## TESTING 'RANDOMNESS' OF RANDOM NUMBER GENERATORS

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## Abstract:

Randomness occurs when events or series of events happen without any pattern, and we cannot predict them in advance. Natural examples of randomness include nuclear decay, and true random processes can only be generated using hardware random generators, which generate random numbers based on physical processes such as quantum-mechanical phenomena. In contrast, everyday computing systems, algorithmic settings, and cryptography rely on pseudo-random numbers. Various methods of PRNGS are used in software, such as Linear Congruential Generator, Mersenne Twister, and Pseudo Random Number Generator, which use mathematical formulas to approximate random numbers. They typically start with a seed value and perform operations to generate sequences of seemingly random numbers. The article then discusses four sources that use different methods to generate random numbers, including the Python-random function, PUF-based code using hardware, Python-secrets, and a true random generator that uses atmospheric noise.

Random number generators find wide application in fields where security is a key requirement. Cybersecurity, cryptography, and IOT are such crucial places where we require a non-predictable key for protection.

## 1. Introduction:

#### 1.1 True Vs. Pseudo-Random Number Generators

Random Number Generators(RNGs) are sequence generators that use a nondeterministic source (i.e., the entropy source), along with some processing function (i.e., the entropy distillation process) to produce randomness. The use of a distillation process is needed to overcome any weakness in the entropy source that results in the production of non-random numbers (e.g., the occurrence of long strings of zeros or ones). The noise in an electrical circuit, the timing of user processes (e.g keystrokes or mouse movements), or the quantum effects in a semiconductor are usually used as entropy sources. This is the working principle of RNGs.

However, in our realization, most RNGs are not truly random. The output of any RNG needs to satisfy strict randomness criteria as measured by statistical tests in order to determine that the physical sources of the RNG inputs appear random. For example, a physical source such as electronic noise may contain a superposition of regular structures, such as waves or other periodic phenomena, which may appear to be random, yet are determined to be non-random using statistical tests. This is a major flaw existing in the world of computer science.

The primary criteria for a sequence to be qualified as 'Truly Random' is-

All elements of the sequence are generated independently of each other, and the value of the next element in the sequence cannot be predicted, regardless of how many elements have already been produced.

Such an expectation is not practically realizable and thus RNGs can only get close to becoming truly random but it is impossible to be truly random.

The main objective of RNGs is to fulfill the property of Forward Unpredictability. It implies that the next output number in the sequence should be unpredictable in spite of any knowledge of previous random numbers in the sequence. The biggest application of RNGs is for cryptographic applications. For example, common cryptosystems employ keys that must be generated in a random fashion. Many cryptographic protocols also require random or pseudo random inputs at various points, e.g., for auxiliary quantities used in generating digital signatures, or for generating challenges in authentication protocols.

Other applications include science experiments, art (such as synthesizers for music production), statistical analysis of data, cryptography, gaming, gambling, and other fields.

#### 1.2 RNGs used:

## (a) Random module from Python:

This module implements a pseudo-random number generator by using the Mersenne Twister algorithm. Mersenne twister has a long period of  $2^{19937}$ -1 before it repeats itself.

The code to generate a random sequence of 0 and 1 bits uses random.choice() which selects either 0 or 1 and repeats this process 1.5 million times to generate a sequence.[1]

#### (b) Secrets module from Python:

The Mersenne Twister is not cryptographically secure [2]. It keeps track of 624 32-bit values, and if an attacker can gather 624 sequential values, then the entire sequence can be predicted. Therefore, python produces a 'secrets' module which provides cryptographically secure sequences.

Similar to the random module, we use secrets.choice() to generate a sequence.

#### (c) PUF-based random number generator:

PUF( Physical Unclonable Functions), are hardware random number generators that extract randomness directly from complex physical systems.[3]

#### (d) Random number database from Random.org:

The bytes generated are from atmospheric noise. The bytes are then concatenated to form a bitstream.

## 1.3 Testing methodology to find 'randomness' of RNGs:

Randomness is a probabilistic property; that is, the properties of a random sequence can be characterized and described in terms of probability. A statistical test is formulated to test a specific null hypothesis (H0), which is that the sequence being tested is random. The alternate hypothesis (Ha) would be that the sequence is not random.

Each test tries to accept or reject the null hypothesis. To do that, each test has a randomness statistic, which under an assumption, gives a distribution of likely values. A critical value is then determined (at 99% significance level for the test). If the test statistic exceeds the critical value, then the null hypothesis is rejected, and otherwise accepted.

# 1.4 Most commonly used terminologies in the tests

(a) Entropy: A measure of the disorder or randomness in a closed system. The entropy of uncertainty of a random variable X with probabilities  $p_1, ..., p_n$  is defined to be

$$H(X) = -\sum_{i=1}^{n} p_i \log p_i$$

(b) Erfc: The complementary error function erfc(z)

$$erfc(z) = \frac{2}{\sqrt{\pi}} \int_{z}^{\infty} e^{-u^{2}} du$$

(c) Incomplete Gamma Function: The incomplete gamma function Q(a,x)

$$P(a,x) = \frac{\gamma(a,x)}{\Gamma(a)} = \frac{1}{\Gamma(a)} \int_{0}^{\infty} t^{z-1} e^{-t} dt$$

- (d) Level of Significance (α): The probability of falsely rejecting the null hypothesis, i.e., the probability of concluding that the null hypothesis is false when the hypothesis is, in fact, true.
- (e) P-value: The probability (under the null hypothesis of randomness) that the chosen test statistic will assume values equal to or worse than the observed test statistic value when considering the null hypothesis.

# 2. Test Implementations and Results:

# Frequency (Monobit) Test:

This test checks to see if a sequence has roughly the same amount of ones and zeros as would be anticipated for a really random sequence. The test determines whether the fraction of ones is near to 0.5, meaning that there should be roughly equal numbers of ones and zeros in a sequence. All subsequent tests are based on the result of this test.

The variables used in the Test:

n - The length of the bit string

*Decision Rule:* P-Value. If P-value < 0.01, the sequence is classified as non-random, and the test FAILS.

#### *Test Implementation:*

1. Conversion to  $\pm 1$ : The zeros and ones of the input sequence are converted to values of -1 and +1 and are added together to produce Sn ('difference' as per the code).

2. Compute the test statistic,

$$s_{obs} = \frac{|S_n|}{\sqrt{n}}$$

3. Calculate the P-value,

$$P$$
-value =  $erfc\left(\frac{s_{obs}}{\sqrt{2}}\right)$ 

```
# Compute score
    score: float = math.erfc(float(difference) /
(math.sqrt(float(bits.size)) * math.sqrt(2.0)))
    # Return result
    if score >= self.significance_value:
        return Result(self.name, True, numpy.array(score))
    return Result(self.name, False, numpy.array(score))
```

#### **Results and Inferences:**

After executing the test for the four random number datasets, the results were :-

PUF Random Dataset - P-value = 0.376, Pass

Python Random Dataset - P-value = 0.407, Pass

Python Secrets Dataset - P-value = 0.02, Pass

True Random Dataset - P-value = 0.417, Pass

Thus, we conclude that all selected datasets passed the Monobit test. All selected datasets can be declared 'Random' according to this test. The proportion of zeros to ones in all datasets is close to 1. In case of a Fail result, this test is a clear indication that the dataset is not random and further tests need not be performed.

# Frequency Test within a Block:

The focus of this test is the proportion of ones within M-bit blocks. The purpose of this test is to determine whether the frequency of ones in an M-bit block is approximately M/2, as would be expected under an assumption of randomness.

The variables used in the Test:

M - The length of each block n - The length of the bit string

*Decision Rule:* P-Value. If P-value < 0.01, the sequence is classified as non-random, and the test FAILS.

#### *Test Implementation:*

1. Partition the input sequence into N =  $\left\lfloor \frac{n}{M} \right\rfloor$  n-overlapping blocks. Discard any unused bits.

Here, we have used 2 functions, one is for defining the size of the block and other is for executing the test. Extremum limits have been applied such that the test always remains within the range of functionality.

```
def __init__(self):
        self. sequence size min: int = 100
        self. default block size: int = 20
        self. blocks number max: int = 100
        self. last bits size: int = -1
        self._block_size: int = -1
        self. blocks number: int = -1
        super(FrequencyWithinBlockTest, self).__init__("Frequency Within")
Block", 0.01)
    def _execute(self,
                 bits: numpy.ndarray) -> Result:
        Overridden method of Test class: check its docstring for further
information.
        if self._last_bits_size == -1 or self._last_bits_size !=
bits.size:
            block_size: int = self._default_block_size
            blocks number: int = int(bits.size // block size)
```

2. Determine the proportion  $\pi_i$  of ones in each M-bit block using the equation

$$\pi_i = \frac{\sum_{j=l}^{M} \varepsilon_{(i-l)M+j}}{M}, \quad \text{for } 1 \le i \le N.$$

3. Compute the  $\chi^2$  tatistic:  $\chi^2(obs) = 4 M \sum_{i=1}^{N} (\pi_i - \frac{1}{2})^2$ .

```
# Compute Chi-square
      chi_square: float = numpy.sum(4.0 * block_size *
      ((block_fractions[:] - 0.5) ** 2))
```

#### 4. Compute P-value = igame

#### **Results and Inferences:**

After executing the test for the four random number datasets, the results were :-

```
PUF Random Dataset - P-value = 0.0, Fail

Python Random Dataset - P-value = 0.472, Pass

Python Secrets Dataset - P-value = 0.851, Pass

True Random Dataset - P-value = 0.843, Pass
```

Thus, we conclude that except for the PUF Random Dataset, all tests passed the test. This reflects that there were large deviations from equal proportions of ones and zeros in many blocks of the PUF Random Dataset. On the other hand, the rest of the datasets do not have any large variations in blocks thus they have passed this test. Thus, we can be sure that except the PUF Random Dataset, all datasets are 'Random' according to this test.

#### Runs Test:

The focus of this test is the total number of runs in the sequence, where a run is an uninterrupted sequence of identical bits. A run of length k consists of exactly k identical bits and is bounded before and after with a bit of the opposite value. The purpose of the runs test is to determine whether the number of runs of ones and zeros of various lengths is as expected for a random sequence. In particular, this test determines whether the oscillation between such zeros and ones is too fast or too slow

The variables used in the Test:

n - The length of the bit string

*Decision Rule:* P-Value. If P-value < 0.01, the sequence is classified as non-random, and the test FAILS.

## *Test Implementation:*

1. Compute the pre-test proportion  $\pi$  of ones in the input sequence:  $\pi = \frac{\sum_{j} \varepsilon_{j}}{n}$ .

$$V_n(obs) = \sum_{k=1}^{n-1} r(k) + 1$$
,

2. Compute the test statistic, where r(k)=0 if  $\varepsilon_k=\varepsilon k+1$ , and r(k)=1 otherwise.

```
# Count the observed runs (list of adjacent equal bits)
  observed_runs: float = 1.0
  for i in range(bits.size - 1):
    if bits[i] != bits[i + 1]:
      observed_runs += 1.0
```

3. Compute P-value = 
$$\operatorname{erfc}\left(\frac{|V_n(obs) - 2n\pi(1-\pi)|}{2\sqrt{2n}\pi(1-\pi)}\right)$$
.

```
# Compute score (P-value)
score: float = math.erfc(abs(observed_runs - (2.0 *
bits.size * proportion * (1.0 - proportion))) / (2.0 *
math.sqrt(2.0 * bits.size) * proportion * (1 - proportion)))
# Return result
if score >= self.significance_value:
    return Result(self.name, True, numpy.array(score))
return Result(self.name, False, numpy.array(score))
```

#### **Results and Inferences:**

After executing the test for the four random number datasets, the results were :-

```
P-value = 0.0, Fail

Python Random Dataset - P-value = 0.201, Pass

Python Secrets Dataset - P-value = 0.64, Pass

True Random Dataset - P-value = 0.738, Pass
```

Thus, we conclude that except for the PUF Random Dataset, all tests passed the test. This reflects that there were too few oscillations (changes from a string of zeros to string of ones or vice versa) in the PUF Random Dataset. On the other hand, the rest of the datasets have very high oscillations leading them to becoming more 'random' in comparison with the PUF Random Dataset. Thus, except for the PUF Random Dataset, all datasets are 'Random' according to this test.

A major real life application of this test is in the field of Trading where Technical traders can use a runs test to analyze statistical trends and help spot profitable trading opportunities.

## **Linear Complexity Test:**

The test divides the bitstream (n) into N blocks of M bits (n=MN). Then, it tries to find the length of the fixed LSFR (Linear Fixed Shift Register), L<sub>i</sub> (i=1,2...N). The randomness of the stream is characterized by a longer LFSR.

The variables used in the Test:

```
M- length of bits in a block n- length of the bitstream \epsilon - Sequence of bits generated by the random number generator(RNG) K - Degrees of freedom (K=6 for the test)
```

*Test Statistic*:  $\chi^2$  (obs), Measure how well the observed occurrences match the expected number under an assumption of randomness.

Decision Rule: P-Value (computed using  $\chi$  and K). If P-value < 0.01, the sequence is classified as non-random, and the test FAILS.

## *Test Implementation:*

1. Partition the n-bit sequence into N-independent blocks of M bits.

```
def __init__(self):
    # Define specific test attributes
    self._sequence_size_min: int = 1000000
    self._pattern_length: int = 512
    self._freedom_degrees: int = 6
    self._probabilities: numpy.ndarray = numpy.array([0.010417,
0.03125,
    0.125,0.5, 0.25, 0.0625, 0.020833])
```

Initialisation and parameters are set [1]

2. Use the Berlekamp-Massey algorithm and determine the linear complexity  $L_i$  (length of the shortest LFSR sequence that generates all bits in the block i) for each N block. Within any  $L_i$  bit sequence, some combination of the bits, when added together modulo 2, produces the next bit in the sequence (bit Ly + 1)

```
def _berlekamp_massey(sequence: numpy.ndarray) -> int:
    """
```

```
Compute the linear complexity of a sequence of bits by the means of the
Berlekamp Massey algorithm.
        :param sequence: the sequence of bits to compute the linear complexity
for
        :return: the int value of the linear complexity
        # Initialize b and c to all zeros with first element one
        b: numpy.ndarray = numpy.zeros(sequence.size, dtype=int)
        c: numpy.ndarray = numpy.zeros(sequence.size, dtype=int)
        b[0] = 1
        c[0] = 1
        # Initialize the generator length
        generator length: int = ∅
        # Initialize variables
        m: int = -1
        n: int = 0
        while n < sequence.size:</pre>
            # Compute discrepancy
            discrepancy = sequence[n]
            for j in range(1, generator_length + 1):
                discrepancy: int = discrepancy ^ (c[j] & sequence[n - j])
            # If discrepancy is not zero, adjust polynomial
            if discrepancy != 0:
                t = c[:]
                for j in range(∅, sequence.size - n + m):
                    c[n - m + j] = c[n - m + j] ^ b[j]
                if generator_length <= n / 2:</pre>
                    generator length = n + 1 - generator length
                    m = n
                    b = t
             n = n + 1
        # Return the length of generator
        return generator_length
```

3. Under an assumption of randomness, calculate the theoretical mean  $\mu$ :

$$\mu = \frac{M}{2} + \frac{\left(0 + \left(-1\right)^{M+1}\right)}{36} - \frac{\left(\frac{M}{3} + \frac{2}{9}\right)}{2^{M}}.$$

Also, For each substring, calculate a value of T<sub>i</sub>

$$T_i = (-1)^M \cdot (L_i - \mu) + \frac{2}{2}$$

```
# Compute mean
self._mu: float = (self._pattern_length / 2.0) + (((-1) **
(self._pattern_length + 1)) + 9.0) / 36.0 - ((self._pattern_length / 3.0) +
(2.0 / 9.0)) / (2 ** self._pattern_length)
```

4. Record the Ti values in v0,..., v6 as follows:

```
Ti \le -2.5 Increment v_0 by 1

-2.5 < Ti \le -1.5 Increment v_1 by 1

-1.5 < Ti \le -0.5 Increment v_2 by 1

-0.5 < Ti \le 0.5 Increment v_3 by 1

0.5 < Ti \le 1.5 Increment v_4 by 1

1.5 < Ti \le 2.5 Increment v_5 by 1

Ti > 2.5 Increment v_6 by 1
```

5. Compute  $\chi^2$  (obs) using,

$$\chi^{2}(obs) = \sum_{i=0}^{K} \frac{(v_{i} - N\pi_{i})^{2}}{N\pi_{i}}$$

Where,  $\pi_0 = 0.010417$ ,  $\pi_1 = 0.03125$ ,  $\pi_2 = 0.125$ ,  $\pi_3 = 0.5$ ,  $\pi_4 = 0.25$ ,  $\pi_5 = 0.0625$ ,  $\pi_6 = 0.020833$ 

And calculate P-Value using,

$$P$$
-value = **igamc**  $\left(\frac{K}{2}, \frac{\chi^2(obs)}{2}\right)$ 

```
chi_square: float = float(numpy.sum(((frequencies[:] -
  (blocks_number * self._probabilities[:])) ** 2.0) / (blocks_number *
  self._probabilities[:])))
    # Compute the score (P-value)
    score: float = scipy.special.gammaincc((self._freedom_degrees
/ 2.0), (chi_square / 2.0))
    # Return result
    if score >= self.significance_value:
        return Result(self.name, True, numpy.array(score))
    return Result(self.name, False, numpy.array(score))
```

#### **Results and Inferences:**

```
Python RNG - FAILED - score: 0.0
Python RNG (Secrets)
```

The result implies that there was a degree of randomness below the accepted level of significance (P>0.01). Also, the test requires n>1,000,000 and 500< M<5000 (M=512 in the test code). Failing the test and score=0 would imply that the observed frequency counts of  $T_i$  stored in the  $v_i$  values were different from the expected values. If score>0.01, then it is expected that the distribution of the frequency of the Ti should be proportional to the computed  $\pi_i$ .

```
\begin{array}{ll} PUF\;RNG & \quad \text{-Not eligible for the test} \\ True\;RNG & \quad \end{array}
```

For PUF RNG, n=30,000 and for True RNG, n=131,076. As the test requires a minimum of n=1,000,000, they are not eligible for the test.

## Serial Test:

The test finds the frequency of all overlapping m-bit patterns across n-bit sequences and determines whether the occurrence of  $2^m m$  bit overlapping patterns is the same as what would be expected for the true random sequence. A random sequence would have uniformity. Every m-bit pattern would have the same chance of occurring as any other m-bit pattern. If m=1, the test becomes the Monobit test.

The variables used in the Test:

```
m- length of bits in a block n- length of the bitstream \epsilon - Sequence of bits generated by the random number generator(RNG)
```

*Test Statistic*:  $\nabla \psi^2_m(obs)$  and  $\nabla^2 \psi^2_m$ . A measure of how well the observed frequencies of m-bit patterns match he expected frequencies of the m-bit patterns.

Decision Rule: If the P-value < 0.01, then the sequence is classified as non-random.

## Test Implementation:

1. Form an augmented sequence,  $\varepsilon$ , by padding first m-1 bits of  $\varepsilon$  at the end.

```
# Pad the sequence
padded_bits: numpy.ndarray =
numpy.concatenate((bits,bits[0:self._pattern_length - 1]))
```

2. Determine the frequency of all possible overlapping m-bit blocks, all possible overlapping (m-1)-bit blocks and all possible overlapping (m-2)-bit blocks.

```
Let, V_{i1..i(m)} denote the frequency of the m-bit pattern i_1...i_m; V_{i1..i(m-1)} denote the frequency of the (m-1)-bit pattern i_1...i_{m-1}; V_{i1..i(m-2)} denotes the frequency of the (m-2)-bit pattern i_1...i_{m-2}.
```

#### 3. Calculate:

$$\psi_{m}^{2} = \frac{2^{m}}{n} \sum_{i_{1} \dots i_{m}} \left( v_{i_{1} \dots i_{m}} - \frac{n}{2^{m}} \right)^{2} = \frac{2^{m}}{n} \sum_{i_{1} \dots i_{m}} v_{i_{1} \dots i_{m}}^{2} - n$$

$$\psi_{m-1}^{2} = \frac{2^{m-1}}{n} \sum_{i_{1} \dots i_{m-1}} \left( v_{i_{1} \dots i_{m-1}} - \frac{n}{2^{m-1}} \right)^{2} = \frac{2^{m-1}}{n} \sum_{i_{1} \dots i_{m-1}} v_{i_{1} \dots i_{m-1}}^{2} - n$$

$$\psi_{m-2}^{2} = \frac{2^{m-2}}{n} \sum_{i_{1} \dots i_{m-2}} \left( v_{i_{1} \dots i_{m-2}} - \frac{n}{2^{m-2}} \right)^{2} = \frac{2^{m-2}}{n} \sum_{i_{1} \dots i_{m-2}} v_{i_{1} \dots i_{m-2}}^{2} - n$$

```
def _psi_sq_mv1(block_size: int, sequence_size: int, padded_sequence:
    numpy.ndarray) -> float:
    """

Compute the Psi-Squared statistics from the NIST paper.
:param block_size: the size of the block
:param sequence_size: the size of the sequence of bits
:param padded_sequence: the original sequence once padded
:return: the float value of Psi-Squared statistics
"""

# Count the patterns
counts: numpy.ndarray = numpy.zeros(2 ** block_size, dtype=int)
for i in range(2 ** block_size):
    pattern: numpy.ndarray = (i >> numpy.arange(block_size, dtype=int)) &
1
```

```
counts[i] = SerialTest._count_pattern(pattern, padded_sequence,
sequence_size)
# Compute Psi-Squared statistics and return it
psi_sq_m: float = numpy.sum(counts[:] ** 2)
psi_sq_m *= (2 ** block_size) / sequence_size
psi_sq_m -= sequence_size
return psi_sq_m
```

```
# Compute Psi-Squared statistics
psi_sq_m_0: float = self._psi_sq_mv1(self._pattern_length, bits.size,
padded_bits)
psi_sq_m_1: float = self._psi_sq_mv1(self._pattern_length - 1,
bits.size, padded_bits)
psi_sq_m_2: float = self._psi_sq_mv1(self._pattern_length - 2,
bits.size, padded_bits)
```

#### 4. Calculate:

$$\nabla \psi_{m}^{2} = \psi_{m}^{2} - \psi_{m-1}^{2}, \text{ and } P-value 1 = igamc} \left(2^{m-2}, \nabla \psi_{m}^{2}\right) \text{ and }$$

$$\nabla^{2} \psi_{m}^{2} = \psi_{m}^{2} - 2\psi_{m-1}^{2} + \psi_{m-2}^{2}. P-value 2 = igamc} \left(2^{m-3}, \nabla^{2} \psi_{m}^{2}\right).$$

```
delta_1: float = psi_sq_m_0 - psi_sq_m_1
delta_2: float = psi_sq_m_0 - (2 * psi_sq_m_1) + psi_sq_m_2
# Compute the scores (P-values)
score_1: float = scipy.special.gammaincc(2 ** (self._pattern_length - 2),
delta_1 / 2.0)
score_2: float = scipy.special.gammaincc(2 ** (self._pattern_length - 3),
delta_2 / 2.0)
# Return result
if score_1 >= self.significance_value and score_2 >= self.significance_value:
    return Result(self.name, True, numpy.array([score_1, score_2]))
return Result(self.name, False, numpy.array([score_1, score_2]))
```

#### **Results and Inferences:**

```
Python RNG

Python RNG (Secrets) - FAILED - score: 0.0
```

True RNG

The result implies that there was a degree of randomness below the accepted level of significance (P>0.01). As we have selected m=4 in the test, the minimum input size n becomes  $floor(log_2^{128})-2=5$ . As all inputs are greater than 128, all of them are eligible for the test. Failing the test and score=0, would imply that uniformity present in the sequence is not as expected in a random sequence. The score=0 does not imply that P-value1=0 and P-value2=0. In step-4 of implementation, result is TRUE iff both P-value > 0.01. If one or both P-values are less than 0.01, result=FALSE and score=0.

P-value would be less if the  $\nabla \psi_m^2$  and  $\nabla^2 \psi_m^2$  had been large. [2]

# Approximate Entropy Test:

Similar to the *Serial Test*, this test finds the frequency of all possible overlapping m-bit patterns across the entire sequence. Instead of comparing occurrence of  $2^m m$  bit overlapping patterns to random sequence, this test compares the frequency of overlapping blocks of consecutive lengths (m & m+1 bits) with the expected result of random sequence.

*The variables used in the Test:* 

m- length of bits in the first block. m+1 is the length of a consecutive block. n- length of the bitstream

 $\varepsilon$  - Sequence of bits generated by the random number generator(RNG)

*Test Statistic*:  $\chi^2$  (obs), calculates how well the value of ApEn(m) (observed) matches the value of the expected value. [ApEn(m) =  $\phi^{(m)} - \phi^{(m+1)}$ , where  $\phi^{(m)} = \sum \pi_i \log \pi_i$ ]

Decision Rule: If the P-value < 0.01, then the sequence is classified as non-random

## *Test Implementation:*

- 1. Form an augmented sequence,  $\varepsilon$ ', by padding first m-1 bits of  $\varepsilon$  at the end.
- 2. A frequency count is made of the *n* overlapping blocks (e.g., if a block containing  $\varepsilon_j$  to  $\varepsilon_{j+m-1}$  is examined at time j, then the block containing  $\varepsilon_{j+1}$  to  $\varepsilon_{j+m}$  is examined at time j+1). Let  $C_m^i$  represent the count of the possible *m-bit* ((*m+1*)-*bit*) values, where *i* is the *m-bit* value.

- 3. Compute  $C_i^m = \frac{\#i}{n}$  for each value of i.
- 4. Calculate  $\phi^{(m)} = \sum_{i=0}^{2^m 1} \pi_i \log \pi_i$ , where,  $\pi_i = C_j^3$  and  $j = \log_2 i$
- 5. Repeat steps 1-4 and by replacing m to m+1.

```
# Define Phi-m statistics list
                              phi m: [] = []
                             for iteration in range(blocks_length, blocks length + 2):
                                             # Compute the padded sequence of bits
                                             padded_bits: numpy.ndarray = numpy.concatenate((bits,
                                             bits[0:iteration - 1]))
                                             # Compute the frequency count
                                             counts: numpy.ndarray = numpy.zeros(2 ** iteration, dtype=int)
                                             for i in range(2 ** iteration):
                                                        count: int = 0
                                                        for j in range(bits.size):
                                                            if self. pattern to int(padded bits[j:j + iteration]) ==i:
                                                                                           count += 1
                                                             counts[i] = count
                                             # Compute C-i as the average of counts on the number of bits
                                             c i: numpy.ndarray = counts[:] / float(bits.size)
                                             # Compute Phi-m based on C-i
                                             phi_m.append(numpy.sum(c_i[c_i > 0.0] * numpy.log((c_i[c_i > 0.0]) * numpy.log((c_i[c_i
```

6. Compute the test statistic:

$$\gamma^2 = 2n\lceil \log 2 - ApEn(m) \rceil$$
, where  $ApEn(m) = \phi^{(m)} - \phi^{(m+1)}$ 

And, compute *P-value* using,

$$P$$
-value = **igamc** $(2^{m-1}, \frac{\chi^2}{2})$ 

```
# Compute Chi-Square from the computed statistics
chi_square: float = 2 * bits.size * (math.log(2) - (phi_m[0] -
phi_m[1]))
# Compute the score (P-value)
score: float = scipy.special.gammaincc(2 ** (blocks_length - 1),
    (chi_square / 2.0))
# Return result
```

```
if score >= self.significance_value:
          return Result(self.name, True, numpy.array(score))
return Result(self.name, False, numpy.array(score))
```

#### **Results and Inferences:**

```
Python RNG

Python RNG (Secrets) - FAILED - score: 0.0

PUF RNG

True RNG
```

For the test to run, minimum input size is given by  $m < floor(log_2n) - 5$ . We have set the m=4 for this test. Therefore n > 1024 which is true for all RNGs. The reason for failing is that the P-values < 0.01 for all tests which can be caused by large values of ApEn(m) as smaller values of ApEn(m) imply strong regularity and thus more random nature [2].

## Cumulative Sums (Cusum) Test

The purpose of the test is to determine whether the cumulative sum of the partial sequences occurring in the tested sequence is too large or too small relative to the expected behavior of that cumulative sum for random sequences. The test is based on the principle of random walks and looks for deviations in the cumulative sum of the binary sequence from what would be expected for a random sequence.

The variables used in the Test:

```
n = The length of the bit string.
```

 $\varepsilon$  = The sequence of bits as generated by the RNG or PRNG being tested; this exists as a global structure at the time of the function call;  $\varepsilon = \varepsilon 1, \varepsilon 2, ..., \varepsilon n$ .

Mode = A switch for applying the test either forward through the input sequence (mode = 0) or backward through the sequence (mode = 1).

Test Statistic: (Z) The largest excursion from the origin of the cumulative sums in the corresponding (-1, +1) sequence

Decision Rule: If the P-value obtained is  $\geq 0.01$  (P-value = 0.615, 0, 0, 0), the conclusion is that the sequence is random and the test passes

## Test Implementation:

1. Form a normalized sequence: The zeros and ones of the input sequence ( $\epsilon$ ) are converted to values Xi of -1 and +1 using Xi =  $2\epsilon i - 1$ .

For example, if  $\varepsilon = 1011010111$ , then X = 1, (-1), 1, 1, (-1), 1, (-1), 1, 1, 1

2. Compute partial sums Si of successively larger subsequences, each starting with X1 (if mode = 0) or Xn (if mode = 1).

Mode = 0 (forward)	Mode = 1 (backward)
S1 = X1	S1 = Xn
S2 = X1 + X2	S2 = Xn + Xn-1
S3 = X1 + X2 + X3	S3 = Xn + Xn-1 + Xn-2
Sk = X1 + X2 + X3 + + Xk	Sk = Xn + Xn-1 + Xn-2 + + Xn-k+1
Sn = X1 + X2 + X3 + + Xk + + Xn	Sn = Xn + Xn-1 + Xn-2 + + Xk-1 + + X1

That is, Sk = Sk-1 + Xk for mode 0, and Sk = Sk-1 + Xn-k+1 for mode 1.

```
# Compute the partial sum with forward (mode 0) and backward
(mode 1) modes and record the largest excursion
forward_sum: int = 0
  backward_sum: int = 0
  backward_max: int = 0
for i in range(bits_copy.size):
forward_sum += bits_copy[i]
backward_sum += bits_copy[bits_copy.size - 1 - i]
forward_max = max(abs(forward_sum), forward_max)
backward_max = max(abs(backward_sum), backward_max)
```

3. Compute the test statistic  $z = \max 1 \le k \le n|Sk|$ , where  $\max 1 \le k \le n|Sk|$  is the largest of the absolute values of the partial sums Sk.

4. Compute 
$$P$$
-value =  $1 - \sum_{k=\left(\frac{-n}{z}-1\right)/4}^{\left(\frac{n}{z}-1\right)/4} \left[ \Phi\left(\frac{(4k+1)z}{\sqrt{n}}\right) - \Phi\left(\frac{(4k-1)z}{\sqrt{n}}\right) \right] + \sum_{k=\left(\frac{-n}{z}-3\right)/4}^{\left(\frac{n}{z}-1\right)/4} \left[ \Phi\left(\frac{(4k+3)z}{\sqrt{n}}\right) - \Phi\left(\frac{(4k+1)z}{\sqrt{n}}\right) \right]$ 

If the computed P-value is < 0.01, then conclude that the sequence is non-random. Otherwise, conclude that the sequence is random.

```
# Compute the scores (P-Values)

score_1: float = self._compute_p_value(bits_copy.size,
forward_max)
    score_2: float = self._compute_p_value(bits_copy.size,
backward_max)

# Return result
    if score_1 >= self.significance_value and score_2 >=
self.significance_value:
        return Result(self.name, True, numpy.array([score_1, score_2]))
return Result(self.name, False, numpy.array([score_1, score_2]))
```

```
@staticmethod
def _compute_p_value(sequence_size: int, max_excursion: int) ->
float:
    """
Compute P-Value given the sequence size and the max excursion.
:param sequence_size: the length of the sequence of bits
:param max_excursion: the max excursion backward or forward
:return: the computed float P-Value
"""
```

## 5. Input Size Recommendation

It is recommended that each sequence to be tested consist of a minimum of 100 bits (i.e.,  $n \ge 100$ ). Those random no. sequences have less than 100 bits then the test will be failed.

```
# Execute first sum
sum_a: float = 0.0
start_k: int = int(math.floor((((float(-sequence_size) /
max_excursion) + 1.0) / 4.0)))
end_k: int = int(math.floor((((float(sequence_size) /
max_excursion) - 1.0) / 4.0)))
for k in range(start_k, end_k + 1):
c: float = 0.5 * math.erfc(-(((4.0 * k) + 1.0) * max_excursion) /
math.sqrt(sequence_size) * math.sqrt(0.5))
d: float = 0.5 * math.erfc(-(((4.0 * k) - 1.0) * max_excursion) /
math.sqrt(sequence_size) * math.sqrt(0.5))
sum_a = sum_a + c - d
# Execute second sum
sum_b: float = 0.0
start_k = int(math.floor((((float(-sequence_size) / max_excursion))
```

```
- 3.0) / 4.0)))
end_k = int(math.floor((((float(sequence_size) / max_excursion) -
1.0) / 4.0)))
for k in range(start_k, end_k + 1):
c: float = 0.5 * math.erfc(-(((4.0 * k) + 3.0) * max_excursion) /
math.sqrt(sequence_size) * math.sqrt(0.5))
d: float = 0.5 * math.erfc(-(((4.0 * k) + 1.0) * max_excursion) /
math.sqrt(sequence_size) * math.sqrt(0.5))
sum_b = sum_b + c - d
# Return value
return 1.0 - sum_a + sum_b
```

#### **Results and Inferences:**

```
Python RNG - FAILED - score: 0.0
Python RNG (Secrets)
```

The result implies that there was a degree of randomness below the accepted level of significance (P>0.01). Also, the test requires  $n \ge 100$  (in the test code).

The excursions of the random walk from zero are near zero, indicating that the tested sequence is likely random. For certain types of non-random sequences, the excursions of this random walk from zero will be large and the test will be failed.

If the Cusum Test fails, it indicates that the binary sequence may be non-random and further investigation is necessary to identify the source of the non-randomness.

```
PUF RNG - PASSED - score: 0.615

True RNG - FAILED - score: 0.0
```

For PUF RNG, n=30,000 and for True RNG, n=131,076. As the test requires a minimum n=1,00 they are eligible for the test and P value should be P>0.01 to pass the test.

## Random Excursions Test

The focus of this test is the number of cycles having exactly K visits in a cumulative sum random walk. The cumulative sum random walk is derived from partial sums after the (0,1) sequence is transferred to the appropriate (-1, +1) sequence.

A cycle of a random walk consists of a sequence of steps of unit length taken at random that begin at and return to the origin. The purpose of this test is to determine if the number of visits to a particular state within a cycle deviates from what one would expect for a random sequence.

The variables used in the Test:

```
n = The length of the bit string.
```

 $\varepsilon$  = The sequence of bits as generated by the RNG or PRNG being tested; this exists as a global structure at the time of the function call;  $\varepsilon = \varepsilon 1, \varepsilon 2, ..., \varepsilon n$ .

Test Statistic:  $\chi$  2 (obs), For a given state x, a measure of how well the observed number of state visits within a cycle match the expected number of state visits within a cycle, under an assumption of randomness.

The reference distribution for the test statistic is the  $\chi$  2 distribution.

Decision Rule: If the computed P-value is < 0.01 (P-value =0.005, 0.683, 0.683, 0.683), then conclude that the sequence is non-random. Otherwise, conclude that the sequence is random and in a random case then the sequence of no. will pass the test.

## Test Implementation:

1. Form a normalized (-1, +1) sequence X: The zeros and ones of the input sequence ( $\epsilon$ ) are changed to values of -1 and +1 via Xi =  $2\epsilon i - 1$ .

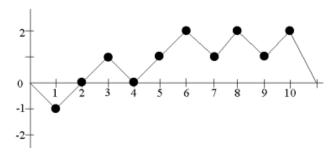
For example, if  $\varepsilon = 0110110101$ , then n = 10 and X = -1, 1, 1, -1, 1, 1, -1, 1, 1.

```
# Copy the bits to a new array
bits_copy: numpy.ndarray = bits.copy()
# Convert all the zeros in the array to -1
bits_copy[bits_copy == 0] = -1
```

2. Compute the partial sums Si of successively larger subsequences, each starting with X1. Form the set  $S = \{Si\}$ .

Form a new sequence S' by attaching zeros before and after the set S. That is, S' = 0, s1, s2, ..., sn, 0

For the example in this section, S' = 0, -1, 0, 1, 0, 1, 2, 1, 2, 1, 2, 0. The resulting random walk is shown below.



Example Random Walk (S')

3. Let J = the total number of zero crossings in S', where a zero crossing is a value of zero in S' that occurs after the starting zero. J is also the number of cycles in S', where a cycle of S' is a subsequence of S'consisting of an occurrence of zero, followed by no-zero values, and ending with another zero. The ending zero in one cycle may be the beginning zero in another cycle. The number of cycles in S' is the number of zero crossings. If J < 500, discontinue the test6.

For example in this section, if  $S' = \{0, -1, 0, 1, 0, 1, 2, 1, 2, 1, 2, 0\}$ , then J = 3 (there are zeros in positions 3, 5 and 12 of S'). The zero crossings are easily observed in the above plot. Since J = 3, there are 3 cycles, consisting of  $\{0, -1, 0\}$ ,  $\{0, 1, 0\}$  and  $\{0, 1, 2, 1, 2, 1, 2, 0\}$ .

4. For each cycle and for each non-zero state value x having values  $-4 \le x \le -1$  and  $1 \le x \le 4$ , compute the frequency of each x within each cycle.

For the example in this section, the first cycle has one occurrence of -1, the second cycle has one occurrence of 1, and the third cycle has three occurrences each of 1 and 2. This can be visualized.

5. For each of the eight states of x, compute vk(x) = the total number of cycles in which state x occurs exactly k times among all cycles, for k = 0, 1, ..., 5 (for k = 5, all frequencies  $\geq 5$  are 5 stored in v5(x)). Note that  $\sum v k(x) = J$ .

```
# Generate the padded cumulative sum of the array of -1, 1
   sum prime: numpy.ndarray = numpy.concatenate((numpy.array([0]),
numpy.cumsum(bits_copy), numpy.array([0]))).astype(int)
# Compute the cycles iterating over each position of S'
(sum prime) and define the first cycle
cycles: [] = []
 cycle: [] = [0]
for index, _ in enumerate(sum_prime[1:]):
# Once a zero crossing is found add all the non zero elements of
S' to the cycle
# Else wrap up the cycle and start a new cycle
if sum prime[index] != 0:
cycle += [sum_prime[index]]
else:
cycle += [0]
cycles.append(cycle)
cycle: [] = [0]
# Append the last cycle
cycles.append(cycle)
    # Compute the size of the cycles list
cycles size: int = len(cycles)
# Setup frequencies table (Vk(x))
frequencies table: dict = {
-4: numpy.zeros(6, dtype=int),
-3: numpy.zeros(6, dtype=int),
-2: numpy.zeros(6, dtype=int),
-1: numpy.zeros(6, dtype=int),
1: numpy.zeros(6, dtype=int),
2: numpy.zeros(6, dtype=int),
3: numpy.zeros(6, dtype=int),
4: numpy.zeros(6, dtype=int),
# Count occurrences
for value in frequencies table.keys():
for k in range(frequencies table[value].size):
  count: int = 0
# Count how many cycles in which x occurs k times
```

```
for cycle in cycles:
# Count how many times the value used as key of the table occurs
in the current cycle
occurrences: int = numpy.count_nonzero(numpy.array(cycle) ==
value)
# If the value occurs k times, increment the cycle count
    if 5 > k == occurrences:
count += 1
elif occurrences >= 5:
count += 1
frequencies_table[value][k] = count
```

6. For each of the eight states of x, compute the test statistic

$$\chi^{2}(obs) = \sum_{k=0}^{5} \frac{(v_{k}(x) - J\pi_{k}(x))^{2}}{J\pi_{k}(x)}$$

7. For each state of x, compute P-value = igamc(5/2,  $\chi 2$  (obs) 2). Eight P-values will be produced. 5 4.333033 For the example when x = 1,

P-value= 
$$igamc\left(\frac{5}{2}, \frac{4.333033}{2}\right) = 0.502529.$$

Since the P-value obtained is  $\ge 0.01$  (P-value = 0.502529), the conclusion is that the sequence is random.

Note that if  $\chi$  2 (obs) were too large, then the sequence would have displayed a deviation from the theoretical distribution for a given state across all cycles.

8. It is recommended that each sequence to be tested consist of a minimum of 1,000,000 bits (i.e.,  $n \ge 10^6$ ). Those random no. sequences have less than 1,000,000 bits then the test will be failed.

```
# Compute the scores (P-values)
scores: [] = []
```

```
for value in frequencies table.keys():
  # Compute Chi-Square for this value
chi square: float = numpy.sum(((frequencies table[value][:] -
(cycles_size * (self._probabilities_xk[abs(value) - 1][:]))) ** 2)
/ (cycles size * self. probabilities xk[abs(value) - 1][:]))
# Compute the P-value for this value
score: float = scipy.special.gammaincc(5.0 / 2.0, chi_square /
2.0)
scores.append(score)
# Return result
if all(score >= self.significance value for score in scores):
return Result(self.name, True, numpy.array(scores))
def is_eligible(self,
bits: numpy.ndarray) -> bool:
Overridden method of Test class: check its docstring for further
information.
# This test is always eligible for any sequence
return True
```

#### **Results and Inferences:**

Python RNG - FAILED - score: 0.683 Python RNG (Secrets) - FAILED - score: 0.683

The result implies that there was a degree of randomness below the accepted level of significance (P>0.01). Also, the test requires  $n \ge 100$  (in the test code).

Random Excursions Test checks whether the number of visits to a particular state within a cycle deviates from what one would expect for a random sequence. Passing the test suggests that the sequence is consistent with being random and exhibits expected behavior of a random walk.

PUF RNG - FAILED - score: 0.005 True RNG - FAILED - score: 0.63

For PUF RNG, n=30,000 and for True RNG, n=131,076. As the test requires a minimum  $n \ge 10^6$  they are eligible for the test and P value should be P>0.01 to pass the test.

## Random Excursions Variant Test

The focus of this test is the total number of times that a particular state is visited (i.e., occurs) in a cumulative sum random walk. The purpose of this test is to detect deviations from the expected number of visits to various states in the random walk. This test is actually a series of eighteen tests (and conclusions), one test and conclusion for each of the states: -9, -8, ..., -1 and +1, +2, ..., +9.

The variables used in the Test:

n =The length of the bit string.

 $\varepsilon$  = The sequence of bits as generated by the RNG or PRNG being tested; this exists as a global structure at the time of the function call;  $\varepsilon = \varepsilon 1, \varepsilon 2, ..., \varepsilon n$ .

*Test Statistic*:  $(\xi)$  For a given state x, the total number of times that the given state is visited during the entire random walk as determined.

If  $\xi$  is distributed as normal, then  $|\xi|$  is distributed as half normal.) If the sequence is random, then the test statistic will be about 0. If there are too many ones or too many zeroes, then the test statistic will be large

Decision Rule: If the computed P-value is < 0.01, then conclude that the sequence is non-random. Otherwise, conclude that the sequence is random.

## *Test Implementation:*

1. Form the normalized (-1, +1) sequence X in which the zeros and ones of the input sequence ( $\epsilon$ ) are converted to values of -1 and +1 via X = X1, X2, ..., Xn, where Xi =  $2\epsilon i - 1$ .

For example, if  $\varepsilon = 0110110101$ , then n = 10 and X = -1, 1, 1, -1, 1, 1, -1, 1, 1.

2. Compute partial sums Si of successively larger subsequences, each starting with x1. Form the set  $S = \{Si\}$ 

```
S1 = X1

S2 = X1 + X2

S3 = X1 + X2 + X3

.

.

Sk = X1 + X2 + X3 + ... + Xk

.

.

.

.

.

.

.

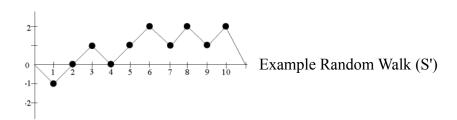
.

.

.
```

3. Form a new sequence S' by attaching zeros before and after the set S. That is, S' = 0, s1, s2, ..., sn, 0. For example,

S' = 0, -1, 0, 1, 0, 1, 2, 1, 2, 1, 2, 0. The resulting random walk is shown below



```
def execute(self,
bits: numpy.ndarray) -> Result:
Overridden method of Test class: check its docstring for further
information.
.....
# Copy the bits to a new array
bits copy: numpy.ndarray = bits.copy()
# Convert all the zeros in the array to -1
bits copy[bits copy == 0] = -1
sum_prime: numpy.ndarray = numpy.concatenate((numpy.array([0]),
numpy.cumsum(bits_copy), numpy.array([0]))).astype(int)
       # Count the number of cycles in S' (sum_prime)
cycles_size: int = numpy.count_nonzero(sum_prime[1:] == 0)
# Generate the counts of offsets
 unique, counts = numpy.unique(sum_prime[abs(sum_prime) < 10],</pre>
return counts=True)
 # Compute the scores (P-values)
scores: [] = []
 for key, value in zip(unique, counts):
```

$$P-value = erfc \left( \frac{|\xi(x) - J|}{\sqrt{2J(4|x|-2)}} \right)$$

4. For each  $\xi(x)$ , compute P-value . Eighteen P-values are computed.

If the computed P-value is < 0.01, then conclude that the sequence is non-random. Otherwise, conclude that the sequence is random.

5. It is recommended that each sequence to be tested consist of a minimum of 1,000,000 bits (i.e.,  $n \ge 106$ )

```
# Compute the scores (P-values)
scores: [] = []
for key, value in zip(unique, counts):
# Compute the P-value for this value (if not zero)
if key != 0:
scores.append(abs(value - cycles_size) / math.sqrt(2.0 *
cycles size * ((4.0 * abs(key)) - 2.0)))
# Return result
if all(score >= self.significance value for score in scores):
return Result(self.name, True, numpy.array(scores))
return Result(self.name, False, numpy.array(scores))
def is eligible(self,
      bits: numpy.ndarray) -> bool:
Overridden method of Test class: check its docstring for further
information.
# This test is always eligible for any sequence
return True
```

#### **Results and Inferences:**

```
Python RNG - PASSED - score: 0.17
Python RNG (Secrets) - PASSED - score: 0.17
```

The result implies that there was a degree of randomness below the accepted level of significance (P>0.01). Also, the test requires  $n \ge 10^6$  (in the test code).

It is important to note that passing the Random Excursions Variant Test does not guarantee that the sequence is completely random, as no test can prove randomness with absolute certainty. Instead, it provides evidence that the sequence is likely to be random and can be used with some level of confidence in applications that require randomness.

```
PUF RNG - PASSED - score: 1.208
True RNG - FAILED - score: 0.0
```

For PUF RNG, n=30,000 and for True RNG, n=131,076. As the test requires a minimum  $n \ge 10^6$  they are eligible for the test and If the computed P-value is < 0.01 (P-value = 1.208, 0.17, 0.17, 0.0), then conclude that the sequence is non-random. Otherwise, conclude that the sequence is random and in a random case then the sequence of no. will pass the test.

# Non-Overlapping Template Test:

We have derived the code from the GitHub repository; the essential parts include checking the eligibility of running the test based on the size of the dataset. The principle involves defining the templates to derive data points (namely, the occurrence of each sequence of m bits in a non-overlapping fashion. Due to this approach, we shift the pointer with m bits every time we find a matching sequence, and if not, then just by 1 bit. The time taken by code thus reduces as repeating patterns are not checked due to omission by m bits. For all the templates of a given length of m bits, we check the occurrences of that template and then code to get the p-value of the provided data.

The variables used in the Test:

n = The length of the bit string.m = size of the template.

#### Step 1:

#### Step 2: checking eligibility

## Step 3:(declaring initial values and base case)

Step 4: Checking for all templates throughout bitstring and incrementing as per the result:

```
block: numpy.ndarray = bits[i * substring_bits_length:(i + 1) * substring_bits_length]
# Define counting variables
position: int = 0

count: int = 0

# Count the matches in the block with the chosen template
while position < (substring_bits_length - b_template.size):
    if (block[position:position + b_template.size] == b_template).all():
        position += b_template.size
        count += 1
    else:
        position += 1

matches[i] = count</pre>
```

Step 5: Statistical computation:

```
sigma: float = substring_bits_length * ((1.0 / float(2 ** b_template.size)) - (float((2 * b_template.size) - 1) / float(2 ** (2 * b_template.size))))

# Compute Chi-square

chi_square: float = float(numpy.sum(((matches[:] - mu) ** 2) / (sigma ** 2)))

# If Chi-square is zero, fail the test

if chi_square != 0:

# Compute the score (P-value)

score: float = scipy.special.gammaincc(self._blocks_number / 2.0, chi_square / 2.0)

# Return result

if score >= self.significance_value:

return Result(self.name, True, numpy.array(score))

return Result(self.name, Ealse_numpy.array(score))
```

### **Results And Inferences**

PYTHON RANDOM- Passed: score-0.063: p-value greater than 0.01.

The test passed, indicating that for blocks of size 2 to 8, we can say with a 99% confidence interval that a particular sequence doesn't repeat often.

PYTHON-SECRETS- Passed:score-0.355:

It can be inferred that python secrets implements a much more random code then that provided by random function.

TRUE RANDOM- Passed:-0.347:

Using the non-overlapping test we can agree with the online source that it generates fairly randomised code.

PUF GENERATOR- Passed: score-0.989: p-value greater than 0.01.

Here it's evident that the PUF code is much more random than the python-random number generator as the score is much larger.

# Overlapping Matching Template Test:

The non-overlapping test function has a very big demerit that it does'nt consider repeating sequences due to large omission but now it takes more time to analyse the dataset as we check all subsequences of m bits in the entire data.

The code involves checking eligibility of the dataset based on size, it consumes a large amount of data but generally we are constrained by availability of space and fast servers. The code involves defining the templates to be used to derive data points (namely occurrence of each sequence of m bits in an overlapping fashion). With this approach, we shift the pointer with 1 bit each and every time irrespective of the result of comparing. The accuracy of code is much better then previous approach as a result of more checks. For all the templates of given length of m bits we check the no of occurrences of that template and then code to get the p-value of the given data. The chi-squared random variable is used to compute the statistics.

Here only the step 4 changes primarily: the pointer

```
matches_distributions: numpy.ndarray = numpy.zeros(self._freedom_degrees + 1, dtype=int)
for i in range(self._blocks_number):
    # Define the block at the current index
    block: numpy.ndarray = bits[i * self._substring_bits_length:(i + 1) * self._substring_bits_length]
    # Define counting variable
    count: int = 0
# Count the matches in the block with respect to the given template
for position in range(self._substring_bits_length - self._template_bits_length):
    if (block[position:position + self._template_bits_length] == b_template).all():
        count += 1
matches distributions[min(count, self. freedom_degrees)] += 1
```

#### Results and inferences

PYTHON RANDOM - Failed: score-0.00: p-value less than 0.01.

We get to know that python generator is not random but surely a PRNG as it has overlapping sequences in a repeating fashion and so does not pass it.

```
PYTHON SECRETS - Failed:-0.00: p-value less than 0.01.
```

Python secrets implement a much more random code than that provided by random function but still it is not completely random as there are repeated sequences which are responsible for failure.

#### PUF-GENERATED ALGORITHM - Not applicable

Not applicable owing to insufficient data set and hardware.

TRUE RANDOM - Not applicable

Not applicable owing to insufficient data set and hardware. Required data size for successful operation of test:

n	L	$Q = 10 \bullet 2^L$
≥ 387,840	6	640
≥ 904,960	7	1280
≥ 2,068,480	8	2560
≥ 4,654,080	9	5120
≥ 10,342,400	10	10240
≥ 22,753,280	11	20480
≥ 49,643,520	12	40960
≥ 107,560,960	13	81920
≥ 231,669,760	14	163840
≥ 496,435,200	15	327680
≥ 1,059,061,760	16	655360

## Maurer's Universal Test:

The number of bits between matching patterns—a statistic related to the length of a compressed sequence—is the test's primary concern. A series that may be compressed significantly is regarded as non-random. One of the best tests follows a semi-log-normal distribution to study the cryptographic use of the device. The steps involved in performing the test are:

A test segment made up of K L-bit non-overlapping blocks, and an initialization segment made up of Q L-bit non-overlapping blocks make up the n-bit sequence. At the end of the row, bits that do not make up a whole L-bit block are discarded.

The test is initialized using the first Q blocks.

A table is generated using the initialization segment for each potential L-bit value.

Determine the number of blocks since the start of the test segment by looking at each of the K blocks.

last instance of the identical L-bit block. Substitute the location of the current block for the value in the table, and add the estimated distance between occurrences using the log2 factor. Finally, computing the p-value with the level of significance of 0.01.

### Step 1: Initialising the test variables

```
ef __init__(self):
  # Define specific test attributes
  # Note: tables from https://static.aminer.org/pdf/PDF/000/120/333/a_universal_statistical_test
  self. sequence size min: int = 387840
  self. default pattern size: int = 6
  self. freedom degrees: int = 5
  self._substring_bits_length: int = 1062
  self. thresholds: [] = [904960, 2068480, 4654080, 10342400, 22753280, 49643520, 107560960, 2316
  self._expected_value_table: [] = [0, 0.73264948, 1.5374383, 2.40160681, 3.31122472, 4.25342659
  self._variance_table: [] = [0, 0.690, 1.338, 1.901, 2.358, 2.705, 2.954, 3.125, 3.238, 3.311,
  # Define cache attributes
  self. last bits size: int = -1
  self._pattern_length: int = -1
  self. blocks number: int = -1
  self._q_blocks: int = -1
  self._k_blocks: int = -1
  # Generate base Test class
  super(MaurersUniversalTest, self).__init__("Maurers Universal", 0.01)
```

## Step 2:

```
if self._last_bits_size == -1 or self._last_bits_size != bits.size:
   # Compute the pattern size
   pattern_length: int = self._default_pattern_size
    for threshold in self._thresholds:
        if bits.size >= threshold:
           pattern_length += 1
    # Split the data into Q and K blocks
   blocks_number: int = int(bits.size // pattern_length)
   q blocks: int = 10 * (2 ** pattern length)
    k_blocks: int = blocks_number - q_blocks
   self._last_bits_size = bits.size
    self._pattern_length = pattern_length
   self._blocks_number = blocks_number
   self. q blocks = q blocks
    self._k_blocks = k_blocks
else:
   pattern_length: int = self._pattern_length
   blocks_number: int = self._blocks_number
   q_blocks: int = self._q_blocks
   k blocks: int = self. k blocks
```

Step 3: Computing statistic for all values of L

```
fn: float = computed_sum / k_blocks

# Compute magnitude

magnitude: float = abs((fn - self._expected_value_table[pattern_length]) / ((math.sqrt(self._variance_table[pattern_length])) * math.sqrt(2)))

# Compute the score (P-value)

score: float = math.erfc(magnitude)

# Return result

if score >= self.significance_value:
    return Result(self.name, True, numpy.array(score))

return Result(self.name, False, numpy.array(score))
```

Step 4: Checking eligibility with size

### Results and inferences

PYTHON RANDOM - Failed: score-0.01 {expected(f(n) > L)}

Failing the Maurer test implies that the code produces the repetition of its bit sequences at regular intervals and so can easily be compressed, that is the distance between repetitions is small enough and so not a good Random generator.

PYTHON SECRETS - Failed:-0.01: p-value is equal to 0.01.

Failing the Maurer test implies that the code produces the repetition of its bit sequences at regular intervals and so can easily be compressed, so we cannot rely on it as a strong RNG.

TRUE RANDOM - Not applicable

Not applicable owing to insufficient data set and hardware.

PUF GENERATED ALGORITH - Not applicable

Not applicable owing to insufficient data set and hardware.

# Longest Run of Ones in a Block Test

A statistical test is used to measure the randomness of a number. It measures the longest consecutive subsequence of ones and compares it with the theoretically expected length of the longest run of ones in a random sequence. The true random sequence should have less deviation from expected results (i.e., larger value  $\chi^2$ ).

### Test Description:

1. The sequence of length n is divided into blocks of length M, i.e., M blocks each of size N(n = N\*M). The values of M and K are selected corresponding to the value of n.

```
# Set the block size depending on the input sequence length block_size: int = 10000
    if bits.size < 6272:
         block_size: int
    elif bits.size < 750000:
    | block_size: int = 128
# Set the block number and K depending on the block size
    blocks_number: int = 75
    if block_size == 8:
         k: int = 3
blocks_number: int = 16
         k: int = 5
blocks_number: int = 49
    # Save in the cache
self._last_bits_size = bits.size
    self._block_size = block_size
    self._blocks_number = blocks_number
    self. k = k
else:
    blocks_number: int = self._blocks_number
    k: int = self._k
```

2. We divide the blocks into k+1 classes based on the length of the longest run of ones and assign each block to a class.

The length of the longest run of ones in observed data against expected data is measured by  $\chi^2$ -distribution with K deg  $X^2 = \sum_{i=0}^{K} \frac{(v_i - N\pi_i)^2}{N\pi_i}$ 

```
chi_square: float = 0.0
for i in range(k + 1):
    chi_square += ((frequencies[i] - blocks_number * self._probabilities(block_size, i)) ** 2)
    (blocks_number * self._probabilities(block_size, i))
```

where  $\pi_i$  (i = 0,..., K) are the probabilities associated with K, and M and has been taken from [1].

The expected probability that  $v \le m$ , given the number of ones in a block, r, is expressed as

$$P(v \le m \mid r) = \frac{1}{\binom{M}{r}} \sum_{j=0}^{U} (-1)^{j} \binom{M-r+1}{j} \binom{M-j(m+1)}{M-r}$$

$$P(v \le m) = \sum_{r=0}^{M} \binom{M}{r} P(v \le m \mid r) \frac{1}{2^{M}}.$$

$$P(v \le m) = \sum_{r=0}^{M} \binom{M}{r} P(v \le m \mid r) \frac{1}{2^{M}}.$$

K=3, M=8

classes	probabilities
{v≤1}	$\pi_0 = 0.2148$
{v=2}	$\pi_1 = 0.3672$
{v=3}	$\pi_2 = 0.2305$
{v≥4}	$\pi_3 = 0.1875$

K=5 M=128

1 3, W 120		
classes	probabilities	
{v≤4}	$\pi_0 = 0.1174$	
{v=5}	$\pi_1 = 0.2430$	
{v=6}	$\pi_2 = 0.2493$	
{v=7}	$\pi_3 = 0.1752$	
{v=8}	$\pi_4 = 0.1027$	
{v≥9}	$\pi_5 = 0.1124$	

K=5, M=512

classes	probabilities
{v≤6}	$\pi_0 = 0.1170$
{ <b>v</b> =7}	$\pi_1 = 0.2460$
{v=8}	$\pi_2 = 0.2523$
{v=9}	$\pi_3 = 0.1755$
{v=10}	$\pi_4 = 0.1027$
{v≥11}	$\pi_5 = 0.1124$

K=5, M=1000

11 0,111 1000		
classes	probabilities	
{ν≤7}	$\pi_0 = 0.1307$	
{v=8}	$\pi_1 = 0.2437$	
{v=9}	$\pi_2 = 0.2452$	
{v=10}	$\pi_3 = 0.1714$	
{v=11}	$\pi_4 = 0.1002$	
{ν≥12}	$\pi_5 = 0.1088$	

4. Compute the p-value using the incomplete gamma function expressed as

$$\frac{\int_{X^{2}(obs)}^{\infty} e^{-u/2} u^{K/2-1} du}{\Gamma(K/2) 2^{K/2}} = igamc \left( \frac{K}{2}, \frac{X^{2}(obs)}{2} \right)$$

```
# Compute score (P-value)
score: float = scipy.special.gammaincc(k / 2.0, chi_square / 2.0)
# Return result
if score >= self.significance_value:
    return Result(self.name, True, numpy.array(score))
return Result(self.name, False, numpy.array(score))
```

5. The result of the test is determined by the P-value, if the P-value is < 0.01, then the test is passed, or otherwise return failed.

```
if size_of_block == 8:
    return [0.1174, 0.2430, 0.2493, 0.1752, 0.1027, 0.1124][index]
elif size_of_block == 128:
    return [0.1174, 0.2430, 0.2493, 0.1752, 0.1027, 0.1124][index]
elif size_of_block == 512:
    return [0.1170, 0.2460, 0.2523, 0.1755, 0.1027, 0.1124][index]
elif size_of_block == 1000:
    return [0.1307, 0.2437, 0.2452, 0.1714, 0.1002, 0.1088][index]
else:
    return [0.0882, 0.2092, 0.2483, 0.1933, 0.1208, 0.0675, 0.0727][index]
```

#### **Results and Inferences:**

```
PUF random test result - FAILED - score: 0.0
```

Failing the test and score=0 implies that the observed frequency  $v_i$  values are different from the expected values. If the score>0.01, then it is expected that the distribution of the frequency  $\pi_i$ . It is possible for a true Random Number Generator to fail some tests due to random bias, but the ratio of successful tests should be considered to term any generator as true-random or pseudo-random.

```
Python_random_ test result - PASSED - score: 0.475

Python_secret_ Test result - PASSED - score: 0.14
```

```
True random Test result - PASSED - score: 0.448
```

The default\_Python\_Generator passed the test. Although it has passed the test, a higher P-value also suggests that the evidence is not strong enough. Thereby pointing to the fact that a single test only looks into a specific property of sequence. If a PRNG passes the given test, it cannot be concluded that the generator is truly random. It only implies that it has a specific property.

# Binary Matrix Rank Test

The binary matrix rank test aims to determine the linear independence of the rows and columns of a binary matrix, i.e., it determines if the rows and columns of a binary matrix can be linearly combined to form all the other rows and columns in the matrix. This is important in many fields, including cryptography, coding theory, and communication systems.

## Test Description:

1. Divide the binary sequence into Q-bit blocks, and collect M such blocks to form M\*Q matrices. There will be N = n / (M\*Q) such matrices.

```
def __init__(self, block: numpy.ndarray, rows_number: int, columns_number: int):
    self._rows = rows_number
    self._columns = columns_number
    self._matrix = block
    self._base_rank = min(self._rows, self._columns)
```

2. Before calculating the binary rank, we perform some transformations on the input binary matrix, including row operations to transform the input binary matrix into an upper triangular binary matrix and then into a reduced row-echelon form.

```
def compute_rank(self) -> int:
   Computes the **binary rank** of the matrix.
    :return: an integer defining binary rank of the matrix.
    # Perform row operations with forward elimination
    i: int = 0
   while i < self. base rank - 1:
        if self. matrix[i][i] == 1:
            self. perform row operations(i, True)
        else:
            found = self. find unit element swap(i, True)
            if found == 1:
                self._perform_row_operations(i, True)
        i += 1
    # Perform row operations without forward elimination
    i = self. base rank - 1
   while i > 0:
        if self. matrix[i][i] == 1:
            self._perform_row_operations(i, False)
        else:
            if self. find unit element swap(i, False) == 1:
                self. perform row operations(i, False)
    # Compute the rank of the transformed matrix
    return self. compute rank()
```

3. Next, calculate the binary rank (Rℓ: ℓ=1, 2, ..., N) for each transformed matrix. This step performs only the rank computation for the transformed matrix already in row-echelon form.

4. Classify matrices in three classes, full\_rank\_matrices, minus\_rank\_matrices and remainder matrices with frequencies:

 $F_M$ = the number of matrices with rank  $R_\ell = M$ ,  $F_{M-1}$  = the number of matrices with R = M-1,  $N - F_M - F_{M-1}$  = the number of matrices remaining.

```
Compute the number of full rank, minus rank and remained
rank matrices
       full_rank_matrices: int = 0
       minus rank matrices: int = 0
       remainder: int = 0
       for i in range(blocks number):
             # Get the bits in the block and reshape them in a
2D array (the matrix)
            block: numpy.ndarray = bits[i * (self._rows_number
   self. cols number):(i + 1) *
                                         (self. rows number
self._cols_number)].reshape((self._rows_number,
self._cols_number))
           # Compute rank of the block matrix
                   matrix: BinaryMatrix = BinaryMatrix(block,
self. rows number, self. cols number)
           rank: int = matrix.compute_rank()
           # Count the result
           if rank == self. rows number:
               full rank matrices += 1
           elif rank == self._rows_number - 1:
               minus rank matrices += 1
           else:
               remainder += 1
```

5. Compute  $\chi^2$  and the P-value of the observations using the expression. If the P-value<0.01, we conclude that the given binary sequence is not random else, return random.

$$\chi^{2}(obs) = \frac{(F_{M} - 0.2888N)^{2}}{0.2888N} + \frac{(F_{M-1} - 0.5776N)^{2}}{0.5776N} + \frac{(N - F_{M} - F_{M-1} - 0.1336N)^{2}}{0.1336N}$$

$$P-value = e^{-\chi^2(obs)/2}$$
  $P-value = igamc\left(1, \frac{\chi^2(obs)}{2}\right)$ 

```
Compute Chi-square
               chi_square: float = (((full_rank_matrices
(self. full rank probability *
                                 blocks number))
                                                        2)
(self. full rank probability
                                        blocks number))
(((minus rank matrices -
                             (self. minus rank probability
blocks_number))
                        __/
                            (self._minus_rank_probability
                     2)
blocks number))
                                      (((remainder
(self._remained_rank_probability * blocks_number))
(self. remained_rank_probability * blocks_number))
# Compute the score (P-value)
       score: float = math.e ** (-chi_square / 2.0)
# Return result
       if score >= self.significance_value:
           return Result(self.name, True, numpy.array(score))
       return Result(self.name, False, numpy.array(score))
```

#### **Results and Inferences:**

```
PUF random test result - Not eligible for the test.
```

For PUF RNG, n=30,000, but as the test requires a minimum of n=38,000, they are not eligible for the test.

```
Python_random_ test result - PASSED - score: 0.30

Python_secret_ Test result - PASSED - score: 0.208

True random Test result - PASSED - score: 0.578
```

The test to randomness passes, suggesting that the data appears to be random, and there is no strong evidence of the presence of any underlying pattern or structure in the data.

It is important to note that passing a single test does not guarantee that the data is completely random. It is always possible for patterns or structures to exist in the data that are not detectable by the particular statistical test being used. Therefore, passing a statistical test for randomness should be viewed as a preliminary indication of randomness rather than definitive proof of randomness.

# Discrete Fourier Transform (Spectral) Test

The Discrete Fourier Transform Test detects any pattern in the input sequence that deviates from the expected theoretical result. We assume that a truly random sequence will have a flat power spectral density while a non-random sequence will have peaks at certain frequencies.

## Test Description:

1. Make the sequence even in length, and convert all 0's and 1's to -1 and +1

```
bits_copy: numpy.ndarray = bits.copy()
if (bits_copy.size % 2) == 1:
    bits_copy = bits_copy[:-1]
# Convert all the zeros in the array to -1
bits_copy[bits_copy == 0] = -1
```

2. Compute the DFT of the sequence of -1 and +1; the output of DFT will be frequency domain sequence in complex variables.

$$f_j = \sum_{k=1}^{n} x_k \exp(2\pi_i(k-1)j/n),$$

where,  $x_k$  is the  $k^{th}$  bit (k = 1, ..., n) of the sequence.

```
# Compute DFT
discrete_fourier_transform = numpy.fft.fft(bits_copy)
```

3. As there is symmetry in the complex numbers, we find the modulus of first-half numbers only. The threshold value is 0.95; thus, 95% of values of mod<sub>j</sub> should be less than h,

$$h = \sqrt{\left(\log \frac{1}{0.05}\right)n}$$

where,

 $mod_j = modulus of the complex number f_j$ h = peak height threshold value

```
# Compute magnitudes of first half of sequence depending on the system type
if sys.version_info > (3, 0):
    magnitudes = abs(discrete_fourier_transform)[:bits_copy.size // 2]
else:
    magnitudes = abs(discrete_fourier_transform)[:bits_copy.size / 2]
# Compute upper threshold
threshold: float = math.sqrt(math.log(1.0 / 0.05) * bits_copy.size)
```

4. Find N<sub>0</sub>, the expected theoretical (95 %) number of peaks, and N<sub>1</sub>, the actual observed number of peaks in M that are less than T.

$$N_0 = (0.95) * n/2$$

```
# Compute the expected number of peaks (N0)
expected_peaks: float = 0.95 * bits_copy.size / 2.0
# Count the peaks above the upper threshold (N1)
counted_peaks: float = float(len(magnitudes[magnitudes < threshold]))</pre>
```

5. Compute the value of d, the normalized difference, find the P-value, and conclude the results.

$$d = \frac{(N_1 - N_0)}{\sqrt{n(.95)(.05)/4}} \qquad P-value = erfc\left(\frac{|d|}{\sqrt{2}}\right)$$

```
# Compute the score (P-value) using the normalized difference
normalized_difference: float = (counted_peaks - expected_peaks) / math.sqrt((bits_copy.size * 0.95 * 0.05) / 4)
score: float = math.erfc(abs(normalized_difference) / math.sqrt(2))
```

```
# Return result
if score >= self.significance_value:
    return Result(self.name, True, numpy.array(score))
return Result(self.name, False, numpy.array(score))
```

### **Results and Inferences:**

PUF random test result - FAILED - score: 0.0

Python random test result - FAILED - score: 0.0

Python secret Test result - FAILED - score: 0.0

True random Test result - FAILED - score: 0.0

The result implies a degree of randomness below the accepted significance level (P<0.01). The test fails, indicating that the binary sequence may be non-random and further investigation is necessary to identify the source of the non-randomness. There could be several reasons why the statistical test of randomness fails. For example, there may be a systematic bias in the data collection process, or the data may have been intentionally manipulated or fabricated. Alternatively, there may be some underlying pattern or structure in the data that is not immediately apparent and requires further investigation.

It is possible for a true Random Number Generator to fail some tests due to random bias, but the ratio of successful tests should be considered to term any generator as true-random or pseudo-random. Here we know that the dataset used is truly random, but it still fails the test. It can be concluded that some bias may be there, even in truly random sequences.

## 4. Conclusion

After performing tests with the 4 generators used in analysis we have come up with the following results:

PUF generator passed 3 tests, true random generator passed 6 tests each whereas both python-secrets and random successfully cleared 7 tests.But this result is not sufficient because we see that we both python generators were eligible for 15 tests owing to large data size whereas we could not perform the 3 tests namely:

- 1. Overlapping-template matching,
- 2. Maurer's universal test,
- 3. Binary matrix rank test.

So, if we see the scores:

PUF had a score of 3/12=0.25 Python secret and random test cleared 7/15=0.46 Whereas the true random passed 6/12=0.5.

Although we are aware of the fact that true random generators must as such pass all the tests but then due to the design of tests in a particular fashion and depending on the constraints of source producing it we don't get a 100% score. If we analyze the test outcomes both the non-overlapping and monobit test pass with all 4 datasets and so they are not of much utility to comment on randomness.

# 5. References:

- [1] https://docs.python.org/3/library/random.html
- [2] https://peps.python.org/pep-0506/
- [3] PUF-based random number generation, CW O'Donnell, GE Suh, S Devadas In MIT CSAIL CSG Technical Memo 481, 2004
- [4] For all tests, the following publication was used.

  NIST Special Publication 800-22. A statistical test suite for random and pseudorandom number generators for cryptographic applications. Information Technology Laboratory of the National Institute of Standards and Technology, May 2000.
- [5] The Python package used to implement the tests is as follows: https://pypi.org/project/nistrng/