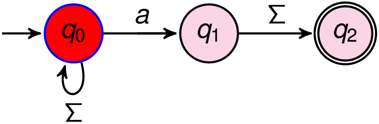
CS 228 : Logic in Computer Science

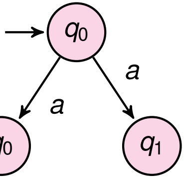
S. Krishna

Nondeterministic Finite Automata(NFA)

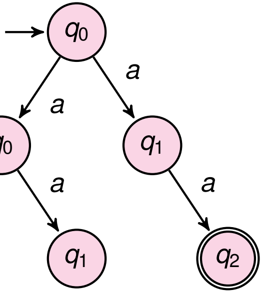


- $N = (Q, \Sigma, \delta, Q_0, F)$
  - Q is a finite set of states
  - $Q_0 \subseteq Q$ is the set of initial states
  - $\delta : Q \times \Sigma \to 2^Q$ is the transition function
  - $F \subseteq Q$ is the set of final states
- Acceptance condition : A word w is accepted iff it has atleast one accepting path

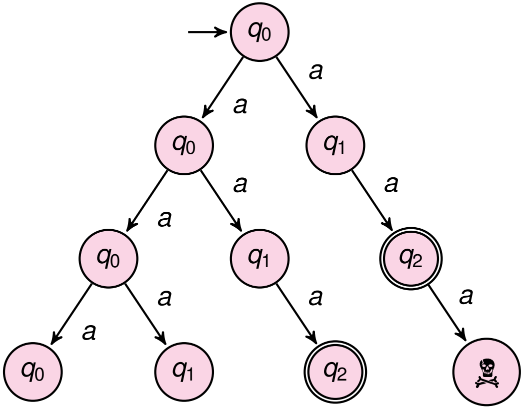
Run Tree of *aaab*



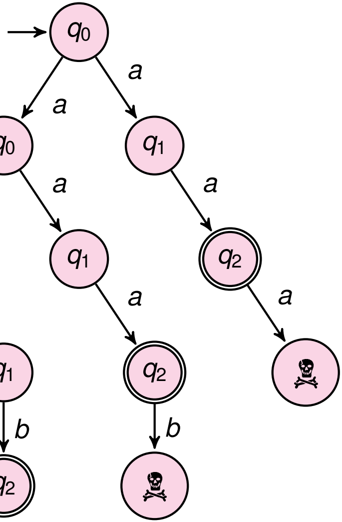
Run Tree of *aaab*



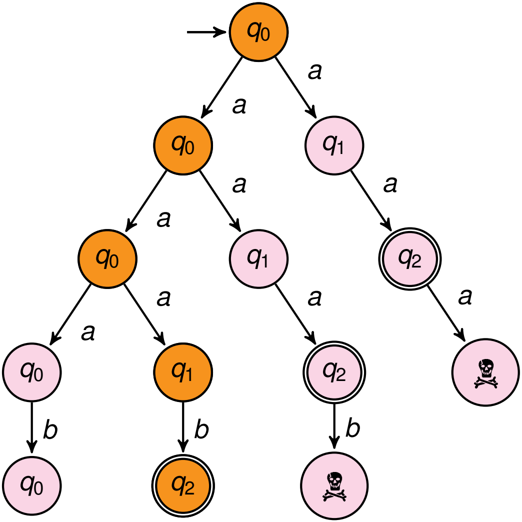
Run Tree of *aaab*



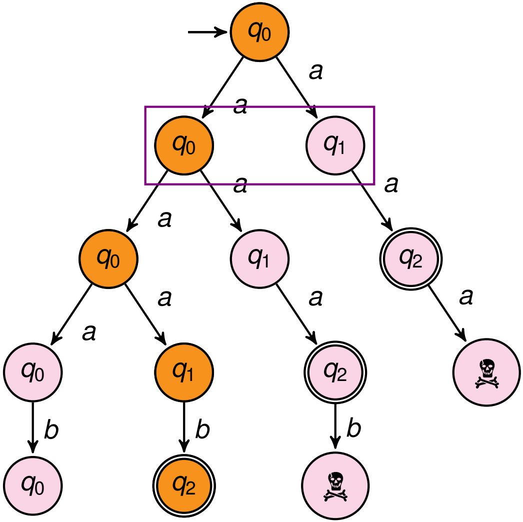
Run Tree of *aaab*



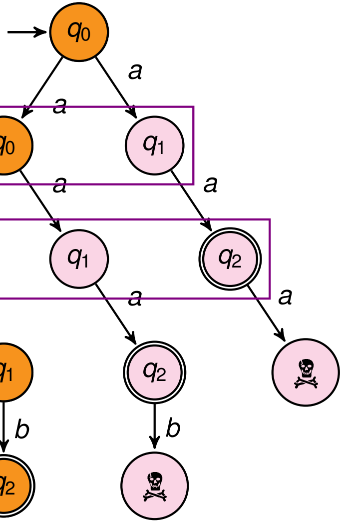
Run Tree of *aaab*



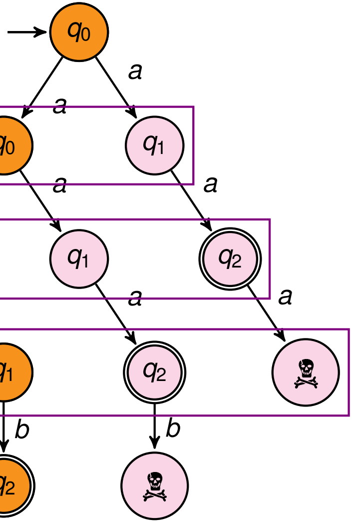
Run Tree of *aaab*



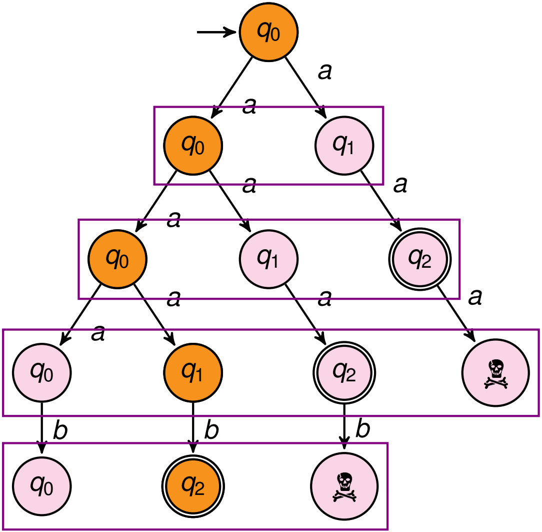
Run Tree of *aaab*



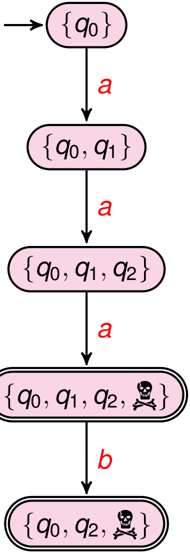
Run Tree of *aaab*



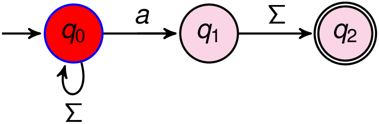
Run Tree of *aaab*

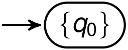


The Single Run

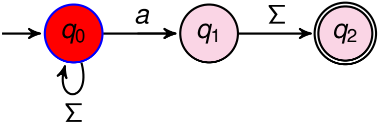


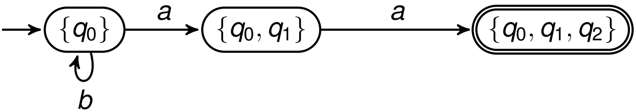
An Example



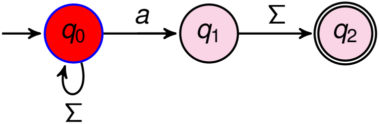


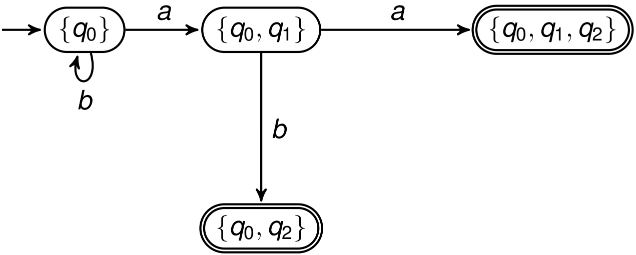
An Example



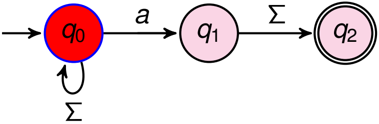


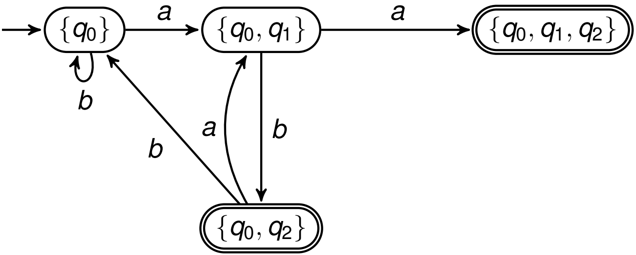
An Example



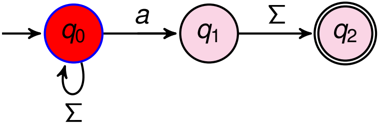


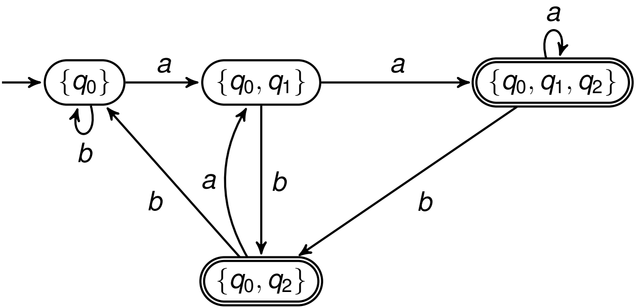
An Example





An Example





NFA and DFA

- Any DFA is also an NFA

NFA and DFA

- Any DFA is also an NFA
- Any NFA can be converted into a language equivalent DFA

NFA and DFA

- Any DFA is also an NFA
- Any NFA can be converted into a language equivalent DFA
  - Combine all the runs of w in the NFA into a single run in the DFA

NFA and DFA

- Any DFA is also an NFA
- Any NFA can be converted into a language equivalent DFA
  - Combine all the runs of w in the NFA into a single run in the DFA
  - Combine states occurring in various runs to obtain a set of states

NFA and DFA

- Any DFA is also an NFA
- Any NFA can be converted into a language equivalent DFA
  - Combine all the runs of w in the NFA into a single run in the DFA
  - Combine states occurring in various runs to obtain a set of states
  - A set of states evolves into another set of states

NFA and DFA

- Any DFA is also an NFA
- Any NFA can be converted into a language equivalent DFA
  - Combine all the runs of w in the NFA into a single run in the DFA
  - Combine states occurring in various runs to obtain a set of states
  - A set of states evolves into another set of states
  - Use $\delta : Q \times \Sigma \rightarrow 2^Q$, obtain $\Delta : 2^Q \times \Sigma \rightarrow 2^Q$

NFA and DFA

- Any DFA is also an NFA
- Any NFA can be converted into a language equivalent DFA
  - Combine all the runs of w in the NFA into a single run in the DFA
  - Combine states occurring in various runs to obtain a set of states
  - A set of states evolves into another set of states
  - Use $\delta : Q \times \Sigma \to 2^Q$, obtain $\Delta : 2^Q \times \Sigma \to 2^Q$
  - Δ is an extension of δ

NFA and DFA

- Any DFA is also an NFA
- Any NFA can be converted into a language equivalent DFA
  - Combine all the runs of w in the NFA into a single run in the DFA
  - Combine states occurring in various runs to obtain a set of states
  - A set of states evolves into another set of states
  - Use $\delta : Q \times \Sigma \to 2^Q$, obtain $\Delta : 2^Q \times \Sigma \to 2^Q$
  - Δ is an extension of δ
  - Accept if the obtained set of states contains a final state

NFA and DFA

Given NFA $N = (Q, \Sigma, Q_0, \delta, F)$, obtain the DFA $D = (2^Q, \Sigma, Q_0, \Delta, F')$

NFA and DFA

Given NFA $N = (Q, \Sigma, Q_0, \delta, F)$, obtain the DFA $D = (2^Q, \Sigma, Q_0, \Delta, F')$

- $\Delta : 2^Q \times \Sigma \to 2^Q$ is defined by $\Delta(A, a) = \bigcup_{q \in A} \delta(q, a)$

NFA and DFA

Given NFA $N = (Q, \Sigma, Q_0, \delta, F)$, obtain the DFA $D = (2^Q, \Sigma, Q_0, \Delta, F')$

- $\Delta : 2^Q \times \Sigma \to 2^Q$ is defined by $\Delta(A, a) = \bigcup_{q \in A} \delta(q, a)$
- $F' = \{S \in 2^Q \mid S \cap F \neq \emptyset\}$

NFA and DFA

Given NFA $N = (Q, \Sigma, Q_0, \delta, F)$, obtain the DFA $D = (2^Q, \Sigma, Q_0, \Delta, F')$

- $\Delta : 2^Q \times \Sigma \to 2^Q$ is defined by $\Delta(A, a) = \bigcup_{q \in A} \delta(q, a)$
- $F' = \{S \in 2^Q \mid S \cap F \neq \emptyset\}$

Note that $\hat{\delta}(A, a) = \bigcup_{q \in A} \delta(q, a) = \Delta(A, a)$

NFA and DFA

Given NFA $N = (Q, \Sigma, Q_0, \delta, F)$, obtain the DFA $D = (2^Q, \Sigma, Q_0, \Delta, F')$

- $\Delta : 2^Q \times \Sigma \to 2^Q$ is defined by $\Delta(A, a) = \bigcup_{q \in A} \delta(q, a)$
- $F' = \{S \in 2^Q \mid S \cap F \neq \emptyset\}$

Note that $\hat{\delta}(A, a) = \bigcup_{q \in A} \delta(q, a) = \Delta(A, a)$
Show that

- $\hat{\Delta} : 2^Q \times \Sigma^* \to 2^Q$ is same as $\hat{\delta} : 2^Q \times \Sigma^* \to 2^Q$ (recall $\delta : Q \times \Sigma \to 2^Q$)

NFA and DFA

Given NFA $N = (Q, \Sigma, Q_0, \delta, F)$, obtain the DFA $D = (2^Q, \Sigma, Q_0, \Delta, F')$

- $\Delta : 2^Q \times \Sigma \to 2^Q$ is defined by $\Delta(A, a) = \bigcup_{q \in A} \delta(q, a)$
- $F' = \{S \in 2^Q \mid S \cap F \neq \emptyset\}$

Note that $\hat{\delta}(A, a) = \bigcup_{q \in A} \delta(q, a) = \Delta(A, a)$
Show that

- $\hat{\Delta} : 2^Q \times \Sigma^* \to 2^Q$ is same as $\hat{\delta} : 2^Q \times \Sigma^* \to 2^Q$ (recall $\delta : Q \times \Sigma \to 2^Q$)
- $\hat{\Delta}(A, xa) = \Delta(\hat{\Delta}(A, x), a) = \bigcup_{q \in \hat{\Delta}(A,x)} \delta(q, a)$

NFA and DFA

Given NFA $N = (Q, \Sigma, Q_0, \delta, F)$, obtain the DFA $D = (2^Q, \Sigma, Q_0, \Delta, F')$

- $\Delta : 2^Q \times \Sigma \rightarrow 2^Q$ is defined by $\Delta(A, a) = \bigcup_{q \in A} \delta(q, a)$
- $F' = \{S \in 2^Q \mid S \cap F \neq \emptyset\}$

Note that $\hat{\delta}(A, a) = \bigcup_{q \in A} \delta(q, a) = \Delta(A, a)$
Show that

- $\hat{\Delta} : 2^Q \times \Sigma^* \rightarrow 2^Q$ is same as $\hat{\delta} : 2^Q \times \Sigma^* \rightarrow 2^Q$ (recall $\delta : Q \times \Sigma \rightarrow 2^Q$)
- $\hat{\Delta}(A, xa) = \Delta(\hat{\Delta}(A, x), a) = \bigcup_{q \in \hat{\Delta}(A,x)} \delta(q, a)$
- $\hat{\delta}(A, xa) = \bigcup_{q \in \hat{\delta}(A,x)} \delta(q, a)$

NFA = DFA

$$x \in L(D) \leftrightarrow \hat{\Delta}(Q_0, x) \in F'$$

$$\leftrightarrow$$

$$\hat{\delta}(Q_0, x) \in F'$$

$$\leftrightarrow$$

$$\hat{\delta}(Q_0, x) \cap F \neq \emptyset$$

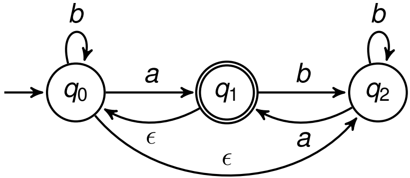
$$\leftrightarrow$$

$$x \in L(N)$$

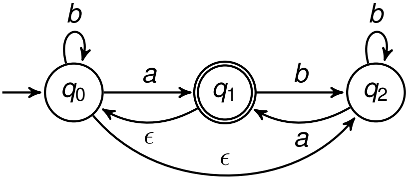
Regularity

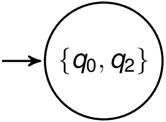
A language L is regular iff there exists an NFA A such that $L = L(A)$

ϵ-NFA

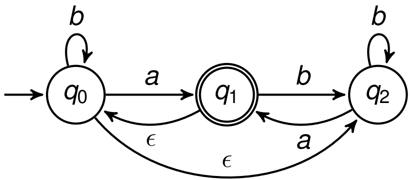


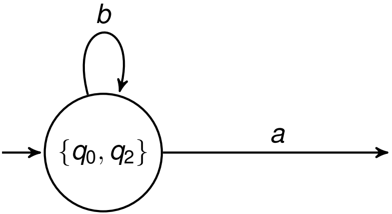
ϵ-NFA



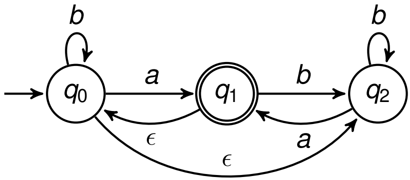


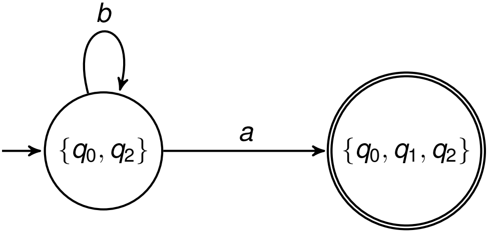
ϵ-NFA



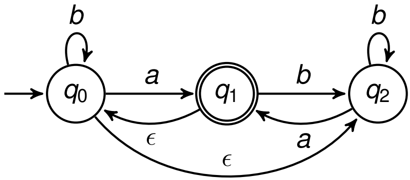


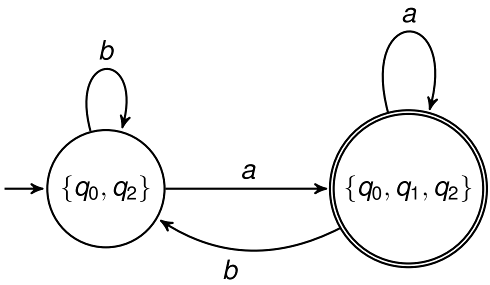
ϵ-NFA





ϵ-NFA



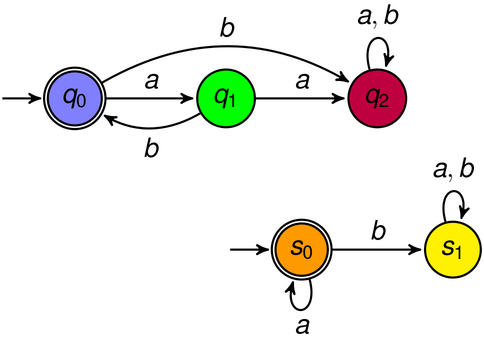


ϵ-NFA and DFA

- ϵ-close the initial states of the ϵ-NFA to obtain initial state of DFA
- From a state S, compute $\Delta(S, a)$ and ϵ-close it
- All states in the DFA are ϵ-closed
- Final states are those which contain a final state of the ϵ-NFA

Closure under Concatenation

- Given regular languages L_1, L_2, is $L_1.L_2$ regular



Closure under Concatenation

- Given regular languages L_1, L_2, is $L_1.L_2$ regular?