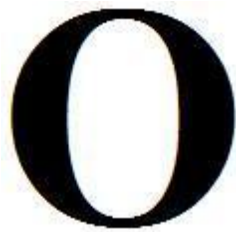


The oark book: <http://code.google.com/p/oark/>



## The Open Source Anti Rootkit

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## Chapter 1

# Call Gates and GDT/LDT

## 1.1 GDT/LDT

The Global Descriptor Table or GDT (extracted from Wikipedia<sup>1</sup>) is a data structure used by Intel x86-family processors in order to define the characteristics of the various memory areas used during program execution, for example the base address, the size and access privileges like executability and writability. These memory areas are called segments in Intel terminology.<sup>2</sup>

The GDT can hold things other than segment descriptors as well. Every 8-byte entry in the GDT is a descriptor, but these can be Task State Segment (or TSS) descriptors, Local Descriptor Table (LDT) descriptors, or Call Gate descriptors. The last one, Call Gates, are particularly important for transferring control between x86 privilege levels although this mechanism is not used on most modern operating systems.

There is also an LDT or Local Descriptor Table. The LDT is supposed to contain memory segments which are private to a specific program, while the GDT is supposed to contain global segments. The x86 processors contain facilities for automatically switching the current LDT on specific machine events, but no facilities for automatically switching the GDT.

Every memory access which a program can perform always goes through a segment. On the 386 processor and later, because of 32-bit segment offsets and limits, it is possible to make segments cover the entire addressable memory, which makes segment-relative addressing transparent to the user.

In order to reference a segment, a program must use its index inside the GDT or the LDT. Such an index is called a segment selector or selector in short. The selector must generally be loaded into a segment register to be used. Apart from the machine instructions which allow one to set/get the position of the GDT (and of the Interrupt Descriptor Table) in memory, every machine instruction referencing memory has an implicit Segment Register, occasionally two. Most of the time this Segment Register can be overridden by adding a Segment Prefix before the instruction.

Loading a selector into a segment register automatically reads the GDT or the LDT and stores the properties of the segment inside the processor itself. Subsequent modifications to the GDT or LDT will not be effective unless the segment register is reloaded.

## 1.2 Call Gates

A Call Gate [Inta] is a mechanism in the Intel x86 architecture to change privilege levels of the CPU when running a predefined function that is called by the instruction CALL/JMP FAR.

A call to a Call Gate allows you to obtain higher privileges than the current, for example we can execute a routine in ring0 using a CALL FAR in ring3. A Call Gate is an entry in

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<sup>1</sup>Wikipedia. “Global Descriptor Table”. In: (2010). URL: [http://en.wikipedia.org/wiki/Global\\_Descriptor\\_Table](http://en.wikipedia.org/wiki/Global_Descriptor_Table).

<sup>2</sup>Intel 64 and IA-32 Architectures, Software Developers Manual, System Programming Guide, Part 1. URL: <http://www.intel.com/design/processor/manuals/253668.pdf>.

the GDT (Global Descriptor Table) or LDT (Local Descriptor Table). There are a GDT for each CORE, and each GDT can have one or more LDTs<sup>3</sup>.

Windows doesn't use Call Gate for anything special, but there are malware, as the worm Gurong.A<sup>4</sup>, that installs a Call Gate via DevicePhysicalMemory to execute code on ring0. An article that talks about it is "Playing with Windows/dev/(k)mem"<sup>5</sup>.

Nowadays we can't easily access to /Device/PhysicalMemory, I recommend reading the presentation by Alex Ionescu at RECON 2006 "Subverting Windows 2003 SP1 Kernel Integrity Protection"<sup>6</sup>. Also, there are examples<sup>7</sup> in the wild that use the API ZwSystemDebugControl<sup>8</sup> to install a Call Gate, but Ionescu's article says that it doesn't work nowadays (although there are techniques to reactivate them).

Gynael and j00ru made a Call-Gate mechanism in kernel/driver exploit development on Windows<sup>9</sup>, or, to be more precise, to use a write-what-where condition to convert a custom LDT entry into a Call-Gate (this can be done by modifying just one byte), and using the Call-Gate to elevate the code privilege from user-land to ring0.

David Reguera Garcia aka Dreg made a Call Gate detector for the free anti-rootkit Rootkit Unhooker<sup>10</sup>. The next releases have new features to detect other new stuff like the 'new LDT Forward to user mode attack' (published also in the Gynael and j00ru's paper.)

An entry in the GDT/LDT looks like this:

---

```
typedef struct _SEG_DESCRIPTOR
{
    WORD size_00_15;
    WORD baseAddress_00_15;
    WORD baseAddress_16_23:8;
    WORD type:4;
    WORD sFlag:1;
    WORD dpl:2;
    WORD pFlag:1;
    WORD size_16_19:4;
    WORD notUsed:1;
    WORD lFlag:1;
    WORD DB:1;
    WORD gFlag:1;
}
```

---

<sup>3</sup>David Reguera Garcia aka Dreg. "Rootkit Arsenal, Installing a Call Gate". In: (2010). URL: <http://www.rootkit.com/blog.php?newsid=992>.

<sup>4</sup>F-secure. "W32/Gurong.A: A worm in e-mails and in Kazaa shared folders. It has a rootkit functionality. This worm appeared on the 21st of March 2006." In: (2006). URL: [http://www.f-secure.com/v-descs/gurong\\_a.shtml](http://www.f-secure.com/v-descs/gurong_a.shtml).

<sup>5</sup>crazylord. "Playing with Windows /dev/(k)mem". In: (2002). URL: <http://www.phrack.com/issues.html?issue=59&id=16>.

<sup>6</sup>Alex Ionescu. "RECON 2006: Subverting Windows 2003 SP1 Kernel Integrity Protection". In: (2006). URL: <http://www.alex-ionescu.com/recon2k6.pdf>.

<sup>7</sup>SACCOPHARYNX. "Call Gates". In: (2006). URL: <http://ricardonarvaja.info/WEB/OTROS/TUTES%20SACCOPHARYNX/#Ring0>.

<sup>8</sup>Undocumented NT Internals. *Function NtSystemDebugControl is used by some low-level debuggers written by Microsoft and available typically in DDK*. URL: <http://undocumented.ntinternals.net/UserMode/Undocumented%20Functions/Debug/NtSystemDebugControl.html>.

<sup>9</sup>Matthew j00ru Jurczyk and Gynael Coldwind. "GDT and LDT in Windows kernel vulnerability exploitation". In: (2010). URL: <http://gynael.coldwind.pl/?id=274>.

<sup>10</sup>DiabloNova aka EP X0FF. "Rootkit Unhooker LE 3.8.386.588 SR1". In: (2010). URL: <http://www.rootkit.com/blog.php?newsid=993>.

```
WORD baseAddress_24_31:8;
} SEG_DESCRIPTOR, *PSEG_DESCRIPTOR;
```

Listing 1.1: GDT/LDT Descriptor structure

A Call Gate is an entry type in the GDT/LDT which has the following appearance:

```
typedef struct _CALL_GATE_DESCRIPTOR
{
    WORD offset_00_15;
    WORD selector;
    WORD argCount:5;
    WORD zeroes:3;
    WORD type:4;
    WORD sFlag:1;
    WORD dpl:2;
    WORD pFlag:1;
    WORD offset_16_31;
} CALL_GATE_DESCRIPTOR, *PCALL_GATE_DESCRIPTOR;
```

Listing 1.2: Call Gate Descriptor structure

- **offset\_00\_15:** is the bottom of the address of the routine to be executed in ring0, **offset\_16\_31** is the top.
- **selector:** specifies the code segment with the value `KGDT_R0_CODE (0x8)`, the routine will run ring0 privileges.
- **argCount:** the number of arguments of the routine in DWORDs.
- **type:** the descriptor type for a 32-bit Call Gate needs the value `0xC`
- **dpl:** minimum privileges that the code must have to call the routine, in this case `0x3`, because it will be called by the routine ring3

To create a Call Gate we can follow the following steps:

1. Build the Call Gate that points to our routine.
2. Set the code only in a core (remember: there are a GDT for each CORE).
3. Read the GDTR register in order to find the GDT address and the size using SGDT instruction:

```
typedef struct _GDTR
{
    WORD nBytes;
    DWORD baseAddress;
} GDTR;
```

Listing 1.3: GDTR register

We can obtain the number of entries (number of GDT descriptors) with `GDTR.nBytes/8`.

4. Find a free entry in the GDT/LDT.

5. Write the Call Gate descriptor.
6. To call the Call Gate is only necessary to make a CALL/JMP FAR to the GDT/LDT selector:
  - ie if we've introduced the Call Gate at the entry **100** of the **GDT**, the user space application must execute a CALL/JMP FAR **0x320:00000000**. 0x320 is in binary 1100100 **0** 00, then the entry is:1100100 (100 in binary is 1100100), **TI=0** (entry is in **GDT**) RPL=00.. The other part of the FAR CALL is not useful but must be in the instruction.
  - ie if we've introduced the Call Gate at the entry **100** of the **LDT**, the user space application must execute a CALL/JMP FAR **0x324:00000000**. 0x324 is in binary 1100100 **1** 00, then the entry is:1100100 (100 in binary is 1100100), **TI=1** (entry is in **LDT**) RPL=00..

## Chapter 2

# PEB Hooking

# Chapter 3

## Introduction

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Pasting code...

### 3.1 Introduction Section

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charge,Permission is hereby granted, free of charge,Permission is hereby granted, free of charge,Permission  
bbbbbb

#### 3.1.1 sub Introduction

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##### sub sub Introduction

a:<sup>1</sup>, This subsection's content..

This subsection's content..This subsection's content..This sub content..This subsection's content..This subsection's content..This subsection's content..This subsection's keyword1  
ggggggggrrrrrr

---

<sup>1</sup>usserman.



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