

Probabilistic Primality Testing

Probabilistic primality testing is a class of algorithms used to determine whether a given number is a prime number or a composite number. Unlike deterministic algorithms that always provide a correct answer, probabilistic algorithms may occasionally produce an incorrect result, but the probability of error is extremely low.

One well-known probabilistic primality testing algorithm is the **Miller-Rabin algorithm**. The Miller-Rabin algorithm is based on the properties of certain types of numbers called "strong pseudoprimes." A strong pseudoprime is a composite number that behaves like a prime number under specific conditions.

Here's a high-level overview of the Miller-Rabin algorithm:

Input: An odd integer n to be tested for primality and a parameter k , which determines the accuracy of the test.

Write n as $2^r * d + 1$ where r is the largest integer such that 2^r divides $n-1$ and d is an odd integer.

Witness Loop: Repeat the following k times:

- a. Choose a random integer a from the range $[2, n-2]$.
- b. Compute $x = a^d \bmod n$.
- c. If x is not congruent to $1 \pmod{n}$ and x is not congruent to $-1 \pmod{n}$, then continue to the next iteration of the loop.
- d. Repeat r times or until x is congruent to $-1 \pmod{n}$: $x = x^2 \bmod n$.
- e. If x is not congruent to $-1 \pmod{n}$ after r iterations, then n is definitely composite. Otherwise, continue to the next iteration of the loop.

If the algorithm has not identified n as composite after k iterations, then n is considered probably prime.

The accuracy of the Miller-Rabin algorithm depends on the chosen value of k . Larger values of k result in a lower probability of error but require more computational effort.