben(.Clk(Clk),.Load(LD_BEN),.reset(reset),.D(ben_input),.Data_Out(BEN));
    always_comb
    begin
    if(BUS == 16'b0)
        NZP = 3'b010;
    else if(BUS[15] == 1'b1)
        NZP = 3'b010;
    else
        NZP = 3'b001;
    end
endmodule
```

Inputs: [15:0] BUS Clk, reset, LD CC, LD BEN[2:0] IR11 9

Outputs: BEN

Description: Based on the value from BUS, this module will have the corresponding conditional code stored in the registers. And this module will generate the corresponding BEN value based on the conditional condes and the control signal and the BEN value will be used for judging whether jump the PC value or not.

Purpose: This module generates Conditional codes and BEN value for BR instruction.

5) datapath.sv

Inputs: Clk, Reset_ah, Run_ah, LD_MAR, LD_MDR, LD_IR, LD_BEN, LD_CC, LD_REG, LD_PC, GatePC, GateMDR, GateALU, GateMARMUX, DRMUX, SR1MUX, SR2MUX, ADDR1MUX, MIO_EN [1:0] PCMUX, ADDR2MUX, ALUK [15:0] MDR_In, S

Outputs: [15:0] MAR, IR, MDR, PC BEN

Description: This is the data path of the whole SLC-3, including the register file, PC-MUX module, instruction decoder, ALU and Conditional Code module. All these components will be connected and be set the control signals through datapath. This module receive all the control signals as input and the output will be the MAR, IR, PC, MDR and BEN.

Purpose: This module connects all the data path components together and set control signals to them.

6) lab6_toplevel.sv

```
module lab6_toplevel( input logic [15:0] S,
                                                       input logic [15:0] S,
input logic Clk, Reset, Run, Continue,
output logic [11:0] LED,
output logic [6:0] HEXO, HEX1, HEX2, HEX3, HEX4, HEX5, HEX6, HEX7,
output logic CE, UB, LB, OE, WE,
output logic [19:0] ADDR,
inout wire [15:0] Data);
slc3 my_slc(.*);
// Even though test memory is instantiated here, it will be synthesized into // a blank module, and will not interfere with the actual SRAM. // Test memory is to play the role of physical SRAM in simulation. test_memory my_test_memory(.Reset(~Reset), .I_O(Data), .A(ADDR), .*);
endmodule
Inputs: [15:0] S Clk, Reset, Run, Continue
```

Inout: [15:0] Data

Outputs: [11:0] LED [6:0] HEX0, HEX1, HEX2, HEX3, HEX4, HEX5, HEX6, HEX7CE, UB, LB,

OE, WE [19:0] ADDR

Descriptions: This is the top level of the SLC-3 processor, including the slc3 module connected with the test memory only for simulation

Purpose: This module works as the top level of lab6

7) Mem2IO.sv

```
( input logic Clk, Reset, input logic [19:0] ADDR, input logic CE, UB, LB, OE, WE, input logic [15:0] Switches, input logic [15:0] Data_from_CPU, Data_from_SRAM, output logic [15:0] Data_to_CPU, Data_to_SRAM, output logic [3:0] HEXO, HEX1, HEX2, HEX3);
Emodule Mem2IO
      logic [15:0] hex_data;
      // Load data from switches when address is xFFFF, and from SRAM otherwise. 
 {\tt always\_comb}
       begin
              Data_to_CPU = 16'd0;
                f (WE && ~OE)
if (ADDR[15:0] == 16'hFFFF)
                    Data_to_CPU = Switches;
                    Data_to_CPU = Data_from_SRAM;
       end
        // Pass data from CPU to SRAM
      assign Data_to_SRAM = Data_from_CPU;
      // Write to LEDs when WE is active and address is xFFFF.
always_ff @ (posedge Clk) begin
  if (Reset)
          hex_data <= 16'd0;
else if ( ~WE & (ADDR[15:0] == 16'hFFFF) )
               hex_data <= Data_from_CPU;
      assign HEX0 = hex_data[3:0];
      assign HEX1 = hex_data[7:4];
assign HEX2 = hex_data[11:8];
      assign HEX3 = hex_data[15:12];
endmodule
```

Input: Clk Reset [19:0] ADDR CE UB LB OE WE [15:0] Switches [15:0] Data from CPU [15:0] Data from SRAM

Output: [15:0] Data to CPU, [15:0] Data to SRAM, [3:0] HEX0, HEX1, HEX2, HEX3

Description: This module uses the data from MAR and input the address which is used for READ and WRITE memory, aided by the signals WE and OE in the ISDU.

Purpose: This module serves as the link between the datapath module and memory contents

module and enables the modification of memory outside of registers.

8)ISDU.sv

```
⊟module ISDU (
                      input logic
                                                clk,
                                     Reset,
                                     Run,
Continue,
                 input logic[3:0]
input logic
input logic
input logic
                                           opcode,
                                           IR_5,
IR_11,
                                           BEN,
                 output logic
                                           LD_MAR,
                                     LD_MDR,
                                     LD_IR,
LD_BEN,
                                     LD_CC,
LD_REG,
                                     LD_LED, // for PAUSE instruction
                 output logic
                                           GatePC,
                                     GateMDR,
                                     GateALU,
                                     GateMARMUX,
                 output logic [1:0] PCMUX, output logic DRMUX,
                                     SR1MUX,
                                     SR2MUX
                                     ADDR1MUX
                 output logic [1:0] ADDR2MUX, ALUK,
                 output logic
                                           Mem_CE,
                                     Mem_UB,
                                     Mem_LB,
                                     Mem OE
      enum logic [4:0] { Halted,
PauseIR1,
                         PauseIR2,
S_18,
                         s_33
                         S_33_2,
S_35,
                         s 32.
                         S_01,
                          s_05,
                         s 09
                         s_06,
                         S_25_1
S_25_2
                         S_27,
S_07,
                         s_16
                         S_21,
S_12,
                         S_0,
S_22}
                                    State, Next_state; // Internal state logic
      always_ff @ (posedge Clk)
(Reset)
              State <= Halted;
             State <= Next_state;
      always_comb
      begin
  // Default next state is staying at current state
  Next_state = State;
```

Inputs: [3:0] Opcode Clk, Reset, Run, Continue, IR_5, IR_11, BEN

Outputs: LD_MAR, LD_MDR, LD_IR, LD_BEN, LD_CC, LD_REG, LD_PC, LD_LED, GatePC,
GateMDR, GateALU, GateMARMUX, DRMUX, SR1MUX, SR2MUX, ADDR1MUX, ALUK,
Mem_CE, Mem_UB, Mem_LB, Mem_OE, Mem_WE [1:0] PCMUX, ADDR2MUX

Description: This module defines our state machine (next state and gate variables for current state),

load signals and mux select signals, etc.

Purpose: This module regulates the states that are specified by the given lc3 state diagram and assign proper values to the gate and mux select signals to make the system function properly, and the data can be stored or loaded in the desired way.

9) tristate.sv

Inputs: Clk, tristate output enable [15:0] Data write,

Inout: [15:0] Data

Outputs: [15:0] Data read

Description: This module is the tristate buffer between Memory IO module and the SRAM. With the control signal from control unit, data can flow along bus. Also it is connected to the SRAM with a bi-direction data.

Purpose: This module works as the tristate buffer that connect the SRAM and the memory io module.

10) synchronizers.sv

```
module sync (
    input logic Clk, d,
    output logic q
);
always_ff @ (posedge Clk)
begin
    q <= d;
end
endmodule</pre>
```

Inputs: Clk, d

Output: q

Description: This module serves as the synchronizer for the input from keys or switches. It is a register to synchronize the input signal and help avoid the potential problems caused by the synchronous signals of inputs

Purpose: This module will synchronize the asynchronous signals to the modules in FPGA.

```
11) reg file.sv
```

Inputs: input logic LD_REG, Reset, Clk, input logic [15:0] in, input logic [2:0] SR1, SR2, DR,

Outputs: output logic [15:0] SR1 OUT, SR2 OUT)

Description:

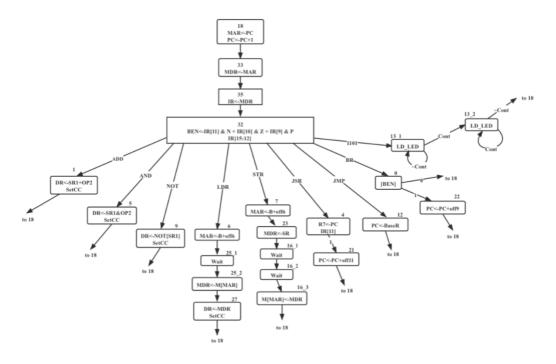
The register file is a unit with eight 16-bit registers that with input logic DR, which will choose which register to load and input logic LD_REG for when to load. SR1, SR2 select which registers to output.

Purpose: For operations in the SLC-3 we need to be able to store "answers" into 8 data registers

e. Description of the Operation of the ISDU

The ISDU module includes signal assignments for all the valid states, controlling many components in the slc3 system with the enable bits that are predefined in each state. The module ISDU. In addition, in this module we define the state transitions for states using various toggle bits to ensure the functionality of the components and the operations of states.

f. State Diagram of ISDU



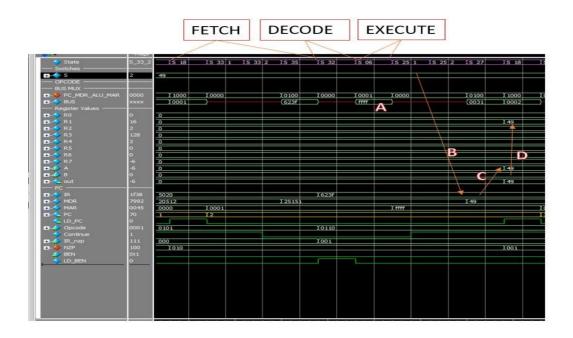
III. Simulation

a. opADD(R1, R1, R1)



At A, we can easily see that 8 from R1 goes into port B of the ALU (from SR2MUX output). Similarly, at **B we can see that 8 from** R1 goes into port A of ALU (from SR1 OUT). At C, we see that 16 is stored back into R1.

b. opLDR(R1, R0, inSW)



At A, we can easily see that xffff in the bus indicates that we will be reading from switches. At B, the value from the switches is loaded into MDR. At C this goes into the ALU (ADD), so 49 + 0 (R0 value) = 49, and at D the output goes into R1

c. opJMP(R1)



At A, we can easily see the value at R1 is 49. When the LD_PC signal goes high right before we see B, PC gets updated the next clock cycle to 49.

IV. Post-Lab Questions

LUT	603
DSP	0
Memory (BRAM)	0
Flip-flop	283
Frequency	85.61
Static Power	98.64mW
Dynamic Power	$0.00 \mathrm{mW}$
Total Power	172.29mW

In week 1, we didn't change the code in testbench, thus our code cannot be compiled. Also, there were too many ports that need to be wired to datapath, we took a lot of time to make it work. It's useful to look at the diagram of LC-3 to do our work.

In week 2, we forgot to use BEN bit in a BR execution, so it always branching. We find this error according to the output result and traced back, that would be easier to find the error.

Q2.1

MEM2IO is used to interface with the SDRAM and the switches, only if the address is xFFFF, it will read from the switches, otherwise, it will read from SDRAM. And we can do reading and writing from SDRAM. MEM2IO interfaces with tristate buffers so data is either from SRAM to CPU or CPU to SRAM, that is, SRAM is bidirectional.

Q2.2

BR is the condition based on NZP bit from the instruction register. JMP is unconditional since it copies memory address from BaseR to PC. The differences is that, JMP can branch the instruction to any indicated memory position while BR can only branch the instruction located at surrounding positions for the current PC.

Q2.3

The R signal in Patt and Patel is a ready signal, which indicate whether LC3 has finished reading and writing from memory. LC3 will wait for R signal to execute the next instruction. However, we don't need R signal in SRAM, so we add the extra states in ISDU, there will be enough time in these states so that the read and write will definitely finish. In LC3, thw R signal is generated form FSM so it's asynchronized, in real situation, R imply that there is an unintended behaviour.

V. Conclusion

Luckily, our design functions well on both analog and FPGA boards. We pass the most difficult XOR tests and SORT tests which is used to determine if all of our operations worked.

We want to thank the lab manual, which is very helpful in explaining the base pipeline about how to design the SLC3 using the SystemVerilog. However, according to our memory, we think that the lab manual doesn't talk much about the memory organization. We use the provided memory code to get familiar with memory organization. We hope it can be specified with more information in the future to help students avoid misunderstanding about this part.