

SF2520 — Applied numerical methods

Lecture 4

Runge–Kutta methods
Multistep methods

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Today's lecture

- Summary of last lecture
- High order methods
 - Runge–Kutta methods
 - Multi-step methods

Absolute stability

- An absolutely stable numerical solution satisfies

$|\tilde{u}_n - u_n|$ is bounded for all n , if $|\tilde{u}_0 - u_0|$ is small enough.

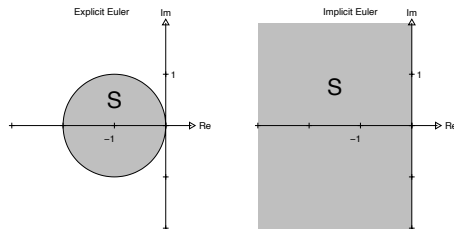
- Absolute stability **necessary in practice**, when exact solution is stable. Ensures that the effect of small perturbations remains small. (Important as small errors/perturbations are always introduced.)
- An absolutely stable solution behaves qualitatively correct.
- The fact that a method is absolutely stable for a certain time step does **not mean that the solution is accurate**. The time step must be chosen such that the method is **both** stable and accurate.

Stability regions

The **region of absolute stability** \mathcal{S} for a method is defined via the "test problem"

$$y' = \lambda y, \quad y(0) = y_0, \quad \text{Real}(\lambda) \leq 0. \quad (1)$$

- Let w_n be a numerical solution to (1). Then \mathcal{S} are the values $h\lambda \in \mathbb{C}$ for which w_n is bounded as $n \rightarrow \infty$.
- Every method has its own stability region \mathcal{S} .



- \mathcal{S} characterizes the absolute stability also for general ODEs:

ODE

$$y' = \lambda y + g$$

$$\mathbf{y}' = \mathbf{A}\mathbf{y} + \mathbf{g}(t)$$

$$\mathbf{y}' = \mathbf{A}(t)\mathbf{y} + \mathbf{g}(t)$$

$$\mathbf{y}' = \mathbf{F}(\mathbf{y})$$

Stability condition

$$h\lambda \in \mathcal{S}$$

$$h\lambda_k \in \mathcal{S} \text{ for all eigenvalues } \lambda_k \text{ of } \mathbf{A}$$

$$h\lambda_k(t) \in \mathcal{S} \text{ for all } t, k$$

$$h\lambda_k(\mathbf{u}_n) \in \mathcal{S} \text{ for all eigenvalues } \lambda_k(\mathbf{u}_n) \text{ of } \mathbf{J}(\mathbf{u}_n), \\ \text{the Jacobian matrix of } \mathbf{F}$$

Stiff ODE and implicit methods

- Explicit methods typically have a *bounded* \mathcal{S} (with the origin on its boundary). This gives a stability limit for h . Since λ is fixed, we must have $h \leq h_{\text{stab}}$.
- Implicit methods typically have an *unbounded* \mathcal{S} . Then there is no stability limit. The method is "unconditionally stable".
(There are examples of useful implicit methods with bounded \mathcal{S} .)
- In *stiff* problems explicit methods require much smaller timestep h than expected based on exact solution and LTE, because of stability requirement: $h \leq h_{\text{stab}} \ll h_{\text{acc}}$.
- Explicit methods become expensive as very many steps are needed, and "unnecessarily" accurate. Implicit methods advantageous.

Higher order methods

Explicit/Implicit Euler are the basic first order methods. To get higher order methods the local truncation error must be made smaller in terms of h .

Two ways to do this:

- 1 Make more function evaluations $f(t, y)$ in each step.
This leads to **Runge–Kutta methods**.

Example: Heun's method (2 evaluations, or "stages")

$$u_{n+1} = u_n + \frac{1}{2}h \left(f(t_n, u_n) + f(t_{n+1}, u_n + hf(t_n, u_n)) \right).$$

- 2 Use several of the previously computed approximations.
This leads to **multi-step methods**.

Example: Midpoint method (2 previous values)

$$u_{n+1} = u_{n-1} + 2hf(t_n, u_n).$$

Runge–Kutta methods

Examples: (for $y' = f(t, y)$)

- Heun's method (2 stages, order 2)

$$\begin{aligned}k_1 &= f(t_n, u_n), \\k_2 &= f(t_n + h, u_n + hk_1), \\u_{n+1} &= u_n + h \left(\frac{1}{2}k_1 + \frac{1}{2}k_2 \right).\end{aligned}$$

- RK4 (4 stages, order 4)

$$\begin{aligned}k_1 &= f(t_n, u_n), \\k_2 &= f\left(t_n + \frac{1}{2}h, u_n + \frac{1}{2}hk_1\right), \\k_3 &= f\left(t_n + \frac{1}{2}h, u_n + \frac{1}{2}hk_2\right), \\k_4 &= f(t_n + h, u_n + hk_3), \\u_{n+1} &= u_n + h \left(\frac{1}{6}k_1 + \frac{1}{3}k_2 + \frac{1}{3}k_3 + \frac{1}{6}k_4 \right).\end{aligned}$$

Runge–Kutta methods

General form: (s stages, explicit RK)

$$k_1 = f(t_n, u_n),$$

$$k_2 = f(t_n + hc_2, u_n + ha_{21}k_1),$$

$$k_3 = f(t_n + hc_3, u_n + ha_{31}k_1 + ha_{32}k_2),$$

$$\vdots$$

$$k_s = f(t_n + hc_s, u_n + ha_{s,1}k_1 + \cdots + ha_{s,s-1}k_{s-1}),$$

$$u_{n+1} = u_n + h(b_1k_1 + b_2k_2 + \cdots + b_sk_s)$$

Intuition: ($y' = f(t, y)$)

$$y(t_{n+1}) = y(t_n) + \int_{t_n}^{t_{n+1}} y'(s) ds = y(t_n) + h \int_0^1 f(t_n + ht, y(t_n + ht)) dt$$

- Approximate the integral by s-point quadrature method

$$y(t_{n+1}) \approx y(t_n) + h \sum_{j=1}^s b_j \underbrace{f(t_n + hc_j, y(t_n + hc_j))}_{\approx k_j}$$

- Approximate $y(t_n + hc_j)$ by an explicit method (note $k_j \approx y'(t_n + hc_j)$)

$$y(t_n + hc_j) \approx y(t_n) + ha_{j,1}k_1 + \cdots + ha_{j,j-1}k_{j-1}$$

Runge–Kutta methods

- Methods often summarized using "Butcher tableaux"

c_1	a_{11}			
c_2	a_{21}	a_{22}		
\vdots	\vdots	\vdots	\ddots	
c_s	$a_{s,1}$	$a_{s,2}$	\cdots	$a_{s,s}$
	b_1	b_2	\cdots	b_s

where $a_{jj} = 0$ for explicit RK-methods, and $c_1 = 0$.

- Example

$$\text{Heun: } \begin{array}{c|cc} 0 & 0 & 0 \\ 1 & 1 & 0 \\ \hline & \frac{1}{2} & \frac{1}{2} \end{array},$$

$$\text{RK4: } \begin{array}{c|cccc} 0 & 0 & & & \\ \frac{1}{2} & \frac{1}{2} & 0 & & \\ \frac{1}{2} & 0 & \frac{1}{2} & 0 & \\ 1 & 0 & 0 & 1 & 0 \\ \hline & \frac{1}{6} & \frac{1}{3} & \frac{1}{3} & \frac{1}{6} \end{array}$$

Runge–Kutta methods – accuracy and convergence

- Runge–Kutta methods are one-step methods (although many *stages*). General theorem for one-step methods applies, so that if $\ell_n = O(h^{p+1})$ then the method is of order p .
- Accuracy determined both by the quadrature (b_j, c_j) and the explicit method $(c_j, a_{j,\ell})$.
- Finding $a_{j,\ell}$, b_j and c_j such that $\ell_n = O(h^{p+1})$ is possible using Taylor expansions, but complicated for p large.
- For $p \leq 4$ it is enough with $s = p$. **Ex.** Heun has $s = p = 2$ and RK4 has $s = p = 4$. (For higher order, s must be larger than p . **Ex.** $p = 5$ requires $s = 6$ and $p = 10$ requires $s = 17$.)
- For consistency, $\ell_n = O(h^2)$, one just needs

$$\sum_{j=1}^s b_j = 1.$$

- There are also **implicit** Runge–Kutta methods. Leads to a $d s \times d s$ nonlinear system of equations in each step. Then $p = 2s$ possible.

Runge–Kutta methods and absolute stability

Example: Heun's method

- Apply to test problem $y' = \lambda y$, $y(0) = 1$,

$$k_1 = f(t_n, w_n) = \lambda w_n,$$

$$k_2 = f(t_n + h, w_n + hk_1) = \lambda(w_n + h\lambda w_n) = (\lambda + h\lambda^2)w_n,$$

$$w_{n+1} = w_n + h \left(\frac{1}{2}k_1 + \frac{1}{2}k_2 \right) = \left(1 + h\lambda + \frac{1}{2}(h\lambda)^2 \right) w_n.$$

- Hence

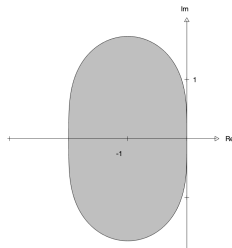
$$w_{n+1} = Q(h\lambda)w_n, \quad Q(z) = 1 + z + \frac{1}{2}z^2.$$

- Since, by induction

$$w_n = Q(h\lambda)^n w_0,$$

we have

$$\mathcal{S} = \{z \in \mathbb{C} : |Q(z)| \leq 1\}.$$



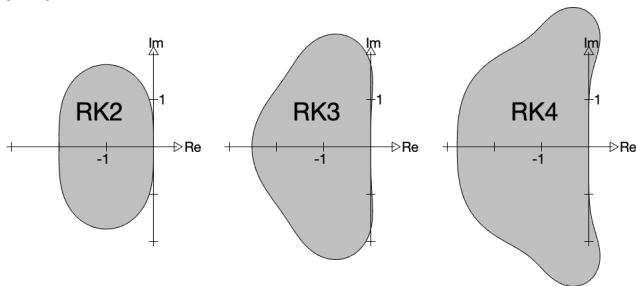
Runge–Kutta methods and absolute stability, cont.

- Applying an explicit R–K method with s stages to the test problem $y' = \lambda y$ always gives an iteration of the type

$$w_{n+1} = Q(h\lambda)w_n, \quad \text{where } Q \text{ is a polynomial of degree } s.$$

- Stability region given by $\mathcal{S} = \{z \in \mathbb{C} : |Q(z)| \leq 1\}$.

Note: Q polynomial $\Rightarrow Q$ unbounded $\Rightarrow \mathcal{S}$ bounded



- In a p -th order method, $Q(z) = e^z + O(z^{p+1})$, since

$$y(t_{n+1}) = e^{h\lambda} y(t_n) \quad \text{and} \quad y(t_{n+1}) = Q(h\lambda)y(t_n) + \ell_{n+1}.$$

Adaptive Runge–Kutta methods

- In adaptive Runge–Kutta methods, the stages chosen such that different quadratures (sets of b_j) give different order of accuracy (= **embedded methods**). (Requires few extra function evaluations.)
- **Example: Dormand–Prince** (as used in Matlab's `ode45` command)
 - Based on a Runge–Kutta-method with 7 stages.
 - First 6 stages give a 5th order approximation:

$$u_{n+1,5} = u_n + h \left(\frac{35}{384} k_1 + \frac{500}{1113} k_3 + \frac{125}{192} k_4 - \frac{2187}{6784} k_5 + \frac{11}{84} k_6 \right).$$

- Then let $k_7 = f(t_n + h, u_{n+1,5})$. (Note: same as k_1 in next step!)
- With k_7 a 4th order approximation is given as

$$u_{n+1,4} = u_n + h \left(\frac{5179}{57600} k_1 + \frac{7571}{16695} k_3 + \frac{393}{640} k_4 - \frac{92097}{339200} k_5 + \frac{187}{2100} k_6 + \frac{1}{40} k_7 \right)$$

- Local truncation error estimated by difference $|u_{n+1,5} - u_{n+1,4}|$.
- Requires 6 function evaluations per step. (Optimal for $p = 5$.)

Linear multistep methods

Example: Adams–Bashforth 3 (for $y' = f(t, y)$)

$$u_{n+1} = u_n + \frac{h}{12}(23f_n - 16f_{n-1} + 5f_{n-2}), \quad f_n = f(t_n, u_n).$$

Uses 3 old values to compute the next. Need u_0, u_1, u_2 to start.

- **General form:** (q steps)

$$u_{n+1} = \alpha_0 u_n + \alpha_1 u_{n-1} + \cdots + \alpha_{q-1} u_{n-q+1} \\ + h(\beta_{-1} f_{n+1} + \beta_0 f_n + \cdots + \beta_{q-1} f_{n-q+1})$$

- Initial data needed for u_0, u_1, \dots, u_{q-1} . Typically obtained from explicit one-step method (e.g. R–K methods).
- If $\beta_{-1} = 0$ the method is explicit, otherwise implicit.
- Very high order of accuracy can be reached in a stable way.
- Multistep methods only need one new function evaluation/time step, but need more memory since old values must be stored.
- Common families of methods: Adams and BDF.

Adams methods

- Adams methods only use previous values of f , (not u)

$$u_{n+1} = u_n + h(\beta_{-1}f_{n+1} + \beta_0f_n + \cdots + \beta_{q-1}f_{n+1-q}).$$

- Explicit ones called Adams–Bashforth (order q).
Implicit ones called Adams–Moulton (order $q + 1$).
- Derivation (Adams–Bashforth):

- Exact solution satisfies

$$y(t_{n+1}) = y(t_n) + \int_{t_n}^{t_{n+1}} y'(t) dt = y(t_n) + \int_{t_n}^{t_{n+1}} f(y(t)) dt.$$

- How to compute integral? Note $f(y(t))$ only approximated in the previous time points $t_{n-q+1}, t_{n-q}, \dots, t_n$.
- Let p be $q - 1$ degree polynomial that approximates $p(t) \approx f(y(t))$ by interpolation of f_n in past points

$$p(t_{n-k}) = f_{n-k}, \quad k = 0, \dots, q - 1.$$

- Then set

$$u_{n+1} = u_n + \int_{t_n}^{t_{n+1}} p(t) dt =: u_n + h(\beta_0f_n + \cdots + \beta_{q-1}f_{n-q+1}).$$

Backward Differentiation Formulae (BDF)

- BDF methods only use previous values of u , (not f)

$$u_{n+1} = \alpha_0 u_n + \alpha_1 u_{n-1} + \cdots + \alpha_{q-1} u_{n-q+1} + h\beta_{-1} f_{n+1}.$$

- Implicit q -step methods of order q .
- Often used for stiff ODEs (e.g. parabolic PDEs).
- Derivation
 - Let p be a degree q polynomial interpolating the *solution* u_n in the next and past points,

$$p(t_{n+1-k}) = u_{n+1-k}, \quad k = 0, \dots, q+1.$$

- Find u_{n+1} such that $p'(t_{n+1}) = f(u_{n+1}) (= f_{n+1})$.
- Leads to

$$u_{n+1} = \alpha_0 u_n + \alpha_1 u_{n-1} + \cdots + \alpha_{q-1} u_{n-q+1} + h\beta_{-1} f_{n+1},$$

for some α_j and β_{-1} .

- **Example: BDF 2**

$$u_{n+1} = \frac{4}{3}u_n - \frac{1}{3}u_{n-1} + h\frac{2}{3}f_{n+1}.$$

Linear multistep methods, summary of methods.

- Adams–Bashforth (explicit, order q)

$$u_{n+1} = u_n + h \left(\beta_0 f_n + \cdots + \beta_{q-1} f_{n-q+1} \right).$$

- **Example: Adams–Bashforth 3** (order 3)

$$u_{n+1} = u_n + \frac{h}{12} (23f_n - 16f_{n-1} + 5f_{n-2}).$$

- Adams–Moulton (implicit, order $q + 1$)

$$u_{n+1} = u_n + h \left(\beta_{-1} f_{n+1} + \beta_0 f_n + \cdots + \beta_{q-1} f_{n-q+1} \right).$$

- **Example: Adams–Moulton 3** (order 4)

$$u_{n+1} = u_n + \frac{h}{24} (9f_{n+1} + 19f_n - 5f_{n-1} + f_{n-2}).$$

- BDF (implicit order q)

$$u_{n+1} = \alpha_0 u_n + \alpha_1 u_{n-1} + \cdots + \alpha_{q-1} u_{n-q+1} + h \beta_{-1} f_{n+1}.$$

- **Example: BDF 2** (order 2)

$$u_{n+1} = \frac{4}{3} u_n - \frac{1}{3} u_{n-1} + h \frac{2}{3} f_{n+1}.$$

Linear multistep methods, accuracy and convergence

- Local truncation error ℓ_n as before ($y_n = y(t_n)$ is exact solution)

$$y_{n+1} = \alpha_0 y_n + \alpha_1 y_{n-1} + \cdots + \alpha_{q-1} y_{n-q+1} \\ + h(\beta_{-1} f(y_{n+1}) + \beta_0 f(y_n) + \cdots + \beta_{q-1} f(y_{n-q+1})) + \ell_{n+1},$$

- Also define the polynomial

$$\rho(z) := \alpha_{q-1} + \alpha_{q-2}z + \cdots + \alpha_0 z^{q-1} - z^q.$$

Theorem (Dahlquist – global error)

Suppose f is Lipschitz and $\ell_n = O(h^{p+1})$. If, in addition, the initial data error in u_0, \dots, u_{q-1} is $O(h^p)$ and ρ satisfies the **root condition**:

All roots z_j of ρ satisfies either $|z_j| < 1$ or, if z_j is a simple root, $|z_j| \leq 1$,

then

$$\max_{0 \leq n \leq N} |y_n - u_n| \leq Ch^p, \quad (Nh = T \text{ fixed, and } C \text{ independent of } h).$$

- Note: For one-step methods $\rho(z) = 1 - z$ (since $\alpha_0 = 1$ then).

Remarks

- Convergence theory more complicated than for one-step methods. Consistent methods ($\ell_n = O(h^2)$) may fail to converge.
- Root condition needed to guarantee convergence (aka as "zero stability"). (Automatically satisfied for one-step methods.)
- Adams–Bashforth and Adams–Moulton are convergent for all q . (They have $\rho(z) = z^{q-1} - z^q$.)
- BDF are only convergent for $q \leq 6$.
- Shortcut available to determine order. Introduce polynomial

$$\sigma(z) := \beta_{q-1} + \beta_{q-2}z + \cdots + \beta_0 z^q + \beta_{-1} z^{q+1}.$$

Then $\ell_n = O(h^{p+1})$ iff $\rho(e^h) + h\sigma(e^h) = O(h^{p+1})$.

- Adaptivity more difficult since changes in time step will change coefficients α_j and β_j .

Linear multistep methods and absolute stability

- Absolute stability analysis leads to study of difference equations.
- For a general multistep method stability governed by the polynomial

$$P(z) = \rho(z) + h\lambda\sigma(z),$$

with ρ, σ defined above.

- Let $z_j(h\lambda)$ be the roots of P . Absolute stability region \mathcal{S} consists of those values $h\lambda \in \mathbb{C}^-$ for which

$$|z_j(h\lambda)| < 1, \quad j = 1, \dots, q.$$

(Or $|z_j(h\lambda)| = 1$ if root is simple.)

- Note that the root condition is the same condition for ρ . This means that root condition (zero stability) equivalent to $0 \in \mathcal{S}$.

Linear multistep methods and absolute stability

- **Example:** Applying AB3 to test problem $y' = \lambda y$ gives

$$w_{n+1} = w_n + \frac{h}{12}(23\lambda w_n - 16\lambda w_{n-1} + 5\lambda w_{n-2}),$$

which can be rewritten as

$$w_{n+1} + r_0(h\lambda)w_n + r_1(h\lambda)w_{n-1} + r_2(h\lambda)w_{n-2} = 0,$$

where $r_0(s) = -(1 + \beta_0 s)$ and for $j \geq 1$ we have $r_j(s) = -\beta_j s$.

- For fixed $h\lambda$ this is a difference equation, whose solution is determined by its characteristic polynomial P ,

$$P(z) = z^3 + r_0 z^2 + r_1 z + r_2.$$

- If z_1 , z_2 and z_3 are simple roots of P , then

$$w_n = c_1 z_1^n + c_2 z_2^n + c_3 z_3^n,$$

where c_j are determined by initial data. (See Appendix A.2 in Edsberg.)

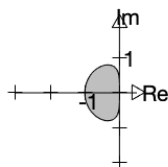
- Absolute stability requires $|z_j| < 1$ for all j . Then $|w_n| \rightarrow 0$.
- Note: r_j , c_j and w_j all depend on $h\lambda$. Also note that

$$P(z) = z^3 - z^2 - \beta_0 h\lambda z^2 - \beta_1 h\lambda z - \beta_2 h\lambda = \rho(z) - h\lambda\sigma(z).$$

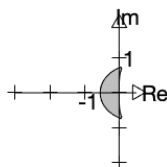
Linear multistep methods and absolute stability

Stability regions, Adams methods

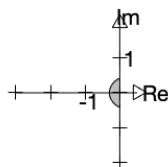
AB2



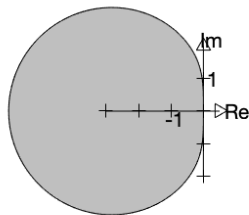
AB3



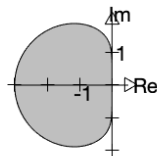
AB4



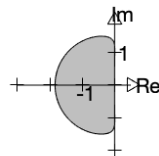
AM2



AM3



AM4



Linear multistep methods and absolute stability

- Stability regions of the implicit Adams–Moulton methods are *bounded* but much larger than those of Adams–Bashforth.
- Adams–Bashforth and Adams–Moulton usually used together in "predictor-corrector" pair: The u_{n+1} value is predicted with explicit AB and corrected with AM. (Fixed point iteration possible.) Used for non-stiff problems to increase stability region of AB.

Linear multistep methods and absolute stability

Stability regions, predictor corrector

AB2



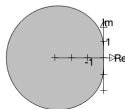
AB3



AB4



AM2



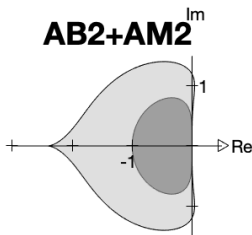
AM3



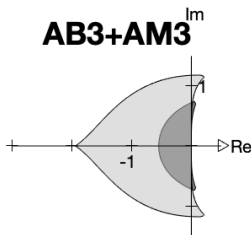
AM4



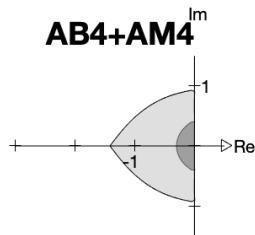
AB2+AM2



AB3+AM3



AB4+AM4



Linear multistep methods and absolute stability

Stability regions, BDF

