CS 228 : Logic in Computer Science

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Recap: Languages, Machines and Logic

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A language $L \subseteq \Sigma^*$ is called FO-definable iff there exists an FO formula φ such that $L = L(\varphi)$.

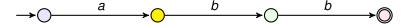
What we plan to show: L is FO-definable $\Rightarrow L$ is regular. Note that the converse is not true.

- $\Sigma = \{a, b\}$. Consider the following languages $L \subseteq \Sigma^*$:
 - ▶ Begins with a, ends with b, and has a pair of consecutive a's
 - Contains a b and ends with aa
 - Contains abb
 - ▶ There are two occurrences of b between which only a's occur
 - ▶ Right before the last position is an a
 - Even length words

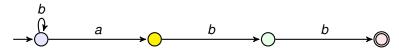
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► Contains abb

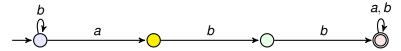
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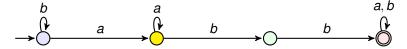
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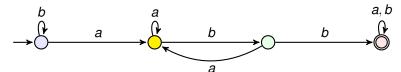


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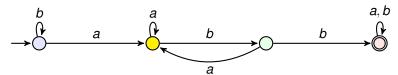
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$$\exists x \exists y \exists z (Q_a(x) \land Q_b(y) \land Q_b(z) \land S(x,y) \land S(y,z))$$

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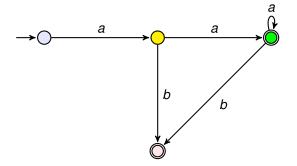
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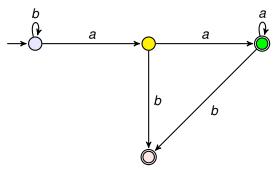
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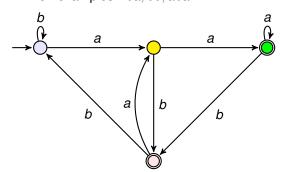
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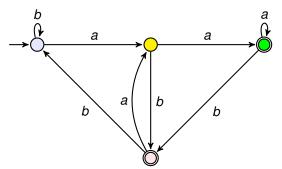
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$$\exists x [Q_a(x) \land \exists y (S(x,y) \land \forall z (z \leqslant y))]$$