Data Structures in Python

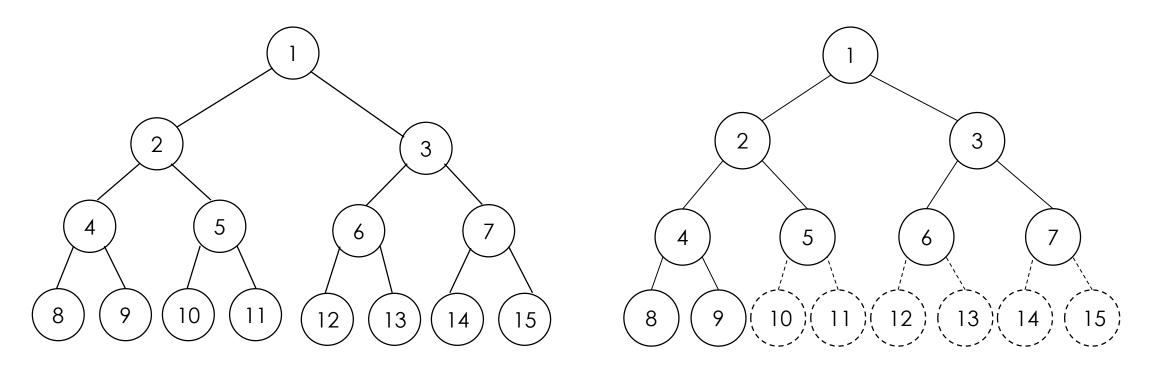
- Heap and Priority Queue
- Heap Coding
- Heap Sort & Min/MaxHeap

Agenda & Readings

- Heap and Priority Queue
 - Complete Binary Tee (Review)
 - Heap and Priority Queue
 - Heap ADT
 - Time Complexity
- Reference:
 - Problem Solving with Algorithms and Data Structures

Binary trees - Properties

- **Definition:** A full binary tree of level k is a binary tree having 2^k 1 nodes, $k \ge 0$.
- **Definition**: A binary tree with n nodes and level k is **complete** if and only if its nodes correspond to the nodes numbered from 1 to n in the full binary tree of level k.

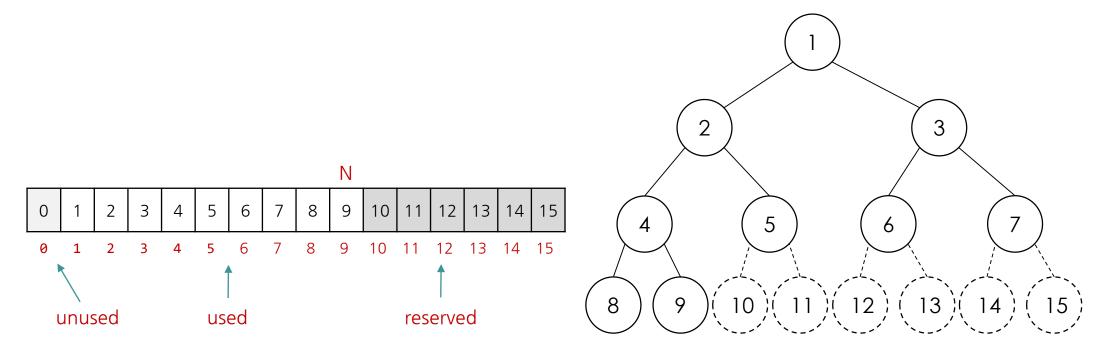


A **full** binary tree

A **complete** binary tree

Binary trees - Array representation

- A complete binary tree with n nodes, any node index i, $1 \le i \le n$, we have
 - parent(i) is at $\lfloor i/2 \rfloor$ If i = 1, i is at the root and has no parent
 - leftChild(i) is at 2i if 2i <= n. If 2i > n, then i has no left child.
 - rightChild(i) is at 2i+1 if 2i+1 <= n. If 2i+1 > n, then i has no right child.



A **complete** binary tree

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 - Because it provides an efficient implementation for priority queues.

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- Priority queues.
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 - Example: A line waiting to be served at a bank and served FIFO except if a senior or a disabled person arrives in the line. They are served first. Seniors and disabled persons have higher priority than others.
- A typical ADT for Priority Queue
 - Get the top priority element (min or max)
 - Insert an element
 - Delete the top priority element
 - Decrease the priority of an element

- O(1)
- O(log n)
- O(log n)
- O(log n)

- Challenge: Find the largest M items in a stream of N items.
- Constraints: Not enough memory to store N items.



Order of insert of finding the largest M in a stream of N items

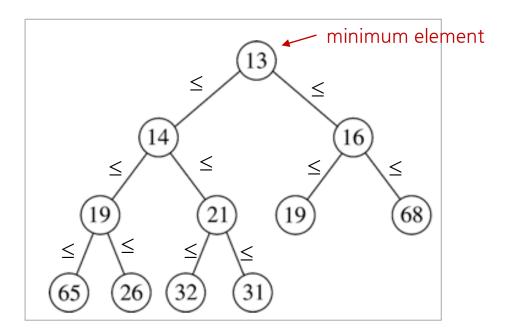
implementation	insert	delete	min/max
unordered array	1	Ν	N
ordered array	Ν	1	N huge
goal	log N	log N	log N
	<u> </u>	<u> </u>	
	Mission	Impossible?	priority

Binary heap

- Binary heap: array representation of a heap-ordered complete binary tree
- Properties:
 - Heap-ordered: Parent's key no smaller than children's keys. [max-heap]
 - Heap-structure:
 A complete binary tree

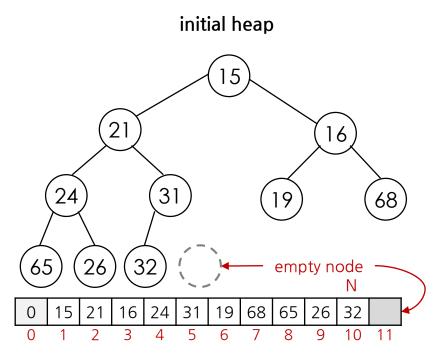
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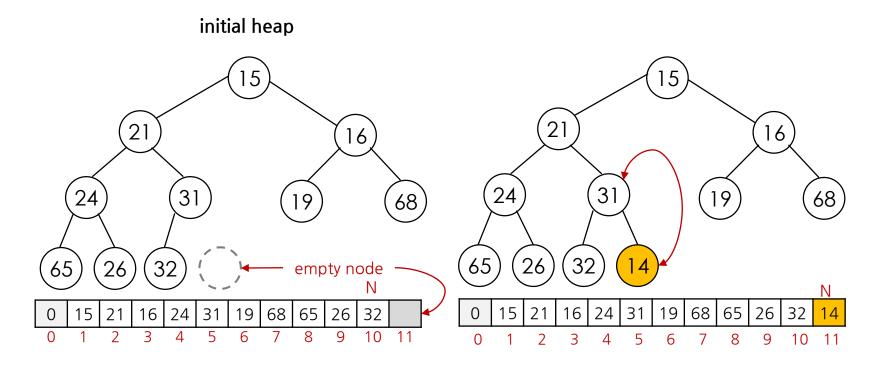


- Duplicates are allowed
- No order implied for elements which do not share ancestor-descendant relationship

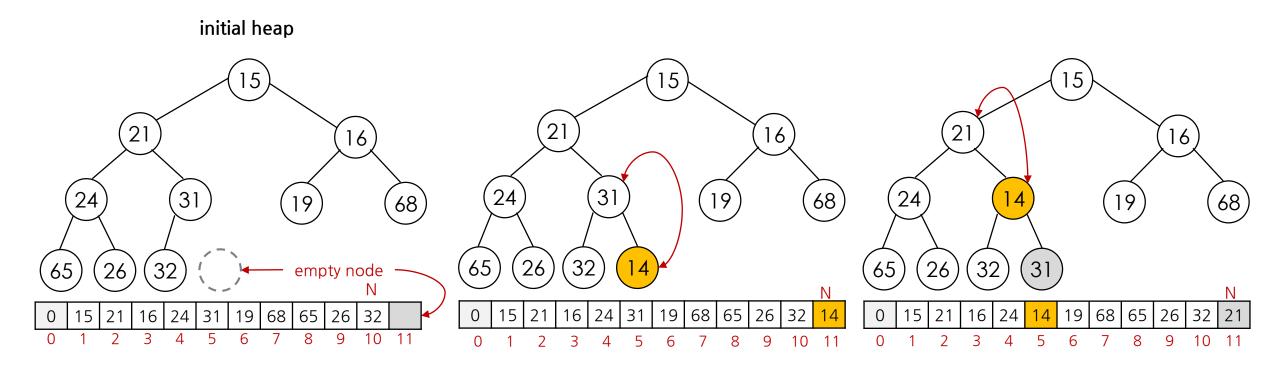
- Insert a new element while maintaining a heap-structure
- Move the element up the heap while not satisfying heap-ordered



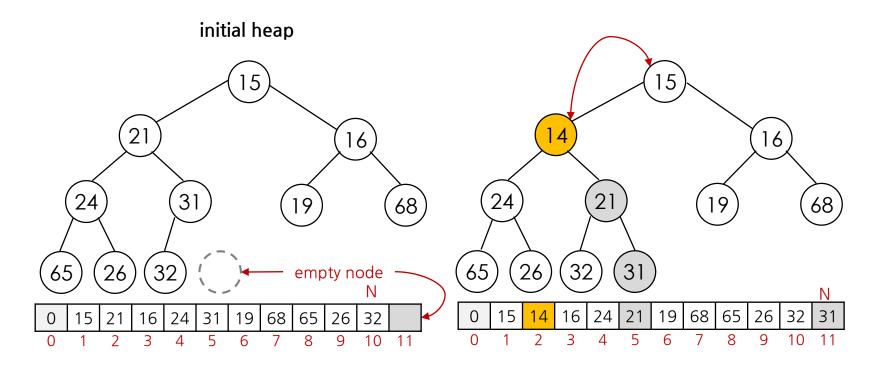
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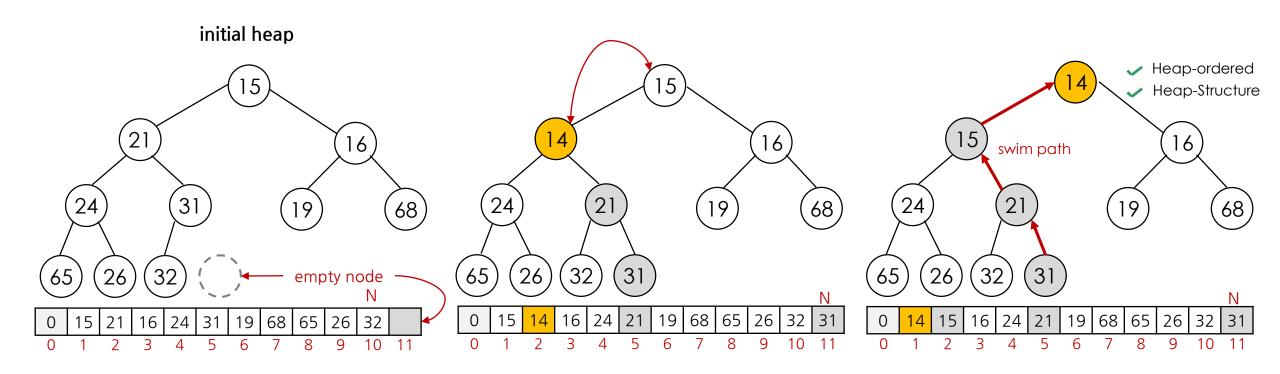
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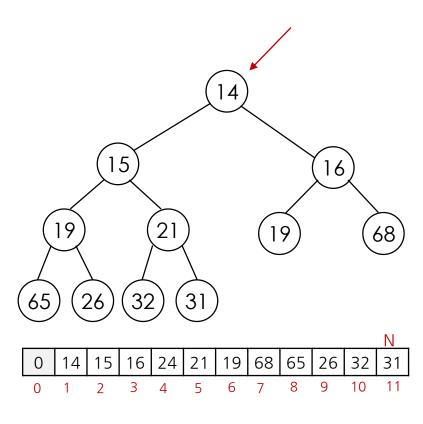


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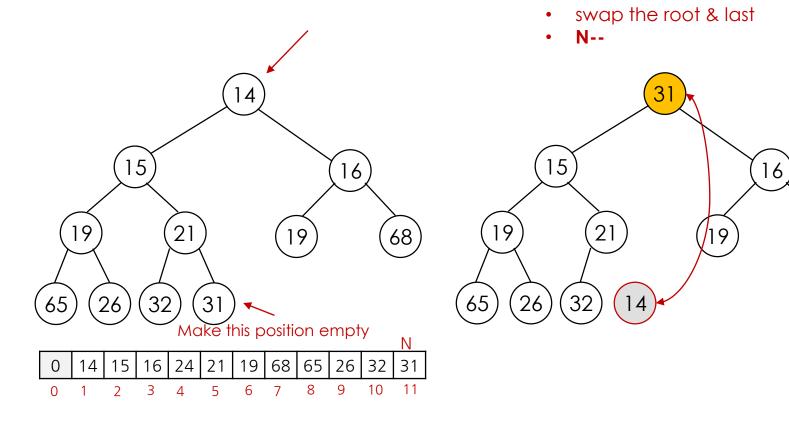


- Swap the root and the last element.
- Heap decreases by one in size.
- Move down (sink) the root while not satisfying heap-ordered.
 - Minimum element is always at the root (by min-heap definition).

Which position of the node will be empty?

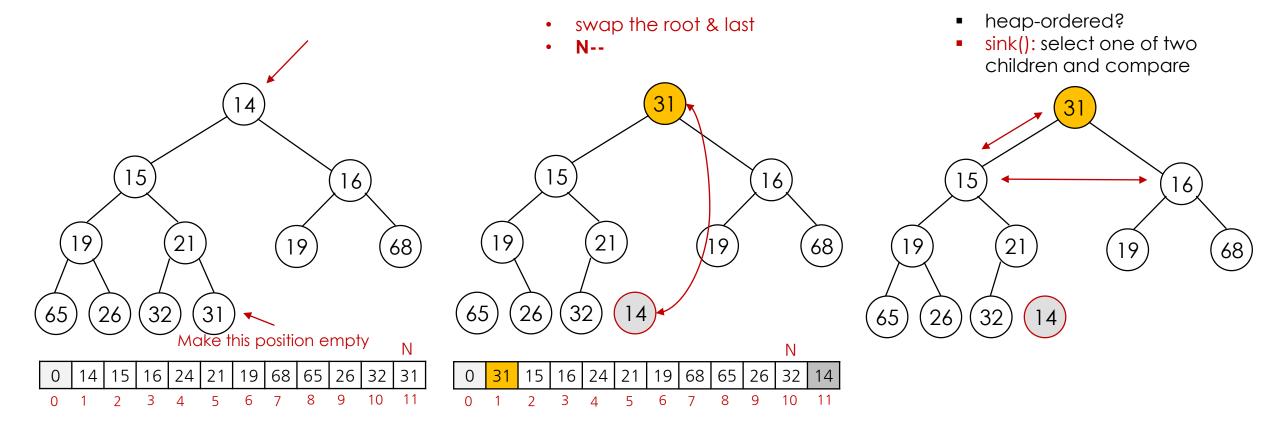


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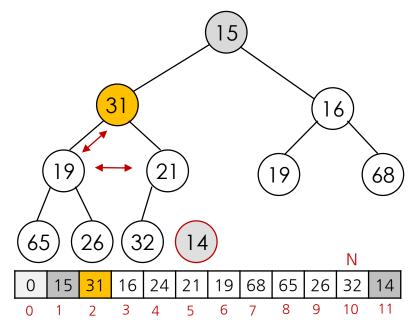


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Which position of the node will be empty?

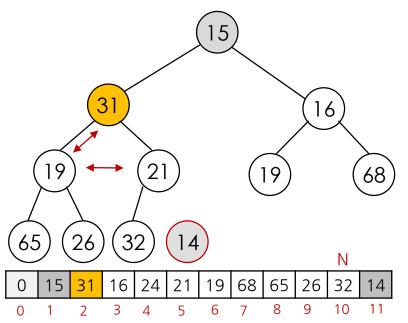


- heap-ordered?
- sink(): select one of two children and compare



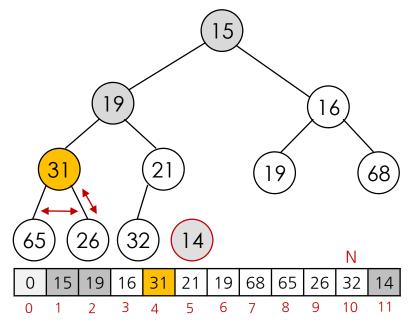
- Is 31 > min(14,16)?
- Yes swap 31 with min(14,16)

- heap-ordered?
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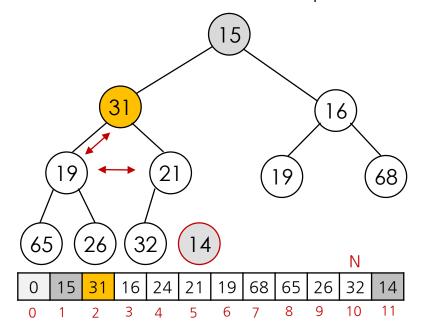
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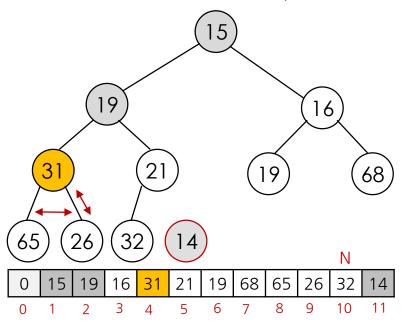
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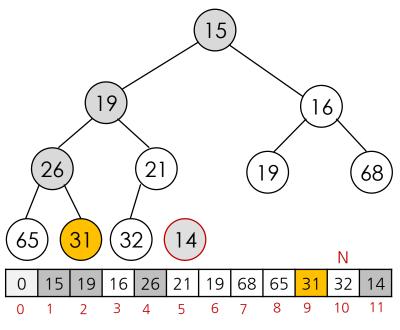
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- ls 31 > min(19,21)?
- Yes swap 31 with min(19,21)

- heap-ordered?
- sink(): select one of two children and compare



- Is 31 > min(65,26)?
- Yes swap 31 with min(65,26)
- Heap-ordered
- ✓ Heap-Structure

Binary heap: Time complexity:

- Level of heap is $\lfloor \log_2 N \rfloor$
- insert: O(log N) for each insert
 - In practice, expect less
- delete: O(log N) // deleting root node or any node
- increase/decrease key: O(log N)

Implementation	Insert	Delete	max
Unordered array	1	N	Ν
Ordered array	Ν	1	1
Binary heap	log N	log N	1

Data Structures in Python

- Heap and Priority Queue
- Heap Coding
- Min/MaxHeap and Heap sort

Proof:

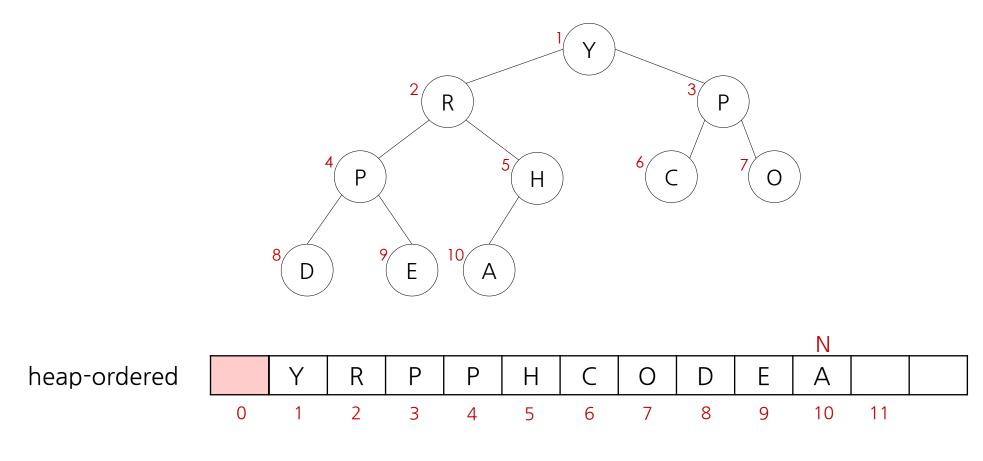
https://stackoverflow.com/questions/9755721/how-can-building-a-heap-be-on-time-complexity https://www.insertingwiththeweb.com/data-structures/binary-heap/build-heap-proof/ https://www.guora.com/How-is-the-time-complexity-of-building-a-heap-is-o-n

References in Korean:

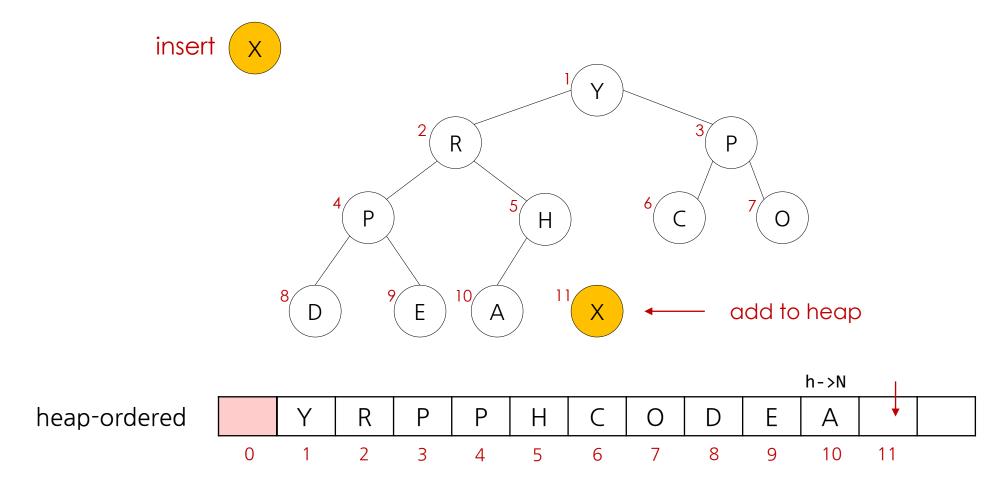
https://ratsgo.github.io/data%20structure&algorithm/2017/09/27/heapsort/https://zeddios.tistory.com/56

Prof. Youngsup Kim, idebtor@gmail.com, CSEE Dept., Grace School Rm204, Handong Global University

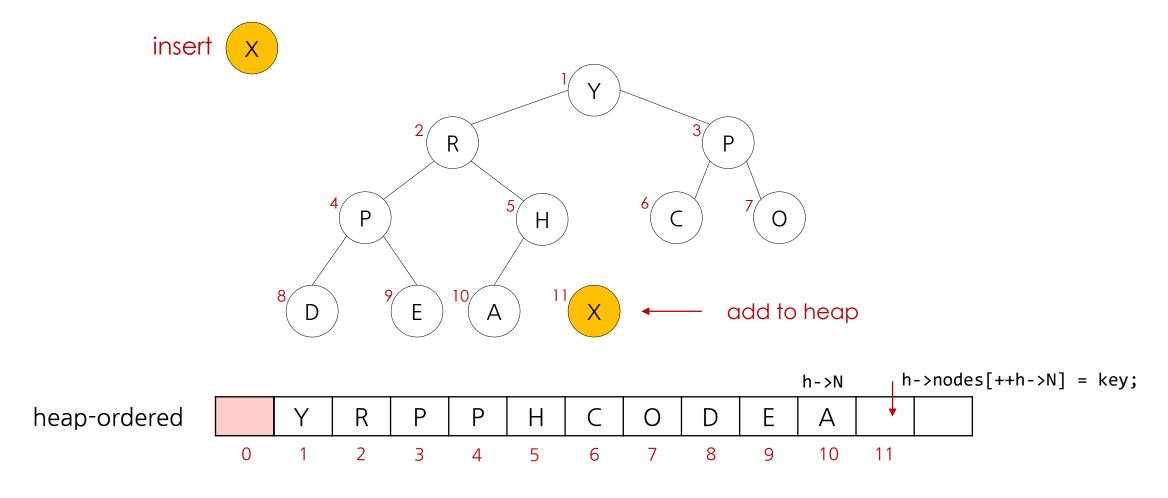
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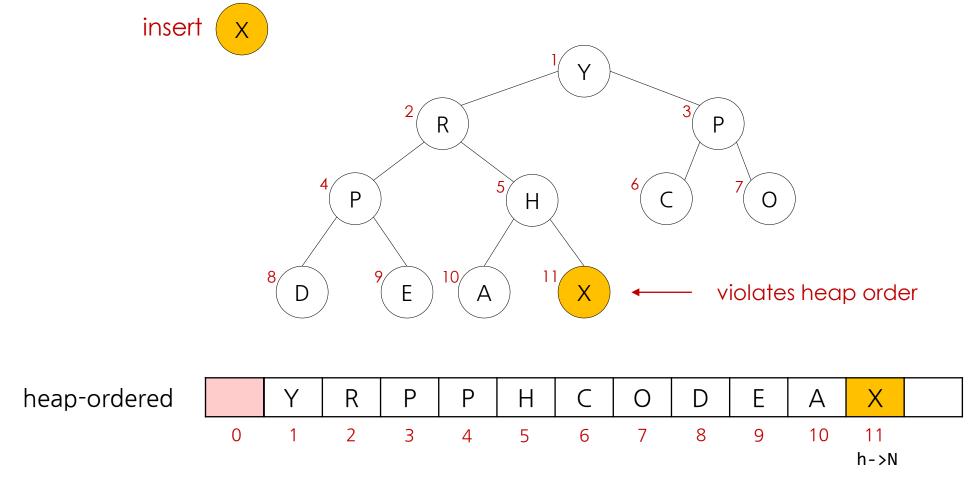
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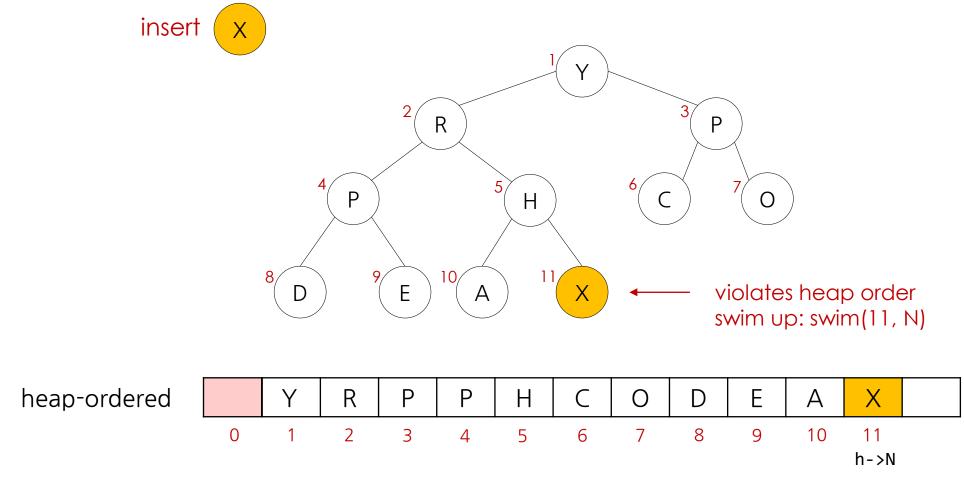
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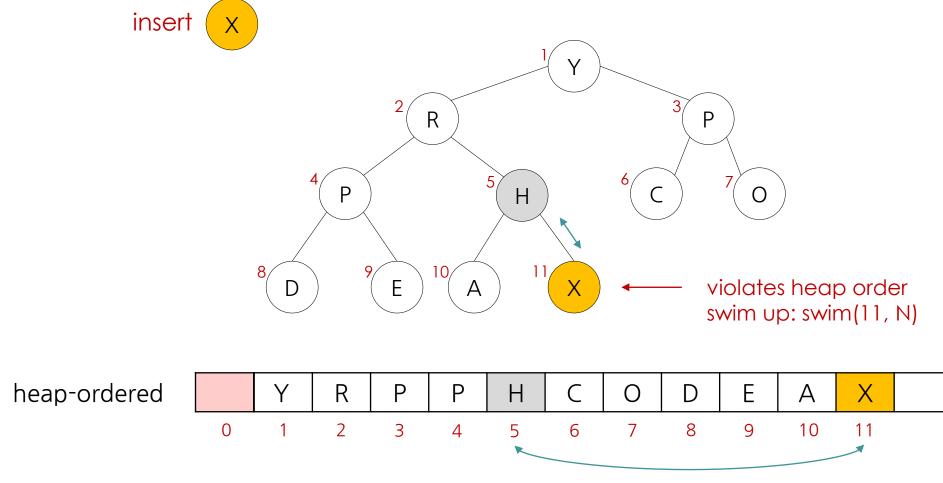
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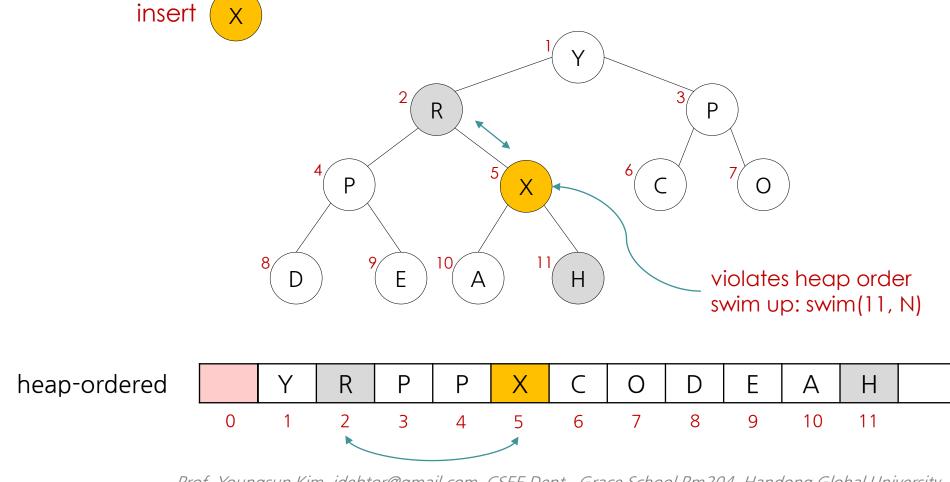
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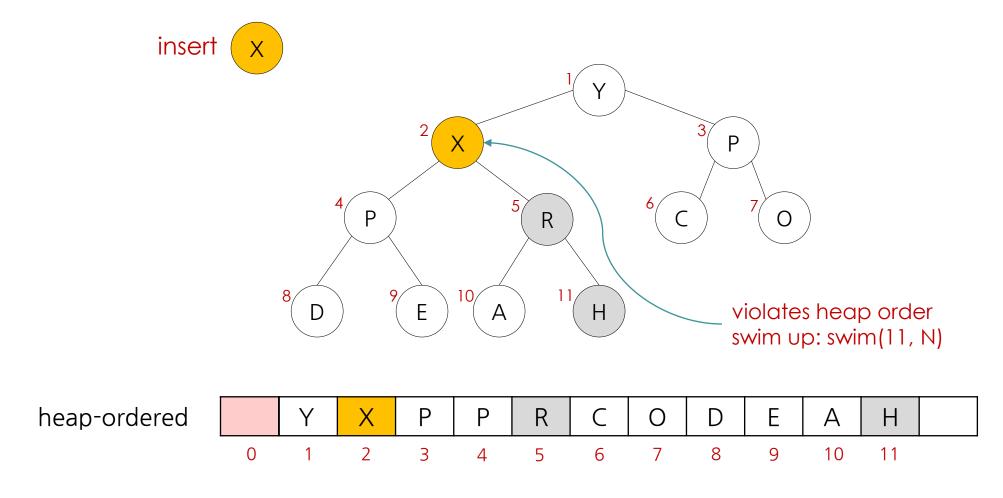
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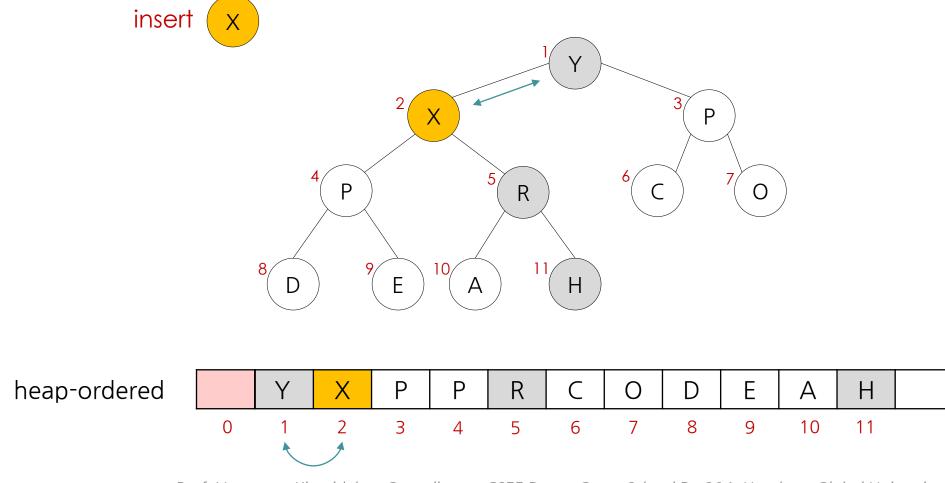
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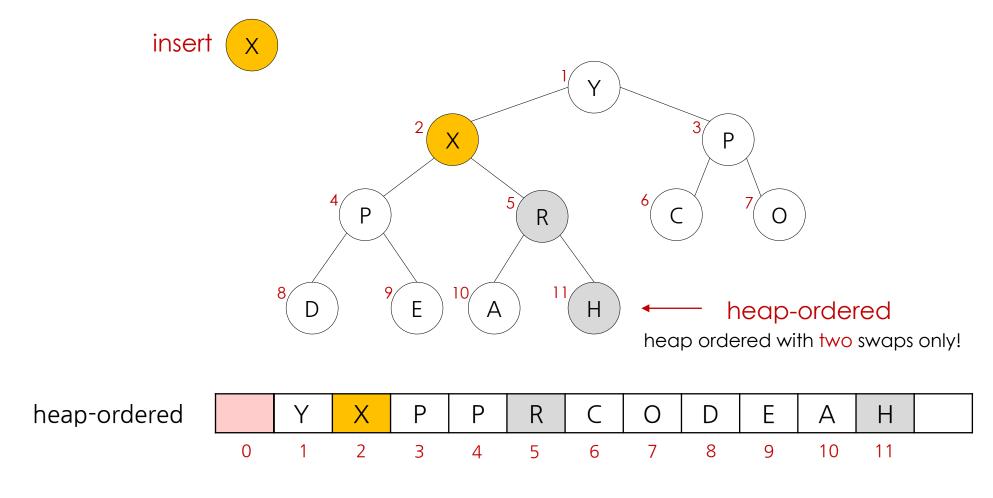
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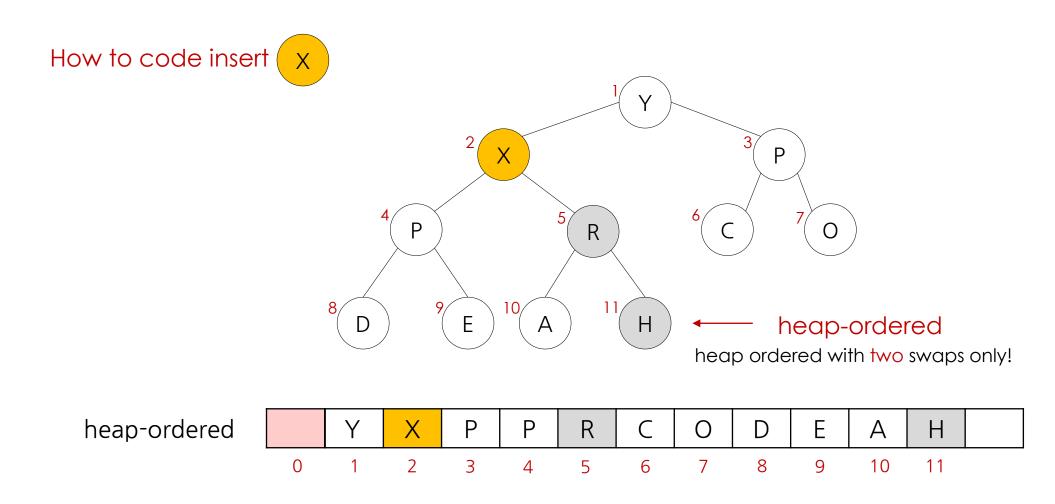
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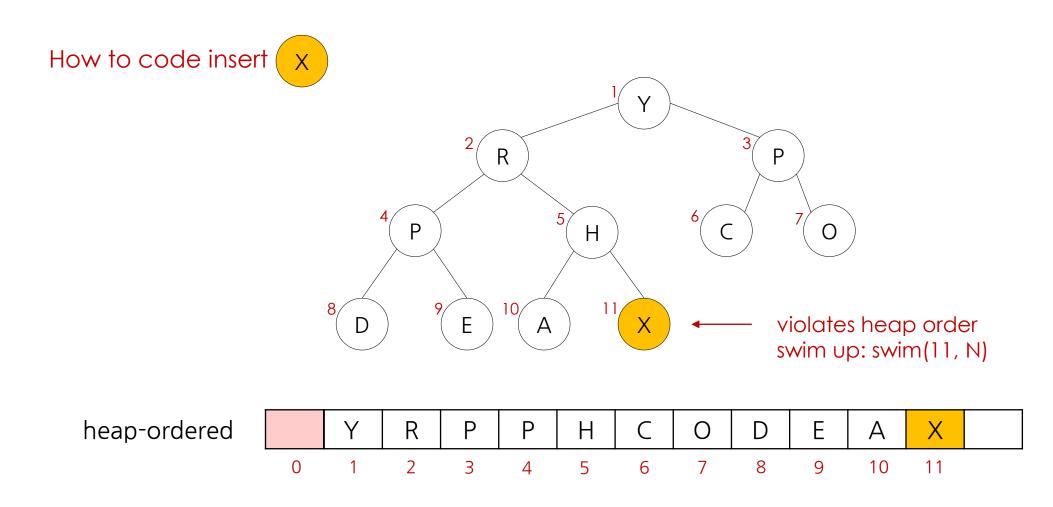
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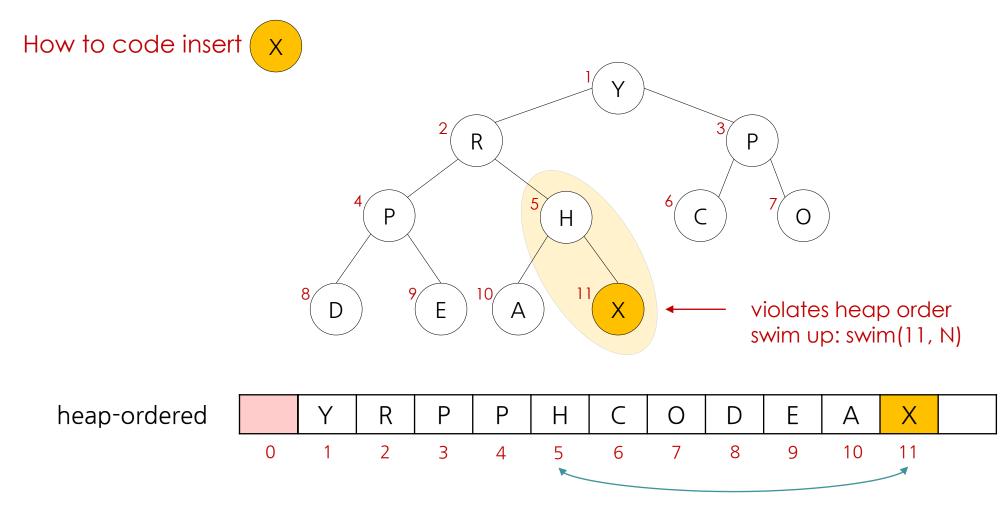
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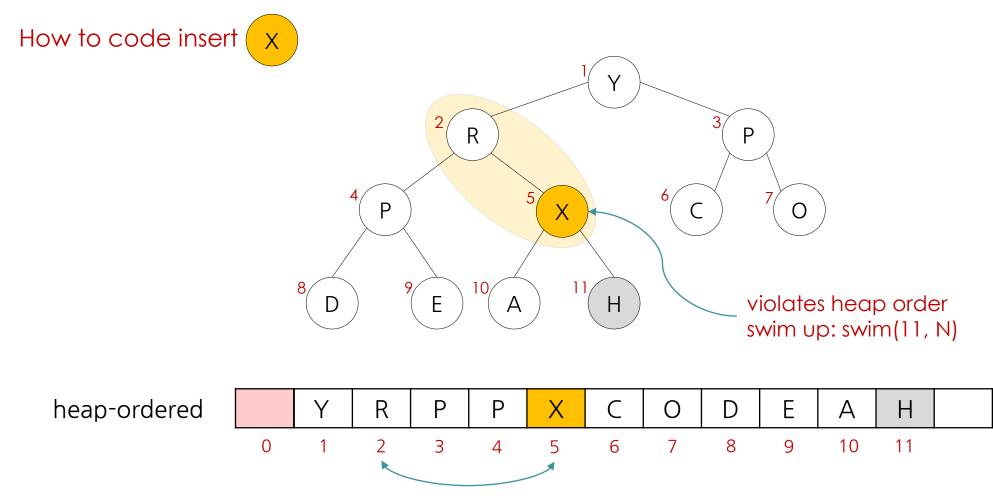
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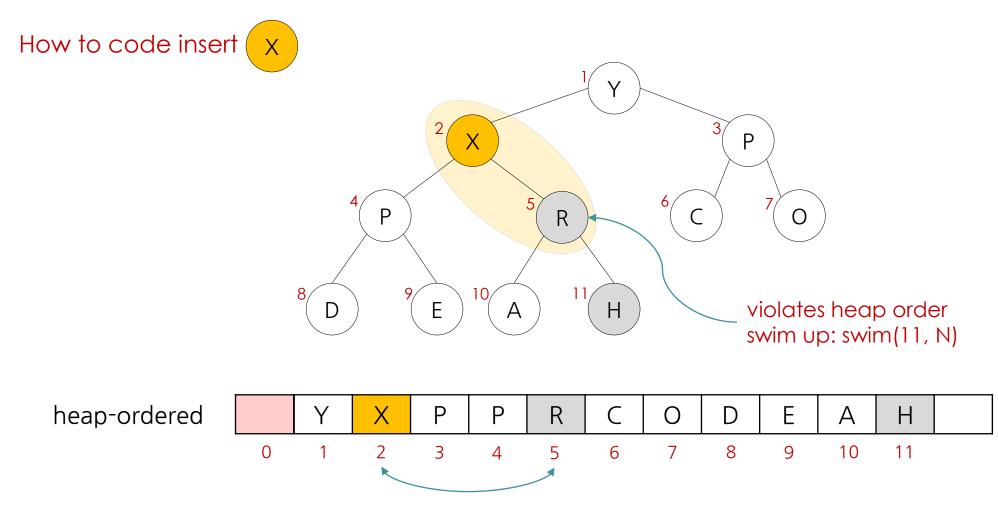
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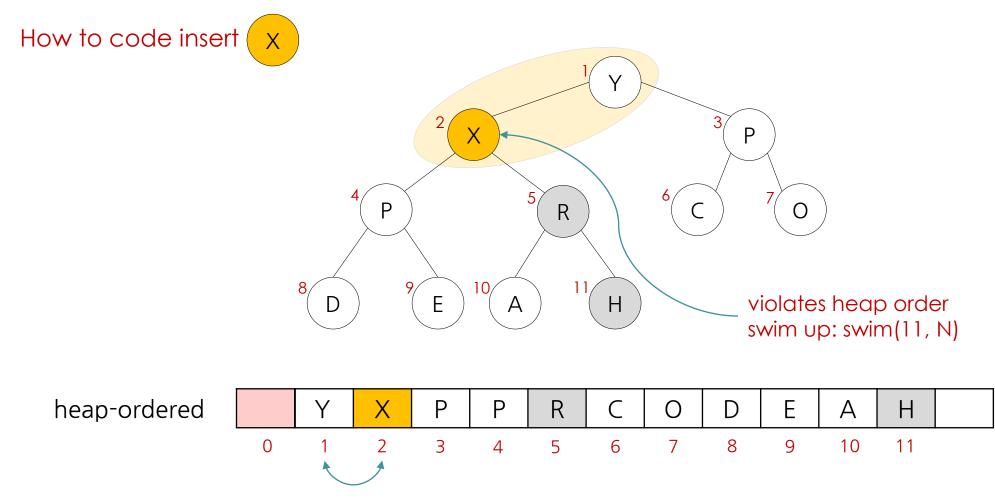
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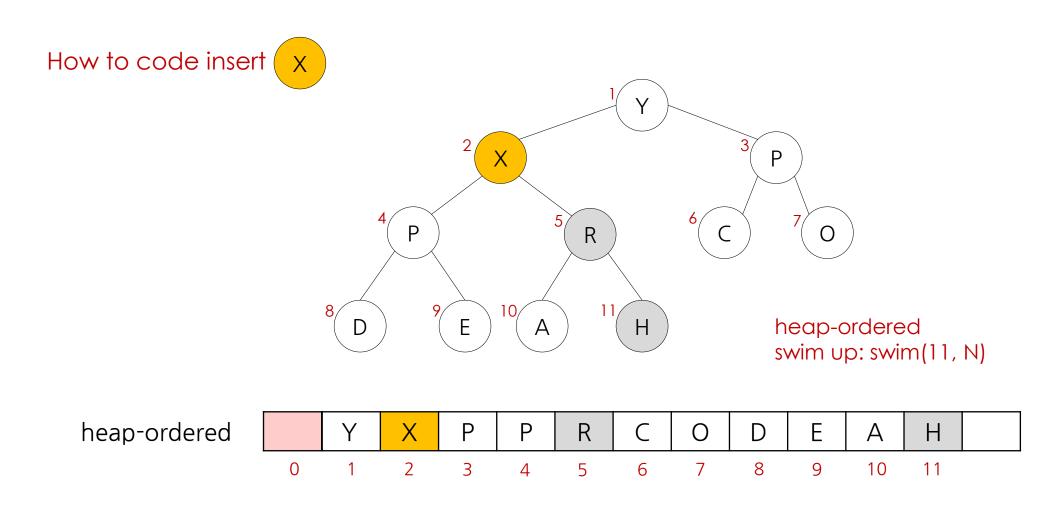
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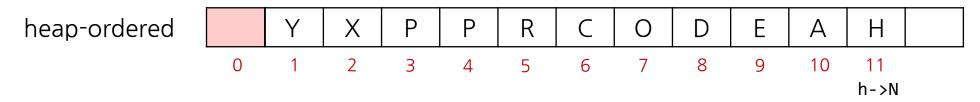


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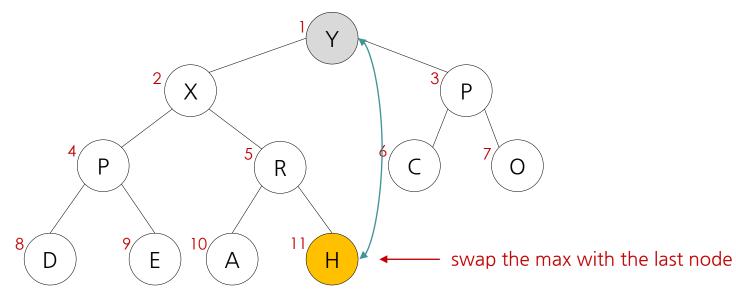
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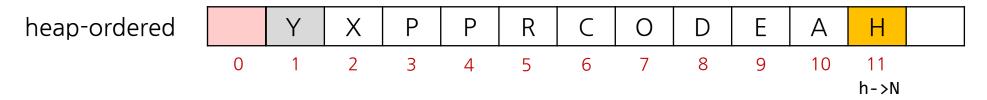
remove the max (root) - pop in priority queue



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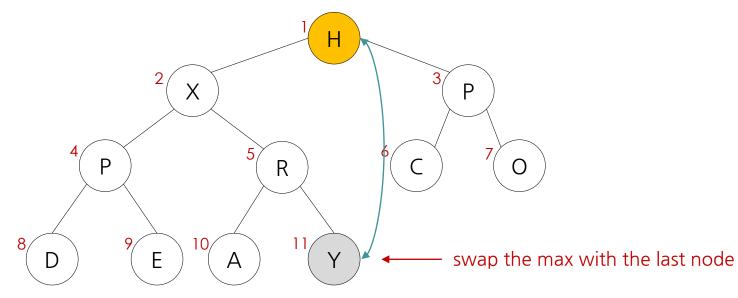
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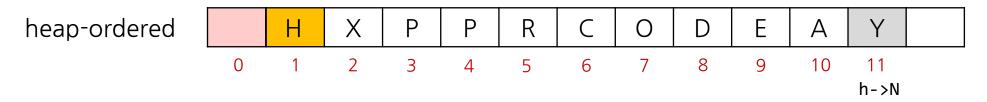




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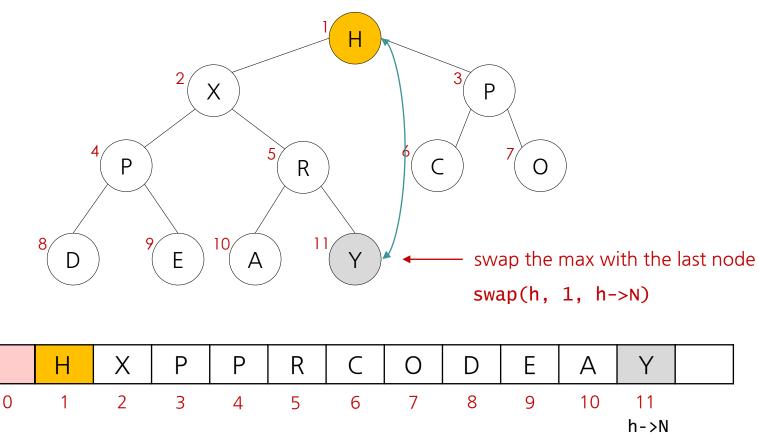




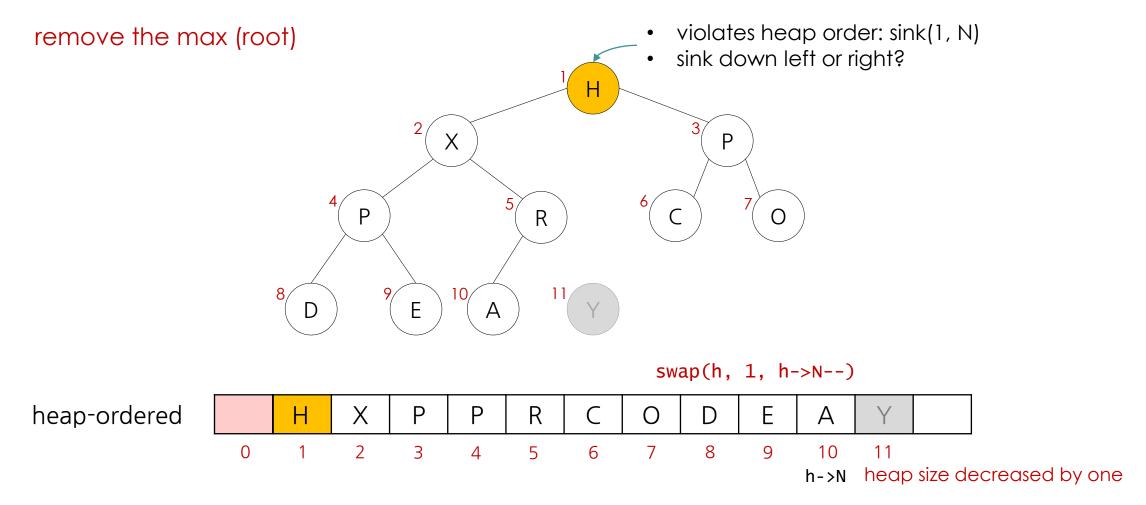
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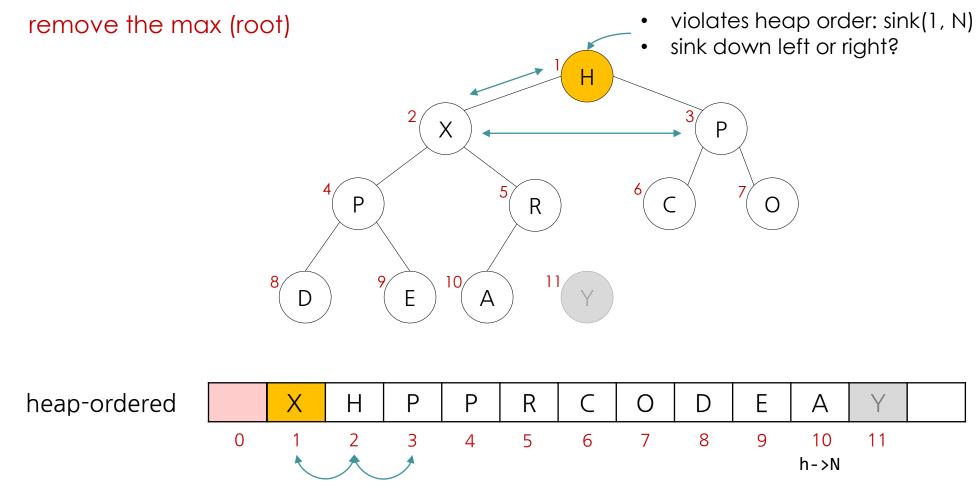
heap-ordered



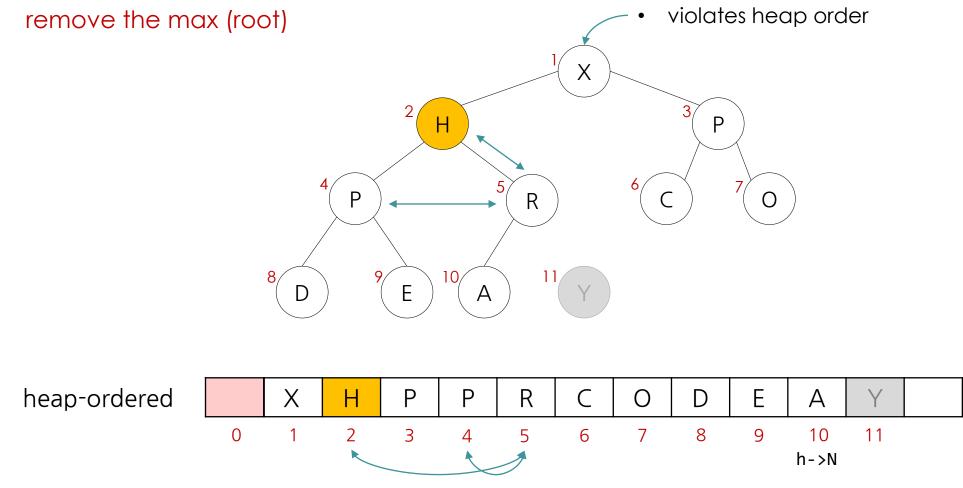
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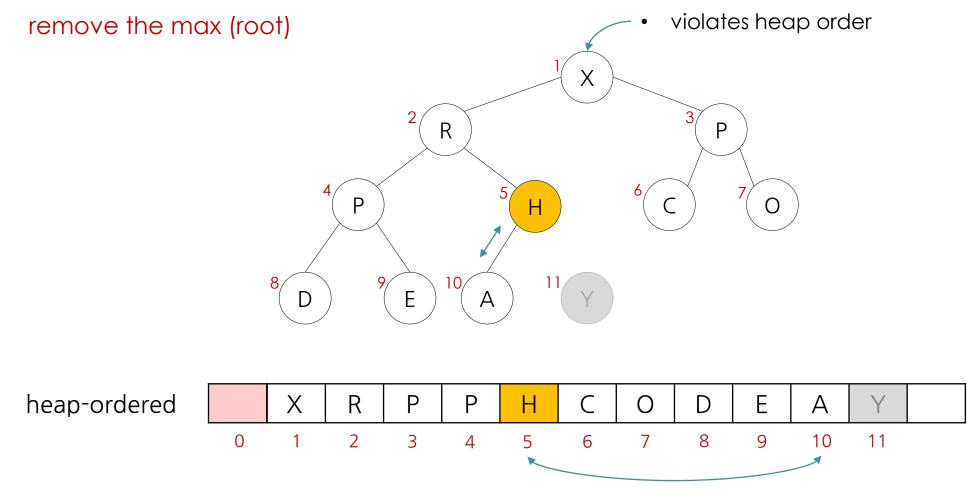
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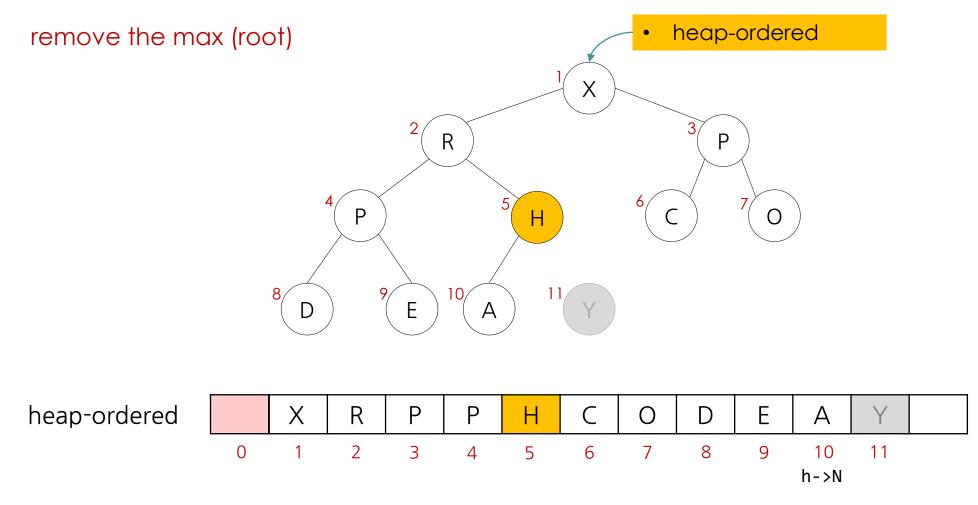
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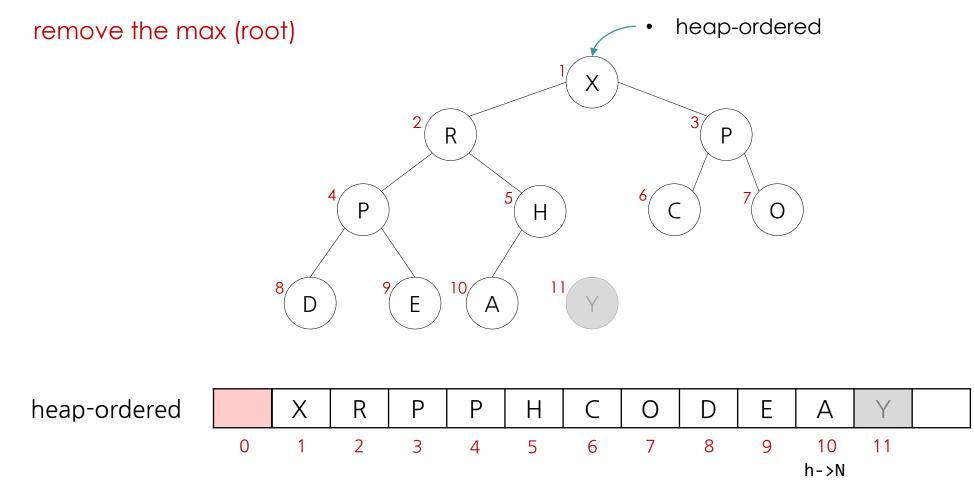
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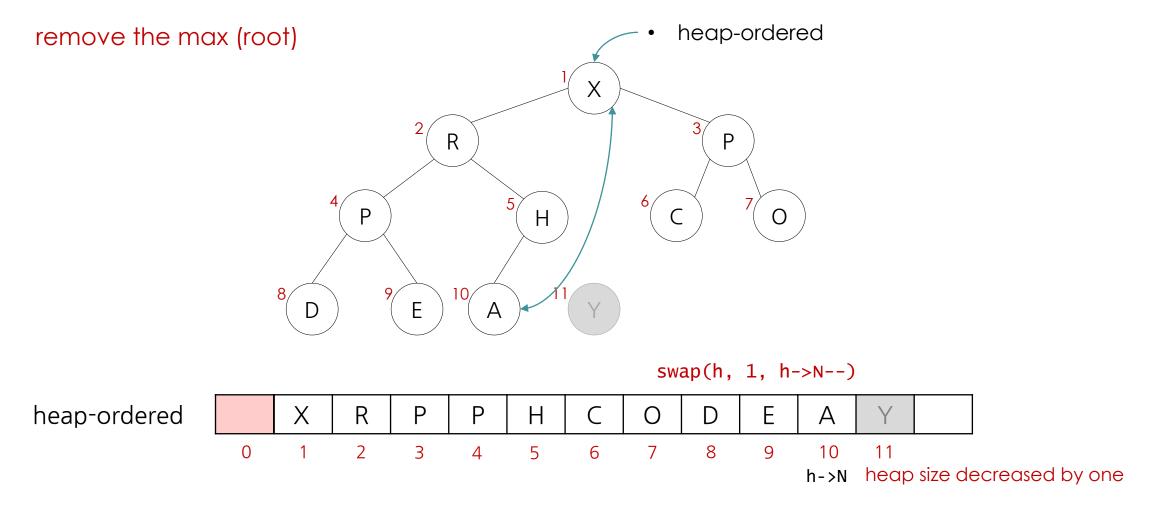
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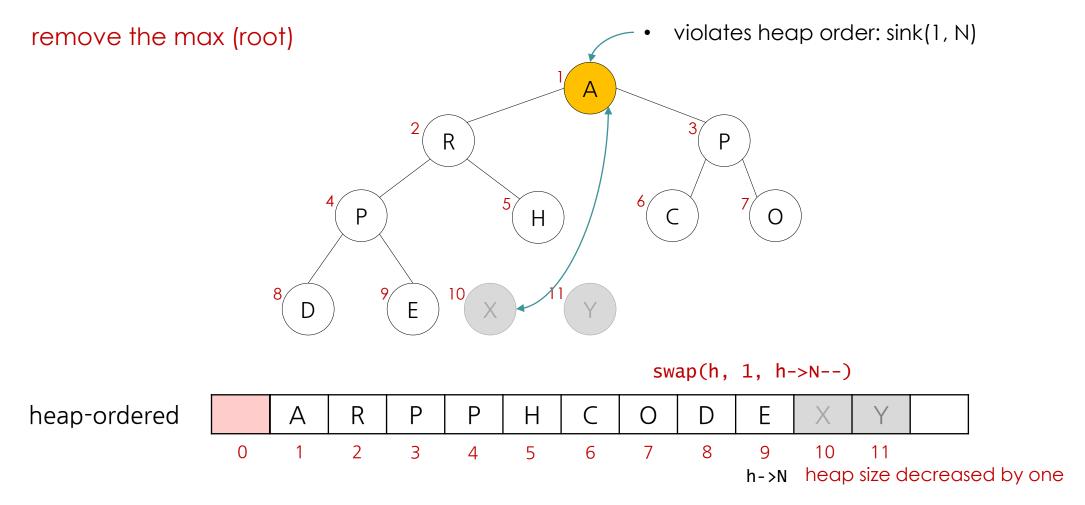
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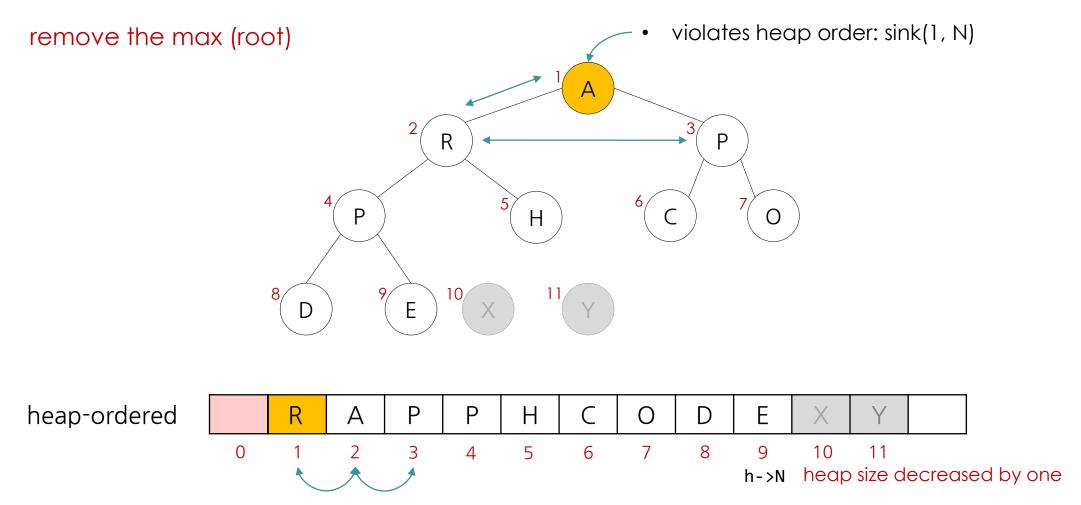
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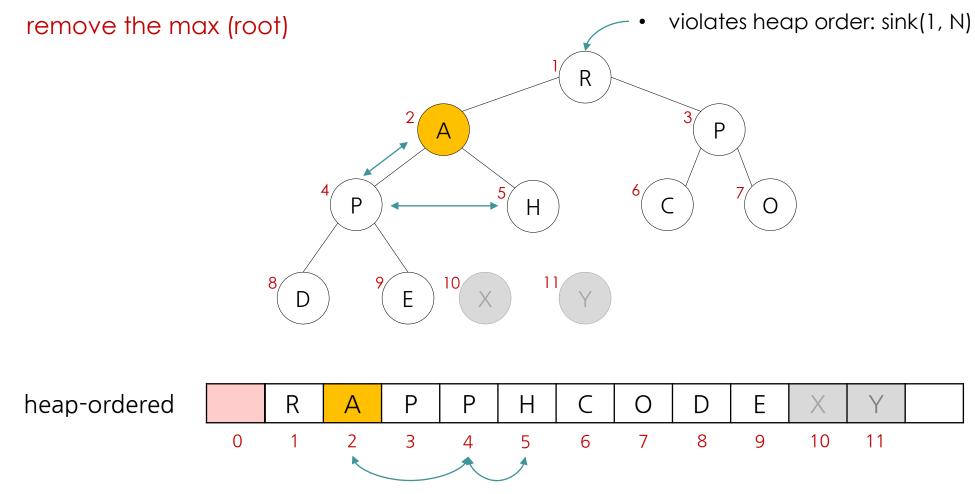
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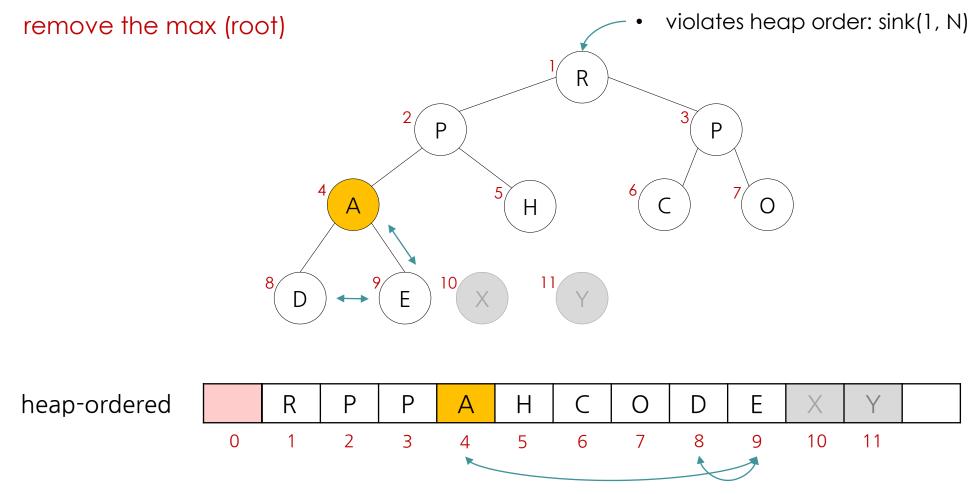
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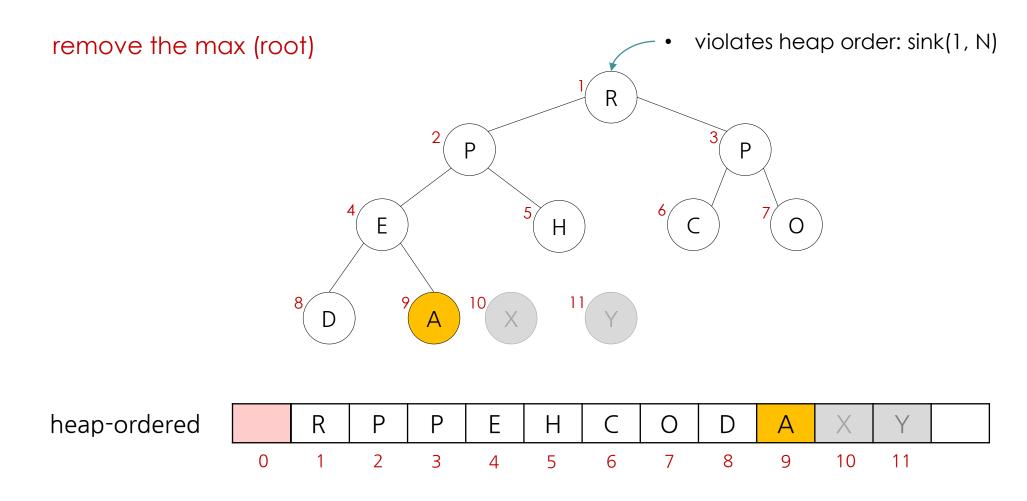
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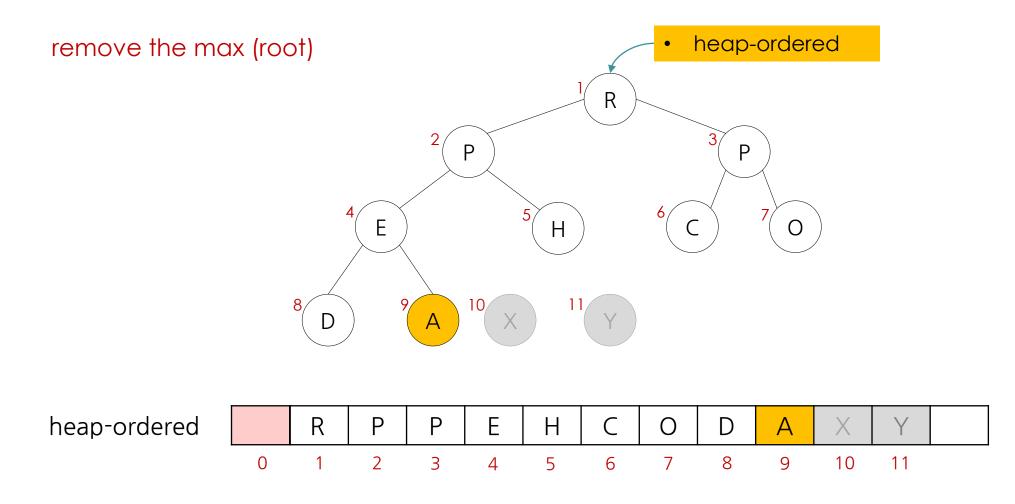
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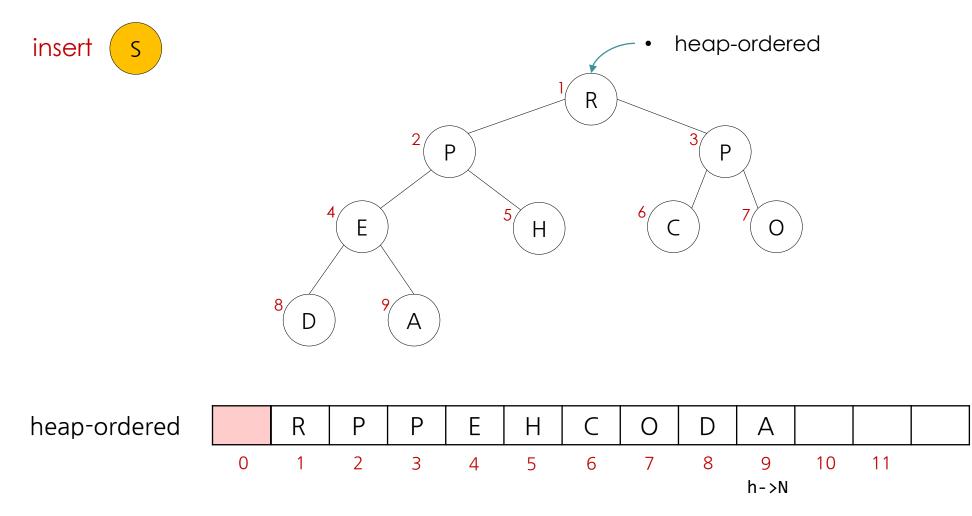
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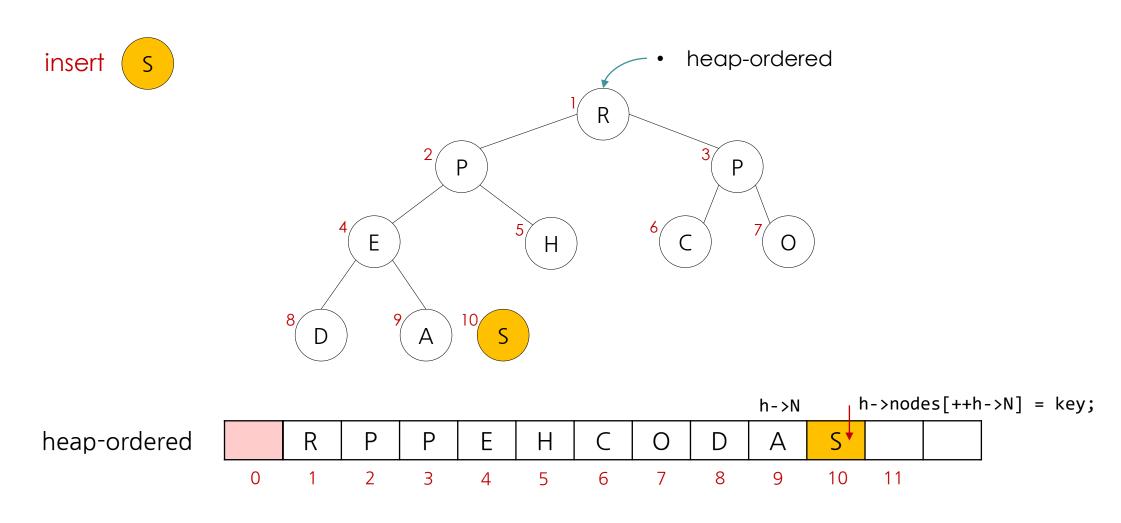
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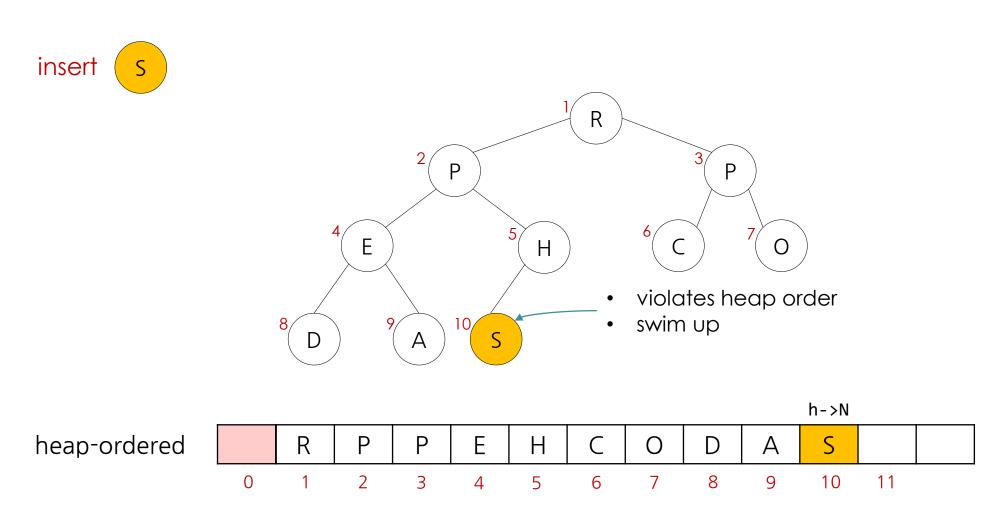
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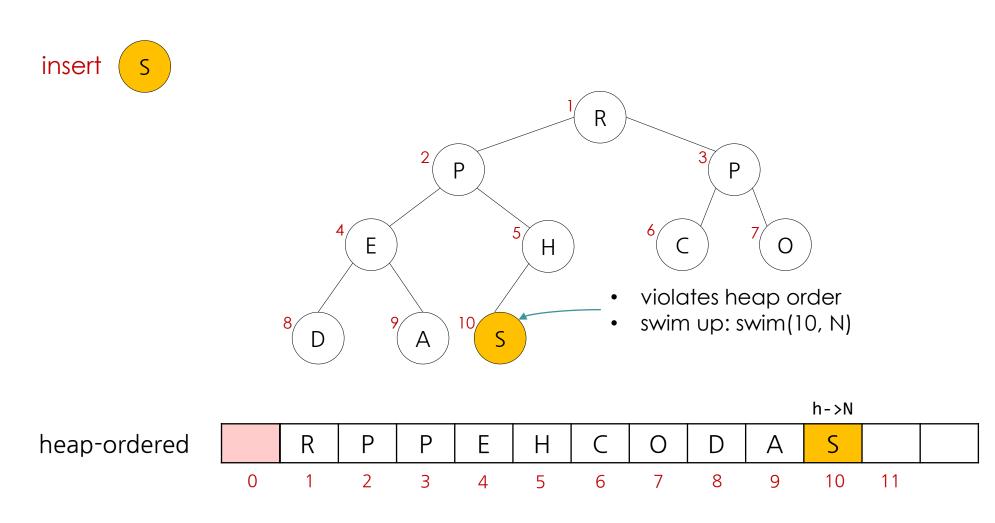
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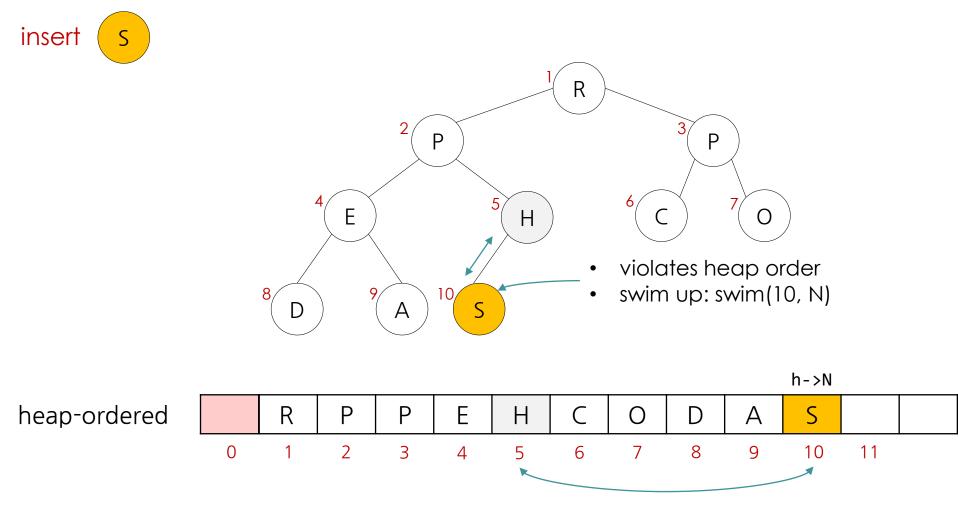
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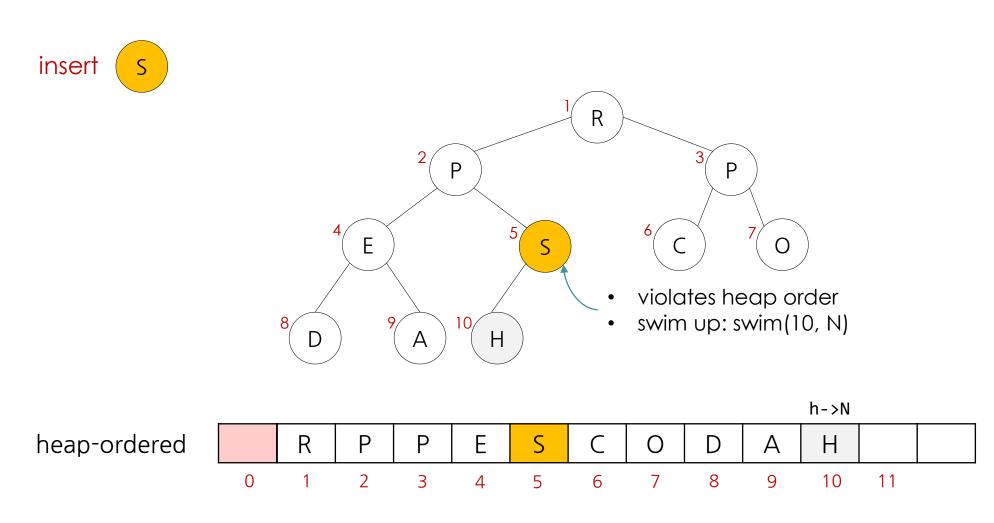
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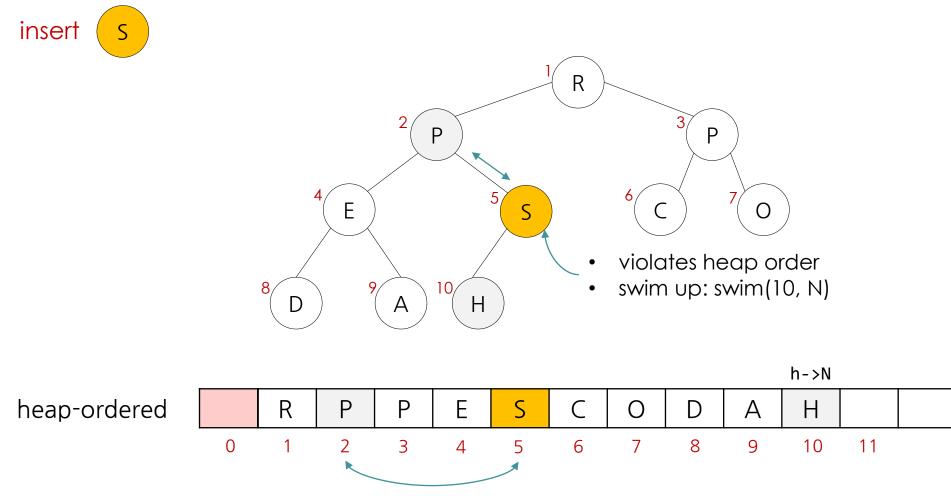
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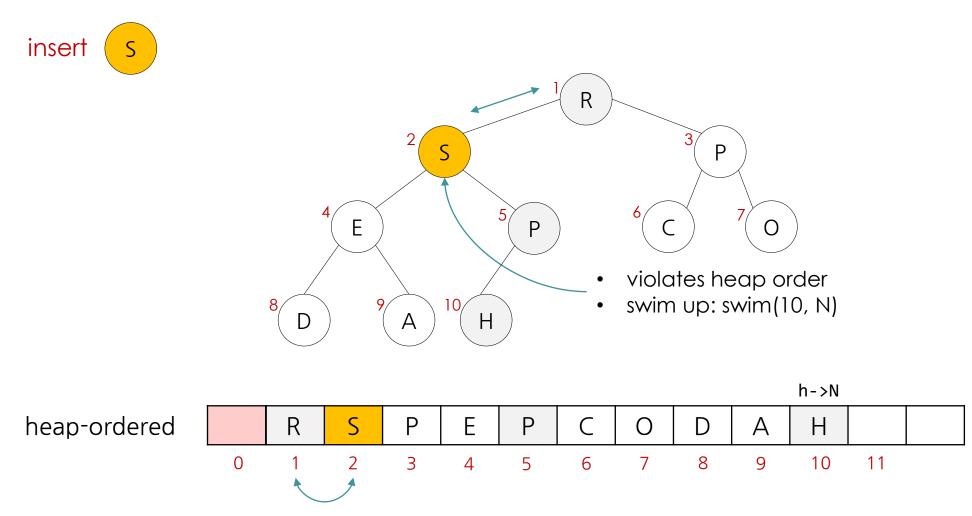
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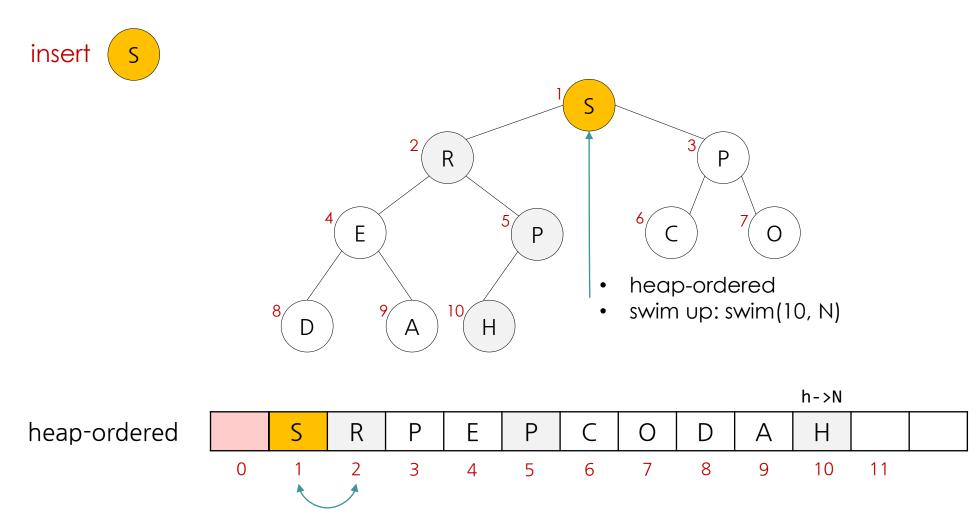
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Binary heap operations time complexity with N items:

- Level of heap is $\lfloor \log_2 N \rfloor$
- insert: O(log N) for each insert
 - In practice, expect less
- delete: O(log N) // deleting root node or any node
- increase/decrease key: O(log N)
- Heapify(): O(N)
- Heapsort(): O(N log N)
- Because O(N) heapify + O(log N) delete = O(N log N)

Proof:

- https://stackoverflow.com/guestions/9755721/how-can-building-a-heap-be-on-time-complexity
- https://www.insertingwiththeweb.com/data-structures/binary-heap/build-heap-proof/
- https://www.guora.com/How-is-the-time-complexity-of-building-a-heap-is-o-n

References in Korean:

- https://ratsgo.github.io/data%20structure&algorithm/2017/09/27/heapsort/
- https://zeddios_tistory.com/56

Data Structures in Python

- Heap and Priority Queue
- Heap Coding
- Heap Sort & Min/MaxHeap