

# Recitation 24: Proofs about Programs (2 of 2)

## Data structures

( $\star\star$  [push  $x$   $s$ ] is the stack  $s$  with  
[ $x$ ] pushed on top  $\star$ )

$(x_1; x_2; \dots; x_n)$   
 $\uparrow$  top

( $\star\star$  [push  $x$   $(x_1; \dots; x_n)$ ] is  $[(x; x_1; \dots; x_n)] \star$ )

val push: 'a  $\rightarrow$  'a t  $\rightarrow$  'a t

peek(push  $x$   $s$ ) =  $x$

## Equational Specifications

empty	1 is_empty empty = true
is_empty	2 is_empty (push $x$ $s$ ) = false
pop	3 pop (push $x$ $s$ ) = $s$
peek	4 peek (push $x$ $s$ ) = $x$
push	

peek(pop(push 3(push 5 empty))))  $\stackrel{(3)}{=}$

peek(push 5 empty)  $\stackrel{(4)}{=}$

5

$(x_1; x_2; \dots; x_n)$   
 $\uparrow$

push  $x_1$  (push  $x_2$  ( ... (push  $x_n$  empty) ... ))

canonical form      generators push, empty

empty } generators  
 push  
 is-empty } queries  
 peek  
 pop } manipulators

Notice

equations are queries manipulators acting on generators

- 1 is-empty empty = true
- 2 is-empty (push x s) = false
- 3 pop (push x s) = s
- 4 peek (push x s) = x

Proofs!

module ListStack = struct

type 'a t = 'a list

let empty = []

let is-empty s = (s = [])

let peek = List.hd

let pop = List.tl

let push = List.cons

end

Show:  $\text{pop}(\text{push } x \ s) = s$

$\text{pop}(\text{push } x \ s)$

$= \text{eval pop, push?}$

$\text{List.tl } x :: s$

$= \text{eval tl?}$

## Queues!

is\_empty  
empty

eng  
front  
deg

is\_empty empty = true

is\_empty (eng x q) = false

front (eng x q) =

$\times$  front q      if is\_empty q  
if not


deg' (eng x q) =

empty      if is\_empty  $\times$  q  
if not  
eng x (deg q)

Simplify

deg (enc 3 (enc 4 (enc 5 empty)))  
= enc 3 (deg (enc 4 (enc 5 empty)))  
= enc 3 (enc 4 (deg (enc 5 empty)))  
= enc 3 (enc 4 (empty))

ListQueue : eval

module TwoListQueue = struct  
 (\* AF:  $(f, b)$  represent the queue  $f$  @ rev  $b$     
 RI: if  $f = []$  then  $b = []$  \*)  
 type 'a t = 'a list \* 'a list  
 let empty = [], []  
 let is\_empty (f, \_) = f = []  
 ...  
end

Proofs use 2 additional techniques

RI( $\underline{q}$ ) : if  $\underline{q} = (f, b)$  and  $f = []$  then  $b = []$

if is\_empty  $\underline{q}$  then  $\underline{q} = [], []$

If  $AF(e) = AF(e') \Rightarrow e' = e$

$([x_1; x_2; \dots x_n], [y])$

$([x_1; x_2; \dots x_n; y], [])$