파이썬으로 배우는 데이터 구조

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Data Structures in Python
Chapter 2 - 2

- Performance Analysis
- Big-O Notation
- Big-O Properties
- **Growth Rates**
- Growth Rates Examples

Agenda & Reading

- **Growth Rate**
  - **Comparison**
  - **Profiling and Prediction**
- Growth Rate Examples
  - Python List & Dictionary

- References:
  - Textbook: Problem Solving with Algorithms and Data Structures
    - Chapter 3. Analysis
  - Textbook: www.github.idebtor/DSpy
    - Chapter 2.1 ~ 3

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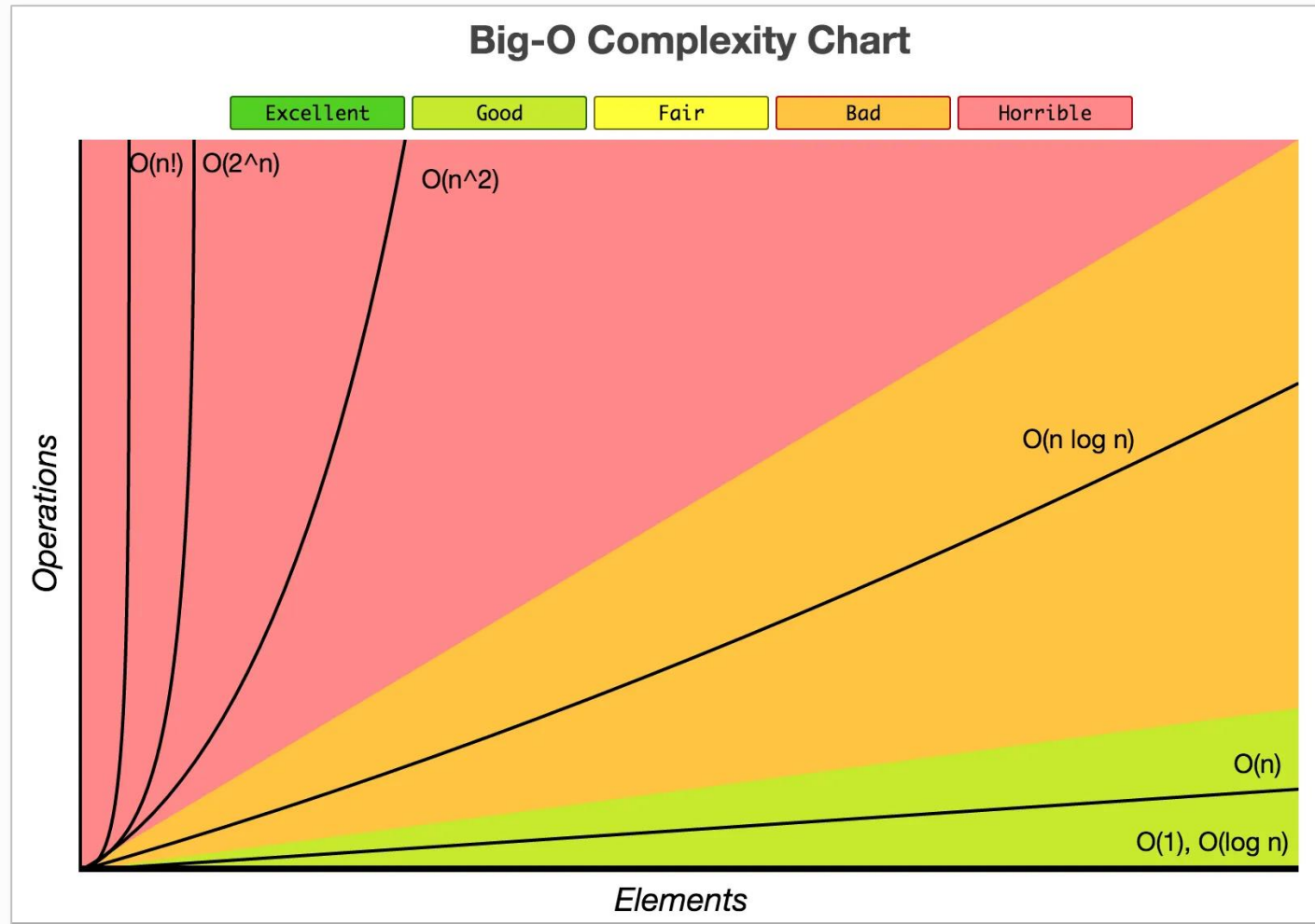3

1 Growth Rate Comparison - Hypothetical Running Time

- The running time on a hypothetical computer that computes 10^6 operations per second for varies problem sizes

$$O(1) \ < \ O(\log n) \ < \ O(n) \ < \ O(n \log n) \ < \ O(n^2) \ < \ O(n^3) \ < \ O(2^n)$$

Notation			n = 10	n = 10^2	n = 10^3	n = 10^4	n = 10^5	n = 10^6
O(1)	Constant	상수	1 µsec	1 µsec	1 µsec	1 µsec	1 µsec	1 µsec
O(log(n))	Logarithmic	대수 함수	3 µsec	7 µsec	10 µsec	13 µsec	17 µsec	20 µsec
O(n)	Linear	선형 함수	10 µsec	100 µsec	1 msec	10 msec	100 msec	1 sec
O(nlog(n))	N log N	선형 대수 함수	33 µsec	664 µsec	10 msec	13.3 msec	1.6 sec	20 sec
O(n²)	Quadratic	2차 함수	100 µsec	10 msec	1 sec	1.7 min	16.7 min	11.6 days
O(n³)	Cubic	3차 함수	1 msec	1 sec	16.7 min	11.6 days	31.7 years	31709 years
O(2ⁿ)	Exponential	지수 함수	10 msec	3e17 years				

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1 Growth Rate Comparison

$$O(1) \; < \; O(\log n) \; < \; O(n) \; < \; O(n \log n) \; < \; O(n^2) \; < \; O(n^3) \; < \; O(2^n)$$

A comparison of growth-rate functions in graphical form

1 Growth Rate Comparison

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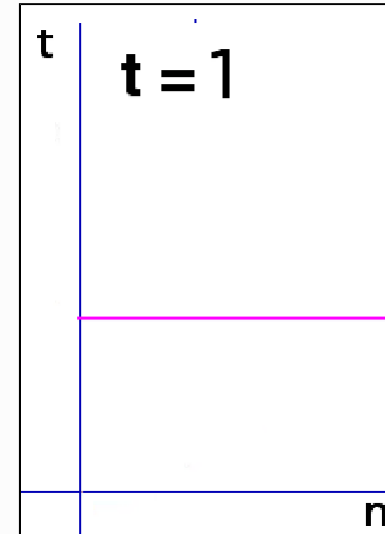
$$T(N) \approx aN^b$$
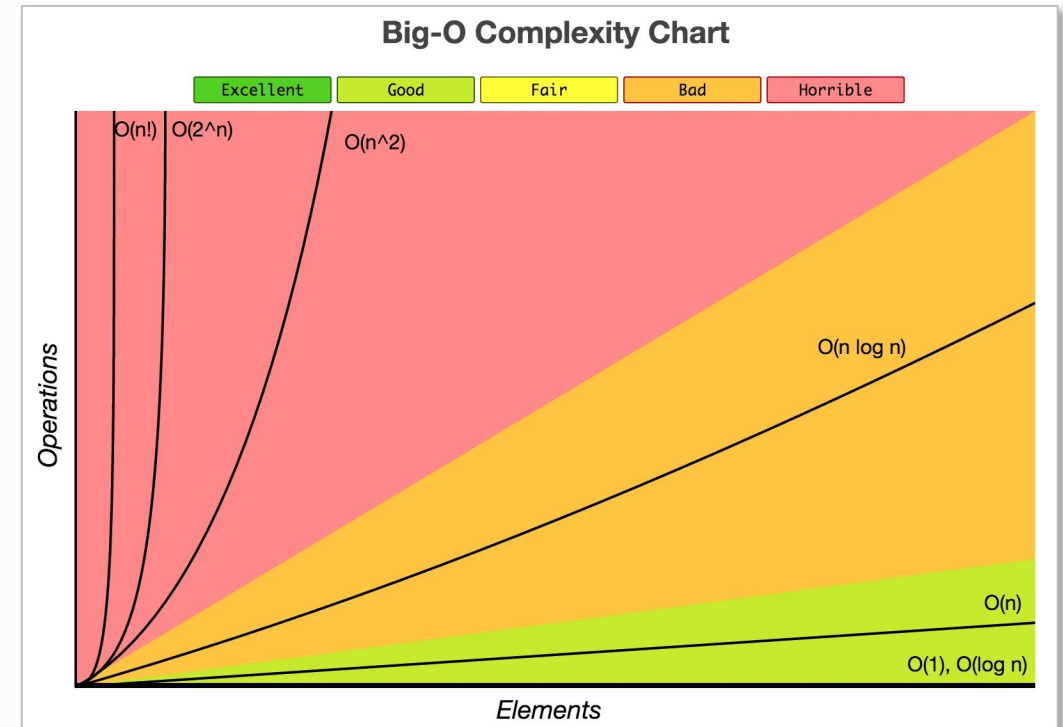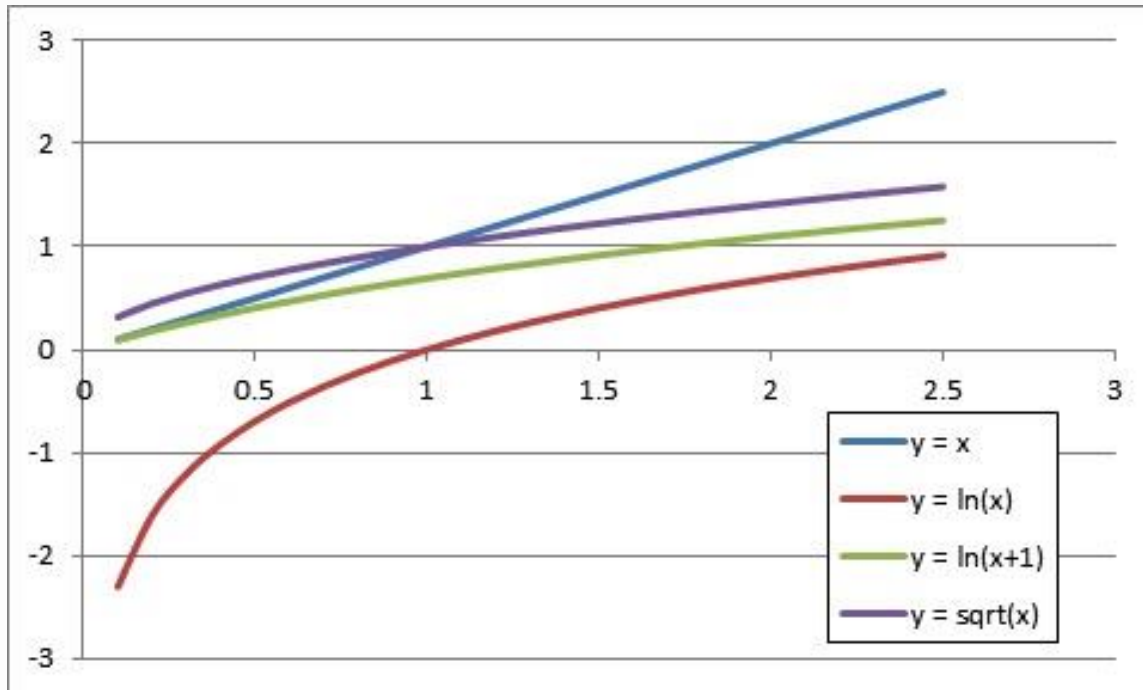
A comparison of growth-rate functions in graphical form

2 Growth Rate Examples – Constant Growth Rate – O(1)

- Time requirement is constant and, therefore, **independent of the problem's size n.**

```python
def rate1(n):
    s = "SWEAR"
    for i in range(25):
        print("I must not ", s)
```

$t = 1$

n	10^1	10^2	10^3	10^4	10^5	10^6
O(1)	1	1	1	1	1	1

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2 Growth Rate Examples - Logarithmic Growth Rate - O(log n)

- Increase **slowly** as the problem size increases.
- If you square the problem size, you only double its time requirement.
- The base of the log does not affect a log growth rate, so you can omit it.

```python
def rate2(n):
    s = "YELL"
    i = 1
    while i < n:
        print("I must not ", s)
        i = i * 2
```
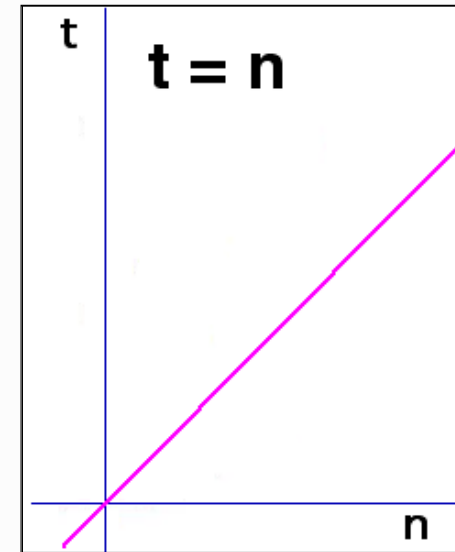
$$t = \log_2(n)$$
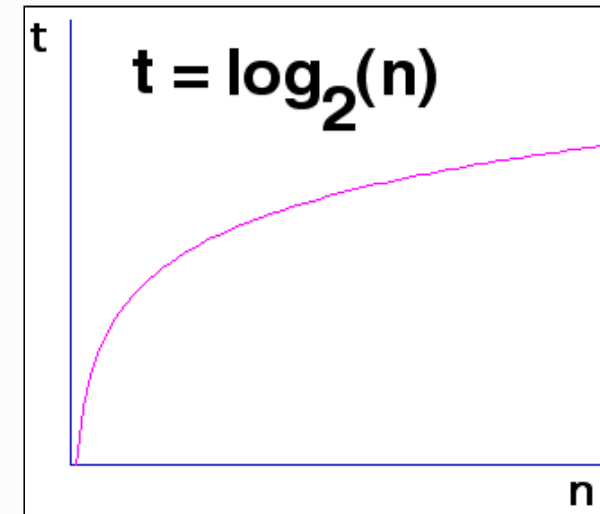
n	10^1	10^2	10^3	10^4	10^5	10^6
$O(\log_2 n)$	3	6	9	13	16	19

2 Growth Rate Examples - Linear Growth Rate - O(n)

- The time increases directly with the sizes of the problem.
- If you square the problem size, you also square its time requirement.

```python
def rate3(n):
    s = "FIGHT"
    for i in range(n):
        print("I must not ", s)
```
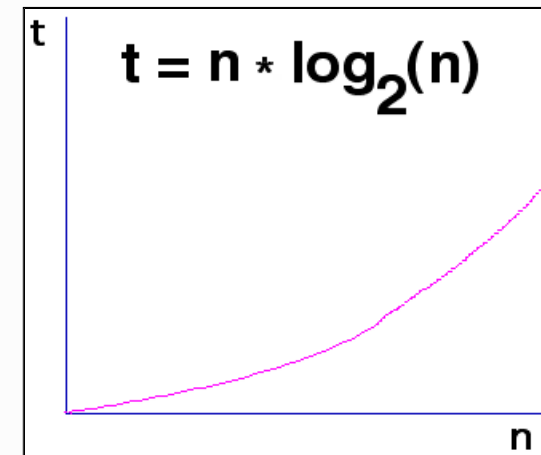
$$t = n$$

n	10^1	10^2	10^3	10^4	10^5	10^6
O(n)	10	10^2	10^3	10^4	10^5	10^6

2 Growth Rate Examples - n* log n Growth Rate - O(n log(n))

- The time requirement increases more rapidly than a linear algorithm.
- Such algorithms usually divide a problem into smaller problem that are each solved separately.

```
def rate4(n):
    s = "HIT"
    for i in range(n):
        j = n
        while j > 1:
            print("I must not ", s)
            j = j // 2
```
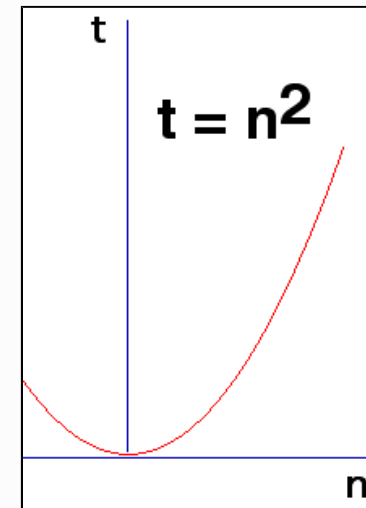
$$t = n * \log_2(n)$$

n	10^1	10^2	10^3	10^4	10^5	10^6
O(nlog(n))	30	664	9965	10^5	10^6	10^7

2 Growth Rate Examples - Quadratic Growth Rate - $O(n^2)$

- The time requirement increases rapidly with the size of the problem.
- Algorithms that use two nested loops are often quadratic.

```
def rate5(n):
    s = "LIE"
    for i in range(n):
        for j in range(n):
            print("I must not ", s)
```

$t = n^2$

n	10^1	10^2	10^3	10^4	10^5	10^6
$O(n^2)$	10^2	10^4	10^6	10^8	10^{10}	10^{12}

2 Growth Rate Examples – Cubic Growth Rate – O(n^3)

- The time requirement increases more rapidly with the size of the problem than the time requirement for a quadratic algorithm.
- Algorithms that use three nested loops are often quadratic and are practical only for small problems.

```
def rate6(n):
    s = "SPACE OUT IN CLASS"
    for i in range(n):
        for j in range(n):
            for k in range(n):
                print("I must not ", s)
```
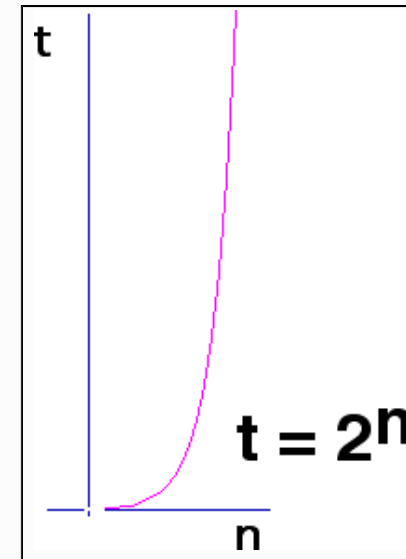

$t = n^3$
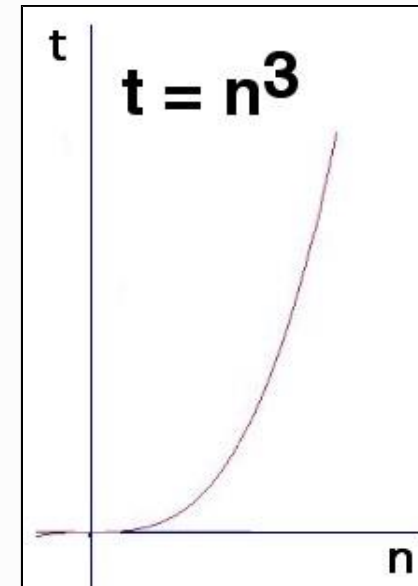
n	10^1	10^2	10^3	10^4	10^5	10^6
O(n^3)	10^3	10^6	10^9	10^{12}	10^{15}	10^{18}

2 Growth Rate Examples - Exponential Growth Rate - $O(2^n)$

- As the size of a problem increases, the time requirement usually increases too rapidly to be practical.

```python
def rate7(n):
    s = "ZONE OUT IN CLASS"
    for i in range(2 ** n):
        print("I must not ", s)
```

$t = 2^n$

n	10^1	10^2	10^3	10^4	10^5	10^6
$O(2^n)$	10^3	10^{30}	10^{301}	10^{3010}	10^{30103}	10^{301030}

Exercise

- What is the Big-O of the following statements?

```
for i in range(n):          executed n times
    for j in range(10):     executed 10 times
        print(i, j)         constant time
```

  - Running time, T(n) = n * 10 * 1 = 10n, Big-O =

- What is the Big-O of the following statements? Big-O =

```
for i in range(n):          executed n times
    for j in range(n):      executed n times
        print(i, j)
for k in range(n):          executed n times
    print(k)
```

- The first set of nested loops is $O(n^2)$ and the second loop is $O(n)$.
  This is $O(\max(n^2,n))$ Big-O =

Exercise

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for i in range(n):          executed n times
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  - Running time, T(n) = n * 10 * 1 = 10n, Big-O = O(n)

- What is the Big-O of the following statements? Big-O =

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for i in range(n):          executed n times
    for j in range(n):      executed n times
        print(i, j)
for k in range(n):          executed n times
    print(k)
```

- The first set of nested loops is $O(n^2)$ and the second loop is $O(n)$.
  This is $O(\max(n^2, n))$ Big-O = $O(n^2)$

Quiz

- What is the Big-O of the following statements?

```
for i in range(n):
    for j in range(i+1, n):
        print(i, j)
```

- When i is 0, the inner loop executes (n - 1) times.
  When i is 1, the inner loop executes (n - 2) times.
  ...
  When i is (n - 2), the inner loop executes once.
- The number of times the inner loop statements execute:
  - Running time, T(n) =
  - Big-O =

3 Profiling: Measuring Growth Rate

Problem: Predict the running time of a big data set (i.e., n = 1 million or 1 billion).

- Most algorithms approximately have the order of growth of the running time:
$$T(N) \approx aN^b$$

- For example, we may compute the constant "a" and the growth rate "b" from data we got from profiling (i.e., performance analysis) as shown below.

N	sec
1000	0.000770
2000	0.002855
3000	0.006579
4000	0.011144
5000	0.014565
6000	0.023295
7000	0.028571
8000	0.036643
9000	0.047810
10000	0.063062

3 Profiling: Measuring Growth Rate

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- **Since T(N) $\approx$ a N^b , T(2N) = a (2N)^b, then**

$$\frac{T(2N)}{T(N)} = \frac{a(2N)^b}{aN^b} = \frac{2^b(N)^b}{N^b} = 2^b$$

Take log both sides

$$log\frac{T(2N)}{T(N)} = \log 2^b$$

$$b = log\frac{T(2N)}{T(N)}$$

3 Profiling: Measuring Growth Rate - Example

- **Example:** let us choose N = 4000 or 2N = 8000, an average case of the insertion sort shown above. Recall that log we use here is **log base 2.**

$$b = \log\frac{T(2N)}{T(N)} = \log\frac{t2(8000)}{t1(4000)} = \log\frac{0.036643}{0.011144} = 1.717$$

- Now, we use this b = 1.717 to solve for a when $N = 4000,\ T(N) = 0.011144$ in the following:

$$T(N) = a\,N^{1.717}$$
$$0.011144 = a\,(4000)^{1.717}$$
$$a = \frac{0.011144}{(4000)^{1.717}}$$
$$a = 7.27\text{x}10^{-9}$$

- Therefore, we have the growth rate b = 1.717, **the constant** $a = 7.27\text{x}10^{-9}$ for the insertion sort average case.

Summary

- Performance analysis or profiling measures an algorithm's time requirement as a function of the problem size **n** by using a growth-rate function.

- The **growth rates** shown below are commonly used :

$$O(1) <\ O(\log n) <\ O(n) <\ O(n \log n) <\ O(n^2) <\ O(n^3) <\ O(2^n)$$

- Generally, **growth rates** can be measured in a form of the following:

$$T(N) \approx\ aN^b$$

Data Structures in Python
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- Big-O Properties
- **Growth Rates**
- Growth Rates Examples