# Design patterns

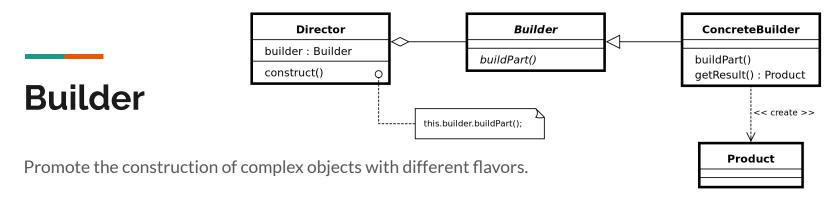
### References

Design Patterns: Elements of Reusable Object-Oriented Software (Erich Gamma, Richard Helm, Ralph Johnson, John Vlissides)

https://github.com/tmrts/go-patterns

en.wikipedia.org

# **Creational patterns**



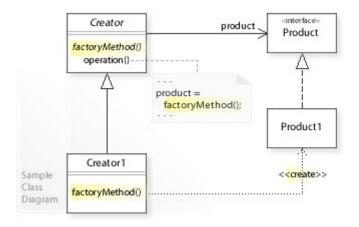
#### Use when:

- Use when you have objects that needs to be built in several steps and want to delay object finalisation until all steps are completed
- You have multiple flavors of a same object to construct.
- You want to decouple the object creation from its constructor.

## **Factory method**

Decouple construction from the object itself.

Use when:



You will only know the concrete type to instantiate at runtime.

# Singleton !



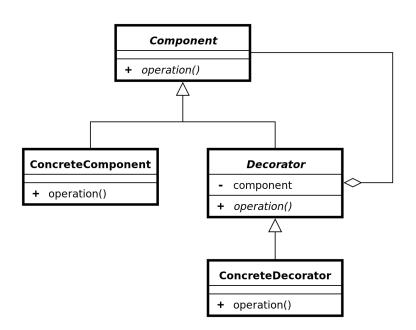
Limits the number of instance to 1.

Often consider to be an anti pattern!

# Structural patterns

### **Decorator**

Allows to dynamically add behavior to an object.

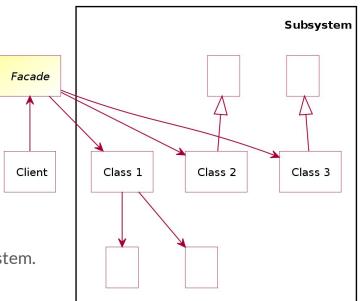


## **Facade**

The Facade presents an interface to a subsystem.

#### Use when:

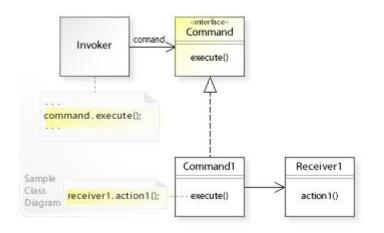
• You need to offer a simplified interface to a complex system.

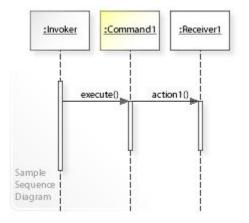


Sample class diagram

# Behavioral patterns

## **Command**





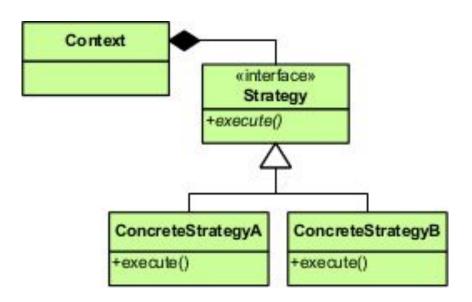
Bundles a command and an argument to be called later.

#### Use when:

• You need to prepare an "action" to be run later.

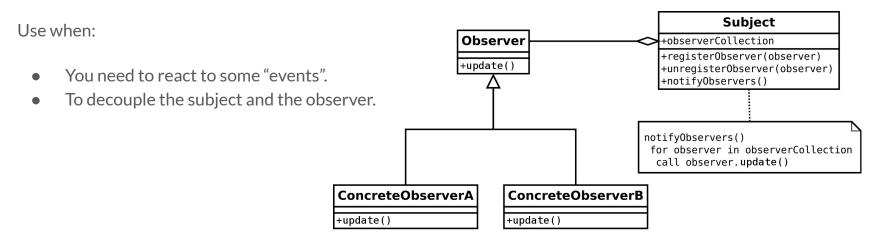
## **Strategy**

Allow to select an algorithm at runtime.



### Observer

A "Subject" notifies registered "Observers" of certain events.



# **Functional patterns**

## Lambda

Simply an anonymous functions.

Because sometimes you don't need to put names on everything, including functions.

Typical use cases: short-live callbacks. (The observer patterns may then become obsolete)

## Closure

A block of code (usually a function) that captures (part of) its surrounding scope (usually variables).

Use when you need to represent a process with its associated state.

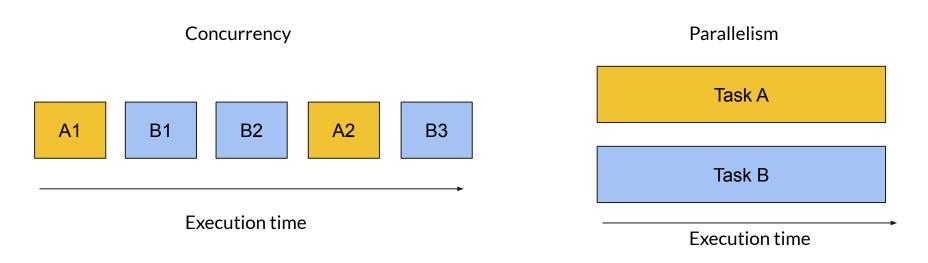
## Memoization (with closure)

Remembering the result of an idempotent function.

Use on expensive function calls. May also be use on things that are not strictly idempotent (Http requests)

# Concurrency patterns

## Concurrency is not parallelism



# Processes, threads and green threads

Terminology might get mixed up! (speaking of processes as a task for example)

- Processes are handled by the OS
- Executions is scheduled on physical threads by OS (hardware might intervene)
- Green threads are scheduled by the language runtime and mapped to OS threads

go someFunction()

# The problem: Data Race

... and also deadlocks

#### Caused by:

- Asynchronicity of your architecture
- Modern CPU caching architecture
- Compiler optimizations

## **Memory Model**

A memory model states what behaviour is guaranteed, or is not guaranteed considering sharing of data between multiple threads.

## Memory Model: "Happens before relationship

If Thread1 states:

a = 1

b = 2

Thread2 might see:

b = 2

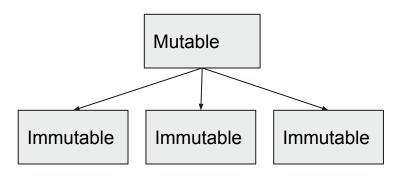
a = 1

The "Happens before" relationship is only guaranteed within the originating thread.

## A property: Immutability

Data races originate from memory writing, aka modification in mutable entities.

Immutable entities are by nature thread safe.

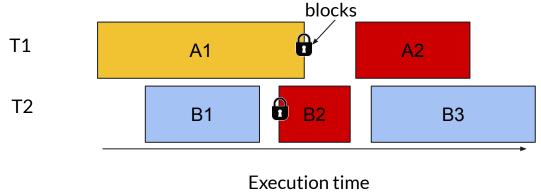


\*Segregate immutability

## Mutex/Locks

A Mutex guards a critical section of your code.

Only 1 thread should execute the critical sections at once.



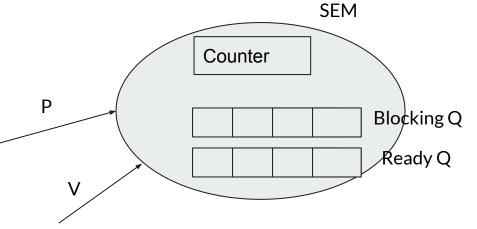


## Semaphore/Signaling

Can be seen of a generalisation of Mutexes.

**P** (wait): decrement counter by 1. If counter < 0, blocks and is added to blocking queue.

V (signal): increments counter by 1. If pre-increment value was negative, transfers the first element in the blocking queue to the ready queue.



## **Producer-Consumer problem**

One process emits items, another process consumes them. Items are put on a shared buffer queue.

- Consumer blocks if buffer queue is empty.
- Producer blocks if buffer queue is full.

#### produce:

```
P(useQueue)
putItemIntoQueue(item)
V(useQueue)
```

#### consume:

```
P(useQueue)
item ← getItemFromQueue()
V(useQueue)
```

Can scale to multiple Producer and Consumers

## **Communication Sequential Process (CSP)**

- Design each task (process) in sequential execution.
- Data is **communicated** through channels. No shared states.

## Channels

"Don't communicate by sharing memory, share memory by communicating"

Build a channel:

c := make(chan int)

Send to a channel: c <- 1

Receive from a channel: msg := <-c

## **Buffered and unbuffered**

Unbuffered Buffered

c := make(chan int, 1)

Sender blocks until receiver consumes. Sender blocks if channel is full.

Receiver blocks until sender sends. Receiver blocks if channel is empty.

## Close channels

To signal receivers not to wait anymore. Most of the time, the sender should be responsible for closing.

```
c := make(chan int)
close(c)
_, ok := <-c
fmt.Println(ok) // Print false</pre>
```

## Select

Analog to switch statement for channels operations.

Choses the first non-blocking case, or default.

## **Avoid blocking on channels**

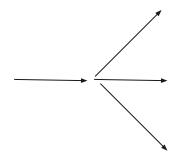
```
func tryReceive(c <-chan int) (data int, more, ok bool) {
    select {
    case data, more = <-c:
        return data, more, true
    default:
        return 0, true, false
    }
}</pre>
```

## Avoid blocking on channels (timeout)

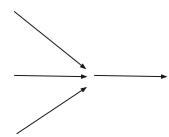
```
func tryReceiveWithTimeout(c <-chan int, duration time.Duration) (data int, more, ok bool) {
    select {
    case data, more = <-c:
        return data, more, true

    case <-time.After(duration): // returns a channel
        return 0, true, false
    }
}</pre>
```

## Shaping the execution flow







Fan-in

## **Fanout**

```
func fanOut(In <-chan int, OutA, OutB chan int) {</pre>
    for data := range In { // receives until closes
        select {
        case OutA <- data:
            //TODO
            break
        case OutB <- data:</pre>
            //TODO
            break
```

## That's it!