

Assignment 03 - MATH 525

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Date: Sunday 5th March, 2023

Contents

Problem 1

Consider any point $y \in S$. We will show that $f(x^*) \leq f(y)$.

We know that x^* is a local minimum, i.e. that $\exists \epsilon > 0$ such that $\forall x \in S, \|x - x^*\| \leq \epsilon$ we have $f(x^*) \leq f(x)$. If $\|y - x^*\| \leq \epsilon$, then by the definition of a local minimum we know $f(x^*) \leq f(y)$. Thus the only remaining case to consider is when $\|y - x^*\| > \epsilon$.

Consider $a = \frac{\epsilon}{\|y - x^*\|} \in \mathbb{R}$. Since $\epsilon > 0$ and $\|y - x^*\| > \epsilon$, it must hold that $0 < a < 1$, or equivalently that $a \in [0, 1]$.

Consider the point $z = ay + (1 - a)x^*$, a convex combination of y and x^* . Since S is a convex set, $z \in S$. We have

$$\begin{aligned}\|z - x^*\| &= \|(ay + (1 - a)x^*) - x^*\| \\ &= \|(ay + x^* - ax^*) - x^*\| \\ &= \|ay - ax^*\| \\ &= a\|y - x^*\| \\ &= \frac{\epsilon}{\|y - x^*\|} \|y - x^*\| \\ &= \epsilon\end{aligned}\tag{1}$$

Since $\|z - x^*\| = \epsilon$ we have $f(x^*) \leq f(z)$. Since f is a convex function, we know that

$$\begin{aligned}f(x^*) \leq f(z) &= f(ay + (1 - a)x^*) \leq af(y) + (1 - a)f(x^*) \\ f(x^*) &\leq af(y) + (1 - a)f(x^*) \\ af(x^*) &\leq af(y) \\ f(x^*) &\leq f(y)\end{aligned}\tag{2}$$

Note that the last step only holds as a must be positive. Thus $\forall y \in S$, we know $f(x^*) \leq f(y)$. By definition, x^* is a global maximum. Q.E.D.

Problem 2

We note that Theorem 2.6 from the lecture notes states that a polyhedron contains an extreme point if and only if it contains a line.

Standard form problems cannot possibly contain a line, as they are confined as a subset of the positive orthant $\{x | x \geq 0\}$. This is not a requirement of general form problems, and as such they can contain a line. For example, the polyhedron $P = \{x \in \mathbb{R}^2 | x_1 \geq 4\}$ comprises a half-space, containing many lines (such as $x_1 = 5$).

Just because we can turn a general form polyhedron into a standard form polyhedron with the same optimal cost, this does not mean that the polyhedron is equivalent. By transforming a polyhedron into standard form, we add dimensions such that we can maintain the constraints of the original problem through a polyhedron in the positive orthant. The fact that the general form polyhedron and the standard form polyhedron (may) lie in different dimensionalities is enough to reason that they may not be equivalent, despite leading to an equivalent LP problem. For this reason, we cannot conclude that every nonempty polyhedron contains an extreme point.

Problem 4

We have $P = \{x \in \mathbb{R}^n : a'_i x = b_i, i = 1 \dots m\}$, where P is bounded and $m \geq n$. We also have $Q = \{x \in P : a'x = b\}$ for some $a \in \mathbb{R}^n, b \in \mathbb{R}$. We show that every extreme point of Q is either an extreme point of P or a convex combination of two adjacent extreme points of P .

Suppose that Q has an extreme point x^* . Since x^* is an extreme point, Theorem 2.3 from the lecture notes tells us that there are n linearly independent constraints active at x^* . We know from the definition of Q that $a'x = b$ is active at x^* , thus the remaining constraints active at x^* must come from the constraints of P . Let $I(x^*) = \{i \in \{1 \dots m\} : a'_i x = b_i\}$ define these remaining constraints from P . Note that a may or may not be linearly independent with this set, so we consider each case:

a is linearly dependent with $\{a_i : i \in I(x^*)\}$

If a is linearly dependent with this set, then the n linearly independent constraints must all lie in $\{a_i : i \in I(x^*)\}$, and x^* is also an extreme point of P .

a is linearly independent with $\{a_i : i \in I(x^*)\}$

If a is linearly independent with this set, then of the n linearly independent constraints active at x^* in Q , $(n-1)$ of them must come from $\{a_i : i \in I(x^*)\}$. These $(n-1)$ linearly independent constraints **do not span** \mathbb{R}^n , and there must exist some nonzero $d \in \mathbb{R}^n$ such that $\forall i \in I(x^*) \quad a'_i d = 0$.

Let $y = x^* + \lambda d$. $\forall \lambda \in \mathbb{R}, i \in I(x^*)$, we have $a'_i y = a'_i x^* + \lambda a'_i d$. Since $a'_i d = 0$, we have $a'_i y = a'_i x^* = b_i$. Thus $I(y) \supseteq I(x^*)$ for any y .

Q is bounded, meaning that if we travel in the direction of any nonzero vector from x^* we must eventually reach a constraint. This is easily shown from the definition of a bounded polyhedron, keeping in mind that we are adding a nonzero vector to a point in Q .

Suppose, going in the direction of d , we find some $\lambda^+ > 0, j \notin I(x^*)$ at which $a'_j(x^* + \lambda^+ d) = b_j$. a_j cannot be expressed as a linear combination of $\{a_i : i \in I(x^*)\}$. This is because $a'_i d = 0$ for all a_i , and

$$a'_j(x^*) > b_j \text{ and } a'_j(x^* + \lambda^+ d) = b_j \Rightarrow a'_j d \neq 0$$

Thus $\{a_i : i \in I(x^*) \cup \{j\}\}$ contains n linearly independent vectors, and we have $y^+ = x^* + \lambda^+ d$ is an extreme point.

Similarly, we can go in the direction $-d$, until, for some $\lambda^- > 0, k \notin I(x^*)$ at which $a'_k(x^* - \lambda^- d) = b_k$, and we find that a_k is also linearly independent from $\{a_i\}$, making $y^- = x^* - \lambda^- d$ is an extreme point. We note that y^+ and y^- share $(n-1)$ active constraints, and that $y^+ \neq y^-$, as

$$y^+ - y^- = x^* + \lambda^+ d - (x^* - \lambda^- d) = d(\lambda^+ + \lambda^-) > 0$$

Thus y^+ and y^- are distinct adjacent basic feasible solutions. Furthermore, since x^* lies on the line segment between y^+ and y^- , it can be expressed as a linear combination of the two, with

$$\begin{aligned}
x^* &= by^+ + (1-b)y^- \\
x^* &= bx^* + b\lambda^+d + x^* - \lambda^-d - bx^* + b\lambda^-d \\
0 &= b\lambda^+d - \lambda^-d + b\lambda^-d \\
\lambda^-d &= b(\lambda^+d + \lambda^-d) \\
b &= \frac{\lambda^-}{\lambda^+ + \lambda^-}
\end{aligned} \tag{3}$$

Since $\lambda^- > 0, \lambda^+ > 0$, we know that $0 < b < 1$, and x^* is actually a convex combination of two adjacent extreme points of P , y^+ and y^- . Q.E.D.

Problem 5

We consider this problem geometrically, by looking at a visual of P :

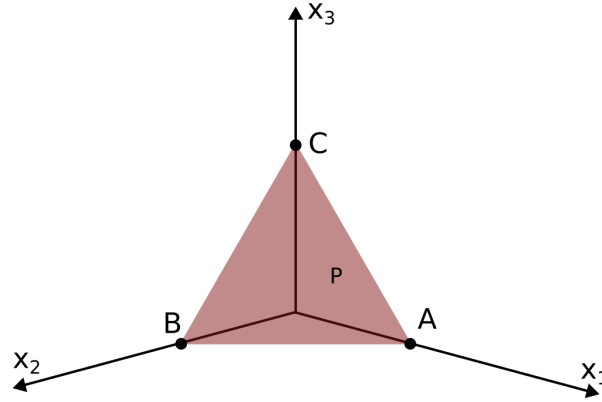


Figure 1: The geometry of P , along with vertices A, B, C

We can see three vertices of P , most notably, $C = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix}$. From C , we can see that feasible directions from C , geometrically speaking, include going towards B , going towards A , or going in a direction in between A and B , as shown in Figure 2.

Arithmetically speaking, any of these directions d can be represented as a convex combination of $(A - C)$ and $(B - C)$:

$$d = \lambda(A - C) + (1 - \lambda)(B - C); \lambda \in [0, 1] \tag{4}$$

Arithmetically speaking, we can only prove that d is a feasible solution at C if $y = C + \theta d \in P$ for some positive θ :

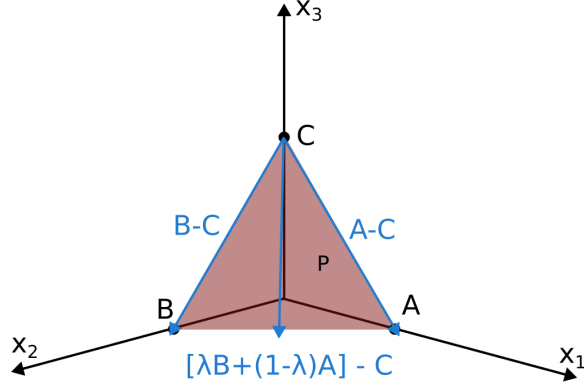


Figure 2: Feasible directions from A

$$\begin{aligned}
 y &= C + \theta d \\
 y &= C + \theta [\lambda(A - C) + (1 - \lambda)(B - C)] \\
 y &= C + \theta \lambda A - \theta \lambda C + \theta(1 - \lambda)B - \theta(1 - \lambda)C \\
 y &= \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} + \begin{bmatrix} \theta \lambda \\ 0 \\ 0 \end{bmatrix} - \begin{bmatrix} 0 \\ 0 \\ \theta \lambda \end{bmatrix} + \begin{bmatrix} 0 \\ \theta(1 - \lambda) \\ 0 \end{bmatrix} - \begin{bmatrix} 0 \\ 0 \\ \theta(1 - \lambda) \end{bmatrix} \\
 y &= \begin{bmatrix} \theta \lambda \\ \theta - \theta \lambda \\ 1 - \theta \end{bmatrix}
 \end{aligned} \tag{5}$$

Suppose $\theta \leq 1$; this means that $y_3 \geq 0$. Since $0 \leq \lambda \leq 1$, $\theta \geq \theta \lambda$ and $y_2 \geq 0$. Finally, since $\theta > 0$, $\lambda \geq 0$ and $y_1 \geq 0$. Finally, we have

$$y_1 + y_2 + y_3 = (\theta \lambda) + (\theta - \theta \lambda) + (1 - \theta) = 1$$

Thus for any $\lambda \in [0, 1]$, $0 < \theta \leq 1$, we have $x + \theta d \in P$, and $d \in \{\lambda(A - C) + (1 - \lambda)(B - C); \lambda \in [0, 1]\}$ is the set of feasible directions at C . Q.E.D.