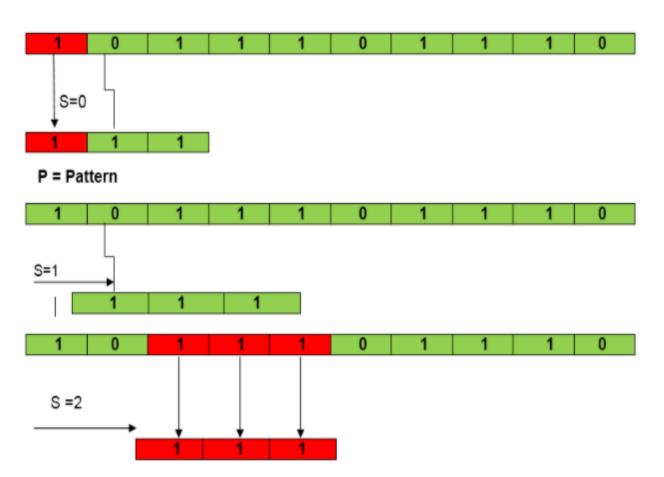
## Naïve String-Matching Algorithm

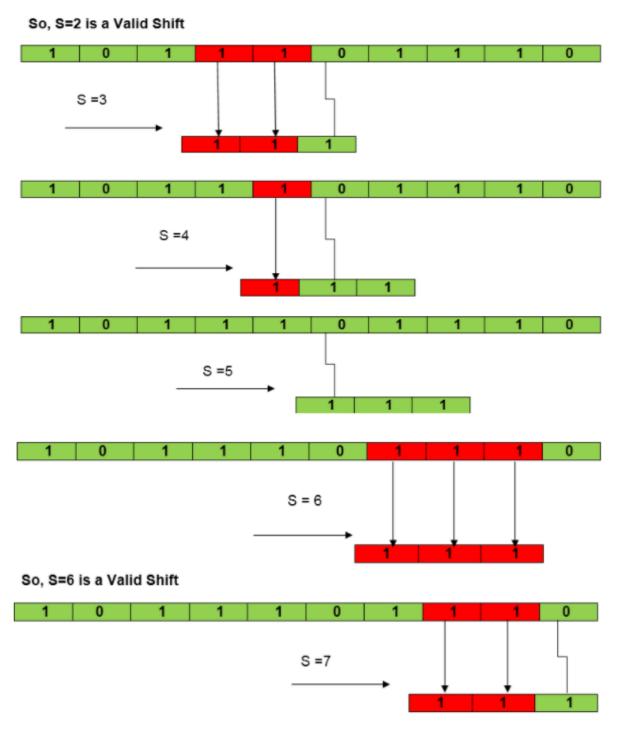
- Naïve pattern searching is the simplest method among other pattern searching algorithms. It checks for all character of the main string to the pattern.
- This algorithm is helpful for smaller texts.
- It does not need any pre-processing phases.
- The naïve approach tests all the possible placement of Pattern P [1.....m] relative to text T [1.....n]. We try shift s = 0, 1.....n-m, successively and for each shift s. Compare T [s+1.....s+m] to P [1.....m].

### Example:

Suppose 
$$T = 1011101110$$
  
 $P = 111$ 







Algorithm:

## **NAIVE-STRING-MATCHER (T, P)**

- 1.  $n \leftarrow length [T]$
- 2.  $m \leftarrow length [P]$
- 3. for  $s \leftarrow 0$  to n m

- 4. do if P[1....m] = T[s + 1...s + m]
- 5. then print "Pattern occurs with shift" s

## Rabin-Karp Algorithm

The Rabin-Karp string matching algorithm calculates a hash value for the pattern, as well as for each M-character subsequence of text to be compared. If the hash values are unequal, the algorithm will determine the hash value for next M-character sequence. If the hash values are equal, the algorithm will analyse the pattern and the M-character sequence. In this way, there is only one comparison per text subsequence, and character matching is only required when the hash values match.

#### Example:

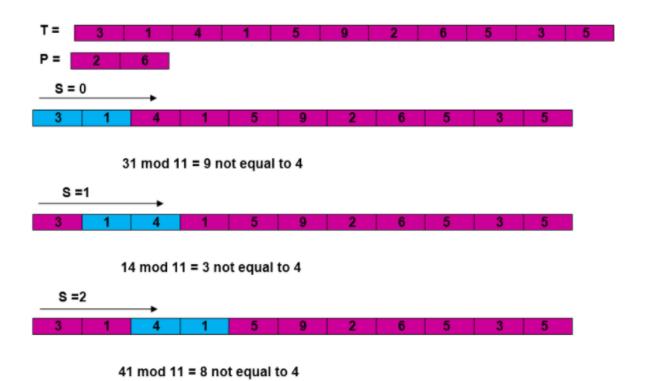
T = 31415926535...

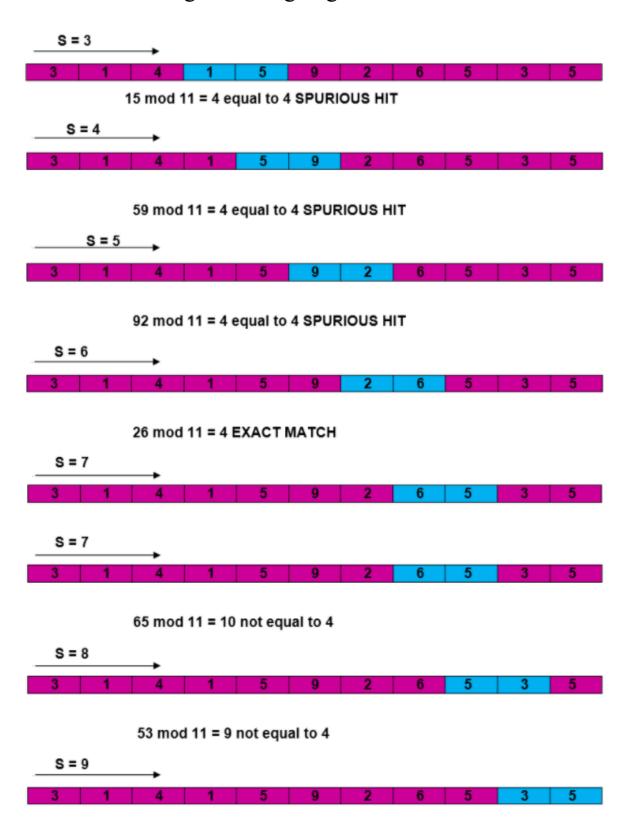
P = 26

Here T. Length = 11 so Q = 11

And P mod  $Q = 26 \mod 11 = 4$ 

Now find the exact match of P mod Q...





35 mod 11 = 2 not equal to 4

The Pattern occurs with shift 6.