

LEBANESE AMERICAN UNIVERSITY  
DEPARTMENT OF COMPUTER SCIENCE AND  
MATHEMATICS



School of  
**Arts and Sciences**

*CSC 430 – Computer Networks*

**Homework III**

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## Table of Contents

Problem 1: .....	3
Problem 2: .....	4
Problem 3: .....	5
Problem 4: .....	6
Problem 5: .....	8
Problem 6: .....	8
Problem 7: .....	9

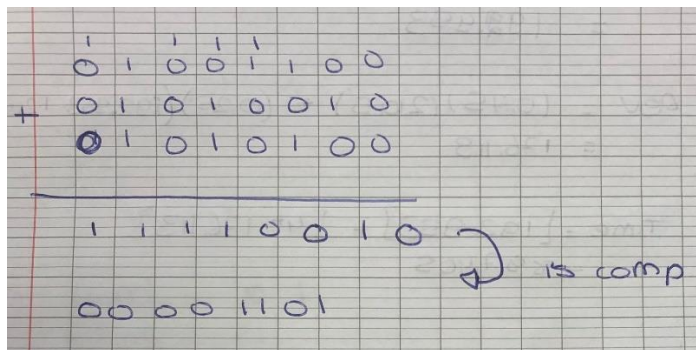
## Problem 1:

What is the 1s complement of the sum of these 8-bit bytes. Show all work.

Why is it that UDP takes the 1s complement of the sum; that is, why not just use the sum? With the 1s complement scheme, how does the receiver detect errors? Is it possible that a 1-bit error will go undetected? How about a 2-bit?

LRT: 01001100    01010010    01010100

Sum: 11110010    1s complement of the sum: 00001101



Finding out why UDP uses 1s complement was confusing, different sources say different things. However, I found an interesting historical analysis:

“In the past, computational resources were scarce, and every machine cycle was valuable. The UDP protocol utilizes one's complement arithmetic for checksum calculation to ensure accuracy and efficiency, independent of system endianness.”

[Why is it that UDP takes the 1s complement of the sum; that is, why not just use the sum? - Quora](#)

Errors will be detected by doing the same calculation and the sender and receiver side and if they don't match then an error is present.

It is not possible that a 1-bit error will go undetected, but it is possible a 2-bit error will go undetected if they result in the same answer in the checksum, example:

example: add two 16-bit integers

	1 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0	0 1
	1 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1	1 0
	<hr/>	
wraparound	1 1 0 1 1 1 0 1 1 1 0 1 1 1 0 1	
sum	1 0 1 1 1 0 1 1 1 0 1 1 1 0 0	
checksum	0 1 0 0 0 1 0 0 0 1 0 0 0 0 1 1	

Even though numbers have changed (bit flips), *no* change in **checksum**!

Taken from book: Kurose, J. and Ross, K. (2020) Computer Networking: A Top-Down Approach. Pearson, 8th Edition.

## Problem 2:

As host A sends data to B at rate higher than what B can handle, the receive buffer @ B becomes full. In response, B will let A know how much it can handle by sending ~~round trip delay~~ ~~remaining window size~~ & A will adjust according to achieve flow control.

### Problem 3: X= 927 ms

$$EstimatedRTT = 0.125 * 927 + 0.875 * 100 = 203.375 \text{ ms}$$

$$DevRTT = 0.25 * |100 - 203.375| + 0.75 * 5 = 29.5 \text{ ms}$$

$$TimeoutInterval = 203.375 + 4 * 29.5 = 354.393 \text{ ms}$$

$$Estimated RTT = 0.125 * 120 + 0.875 * 203.375 = 192.953 \text{ ms}$$

$$DevRTT = 0.25 * |120 - 192.953| + 0.75 * 29.5 = 40.36 \text{ ms}$$

$$TimeoutInterval = 192.953 + 4 * 40.36 = 354.393 \text{ ms}$$

$$Estimated RTT = 0.125 * 140 + 0.875 * 192.953 = 186.3 \text{ ms}$$

$$DevRTT = 0.25 * |140 - 186.3| + 0.75 * 40.36 = 41.845 \text{ ms}$$

$$TimeoutInterval = 186.3 + 4 * 41.845 = 353.68 \text{ ms}$$

$$Estimated RTT = 0.125 * 90 + 0.875 * 186.3 = 174.26 \text{ ms}$$

$$DevRTT = 0.25 * |90 - 174.26| + 0.75 * 41.845 = 52.44 \text{ ms}$$

$$TimeoutInterval = 174.26 + 4 * 52.44 = 384.02 \text{ ms}$$

$$Estimated RTT = 0.125 * 115 + 0.875 * 174.26 = 166.85 \text{ ms}$$

$$DevRTT = 0.25 * |115 - 166.85| + 0.75 * 52.44 = 52.29 \text{ ms}$$

$$TimeoutInterval = 166.85 + 4 * 52.29 = 376.01 \text{ ms}$$

#### Problem 4:

a) slow  $\Rightarrow$  before it starts inc by 1 only  
 $\Rightarrow [1, 6] \cup [23, 26]$

b) inc by 1:  
 $[6, 16] \cup [17, 22]$

c) Triple dup ACK as window size drops to  $1/2$  cwnd.

d) Dropped  $\rightarrow$  1 MSS  $\rightarrow$  timeout

e) 32 (where slow start stops)

f)  $42/2 = 21$

g)  $29/2 = 14.5$



h) 1R : 1  
 2R : 2-3  
 3R : 4-7  
 4R : 8-15  
 5R : 16-31  
 6R : 32-63  
 7R : 64-96  $\Rightarrow$  70 here

70th packet in round 7

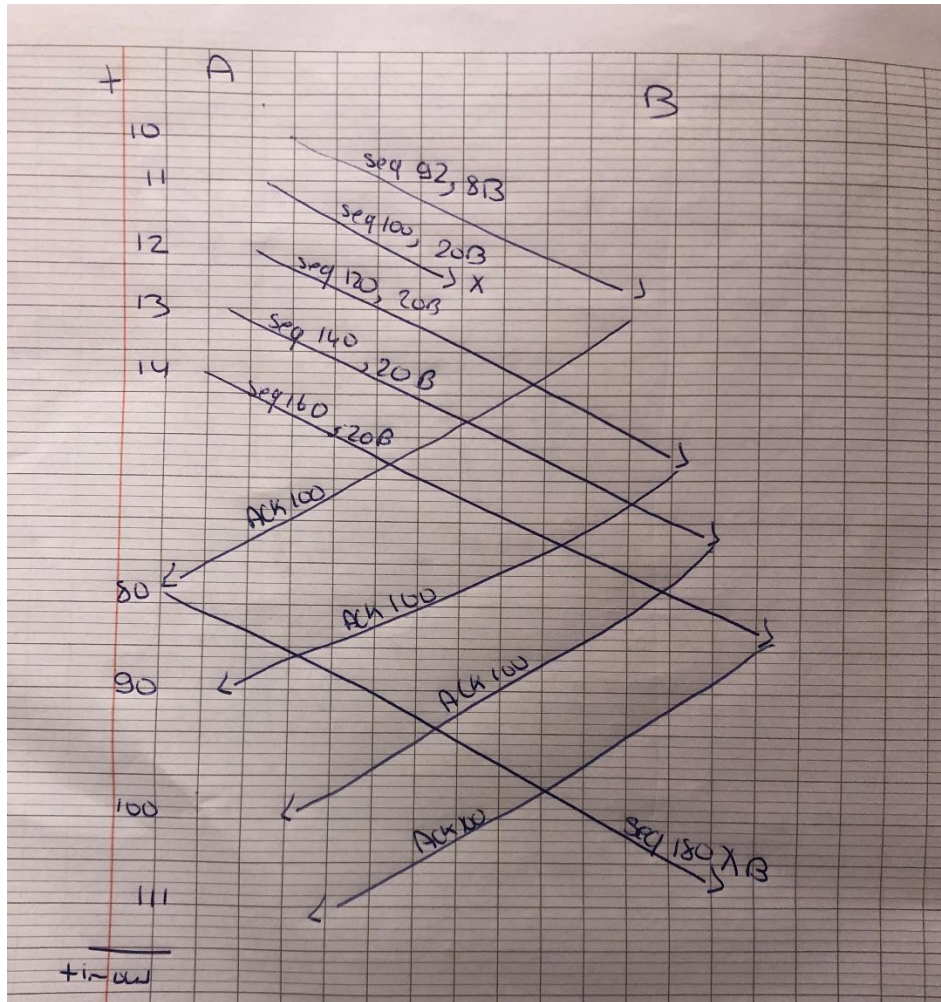
i) cwnd :  $8/2 = 4$   
 ssh = 4

j) ssh = 21 cwnd 8.  
 @ 16 cwnd : 1  
 @ 17 cwnd : 2  
 @ 18 : 4  $\Rightarrow$  @ 19 : cwnd 8

k) 17 : 1  
 18 : 2  
 19 : 4  
 20 : 8  
 21 : 16  
 22 : 21 (ssh)

} total = 52 packets

### Problem 5:





### **Problem 6:**

a)

Go-Back-N:

A sends a total of 9 segments. Initially sent as segments 1 - 5, and then retransmitted as segments 2 - 5. B replies with 8 ACKs: 4 ACKs with ACK number 1 and 4 ACKs with ACK numbers 2 - 5.

Selective Repeat:

A sends a total of 6 segments. Initially sent as segments 1 – 5 and then segment 2 is retransmitted. B acknowledges with 5 ACKs: 4 ACKs with ACK numbers 1, 3, 4, and 5, and 1 ACK with ACK number 2.

TCP:

A sends a total of 6 segments. Initially sent as segments 1 - 5 and then segment 2 is retransmitted. B acknowledges with 5 ACKs: 4 ACKs with ACK number 2 and 1 ACK with sequence number 6.

b) TCP due to its fast retransmit capability.

### Problem 7:

Problem 7 : for Buffered GBN & SR.  
The sequence number space must be able to accommodate to fit both receiver & sender window without overlapping.  
Example! Recv window :  $[m, m+w-1]$  received & acted everything before  $m$   
sender didn't receive  $w$  acks  $\Rightarrow$  window :  $[\underline{m-w}, m-1]$   
for the underlined edge not to overlap we need  $2w$  seq #s  $\Rightarrow K > 2w$

As for non-buffered GBN  $k-1$ .