Chapter 3

Processes (Part 2)

Yunmin Go

School of CSEE



Agenda

- Process Concept
- Process State and Scheduling
- Process API
- Signal
- Inter Process Communication





Signal

- A signal is a small message that notifies a process that an event of some type has occurred in the system.
 - Kernel abstraction for exceptions and interrupts.
 - Sent from the kernel (sometimes at the request of another process) to a process.
 - Different signals are identified by small integer ID's
 - The only information in a signal is its ID and the fact that it arrived.
 - Each signal has a current disposition (action associated with a signal), which determines how the process behaves when it is delivered the signal.



Signal Types

ID	Name	Default Action	Corresponding Event
2	SIGINT	Terminate	Interrupt from keyboard (ctrl-c)
8	SIGFPE	Terminate & Dump	Erroneous arithmetic operation (ex. divide by zero)
9	SIGKILL	Terminate	Kill program (cannot override or ignore)
11	SIGSEGV	Terminate & Dump	Segmentation violation
14	SIGALRM	Terminate	Timer signal
17	SIGCHLD	Ignore	Child stopped or terminated

yunmin@yunmin:~/workspace/ch3\$ kill -l								
1)	SIGHUP	<pre>2) SIGINT</pre>	SIGQUIT	4) SIGILL	5) SIGTRAP			
6)	SIGABRT	7) SIGBUS	8) SIGFPE	9) SIGKILL	10) SIGUSR1			
11)	SIGSEGV	12) SIGUSR2	13) SIGPIPE	14) SIGALRM	15) SIGTERM			
16)	SIGSTKFLT	17) SIGCHLD	18) SIGCONT	19) SIGSTOP	20) SIGTSTP			
21)	SIGTTIN	22) SIGTTOU	23) SIGURG	24) SIGXCPU	25) SIGXFSZ			
26)	SIGVTALRM	27) SIGPROF	28) SIGWINCH	29) SIGIO	30) SIGPWR			



Sending a Signal

- Kernel sends (delivers) a signal to a destination process by updating some state in the context of the destination process.
- Kernel sends a signal for one of the following reasons:
 - Kernel has detected a system event such as divide-by-zero (SIGFPE) or the termination of a child process (SIGCHLD)
 - Another process has invoked the kill() system call to explicitly request the kernel to send a signal to the destination process.



Receiving a Signal

- A destination process receives a signal when it is forced by the kernel to react in some way to the delivery of the signal.
- Signal disposition: Each signal has a default disposition, which determines how the process behaves when it is delivered the signal
 - Terminate the process
 - Core dump and terminate the process
 - Ignore the signal
 - Stop the process
 - Continue the process if it is currently stopped
- Also a process can change the disposition of a signal
 - Process executes a user-level function called a signal handler when the specified signal occurs.

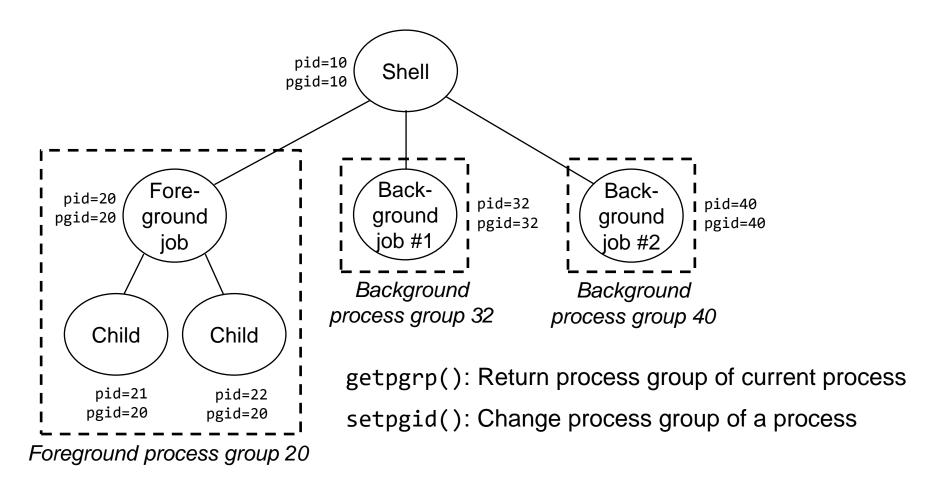
Signal Pending and Blocked

- A signal is pending if it has been sent but not yet received.
 - There can be at most one pending signal of any particular type.
 - Important: Signals are not queued
 - If a process has a pending signal of type k, then subsequent signals of type k that are sent to that process are discarded.
- A process can block the receipt of certain signals.
 - Blocked signals can be delivered, but will not be received until the signal is unblocked.
- A pending signal is received at most once.
- Kernel maintains pending and blocked bit vectors in the context of each process.



Process Groups

Every process belongs to exactly one process group





Kill Program

kill program sends arbitrary signal to a process or process group

- \$ kill -9 4009
 - Send SIGKILL to process 4009
- \$ kill -9 -4224
 - Send SIGKILL to every process in process group 4224

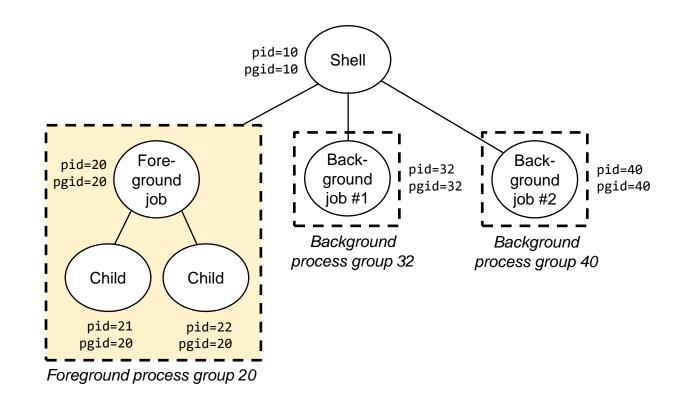
```
yunmin@yunmin:~/ch3 2$ ./killproc1
Child1: pid=4009 pgrp=4008
Child2: pid=4010 pgrp=4008
yunmin@yunmin:~/ch3 2$ ps
    PID TTY
                     TIME CMD
   3605 pts/0
                 00:00:00 bash
                 00:00:01 killproc1
   4009 pts/0
                00:00:01 killproc1
   4010 pts/0
                00:00:00 ps
   4023 pts/0
yunmin@yunmin:~/ch3 2$ kill -9 4009
yunmin@yunmin:~/ch3 2$ ps
    PID TTY
                     TIME CMD
   3605 pts/0
                 00:00:00 bash
                 00:00:10 killproc1
   4010 pts/0
   4069 pts/0
                 00:00:00 ps
yunmin@yunmin:~/ch3_2$ kill -9 4010
yunmin@yunmin:~/ch3 2$ ps
    PID TTY
                     TIME CMD
   3605 pts/0
                 00:00:00 bash
   4086 pts/0
                 00:00:00 ps
```

```
yunmin@yunmin:~/ch3 2$ ./killproc1
Child1: pid=4225 pgrp=4224
Child2: pid=4226 pgrp=4224
yunmin@yunmin:~/ch3 2$ ps
    PID TTY
                     TIME CMD
   3605 pts/0
                 00:00:00 bash
   4225 pts/0
                 00:00:02 killproc1
                 00:00:02 killproc1
   4226 pts/0
   4238 pts/0
                 00:00:00 ps
yunmin@yunmin:~/ch3 2$ kill -9 -4224
yunmin@yunmin:~/ch3_2$ ps
    PID TTY
                     TIME CMD
                 00:00:00 bash
   3605 pts/0
   4291 pts/0
                 00:00:00 ps
```



Signals from Keyboard

- Typing ctrl-c (ctrl-z) sends a SIGTERM (SIGTSTP) to every job in the foreground process group.
 - SIGTERM: default action is to terminate each process
 - SIGTSTP: default action is to stop (suspend) each process





Signals from Keyboard: Example

Example of ctrl-c and ctrl-z

```
int main()
    if (fork() == 0) {
        printf("Child: pid=%d pgrp=%d\n",
                getpid(), getpgrp());
    } else {
        printf("Parent: pid=%d pgrp=%d\n",
                getpid(), getpgrp());
   while(1);
    return 0;
```

- fg: By default, this brings the last background job to the foreground.
- fg %job number: Brings a specific job (by job number) to the foreground.

```
yunmin@yunmin:~/ch3 2$ ./killproc2
Parent: pid=5339 pgrp=5339
Child: pid=5340 pgrp=5339
[1]+ Stopped
                             ./killproc2
yunmin@yunmin:~/ch3_2$ ps -o pid,tty,time,stat,cmd
                    TIME STAT CMD
   PID TT
  3605 pts/0 00:00:00 Ss /bin/bash --init-file /home/
  5339 pts/0
                             ./killproc2
               00:00:03 T
  5340 pts/0 00:00:03 T
                            ./killproc2
   5390 pts/0
                00:00:00 R+
                             ps -o pid,tty,time,stat,cmd
yunmin@yunmin:~/ch3_2$ fg
 /killproc2
yunmin@yunmin:~/ch3_2$ ps -o pid,tty,time,stat,cmd
   PID TT
                    TIME STAT CMD
   3605 pts/0
               00:00:00 Ss /bin/bash --init-file /home/y
   5455 pts/0
                00:00:00 R+ ps -o pid,tty,time,stat,cmd
```



kill()

kill(): Send signal to a process

```
#include <signal.h>
int kill(pid_t pid, int sig);
```

- The kill() system call can be used to send any signal to any process group or process.
- pid>0: sig is sent to the process with specified by pid
- pid=0: sig is sent to every process in the process group of the calling process
- pid=-1: sig is sent to every process for which the calling process has permission to send signals, except for process 1 (init)
- pid<-1: sig is sent to every process in the process group whose ID is -pid.
- Return value
 - Success: zero is returned
 - Error: -1 is returned



kill(): Example

```
void main()
                                                           killproc3.c
    pid t pid[N];
    int i, child status;
    for (i = 0; i < N; i++)
         if ((pid[i] = fork()) == 0)
             while(1); /* Child infinite loop */
    /* Parent terminates the child processes */
    for (i = 0; i < N; i++) {
         printf("Killing process %d\n", pid[i]);
         kill(pid[i], SIGINT);
    /* Parent reaps terminated children */
    for (i = 0; i < N; i++) {
         pid t wpid = wait(&child status);
         if (WIFEXITED(child status))
             printf("Child %d terminated with exit status %d\n",
                     wpid, WEXITSTATUS(child status));
         else
             printf("Child %d terminated abnormally\n", wpid);
```

```
yunmin@yunmin:~/ch3_2$ ./killproc3
Killing process 5703
Killing process 5704
Killing process 5705
Killing process 5706
Killing process 5707
Child 5705 terminated abnormally
Child 5704 terminated abnormally
Child 5706 terminated abnormally
Child 5707 terminated abnormally
Child 5703 terminated abnormally
```

signal()

signal(): Define the action associated with a signal

```
#include <signal.h>
typedef void (*sighandler t)(int);
sighandler_t signal(int signum, sighandler_t handler);
```

- Sets the disposition of the signal signum to handler
- signum: the name of signal
- handler: could be one among three possible reaction to the signal signum
 - SIG IGN: ignore the signals
 - SIG_DFL: use default action
 - Address of signal handler: use a programmer-defined function

- Return value
 - Success: the previous value of the signal handler
 - Error: SIG ERR



HANDONG GLOBAL Installing the handler

signal(): Example

```
// SIGINT handler
                                         signal1.c
void int handler(int sig)
    printf("Process %d received signal %d\n",
            getpid(), sig);
    exit(0);
void main()
    pid t pid[N];
    int i;
    int child status;
    signal(SIGINT, int handler);
    for (i = 0; i < N; i++) {
        if ((pid[i] = fork()) == 0)
            while(1); /* Child: Infinite Loop */
    for (i = 0; i < N; i++) {
        printf("Killing process %d\n", pid[i]);
        kill(pid[i], SIGINT);
```

```
/* Parent reaps terminated children */
    for (i = 0; i < N; i++) {
        pid_t wpid = wait(&child_status);
        if (WIFEXITED(child_status))
            printf("Child %d terminated with exit status %d\n",
                 wpid, WEXITSTATUS(child_status));
        else
            printf("Child %d terminated abnormally\n", wpid);
    }
}</pre>
```

```
yunmin@yunmin:~/ch3_2$ ./signal1
Killing process 2597
Killing process 2598
Killing process 2599
Killing process 2600
Process 2600 received signal 2
Process 2599 received signal 2
Killing process 2601
Process 2601 received signal 2
Child 2599 terminated with exit status 0
Child 2600 terminated with exit status 0
Child 2601 terminated with exit status 0
Process 2598 received signal 2
Child 2598 terminated with exit status 0
Process 2597 received signal 2
Child 2597 terminated with exit status 0
```



Signal Handler Funkiness

```
int ccount = 0;
                                                      signal2.c
// SIGCHLD handler that reaps one terminated child
void child handler(int sig)
   int child status;
    pid t pid = wait(&child status);
    ccount--;
    printf("Received signal %d from process %d\n", sig, pid);
// Signal funkiness: Pending signals are not queued
void main()
    pid_t pid[N];
   int i;
   ccount = N;
    signal(SIGCHLD, child handler);
   for (i = 0; i < N; i++)
        if ((pid[i] = fork()) == 0) {
            exit(0); /* Child: Exit */
    while (ccount > 0)
        pause(); /* Suspend until signal occurs */
```

- Pending signals are not queued
 - For each signal type, just have single bit indicating whether or not signal is pending even if multiple processes have sent this signal

```
yunmin@peace:~/ch10$ ./signal2
Received signal 17 from process 34372
Received signal 17 from process 34373
Received signal 17 from process 34374
Received signal 17 from process 34375
```

Program does not exit



Living With Nonqueuing Signals

```
int ccount = 0;
                                                       signal3.c
// SIGCHLD handler that reaps all terminated children
void child handler2(int sig)
    int child status;
    pid t pid;
    while ((pid = wait(&child status)) > 0) {
        ccount--;
        printf("Received signal %d from process %d\n", sig, pid);
// Using a handler that reaps multiple children
void main()
    pid_t pid[N];
    int i;
    ccount = N;
    signal(SIGCHLD, child handler2);
    for (i = 0; i < N; i++)
        if ((pid[i] = fork()) == 0) {
            sleep(1);
            exit(0); /* Child exits */
    while (ccount > 0); /* Parent spins */
```

- Must check for all terminated jobs
 - Typically loop with wait

```
yunmin@yunmin:~/ch3_2$ ./signal3
Received signal 17 from process 2875
Received signal 17 from process 2876
Received signal 17 from process 2877
Received signal 17 from process 2879
Received signal 17 from process 2878
```

Example for SIGINT

A program that reacts to externally generated events (ctrl-c)

```
#include <stdlib.h>
                                                       sigint.c
#include <stdio.h>
#include <unistd.h>
#include <signal.h>
void handler(int sig) {
    printf("You think hitting ctrl-c will stop the bomb?\n");
    sleep(2);
    printf("Well...");
    fflush(stdout);
    sleep(1);
    printf("OK\n");
    exit(0);
void main() {
    signal(SIGINT, handler); /* installs ctl-c handler */
    while(1) {
```

```
Typed ctrl-c

yummin@yunmin:~/ch3_2$ ./sigint

^CYou think hitting ctrl-c will stop the bomb?

Well...OK
```



Example for SIGALRM

A program that reacts to internally generated events

```
int beeps = 0;
                                                        sigalrm.c
/* SIGALRM handler */
void handler(int sig) {
    printf("BEEP\n");
    fflush(stdout);
    if (++beeps < 5)
        alarm(1);
    else {
        printf("BOOM!\n");
        exit(0);
void main() {
    signal(SIGALRM, handler);
    alarm(1); /* send SIGALRM in 1 second */
    while (1) {
        /* handler returns here */
```

```
yunmin@yunmin:~/ch3_2$ ./sigalrm

BEEP

BEEP

BEEP

BEEP

BEEP

BOOM!
```

sigaction()

sigaction(): Examine and change a signal action

- Used to change the action taken by a process on receipt of a specific signal.
- signum: the name of signal
- act: the new action
- oldact: the previous actions
- Return value
 - Success: 0
 - Error: -1

struct sigaction

```
struct sigaction {
   void (*sa_handler)(int);
   void (*sa_sigaction)(int, siginfo_t *, void *);
   sigset_t sa_mask;
   int sa_flags;
   void (*sa_restorer)(void);
};
```

- sa_handler: the address of a function to be called when the signal occurs
- sa_mask: specifies a set of signals that will be blocked when the signal handler is called
- sa_flags: options. Basically 0



sigaction(): Example

```
#include <stdio.h>
#include <unistd.h>
#include <signal.h>

// Signal handler for timeout
void timeout(int sig)
{
   if (sig == SIGALRM)
        puts("Time out!");

   // generate SIGALRM after 2 seconds
   alarm(2);
}
```

```
yunmin@yunmin:~/ch3_2$ ./sigaction
wait...
Time out!
wait...
Time out!
wait...
Time out!
```

```
int main(int argc, char *argv[])
    int i;
    struct sigaction act;
    act.sa_handler = timeout;
    sigemptyset(&act.sa_mask);
    act.sa_flags = 0;
    // specify signal type and signal handler
    sigaction(SIGALRM, &act, 0);
    alarm(2);
    for (i = 0; i < 3; i++)
        puts("wait...");
        sleep(10);
    return 0;
```



Agenda

- Process Concept
- Process State and Scheduling
- Process API
- Signal
- Inter Process Communication





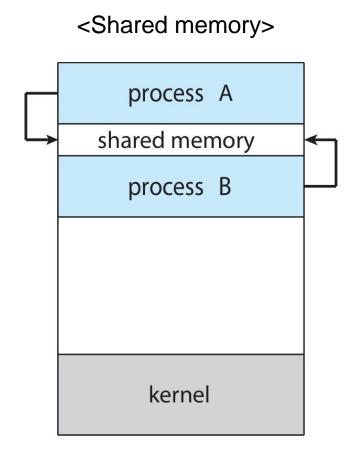
Inter Process Communication

- Processes within a system may be independent or cooperating
 - Cooperating process can affect or be affected by other processes, including sharing data
- Reasons for cooperating processes
 - Information sharing
 - Computation speedup
 - Modularity
- Cooperating processes need inter process communication (IPC)
- Two models of IPC
 - Shared memory
 - Message passing



Shared Memory

- Shared memory: an area of memory shared among the processes that wish to communicate
- The communication is under the control of the users processes not the operating system
- Major issues is to provide mechanism that will allow the user processes to synchronize their actions when they access shared memory
 - Synchronization is discussed in great details in Chapter 5.





Shared Memory

- POSIX shared memory
 - Communication through memory-mapped file.
 - Process creates or open shared memory segment

```
shm_fd = shm_open(name, O_CREAT | O_RDWR, 0666);
```

- O_CREAT: create if it does not yet exist / O_RDWR: open for reading and writing
- Set the size of the object

```
ftruncate(shm_fd, 4096);
```

Map shared memory segment to process address space

```
ptr = mmap(0, 4096, PROT_WRITE, MAP_SHARED, shm_fd, 0);
```

Now the process could write to the shared memory

```
sprintf(ptr, "Writing to shared memory");
```



Shared Memory: Example - Producer

```
#include <stdio.h>
                                             shm pro.c
#include <stdlib.h>
#include <unistd.h>
#include <string.h>
#include <fcntl.h>
#include <sys/shm.h>
#include <sys/stat.h>
#include <sys/mman.h>
#include <sys/types.h>
int main() {
    const int SIZE = 4096;
    const char *name = "OS";
    const char *message0 = "Studying ";
    const char *message1 = "Operating Systems ";
    const char *message2 = "Is Fun!\n";
    int shm fd;
    void *ptr;
    // create the shared memory segment
    shm fd = shm open(name, O CREAT | O RDWR, 0666);
```

```
// configure the size of the shared memory segment
ftruncate(shm fd,SIZE);
// now map the shared memory segment
// in the address space of the process
ptr = (char*) mmap(0,SIZE, PROT READ | PROT WRITE,
                   MAP SHARED, shm fd, 0);
if (ptr == MAP FAILED) {
    printf("Map failed\n");
    return -1;
// Now write to the shared memory region.
// Note we must increment the value of ptr
// after each write.
sprintf(ptr, "%s", message0);
ptr += strlen(message0);
sprintf(ptr, "%s", message1);
ptr += strlen(message1);
sprintf(ptr, "%s", message2);
ptr += strlen(message2);
return 0;
```



Shared Memory: Example - Consumer

```
#include <stdio.h>
                                             shm con.c
#include <stdlib.h>
#include <unistd.h>
#include <fcntl.h>
#include <sys/shm.h>
#include <sys/stat.h>
#include <sys/mman.h>
int main()
    const char *name = "OS";
    const int SIZE = 4096;
   int shm fd;
   void *ptr;
    int i;
    // open the shared memory segment
    shm_fd = shm_open(name, O_RDONLY, 0666);
    if (shm fd == -1) {
        printf("shared memory failed\n");
        exit(-1);
```

```
// now map the shared memory segment
// in the address space of the process
ptr = (char*) mmap(0,SIZE, PROT READ, MAP SHARED,
                    shm fd, 0);
if (ptr == MAP FAILED) {
    printf("Map failed\n");
    exit(-1);
// now read from the shared memory region
printf("%s\n",(char *)ptr);
// remove the shared memory segment
if (shm unlink(name) == -1) {
    printf("Error removing %s\n", name);
    exit(-1);
return 0;
  yunmin@yunmin:~/ch3 2$ gcc shm pro.c -o shm pro -lrt
  yunmin@yunmin:~/ch3 2$ gcc shm con.c -o shm con -lrt
  yunmin@yunmin:~/ch3 2$ ./shm_pro
  yunmin@yunmin:~/ch3 2$ ./shm con
  Studying Operating Systems Is Fun!
  yunmin@yunmin:~/ch3_2$ ./shm_con
  shared memory failed
```



Shared Memory: Reference

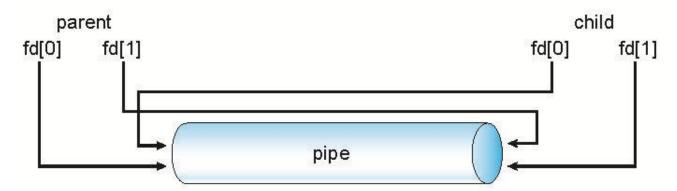
Search the following topics from Internet and study them

- POSIX shared memory (requires compilation option '-Irt')
 - shm_open(): http://man7.org/linux/man-pages/man3/shm_open.3.html
 - shm_unlink(): http://man7.org/linux/man-pages/man3/shm_unlink.3p.html
 - ftruncate(): http://man7.org/linux/man-pages/man3/ftruncate.3p.html
 - mmap(): http://man7.org/linux/man-pages/man2/mmap.2.html
 - munmap(): http://man7.org/linux/man-pages/man3/munmap.3p.html



Pipes

Pipes acts as a conduit allowing two processes to communicate

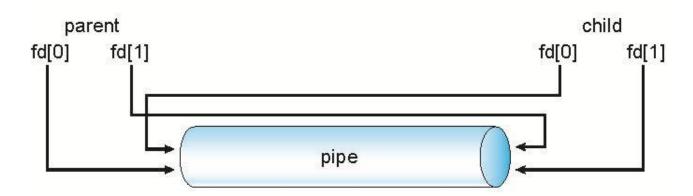


- Ordinary pipes
 - Unidirectional communication between parent and child
 - Typically, a parent process creates a pipe and uses it to communicate with a child process that it created
- Named pipes
 - Can be bidirectional and no parent-child relationship is required



Ordinary Pipes

- Ordinary pipes allow unidirectional communication in standard producer-consumer style
 - Producer writes to one end (the write-end of the pipe)
 - Consumer reads from the other end (the read-end of the pipe)
 - Require parent-child relationship between communicating processes



Windows calls these anonymous pipes



Ordinary Pipes: Example

```
#include <sys/types.h>
                                                        pipe.c
#include <stdio.h>
#include <string.h>
#include <unistd.h>
#define BUFFER SIZE 25
#define READ END
#define WRITE END 1
int main(void)
    char write msg[BUFFER SIZE] = "Greetings";
    char read msg[BUFFER SIZE];
    pid t pid;
    int fd[2];
    /* create the pipe */
    if (pipe(fd) == -1) {
        fprintf(stderr, "Pipe failed");
        return 1;
    /* now fork a child process */
    pid = fork();
    if (pid < 0) {
        fprintf(stderr, "Fork failed");
        return 1;
```

```
if (pid > 0) { /* parent process */
       /* close the unused end of the pipe */
       close(fd[READ END]);
       /* write to the pipe */
       write(fd[WRITE END], write msg, strlen(write msg)+1);
       /* close the write end of the pipe */
       close(fd[WRITE END]);
   else { /* child process */
       /* close the unused end of the pipe */
       close(fd[WRITE END]);
       /* read from the pipe */
       read(fd[READ END], read msg, BUFFER SIZE);
       printf("child read %s\n",read msg);
       /* close the write end of the pipe */
       close(fd[READ END]);
   return 0;
```

```
yunmin@yunmin:~/ch3_2$ gcc pipe.c -o pipe
yunmin@yunmin:~/ch3_2$ ./pipe
child read Greetings
```

Named Pipes

- Named pipes are more powerful than ordinary pipes
 - Communication is bidirectional
 - No parent-child relationship is necessary between the communicating processes
 - Several processes can use the named pipe for communication
 - Provided on both UNIX and Windows systems
 - Named pipes are referred to as FIFOs in UNIX systems



Named Pipes: Example

```
// This side writes first, then reads
                                                    npipe1.c
int main() {
   int fd;
    char *myfifo = "/tmp/myfifo";
                                            // FIFO file path
    // Creating the named file(FIFO)
    mkfifo(myfifo, 0666);
    char arr1[80], arr2[80];
    while (1) {
        // Open FIFO for write only
        fd = open(myfifo, O WRONLY);
        // Take an input arr2 from user. 80 is maximum length
        fgets(arr2, 80, stdin);
        // Write the input arr2 on FIFO and close it
        write(fd, arr2, strlen(arr2)+1);
        close(fd);
        // Open FIFO for Read only
        fd = open(myfifo, O RDONLY);
        // Read from FIFO
        read(fd, arr1, sizeof(arr1));
        // Print the read message
        printf("User2: %s\n", arr1);
        close(fd);
    return 0;
```

```
// This side reads first, then reads
                                                            npipe2.c
      int main() {
          int fd1;
          char *myfifo = "/tmp/myfifo";
                                                   // FIFO file path
          // Creating the named file(FIFO)
          mkfifo(myfifo, 0666);
          char str1[80], str2[80];
          while (1) {
              // First open in read only and read
              fd1 = open(myfifo, O RDONLY);
              read(fd1, str1, 80);
              // Print the read string and close
              printf("User1: %s\n", str1);
              close(fd1);
              // Now open in write mode and write string taken from user.
              fd1 = open(myfifo, O WRONLY);
              fgets(str2, 80, stdin);
              write(fd1, str2, strlen(str2)+1);
              close(fd1);
          return 0;
yunmin@yunmin:~/ch3 2$ ./npipe1
                                       yunmin@yunmin:~/ch3_2$ ./npipe2
                                       User1: Hello!
```

```
yunmin@yunmin:~/ch3_2$ ./npipe1

Hello!

User1: Hello!

User2: Handong!

Operating Systems

User2: University

University

University

Vunmin@yunmin:~/ch3_2$ ./npipe2

User1: Hello!

User1: Operating Systems

University
```

Message Passing: Reference

Search the following topics from Internet and study them

- POSIX message passing
 - msgget(): create a message queue
 - https://man7.org/linux/man-pages/man2/msgget.2.html
 - msgsnd(): send a message to a message queue
 - https://man7.org/linux/man-pages/man3/msgsnd.3p.html
 - msgrcv(): receive a message from a message queue
 - https://man7.org/linux/man-pages/man3/msgrcv.3p.html
 - msgctl(): control/deallocate message queue
 - https://man7.org/linux/man-pages/man2/msgctl.2.html

