



CS 228 : Logic in Computer Science

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GNBA

- ▶ Generalized NBA, a variant of NBA
- ▶ Only difference is in acceptance condition
- ▶ Acceptance condition in GNBA is a set $\mathcal{F} = \{F_1, \dots, F_k\}$, each $F_i \subseteq Q$
- ▶ An infinite run ρ is accepting in a GNBA iff

$$\forall F_i \in \mathcal{F}, \text{Inf}(\rho) \cap F_i \neq \emptyset$$

- ▶ Note that when $\mathcal{F} = \emptyset$, all infinite runs are accepting
- ▶ GNBA and NBA are equivalent in expressive power.

LTL to GNBA

- ▶ Let $\varphi = a \text{ Ub}$.
- ▶ Subformulae of $\varphi : \{a, b, a \text{ Ub}\}$. Let $B = \{a, \neg a, b, \neg b, a \text{ Ub}, \neg(a \text{ Ub})\}$.
- ▶ Possibilities at each state : maximally **consistent** subsets of B
 - ▶ $\{a, \neg b, a \text{ Ub}\}$
 - ▶ $\{\neg a, b, a \text{ Ub}\}$
 - ▶ $\{a, b, a \text{ Ub}\}$
 - ▶ $\{a, \neg b, \neg(a \text{ Ub})\}$
 - ▶ $\{\neg a, \neg b, \neg(a \text{ Ub})\}$
- ▶ Our initial state(s) must guarantee truth of $a \text{ Ub}$. Thus, initial states: $\{a, b, a \text{ Ub}\}$ and $\{\neg a, b, a \text{ Ub}\}$ and $\{a, \neg b, a \text{ Ub}\}$.

LTL to GNBA

→ $\{a, b, a \cup b\}$

$\{a, \neg b, \neg(a \cup b)\}$

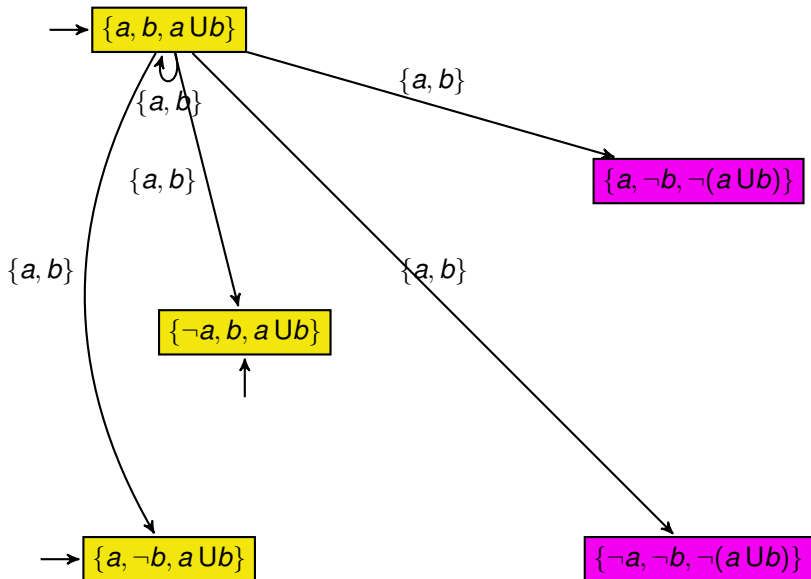
$\{\neg a, b, a \cup b\}$



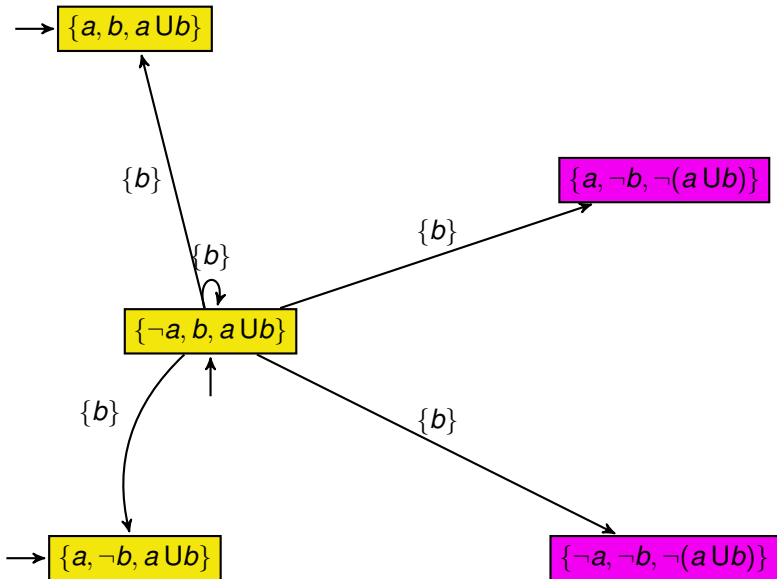
→ $\{a, \neg b, a \cup b\}$

$\{\neg a, \neg b, \neg(a \cup b)\}$

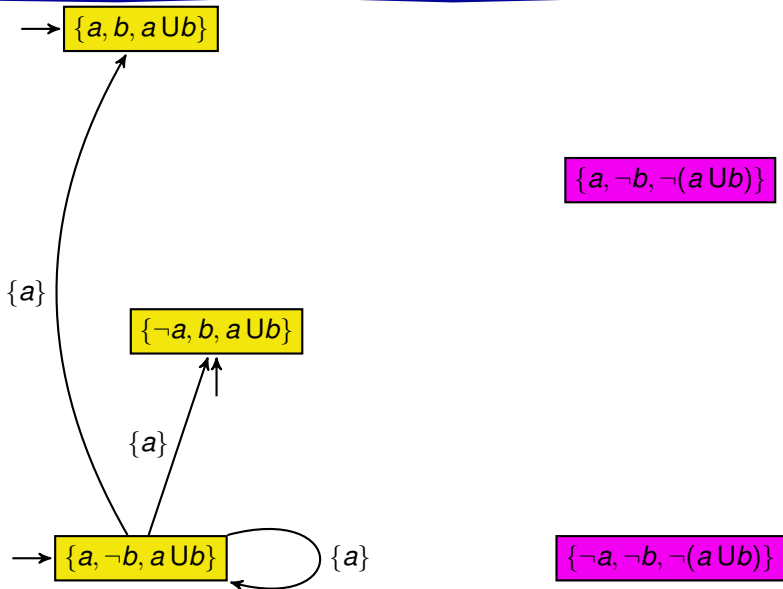
LTL to GNBA



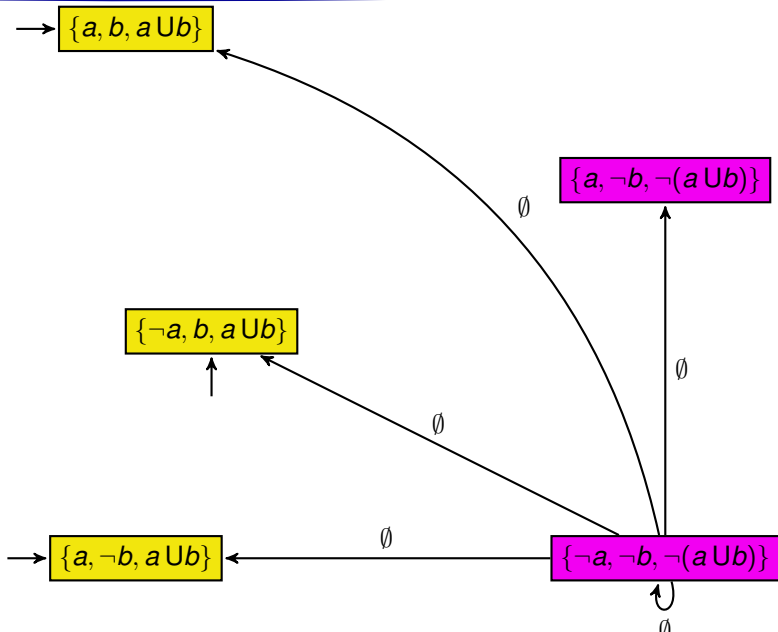
LTL to GNBA



LTL to GNBA



LTL to GNBA

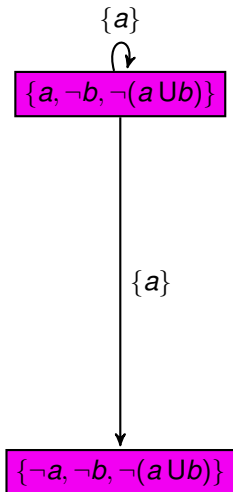


LTL to GNBA

→ $\{a, b, a \cup b\}$

$\{\neg a, b, a \cup b\}$
↑

→ $\{a, \neg b, a \cup b\}$



LTL to GNBA : Accepting States

→ $\{a, b, a \cup b\}$

$\{a, \neg b, \neg(a \cup b)\}$

$\{\neg a, b, a \cup b\}$



→ $\{a, \neg b, a \cup b\}$

$\{\neg a, \neg b, \neg(a \cup b)\}$

LTL to GNBA

Construct GNBA for $\neg(a \text{ U } b)$.

LTL to GNBA

- ▶ Let $\varphi = a \cup (\neg a \cup c)$. Let $\psi = \neg a \cup c$
- ▶ Subformulae of φ : $\{a, \neg a, c, \psi, \varphi\}$. Let $B = \{a, \neg a, c, \neg c, \psi, \neg \psi, \varphi, \neg \varphi\}$.
- ▶ Possibilities at each state : some **consistent** subset of B holds
 - ▶ $\{a, c, \psi, \varphi\}$
 - ▶ $\{\neg a, c, \psi, \varphi\}$
 - ▶ $\{a, \neg c, \neg \psi, \varphi\}$
 - ▶ $\{a, \neg c, \neg \psi, \neg \varphi\}$
 - ▶ $\{\neg a, \neg c, \psi, \varphi\}$
 - ▶ $\{\neg a, \neg c, \neg \psi, \neg \varphi\}$

LTL to GNBA

→ $\{a, c, \psi, \varphi\}$

$\{\neg a, \neg c, \psi, \varphi\}$ ←

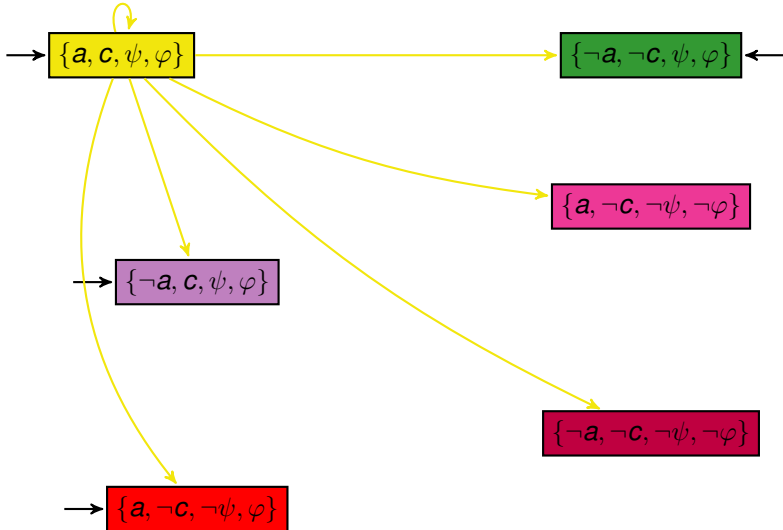
→ $\{\neg a, c, \psi, \varphi\}$

$\{a, \neg c, \neg \psi, \neg \varphi\}$

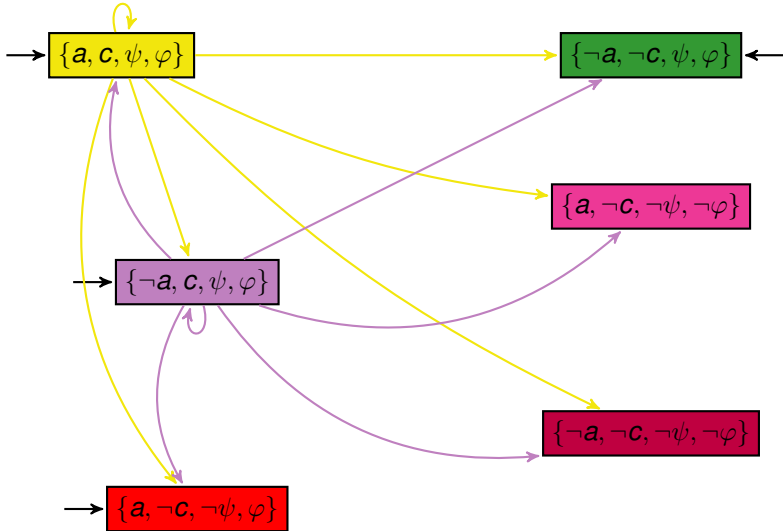
→ $\{a, \neg c, \neg \psi, \varphi\}$

$\{\neg a, \neg c, \neg \psi, \neg \varphi\}$

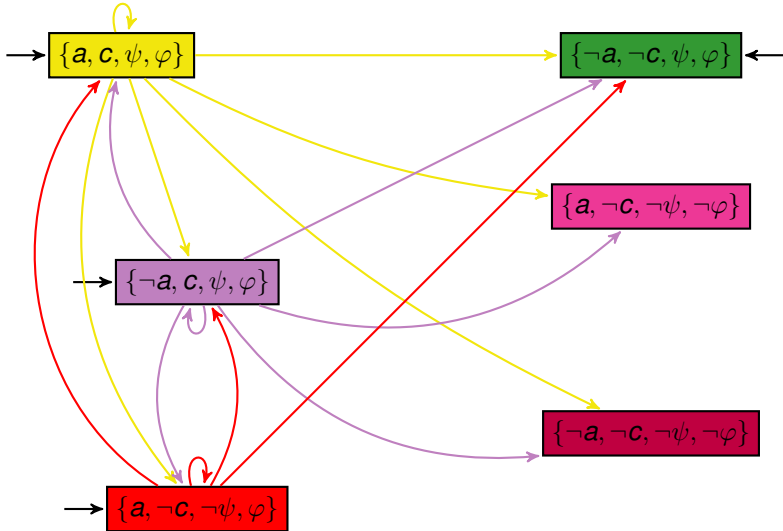
LTL to GNBA



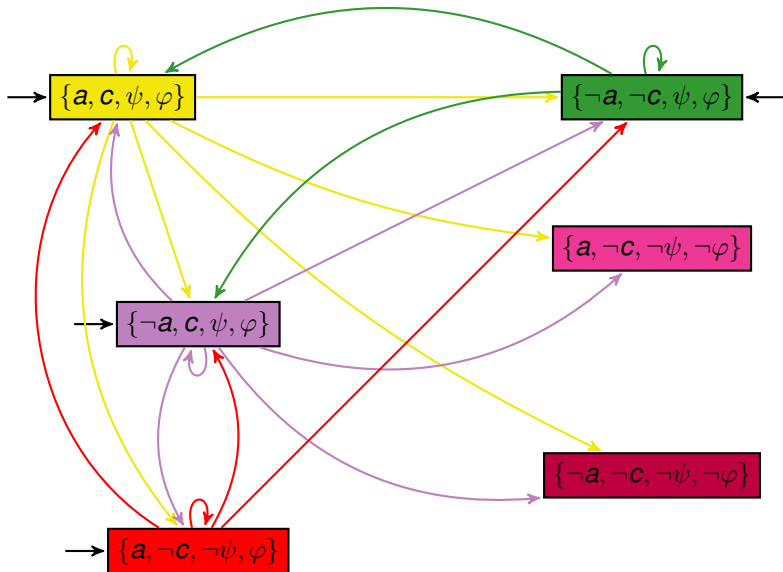
LTL to GNBA



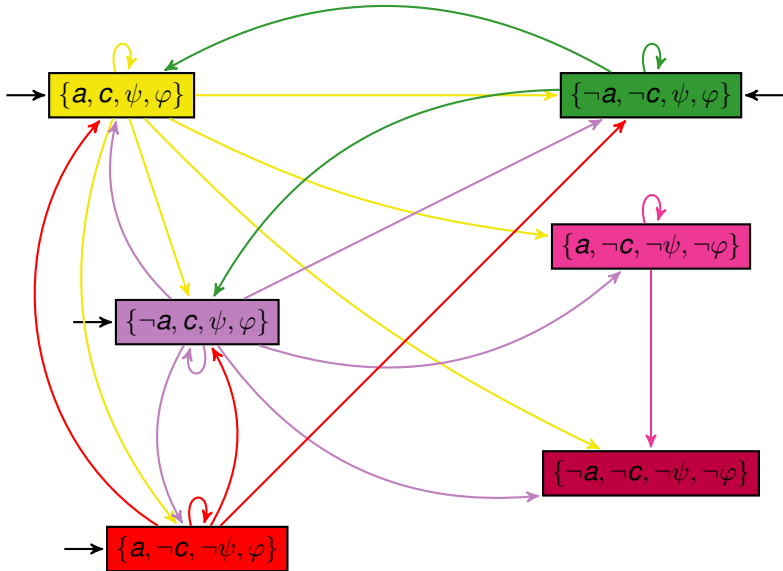
LTL to GNBA



LTL to GNBA



LTL to GNBA



LTL to GNBA

