

# Dynamic Programming

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# Order Isomorphism

# Order Isomorphism

## Definition 1

A surjective map  $F$  from poset  $(V, \preceq)$  to poset  $(\hat{V}, \leq)$  is called an

- **Order isomorphism** if  $v \preceq w \iff Fv \leq Fw$
- **Order anti-isomorphism** if  $v \preceq w \iff Fw \leq Fv$

**Comment:**  $F$  under this definition is bijective.

### Exercise 3.1.1.

Given  $h \in \mathbb{R}^X$ , let  $Fh = \exp(\theta h)$ . Show that  $F$  is an order isomorphism from  $\mathbb{R}^X$  to  $(0, \infty)^X$  whenever  $\theta > 0$ .

Proof.

Fix  $\theta > 0$ . We know that  $\exp(\theta x)$  is a bijective function. Let  $h_1, h_2 \in \mathbb{R}^X$  such that  $h_1 \leq h_2$ . This implies

$$\theta h_1 \leq \theta h_2$$

As  $\exp(\cdot)$  is order preserving, hence, we have

$$Fh_1 = \exp(\theta h_1) \leq \exp(\theta h_2) = Fh_2$$

Let  $k_1, k_2 \in (0, \infty)^X$ ,  $k_1 \leq k_2$  by surjectivity, we have

$$k_1 = F(q_1) = \exp(\theta q_1), k_2 = F(q_2) = \exp(\theta q_2), q_1, q_2 \in \mathbb{R}^X$$

□

## Exercise 3.1.1. Continue

Proof.

We have

$$q_1 = \frac{\ln k_1}{\theta} \leq \frac{\ln k_2}{\theta} = q_2$$

as  $\ln$  is order preserving. Therefore, by definition,  $F$  is an order isomorphism from  $\mathbb{R}^X$  to  $(0, \infty)^X$  □

## Exercise 3.1.2.

Let  $V = M^X$  and  $\hat{V} = \hat{M}^X$ ,  $M, \hat{M} \subset \mathbb{R}$ . Let  $\varphi$  be a map from  $M$  onto  $\hat{M}$  and let  $Fv = \varphi \circ v$ . Prove if  $\varphi$  is an order isomorphism from  $M$  to  $\hat{M}$ , then  $F$  is an order isomorphism from  $V$  to  $\hat{V}$ .

Proof.

$\varphi$  is order isomorphism then  $\varphi$  is bijective, order preserving with order preserving inverse. Hence apply this  $\dim X$  times, we get  $F$  is bijective, order preserving with order preserving inverse. □

### Exercise 3.1.3

Let  $V, \hat{V}$  be posets. Show that every order isomorphism  $F$  is a bijection. Show that every order anti-isomorphism is also a bijection.

Proof.

Let  $v_1, v_2 \in \hat{V}$  such that  $v_1 = v_2$ . By surjectivity, we have

$$v_1 = F(w_1), \quad v_2 = F(w_2), \quad w_1, w_2 \in V$$

Hence, we have

$$F(w_1) \leq F(w_2) \implies w_1 \preceq w_2$$

$$F(w_2) \leq F(w_1) \implies w_2 \preceq w_1$$

Hence,  $w_1 = w_2$ . This proves that  $F$  is injective. □



## Exercise 3.1.4.

Let  $F$  be a bijection from  $(V, \preceq)$  to  $(\hat{V}, \leq)$ . Show that

- ①  $F$  is an order isomorphism if and only if  $F$  and  $F^{-1}$  are order preserving
- ②  $F$  is an order anti-isomorphism if and only if  $F$  and  $F^{-1}$  are order reversing.

Proof.

Skip



### Lemma 3.1.1.

Let  $F$  be an order isomorphism from  $(V, \preceq)$  to  $(\hat{V}, \leq)$ . If the supremum of  $\{v_\alpha\}_{\alpha \in \Lambda} \subset V$  exists in  $V$ , then

$$\bigvee_{\alpha} Fv_{\alpha} \text{ exists in } \hat{V} \text{ and } \bigvee_{\alpha} Fv_{\alpha} = F \bigvee_{\alpha} v_{\alpha}$$

Proof.

Let  $v := \bigvee_{\alpha} v_{\alpha} \in V$ . Let  $\hat{w}$  be any upper bound of  $\{Fv_{\alpha}\}$ , i.e.,  $Fv_{\alpha} \leq \hat{w}$  for all  $\alpha \in \Lambda$ . By surjectivity, we let  $\hat{w} = F(w)$ , and by order isomorphism, we have

$$v_{\alpha} \preceq w \quad \text{for all } \alpha \in \Lambda$$

Hence,  $w$  is an upper bound of  $\{v_{\alpha}\}$ , this implies  $v \preceq w$ . Hence,

$$F(v) \leq F(w) = \hat{w}$$

This implies  $F(v) = F \bigvee_{\alpha} v_{\alpha}$  is the least upper bound of  $\{Fv_{\alpha}\}$ . □

## Exercise 3.1.6

Let  $V, \hat{V}$  be posets and let  $(v_n)$  be a sequence in  $V$ . And let  $F$  be a map from  $V$  to  $\hat{V}$ . Prove the following

- 1 If  $F$  is an order isomorphism, then  $v_n \uparrow v$  if and only if  $Fv_n \uparrow Fv$  in  $\hat{V}$
- 2 If  $F$  is an order anti-isomorphism, then  $v_n \uparrow v$  if and only if  $Fv_n \downarrow Fv$  in  $\hat{V}$ .

Proof.

$v_n \uparrow v \implies v_1 \leq v_2 \leq \dots \leq v$  and  $v = \bigvee_n v_n$ .

Hence, by order isomorphism, we have

$$Fv_1 \leq Fv_2 \leq \dots \leq Fv$$

and  $Fv = \bigvee_n Fv_n$ . Moreover,  $F$  is order isomorphism implies  $F^{-1}$  is order isomorphism. Hence, the other direction follows. □

## Exercise 3.1.7

Prove

- If  $V, \hat{V}$  are order isomorphic, then  $V$  is totally ordered if and only if  $\hat{V}$  is totally ordered
- $F$  is an order anti-isomorphism from  $V$  to  $\hat{V}$  if and only if  $F$  is an order isomorphism from  $V$  to its dual  $\hat{V}^\partial$

Proof.

Skip.



- We start with the definition of conjugacy between dynamical systems  $((V, T_\sigma))$  is a dynamical system with state space  $V$  and evolution  $T_\sigma$ ).
- Then, we go to the most basic structure,  $V$  as a poset, and upgrade conjugacy to order conjugacy.
- This prepares for the later upgrade from dynamical system to ADP

## Definition 2

We call a **discrete time dynamical system** is a pair  $(V, S)$ , where  $V$  is any set, and  $S$  is a self-map on  $V$ .

## Definition 3

Two dynamical systems  $(V, S)$  and  $(\hat{V}, \hat{S})$  are said to be **conjugate under**  $F$  if

$F$  is a bijection from  $V$  into  $\hat{V}$  and  $F \circ S = \hat{S} \circ F$  on  $V$

or we can write it as

$$S = F^{-1} \circ \hat{S} \circ F$$

## Proposition 3.1.2.

If  $(V, S)$  and  $(\hat{V}, \hat{S})$  are conjugate, then

- ①  $S^n = F^{-1}\hat{S}^n F$  for all  $n \in \mathbb{N}$
- ②  $v$  is a fixed point of  $S$  if and only if  $Fv$  is a fixed point of  $\hat{S}$
- ③  $\hat{v}$  is fixed point of  $\hat{S}$  if and only if  $F^{-1}\hat{v}$  is a fixed point of  $S$
- ④  $v$  is the unique fixed point of  $S$  in  $V$  if and only if  $Fv$  is the unique fixed point of  $\hat{S}$  in  $\hat{V}$ .

### Proof.

Let  $v$  be the unique fixed point of  $S$ , i.e.,  $Sv = v$ . Hence,

$$F(v) = F(Sv) = \hat{S}(Fv)$$

Hence,  $F(v)$  is a fixed point of  $\hat{S}$ . Let  $\hat{w} = F(w)$  be a fixed point  $\hat{S}$ , then by part 2,  $w$  is the fixed point of  $S$ . Hence,  $w = v$ , and this implies  $F(w) = F(v)$ . □

## Example 3.1.2

### Conjugate Dynamical Systems

Example 3.1.2

System 1:  $(\mathbb{R}, S)$

$$S(x) = ax + b$$

Example:  $S(x) = 2x + 1$

Fixed point:  $x^* = -1$

**Unstable**

$$F(x) = e^x$$

System 2:  $((0, \infty), \hat{S})$

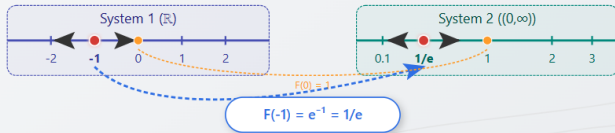
$$\hat{S}(y) = e^b \cdot y^a$$

Example:  $\hat{S}(y) = e \cdot y^2$

Fixed point:  $y^* = 1/e \approx 0.368$

**Unstable**

### Phase Lines and Dynamics



Conjugacy Verification:  $F \circ S = \hat{S} \circ F$

$$F(S(x)) = e^{2x+1} = e^1 \cdot e^{2x} = e \cdot (e^x)^2 = \hat{S}(F(x))$$



## Definition 4

Consider two dynamical systems  $(V, S)$  and  $(\hat{V}, \hat{S})$ , where  $V, \hat{V}$  are posets. We call these systems **order conjugate under  $F$**  if they are conjugate under  $F$ , and,  $F$  is an order isomorphism.

### Exercise 3.1.9.

Prove that order conjugacy is an equivalence relation on the set of dynamical systems over partially ordered set.

Proof.

We denote  $(V, S) \sim (\hat{V}, \hat{S})$  if they are order conjugate. We need to show this relation is reflexive, symmetric and transitive.

- (Reflexivity) Let  $F = Id$  which is a bijection. We have

$$F \circ S = S = S \circ F$$

Moreover, we have  $F = F^{-1}$  is order preserving. Hence,

$$(V, S) \sim (V, S)$$



## Exercise 3.1.9 Continue

### Proof.

- (Symmetry) Let  $(V, S) \sim (\hat{V}, \hat{S})$  under  $F$ .
  - $F$  is bijection implies  $F^{-1}$  is bijection
  - $F \circ S = \hat{S} \circ F \implies F^{-1} \circ \hat{S} = S \circ F^{-1}$

Hence,  $(V, S)$  and  $(\hat{V}, \hat{S})$  are conjugate under  $F^{-1}$ .

- $F$  is order preserving with order preserving inverse implies  $F^{-1}$  is order preserving with order preserving inverse

Hence,  $F^{-1}$  is order isomorphism. Hence  $(\hat{V}, \hat{S}) \sim (V, S)$



## Exercise 3.1.9 Continue

### Proof.

- (Transitive) Let  $(V_1, S_1) \sim (V_2, S_2)$  under  $F$  and  $(V_2, S_2) \sim (V_3, S_3)$  under  $G$ .
  - $F, G$  are bijective implies  $H := (G \circ F)$  is bijective
  - $F \circ S_1 = S_2 \circ F, G \circ S_2 = S_3 \circ G \implies (G \circ F) \circ S_1 = G \circ S_2 \circ F = S_3 \circ (G \circ F)$

Hence,  $(V_1, S_1)$  and  $(V_3, S_3)$  are conjugate under  $H$ .

- $F, G$  are order preserving with order perserving inverses
- $G \circ F$  are order preserving with order preserving inverses

Hence,  $(V_1, S_1) \sim (V_3, S_3)$  under  $(G \circ F)$ .



### Lemma 3.1.3.

If  $(V, S)$  and  $(\hat{V}, \hat{S})$  are order conjugate under  $F$ , then  $S$  is order stable on  $V$  if and only if  $\hat{S}$  is order stable on  $\hat{V}$ .

Proof.

( $\implies$ ) Suppose  $S$  is order stable on  $V$ . This implies

(S1)  $S$  has a unique fixed point  $v^* \in V$

(S2)  $v \in V, v \preceq v^* \implies v \preceq Sv$  and  $v \in V, v^* \preceq v \implies Sv \preceq v$

(S1) implies  $\hat{S}$  has a unique fixed point  $\hat{v}^* := F(v^*) \in \hat{V}$  by Proposition 3.1.2. Moreover we have

$$\bullet \text{ For } \hat{v} := F(v) \in \hat{V}, \hat{v} \leq \hat{v}^* \underbrace{\implies}_{o.i.} v \preceq v^* \underbrace{\implies}_{S2} v \preceq Sv \underbrace{\implies}_{o.i.} Fv \leq \underbrace{FSv}_{conjugate} = \hat{S}\hat{v}$$



In this section, we will see how one ADP can be transformed into another ADP (or ADPs are equivalent up to a transformation).

Such transformation will result

- Simpler form of Bellman equation
- Tailored to solve for some problems (Exponential BE)
- etc.

## Definition 5

Let  $(V, \mathbb{T})$  and  $(\hat{V}, \hat{\mathbb{T}})$  be ADPs with policy sets  $\mathbb{T} := \{T_\sigma : \sigma \in \Sigma\}$  and  $\hat{\mathbb{T}} := \{\hat{T}_\sigma : \sigma \in \Sigma\}$ . We call these ADPs **isomorphic** under  $F$  if

- ①  $F$  is an order isomorphism from  $V$  to  $\hat{V}$
- ② these two ADPs have the same policy set  $\Sigma$
- ③  $(V, T_\sigma)$  and  $(\hat{V}, \hat{T}_\sigma)$  are order conjugate under  $F$  for all  $\sigma \in \Sigma$ .

### Example 3.1.3. Fei et al. (2021) Exponential Bellman Equation

Exponential risk-sensitive Q-factor Bellman equation (ADP:  $((0, \infty)^G, \mathbb{M})$ )

$$M_\sigma h = \exp(\theta r + \beta \ln P_\sigma h), \quad P_\sigma h(x, a) := \sum_{x'} h(x', \sigma(x')) P(x, a, x')$$

Risk-sensitive Q-factor policy operator (ADP:  $(\mathbb{R}^G, \mathbb{T})$ )

$$T_\sigma f = r + \frac{\beta}{\theta} \ln \left[ P_\sigma \exp(\theta f) \right], \quad P_\sigma \exp(\theta f)(x, a) := \sum_{x'} \exp(\theta f(x', \sigma(x'))) P(x, a, x')$$



### Example 3.1.3. Continue

Let  $\theta > 0$ , and

$$Fh = \exp(\theta h)$$

is an order isomorphism from  $\mathbb{R}^G$  to  $(0, \infty)^G$ .

For conjugacy, we have

$$\begin{aligned}(F \circ T_\sigma)(h) &= \exp \left( \theta \left( r + \frac{\beta}{\theta} \right) \ln \left[ P_\sigma \exp(\theta h) \right] \right) \\ &= \exp \left( \theta r + \beta \ln P_\sigma(Fh) \right) \\ &= (M_\sigma \circ F)(h)\end{aligned}$$

Hence,  $((0, \infty)^G, \mathbb{M})$  and  $(\mathbb{R}^G, \mathbb{T})$  are isomorphic

### Example 3.1.4. RDP

Let  $(\Gamma, V, B)$  and  $(\Gamma, \hat{V}, \hat{B})$  be two RDPs with identical state space  $X$ , action space  $A$ , and feasible correspondence  $\Gamma$ . Let  $V = M^X$ ,  $\hat{V} = \hat{M}^X$ , where  $M, \hat{M} \subset \mathbb{R}$ . If there exists an order isomorphism  $\varphi$  from  $M$  to  $\hat{M}$  such that

$$B(x, a, v) = \varphi^{-1}[\hat{B}(x, a, \varphi \circ v)] \quad \text{for all } v \in V \text{ and } (x, a) \in G$$

then  $(V, \mathbb{T})$  and  $(\hat{V}, \hat{\mathbb{T}})$  are isomorphic. From exercise 3.1.2,  $F$  is an order isomorphism from  $V$  to  $\hat{V}$ , and

$$T_\sigma = F^{-1} \circ \hat{T}_\sigma \circ F$$

## Lemma 3.1.4.

### Lemma 6

*Isomorphism between ADPs is an equivalence relation on the set of ADPs.*

### Proof.

Let  $\mathbb{A}$  be the set of ADPs. We denote  $(V_1, \mathbb{T}_1) \cong (V_2, \mathbb{T}_2)$  if there are isomorphic. We need to prove that  $\sim$  is reflexive, symmetric and transitive.

- (Reflexivity) Let  $(V, \mathbb{T}) \in \mathbb{A}$ , as the ADP has the same policy set as itself and by Exercise 3.1.9, we get reflexivity.
- (Symmetry) Let  $(V_1, \mathbb{T}_1) \cong (V_2, \mathbb{T}_2)$ , then they have the same policy set. We use Exercise 3.1.9 get symmetry
- (Transitivity) Let  $(V_1, \mathbb{T}_1) \cong (V_2, \mathbb{T}_2)$  and  $(V_2, \mathbb{T}_2) \cong (V_3, \mathbb{T}_3)$ , hence these three ADPs have the same policy set. We use Exercise 3.1.9. get transitivity.





We take

- $(V, \mathbb{T})$  and  $(\hat{V}, \hat{\mathbb{T}})$  be two ADPs
- $\mathbb{T} := \{T_\sigma : \sigma \in \Sigma\}$ ;  $\hat{\mathbb{T}} := \{\hat{T}_\sigma : \sigma \in \Sigma\}$
- $v_\sigma$  (resp.  $\hat{v}_\sigma$ ) be the unique fixed point of  $T_\sigma$  (resp.  $\hat{T}_\sigma$ )
- $T$  (resp.  $\hat{T}$ ) be the Bellman operator of  $(V, \mathbb{T})$  (resp.  $(\hat{V}, \hat{\mathbb{T}})$ )
- $v^*$  (resp.  $\hat{v}^*$ ) be the value function of  $(V, \mathbb{T})$  (resp.  $(\hat{V}, \hat{\mathbb{T}})$ )

## Theorem 3.1.5.

### Theorem 7

*If  $(V, \mathbb{T})$  and  $(\hat{V}, \hat{\mathbb{T}})$  are isomorphic under  $F$ , then*

- ❶  *$\sigma$  is  $v$ -greedy for  $(V, \mathbb{T}) \iff \sigma$  is  $Fv$ -greedy for  $(\hat{V}, \hat{\mathbb{T}})$*
- ❷  *$\sigma$  is optimal for  $(V, \mathbb{T}) \iff \sigma$  is optimal for  $(\hat{V}, \hat{\mathbb{T}})$*
- ❸ *Regularity, well-posedness, and order stability is preserved under isomorphism.*