# Microprocessor Technologies (102045610)

# MODULE 4 PROGRAMMING TECHNIQUES

#### **Module 4**

- Stack & Subroutines
- Developing Counters and Time Delay Routines
- Code Conversion
- BCD Arithmetic and 16-Bit Data operations

#### Stack

- Stack is a group of memory location in the R/W memory that is used for temporary storage of binary information during execution of a program.
- The starting memory location of the stack can be defined in program and space is reserved usually at the high end of memory map.

E.g.: LXI SP,FFF8; loads 16-bit memory address in stack pointer

## Stack

#### Instruction necessary for stack are as follows:

LXI SP, 2095	Load the stack pointer register with a 16-bit
PUSH B/D/H	address contents of register pair on the stack
PUSH PSW	Operand PSW represents Program status word meaning. i.e. content of accumulator
POP B/D/H	and flags content of top two memory locations of the stack in to specified register pair
POP PSW	It copies content of top two memory locations of the stack in to accumulator and flags respectively.

- A subroutine is a group of instruction that performs a subtask of repeated occurrence.
- A subroutine can be used repeatedly in different locations of the program.

#### **Advantage of using Subroutine**

• Rather than repeat the same instructions several times, they can be grouped into a subroutine that is called from the different locations.

#### Where to write Subroutine?

- In Assembly language, a subroutine can exist anywhere in the code.
- However, it is customary to place subroutines separately from the main program.

- •The 8085 has two instructions for dealing with subroutines.
- The CALL instruction is used to redirect program execution to the subroutine.
- The **RET** instruction is used to return.

#### 1. CALL 16 bit memory

- Call subroutine unconditionally.
- 3 byte instruction.
- Saves the contents of program counter on the stack pointer. Loads the PC by jump
   address (16 bit memory) and executes the subroutine.

#### 2. RET

- Returns from the subroutine unconditionally.
- 1 byte instruction
- Inserts the contents of stack pointer to program counter.
- 3. CC, CNC, CZ, CNZ, CP, CM, CPE, CPO
  - Call subroutine conditionally.
- Same as CALL except that it executes on the basis of flag conditions.
- 4. RC, RNC, RZ, RNZ, RP, RM, RPE, RPO
- Return subroutine conditionally.
- Same as RET except that if executes on the basis of flag conditions.

#### Things to be considered in Subroutine

• Number of **PUSH** and **POP** instruction used in the subroutine must be same, otherwise, **RET** instruction will pick wrong value of the return address from the stack and program will fail.

#### E.g. Write an ALP to add two numbers using subroutines.

2000 MVI B, 4AH

3000 MOV A, B

2002 MVI C, A0H

3001 ADD C

2004 CALL 3000H

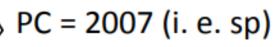
3002 RET

2007 MOV B, A

2008 HLT

SP = 2007 (i.e. PC)

PC = 3000 (i.e. 16bit)



Counters and Time Delays are important techniques.

#### **Applications of Counters and Time Delays**

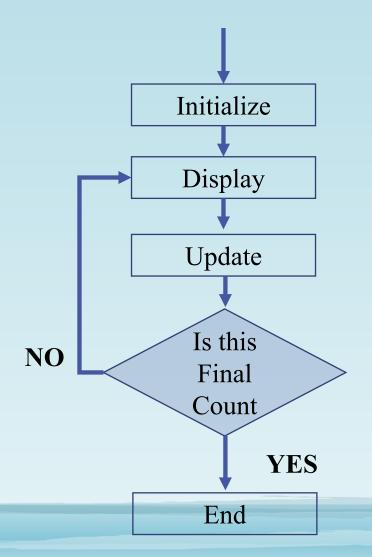
They are commonly used in

- 1. Traffic Signal
- 2. Digital Clocks
- 3. Process Control
- 4. Serial data transfer

- A counter is designed simply by loading appropriate number into one of the registers and using INR or DCR instructions.
- Loop is established to update the count.
- Each count is checked to determine whether it has reached final number; if not, the loop is repeated.

#### Counters

• How to create counters?



- 1. MVI C,05
- 2. LOOP: MOV A,C
- 3. OUT 01
- 4. DCR C
- 5. JNZ LOOP
- **6.** HLT

#### Counters

Label	Opcod e	Operand	Comment	T-states
	MVI	C,o5h	; Load Counter	7
LOOP:	DCR	C	; Decrement Counter	4
МИСО	JNZ	LOOP	; Jump back to Decr. C	10/7

MVI C, o5h

Machine Cycle: F + R = 2

T-States: 4T + 3T = 7T

DCR C

Machine Cycle: **F** =

1

T-States: =

JNZ LOOP (false)

Machine Cycle: F + R = 2

T-States: 4T + 3T = 7T

JNZ LOOP (true)

Machine Cycle: F + R + R = 3

T-States: 4T + 3T + 3T = 10T

- Instruction MVI C, 05h requires 7 T-States to execute.
- Assuming, 8085 Microprocessor with **2MHz** clock frequency.
- How much time it will take to execute above instruction?

Clock frequency of the system (f)= 2 MHz

Clock period (T) = 
$$1/f = \frac{1}{2} * 10^{-6} = 0.5 \mu s$$

Time to execute MVI = 7 T-states \* 
$$0.5 \mu s$$

$$= 3.5 \mu s$$

How much time it will take to execute above instruction with 1 MHz clock frequency?

Label	Opcode	Operand	Comment	T-states
	MVI	C,05h	; Load Counter	7
LOOP:	DCR	C	; Decrement Counter	4
	JNZ	LOOP	; Jump back to Decr. C	10/7

- Now to calculate time delay in loop, we must account for the T-states required for each instruction, and for the number of times instructions are executed in the loop.
- The for the next two instructions:

DCR: 4 T-States

JNZ: + 10 T-States

14 T-States

• Here, the loop is repeated for 5 times.

#### How to calculate time delay for given loop?

• Time delay in loop T<sub>L</sub> with 2MHz clock frequency is calculated as:

$$T_L = T * Loop T-states * N_{10} ----(1)$$

T<sub>L</sub> : Time Delay in Loop

T : Clock Period

N<sub>10</sub>: Equivalent decimal number of hexadecimal count loaded in the delay register

Substituting value in equation (1)

$$T_L = (0.5 * 10^{-6} * 14 * 5)$$
  
= 35  $\mu$ s

• If we want to calculate delay more accurately, we need to accurately calculate execution of JNZ instruction

i.e

If JNZ = true, then T-States = 10

Else if JNZ = false, then T-States = 7

• Now, according to our program:

1. MVI C,05

2. LOOP: DCR C

JNZ LOOP

4. HLT

Here, the last cycle will be executed in **7 T-States**; when **JNZ** = **false** 

Therefore, there is difference of (10T - 7T) 3T-states:

Therefore, there is difference of (10T - 7T) 3T-states:

• Delay generated by last clock cycle:

= 3T \* Clock Period  
= 3T \* 
$$(1/2 * 10^{-6})$$
  
= 1.5  $\mu$ s

• Now, the accurate loop delay is:

$$T_{LA}=T_L$$
 - Delay generated by last clock cycle

= 35 
$$\mu$$
s - 1.5  $\mu$ s  
= 33.5  $\mu$ s



Now, to calculate total time delay

Total Delay = Time taken to execute instruction outside loop

+

Time taken to execute loop instructions

$$T_D$$
 =  $T_O + T_{LA}$   
=  $(7 * 0.5 \mu s) + 33.5 \mu s$   
=  $3.5 \mu s + 33.5 \mu s$   
=  $37 \mu s$ 

Calculate time delay and accurate time delay for given loop with

Counter value =255 (FF h) and

**Clock frequency =2MHz** 

$$T_L$$
= T \* Loop T-states \*  $N_{10}$  -----(1)  
= 0.5 \*10<sup>-6</sup>\* 14\* 255  
= 1785  $\mu$ s = 1.8 ms

 $T_{LA}$ = Time to execute loop instructions

= 
$$T_L - (3T \text{ states*clock period})$$
  
=  $1785 - (3 * \frac{1}{2} * 10^{-6})$   
=  $1785 - 1.5 = 1783.5 \mu s$ 

#### Exercise

1. How much time the 8085 microprocessor will take to execute the MOV B, A instruction, if the **crystal frequency** is 4MHz?

ANS: 1µs

1. How much time will be required to execute the STAX B instruction if the **clock frequency** is 4 MHz?

2. How much time will be required to execute the MVI M,25h instruction if the **clock frequency** is 6 MHz?

22

- Time delay can be considerably increased by setting a loop and using a register pair with a 16-bit number (FFFF h).
- A 16-bit is decremented by using DCX instruction.

#### **Problem with DCX instruction**

- DCX instruction doesn't set Zero flag.
- Without test flag, Jump instruction can't check desired conditions.
- Additional technique must be used to set Zero flag.

23

23

Label	Opcode	Operand	Comment	T-states
	LXI	B,2384 h	; Load BC with 16-bit counter	10
LOOP:	DCX	В	; Decrement BC by 1	6
	MOV	A, C	; Place contents of C in A	4
	ORA	В	; OR B with C to chk Zero flag	4
	JNZ	LOOP	; if result not equal to 0, 10/7 jump back to loop	10/7

Here the loop includes four instructions:

- The loop is repeated for 2384 h times.
- Converting  $(2384)_{16}$  (\_\_\_\_\_\_)<sub>10</sub>  $2384 \text{ h} = (2 * 16^3) + (3 * 16^2) + (8 * 16^1) + (4 * 16^0)$  = 8192 + 768 + 128 + 4 = 9092
- Clock frequency of the system (f)= 2 MHz
- Clock period (T) =  $1/f = \frac{1}{2} * 10^{-6} = 0.5 \mu s$

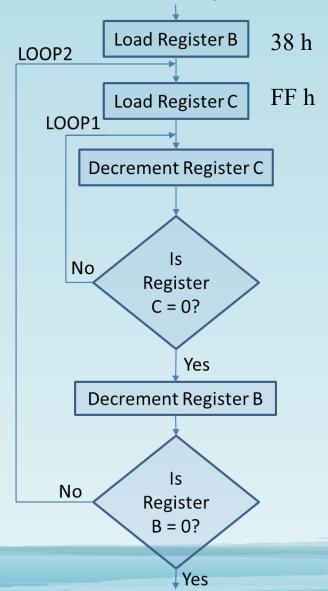
Now, to find delay in the loop

```
T_L= T * Loop T-sates * N_{10}
= 0.5 *10<sup>-6</sup>* 24 * 9092
= 109104 \mus
```

= **109 ms** (without adjusting last cycle)

26

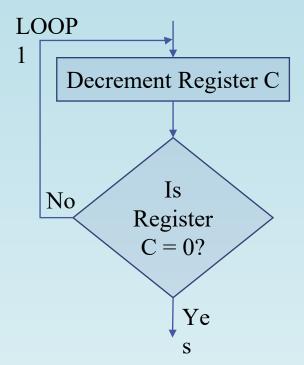
# Time Delay using a LOOP within a LOOP



Label	Opcode	Operand	T-states
	MVI	B,38h	7T
LOOP2:	MVI	C,FFh	7T
LOOP1:	DCR	С	4T
	JNZ	LOOP1	10/7 T
	DCR	В	4T
	JNZ	LOOP2	10/7 T

# Time Delay using a LOOP within a LOOP

Calculating delay of inner LOOP1: T<sub>L1</sub>



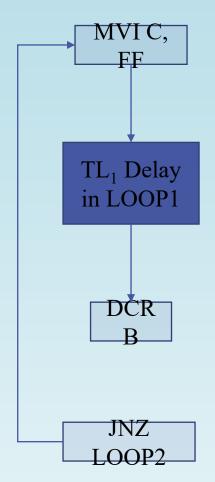
Label	Opcode	Operand	T-states
LOOP	DCR	С	4T
1:	JNZ	LOOP1	10/7 T

$$T_L$$
= T \* Loop T-states \*  $N_{10}$   
= 0.5 \*10<sup>-6</sup>\* 14\* 255  
= 1785  $\mu$ s = 1.8 ms  
 $T_{L1}$ = TL - (3T states\* clock period)  
= 1785 - (3 \*  $\frac{1}{2}$  \* 10<sup>-6</sup>)  
= 1785-1.5=1783.5  $\mu$ s

Delay of Loop1  $T_{L1}$ = 1783.5  $\mu$ s

# Time Delay using a LOOP within a LOOP

• Now, Calculating delay of outer LOOP2: T<sub>L2</sub>



Label	Opcode	Operand	T-states
	MVI	B,38h	7T
LOOP2	MVI	C,FFh	7T
Delay	of Lo	op1 T <sub>L1</sub> =	1783.5 µs
	DCR	В	4T
	JNZ	LOOP2	10/7 T

Counter B : 
$$(38)_{16} = (56)_{10}$$
  
Loop2 is executed for 56 times  
T-States = 7 + 4 + 10 = 21 T-States  
 $T_{L2} = 56 (T_{L1} + 21 T-States * 0.5)$   
= 56( 1783.5 µs + 10.5)  
= 100464 µs  
 $T_{L2} = 100.46 ms$ 

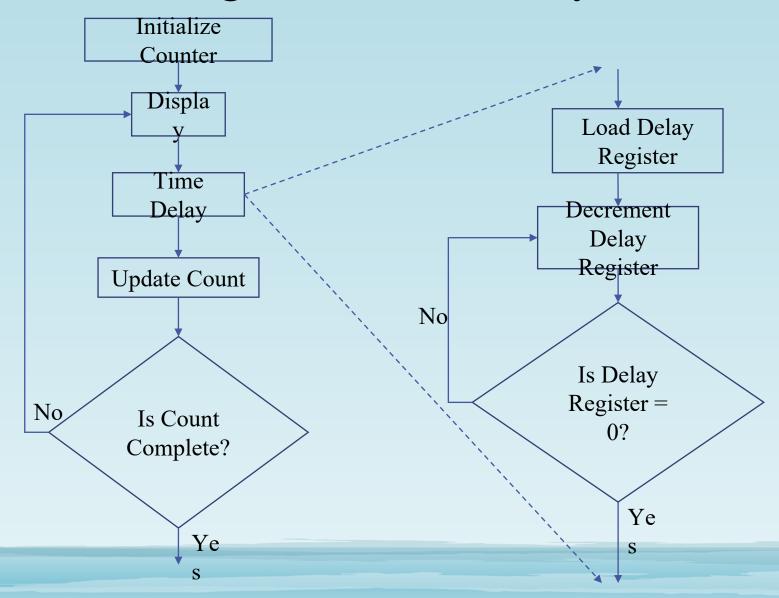
# Disadvantage of using software delay

- Accuracy of time delay depends on the accuracy of system clock.
- Microprocessor is occupied simply in a waiting loop; otherwise, it could be employed to perform other functions.
- The task of calculating accurate time delays is tedious.
- In real time applications timers are commonly used.
- Intel 8254 is a programmable timer chip, that can be interfaced with microprocessor to provide timing accuracy.
- The disadvantage of using hardware chip include the additional expense and the need for extra chip in the system.

30

30

# Counter design with time delay



Write a program to count continuously in hexadecimal from **FFh** to **00h** with **0.5**  $\mu$ s clock period. Use **register** C to set up **1ms** delay between each count and display the number at one of the output port.

#### Given:

Counter= FF h

Clock Period T=0.5 μs

Total Delay = 1 ms

#### **Output:**

To find value of delay counter

32

Program

- 1. MVI B,FF
- 2. LOOP:MOV A,B
- 3. OUT 01
- 4. MVI C, COUNT; need to calculate delay count
- 5. DELAY: DCR C
- 6. JNZ DELAY
- 7. **DCR B**
- 8. JNZ LOOP

#### **Delay Calculations**

1. MVI B, FF;

2. LOOP: MOV A, B

3. OUT 01

4. MVI C, COUNT

5. DELAY: DCR C

6. JNZ DELAY

7. DCR B

Instruction	T-States
DCR C	4
JNZ DELAY	10
Total	14 T

#### **Calculate Delay for Internal Loop**

T<sub>I</sub> = T-States \* Clock Period \* COUNT

$$T_I = (7.0 * 10^{-6}) * COUNT$$

#### **Calculate Delay for Outer Loop:**

$$T_O = T$$
-States \* Clock Period  
= 35 \* 0.5 \*  $10^{-6}$ 

$$T_{\rm O} = 17.5 \; \mu s$$

#### **Calculate Total Time Delay:**

$$T_{D} = T_{O} + T_{L}$$

1 ms = 
$$17.5 * 10^{-6} + (7.0 * 10^{-6}) * COUNT$$

$$1 * 10^{-3} = 17.5 * 10^{-6} + (7.0 * 10^{-6}) * COUNT$$

COUNT= 
$$\frac{1 * 10^{-3} - 17.5 * 10^{-6}}{7.0 * 10^{-6}} \cong (140)_{10} = (8C)_{16}$$

Instruction	T-States
MOV A,B	4
OUT 01H	10
MVI C,	7
<b>BORIST</b>	4
JNZ LOOP	10
TOTAL	35 T

## 0-9 up/down counter program

Write an 8085 assembly language program to generate a decimal counter (which counts 0 to 9 continuously) with a one second delay in between. The counter should reset itself to zero and repeat continuously. Assume a **Clock frequency** of 1MHz.

## 0-9 up counter program

#### Program

1. START: MVI B,00H

2. DISPLAY: OUT 01

3. LXI H, COUNT

4. LOOP: DCX H

5. MOV A,L

6. ORA H

JNZ LOOP

8. INR B

9. **MOV A,B** 

10. **CPI 0A** 

11. JNZ DISPLAY

12. JZ START

Instruction	<b>T-States</b>
DCX	6
MOV A,L	4
ORA H	4
JNZ	10
TOTAL	24 T

## 0-9 up counter program

### **Delay Calculation:**

As show in previous program of register pair, assuming

T-states = 24

Loop Delay  $T_L = 1 \text{ sec}$ 

Clock Period T =  $1 * 10^{-6}$  sec

#### **Find Count**

 $T_L = T * Loop T-states * Count$ 

 $1 \sec = (1.0 * 10^{-6} \sec) * 24 * Count$ 

**Count** = 
$$\frac{1}{24 * 10^{-6}}$$

$$=(41666)_{10}$$

## Exercise

Calculate delay in following loop, assuming clock period =  $0.33\mu$ s

Label	Instruction	T-states
	LXI B, 12FF H	10
DELAY:	DCX B	6
	XTHL	16
	XTHL	16
	NOP	4
	NOP	4
	MOV A,C	4
	ORA B	4
	JNZ DELAY	10/7

TL=T \* T-states \* Count Ans=102ms

### Exercise

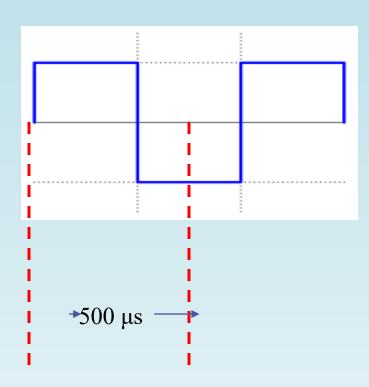
Write a set of 8085 assembly language instructions to generate a 1 second delay, if the **crystal frequency** is 6 MHz.

#### *Note*:

Clock Frequency(Operating Frequency)=Crystal Frequency/2

Clock Period = 1/Clock Frequency

Write a program to generate a continuous square wave with the period of 500  $\mu$ s with clock period of 325ns, and use bit  $D_0$  to output square wave.



### **Problem Analysis**

- In this problem, the period of square wave is 500  $\mu$ s; therefore, the pulse should be ON(logic 1) for 250  $\mu$ s and OFF(logic 0) for remaining 250  $\mu$ s.
- Therefore, the alternate pattern of 0/1 bits can be provided by loading Accumulator with AA h (1010 1010).
- Now rotating the pattern once through each delay loop.
- Bit  $D_0$  of output port is used to provide logic 0 and 1.
- The delay of 250µs can be easily obtained with an 8-bit delay count and one register.

Program		Logic 1	
1. MVI D, 2	<b>AA</b>	(A)	1 0 1 0 1 0 1 0
2. ROTATE:	MOV A,D	After RLC	0 1 0 1 0 1 0 1
3.	RLC	ANI 01h	0 0 0 0 0 0 0 1
4.	MOV D, A	After AND	00000001
5.	ANI 01		
6.	OUT 01	Logic o	
7.	MVI B, COUNT	(A)	0 1 0 1 0 1 0 1
8. DELAY:	DCR B	After RLC	1 0 1 0 1 0 1 0
9.	JNZ DELAY	ANI 01h	0 0 0 0 0 0 0 1
10.	JMP ROTATE	After AND	00000000

### **Delay Calculation:**

- In this problem, the pulse width is relatively small (250  $\mu$ s); therefore, to obtain accurate output pulse, we should take into account for all the T-states.
- The total delay should include the delay in the loop and execution time of the instruction outside the loop.

### **Delay Calculation:**

- No. of instruction outside the loop is seven.
  - Delay outside the loop  $T_0$ =46 T-states \* 325 ns =14.95 $\mu$ s
- Delay Loop includes two instruction, with 14 T-States, except for last cycle 11 T-States

Loop Delay 
$$T_L$$
= 14 T \* 325ns \* (**COUNT -**1) + [ 11 T \*325ns ]

$$T_L = 4.5 \,\mu s \,(COUNT - 1) + 3.575 \,\mu s$$

• Total delay required = 250 μs

$$T_D = T_O + T_L$$
  
250 µs = 14.95µs + 4.5 µs (COUNT -1) + 3.575 µs  
Count = (52)<sub>10</sub>  
= (34)<sub>16</sub>

### Exercise

Write a program to generate a square wave with the period of 400 $\mu$ s with clock period of 325ns. Use bit  $D_0$  to output square wave.

ANS = 
$$(42)_{10}$$

46

### To convert a given decimal number to hexadecimal

LXI H,4300H

MOV A,M

MOV B,A

ANI 0FH

MOV C,A

MOV A,B

RRC

RRC

RRC

RRC

ANI 0FH

MOV B,A

XRA A

CMP B

JZ LAST

BACK:ADI 0AH

DCR B

JNZ BACK

LAST:ADD C

INX H

MOV M,A

#### To convert a given HEXADECIMAL TO DECIMAL.

```
LXI H, 4150H
                     : Point to data
                     ; Initialize hundreds= 0, Tens=0
LXI B, 0000H
MOV A, M
                   ; Get hex data to A
LOOP: SUI 64H
  JC LOOP1
   INR B
                      ; hundreds=hundreds+1
   JMP LOOP
LOOP1: ADI 64H
                         ; if subtracted extra, add it clear carry flag
LOOP2: SUI 0AH
     JC LOOP3
     INR C
                   ; Tens=tens+1
     JMP LOOP2
LOOP3: ADI 0AH
                         ; If subtracted extra, add it again
    INX H
                       ; A = Units
    MOV M, B
                      ; store hundreds
   MOV B, A
                      ; Combine Tens in C &
    MOV A, C
                      ; Units in A to form a
                     ; Single 8-bit number
    RLC
    RLC
    RLC
    RLC
    ADD B
    INX H
    MOV M, A
                       ; Store tens & Units
   HLT
```

Note: In this experiment the number is converted to its equivalent decimal number using the following logic. First count the number of hundreds, the number of tens & units present in that hex number. Then add up to get the equivalent decimal number.

```
Converting A9 we get:
```

```
A9 /64=45 Hundreds = 01
```

Since 64(100 decimal) cannot be subtracted from 45 no. of hundreds = 01. Now count tens

```
45/0A = 3B Tens = 01
```

Now from 09, 0A cannot be subtracted. Hence tens = 06 the decimal equivalent of A9 is

#### PROGRAM 1: Convert BCD TO BINARY OR BCD TO HEX.

		1 0					
Note: Use	nrincing	al of i	nositional	weighing	T 111	ouven	nΩ
Tiote. Osc	princip	աւ Օւ լ	positionai	WCIgiiiiig	5 111	grven	110.

For example: 60=6 x 0A+00

72 = 7x0A + 02

=3C+00

=48

LDA D000H

MOV B,A

ANI 0FH

MOV C,A

MOV A,B

ANI F0H

RRC

RRC

RRC

RRC

MOV B,A

XRA A

MVI D,0AH

L1:ADD D

DCR B

JNZ L1

ADD C

STA D100H

#### Binary to BCD conversion or HEX to BCD CONVERSION

Note: Binary to BCD is done by dividing no. by power often.

Step-1if no. is equal to or greater than 100, divide the no. By 100(i.e. subtract 100 repeatedly till remainder is less than 100) the quotient gives the MSB, DIGIT 2 of BCD no. if no.<100 follow step.2

STEP -2 IF NO. i.e. remainder of 1st division is equal to or greater than 10 divide no. by 10 repeatedly until remainder is less than 10 quotient, the digit 1, if no. is less than 10, go to step 3.

Step-3 Remainder it gives digit 3

LDA 3040H

MVI B,64H

MVI C,0AH

MVI D,00H

MVI E,00H

L1:CMP B

JC L2

SUB B

INR E

JMP L1

L2:CMP C

JC L3

SUB C

INR D

JMP L2

L3: STA 3041H

MOV A,D

STA 3042H

MOV A,E

STA 3043H

ı	Decimal	Hex	ASCII	Decimal	Hex	ASCII	Decimal	Hex	ASCII	Decimal	Hex	ASCII
н	0	00	NUL	32	20	(blank)	64	40	@ A	96	60	
ш	1	01	SOH	33	21	!	65	41		97	61	a
ш	2	02	STX	34	22	-	66	42	В	98	62	b
ш	3	03	ETX	35	23	#	67	43	C	99	63	С
ш	4	04	EOT	36	24	\$	68	44	D	100	64	d
ш	5	05	ENQ	37	25	%	69	45	E	101	65	e
ш	6	06	ACK	38	26	&	70	46	F	102	66	f
ш	7	07	BEL	39	27		71	47	G	103	67	g
ш	8	80	BS	40	28	(	72	48	Н	104	68	h
ш	9	09	HT	41	29	)	73	49	- 1	105	69	i
ш	10	OA.	LF	42	2A	•	74	4A	J	106	6A	j
ш	11	0B	VT	43	2B	+	75	4B	K	107	6B	k
ш	12	OC.	FF	44	2C		76	4C	L	108	6C	1
ш	13	0D	CR	45	2D	-	77	4D	M	109	6D	m
ш	14	0E	SO	46	2E		78	4E	N	110	6E	n
ш	15	0F	SI	47	2F	1	79	4F	0	111	6F	0
ш	16	10	DLE	48	30	0	80	50	P	112	70	p
ш	17	11	DC1	49	31	1	81	51	Q	113	71	q
ш	18	12	DC2	50	32	2	82	52	R	114	72	r
ш	19	13	DC3	51	33	3	83	53	S	115	73	S
ш	20	14	DC4	52	34	4 5	84	54	T	116	74	t
ш	21	15	NAK	53	35	5	85	55	U	117	75	u
ш	22	16	SYN	54	36	6	86	56	V	118	76	v
ш	23	17	ETB	55	37	7	87	57	W	119	77	w
П	24	18	CAN	56	38	8	88	58	X	120	78	×
П	25	19	EM	57	39	9	89	59	Y	121	79	У
	26	1A	SUB	58	3A	2	90	5A	Z	122	7A	Z
	27	1B	ESC	59	3B	;	91	5B	[	123	7B	{
	28	1C	FS	60	3C	<	92	5C	1	124	7C	1
	29	1D	GS	61	3D	=	93	5D	]	125	7D	}
	30	1E	RS	62	3E	>	94	5E	٨	126	7E	~
	31	1F	US	63	3F	?	95	5F		127	7F	(delete)

# **ASCII Chart**

### HEX or Binary to ASCII HEX Code

LXI H,0050H LXI D,0052H MOV A,M MOV B,A RRC RRC RRC RRC CALL ASCII STAX D INX D MOV A,B **CALL ASCII** STAX D HLT ASCII: ANI oFH CPI oAH JC CODE ADI 07H CODE: ADI 30H **RET** 

Write 8085 Assembly language program to convert ASCII to Hexadecimal character values.

LXI H,8000H MOV A, M CPI 58H JNC NUM SUI 37H JMP STORE NUM:SUI 30H STORE:INX H

MOV M, A

To find the square of the number from 0 to 9 using a Table of Square.

#### ALGORITHM:

- 1. Initialize HL pair to point Look up table
- Get the data .
- 3. Check whether the given input is less than 9.
- 4. If yes go to next step else halt the program
- 5. Add the desired address with the accumulator content
- Store the result

#### PROGRAM:

AFTER:

LXI	H,4125	Initialsie Look up table address
LDA	1150	Get the data
CPI	OA	Check input > 9
JC	AFTER	if yes error
MVI	A,FF	Error Indication
STA	4151	
HLT		
MOV	C,A Add	the desired Address
MVI	B,00	
DAD	В	
MOV	A,M	
STA	4151 Store	the result

Terminate the program

#### LOOKUP TABLE:

4125	01
4126	04
4127	09
4128	16
4129	25
4130	36
4131	49
4132	64
4133	81

HLT

#### OBSERVATION:

Input:	4150:	05
Output:	4151	25 (Square)

	Impact	=	4150:	1 1
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### Add two 8 bit BCD numbers stored at consecutive memory locations.

```
C, ooH
START: MVI
           H, 4500H
      LXI
      MOV A, M
      INX
          Η
      ADD M
      DAA
      JNC L1
      INR C
L1:INX H
           M, A
     MOV
           Η
     INX
     MOV
          M, C
     HLT
```