

Topic: Connectivity    ① Paths    ② Cycles \*    ③ Sorting on graph.

① Paths.    Q. graph  $\rightarrow v_1$  and  $v_2 \Rightarrow$  if there is a path bet. these 2 nodes or not.  
if yes, return me path.  
\* undirected    \* directed    logic: generic  
easy    tricky

bool getPath (g, v<sub>1</sub>, v<sub>2</sub>, visited, path):

path.push(v<sub>1</sub>)

visited[v<sub>1</sub>] = true

\* if v<sub>1</sub> == v<sub>2</sub> : return true

for each neighbour v in g[v<sub>1</sub>]:

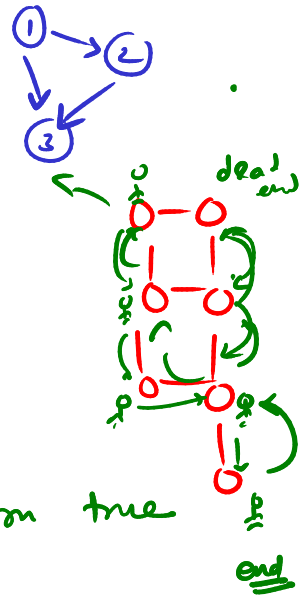
if !visited[v]:

\* if getPath (g, v, v<sub>2</sub>, visited, path) : return true

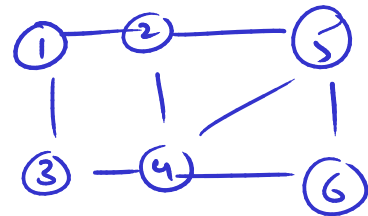
path.pop()

return false

this line.



Why am I getting only 1 path?



1 2 5 4  
1 2 4  
1 3 4  
1 2 5 6 4

void printAllPaths (g, s, d):

visited[n] = {false}

all\_paths = []

path = []

dfsVisit (g, s, d, path, all\_paths, visited)

dfsVisit (g, s, d, path, all\_paths, visited) :

path.push(s) ; visited[s] = true ;

if s == d : all\_paths.push(path) return

for all v in g[s] : if !visited[v] : dfsVisit (g, v, d, path, all\_paths, visited)

path.pop() ; visited[s] = false ;

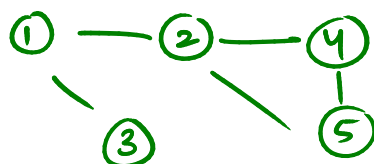
## ② Cycles.

a) undirected

undirected grapho:

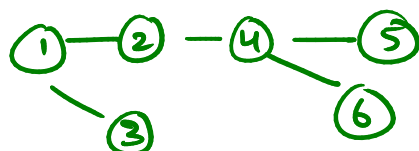
$n$  nodes

$n-1$  edges.



cycle.? Yes

no cycles : tree-like.

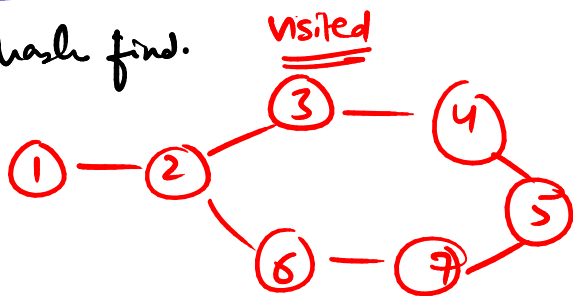


# components = 1

# edges ==  $n-1$

recursion stack. set find.

$O(1)$  hash find.



rec stack

[ 1 2 ~~3~~ ~~4~~ ~~5~~ 6 7 ]

visited

[ 1 2 3 4 5 6 7 ]

Undirected : DFS

bool hasCycle(g):

visited[n] = {false}

for  $i=0$  to  $n-1$   $i++$

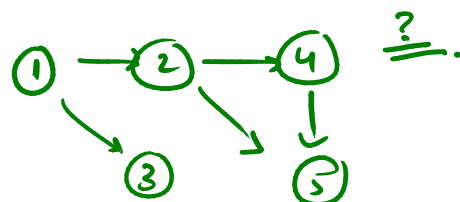
if !visited[i]:

if hasCycleVisit(g, i, visited)

return true

return false

b) directed



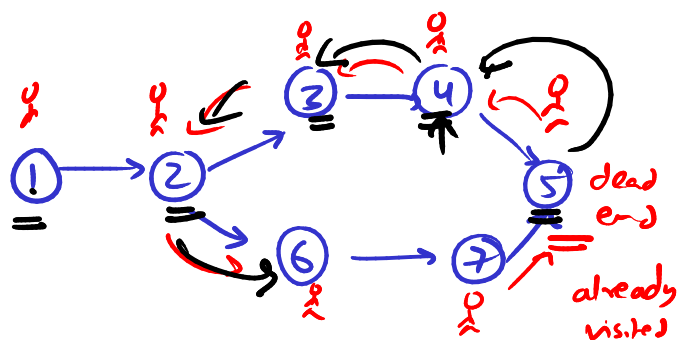
cycle? No

Debate: DFS if I come across a node visited, cycle is there.

Yes.

undirected

Q. directed? same? or not?



cycle  $\Rightarrow$  ?

No.

Yes. cycle

cycle: if  $v$  is in stack: cycle | Backedge

bool hasCycleVisit(g, u, visited):

visited[u] = true

for each  $v$  in  $g[u]$ :

if !visited[v]:

if hasCycleVisit(g, v, visited)

return true

else:

return true

return false

directed :

bool hasCycle(g):

visited[n] = {false}

rec\_stack = set()

for i=0 to n-1

if hasCycleUtil(g, i, visited, rec\_stack)

return true

return false.

bool hasCycleUtil(g, u, visited, rec\_stack)

if !visited[u]:

visited[u] = true rec\_stack.insert(u)

for each v in g[u]:

if !visited[v] and  
hasCycleUtil(g, v, vis, rec)

return true

else if v in rec\_stack:

return true

rec\_stack.remove(u)

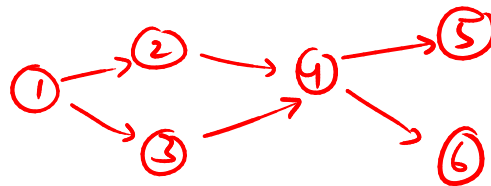
return false

③ sorting on graphs.

Real world : DAGs → Directed acyclic graphs.

Why?

① Chronology/Dependency.



package managers : JS/TS → npm, yarn

Python → pip

Java → maven

Rust → cargo

what is this?

How they operate?

npm i react

node\_modules/

pip i tensorflow

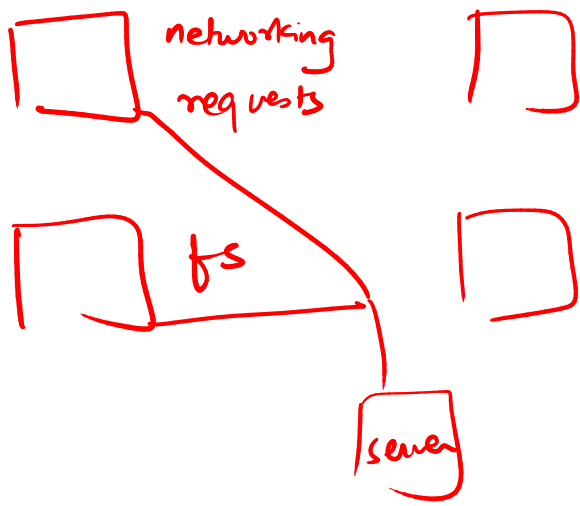
pip freeze > requirements.txt

Can I use my lang for dev? → No!!

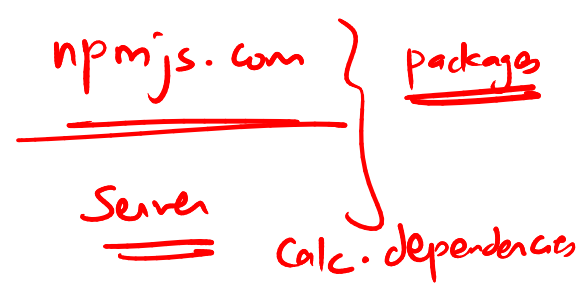
Community.

py

lang



registries

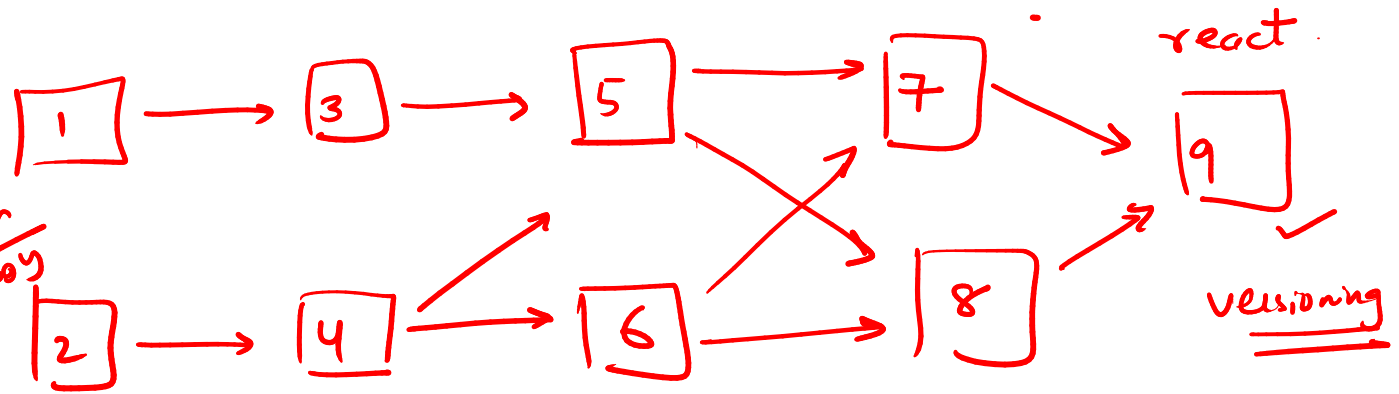


python requirements.txt

js package.json

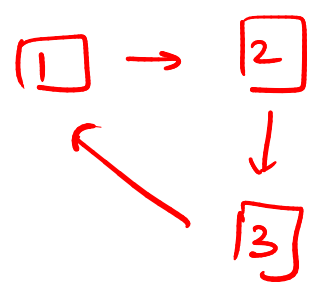
DAG

illustration  
chronology



npm i react.

DFS to construct dependency graph.



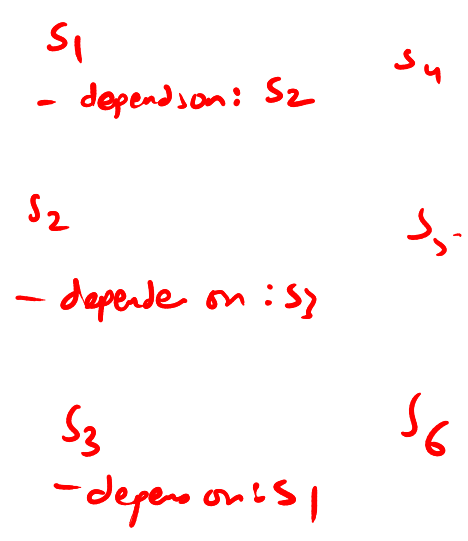
chicken-egg problem / cyclic dependency

Docker containers

docker-compose.yml

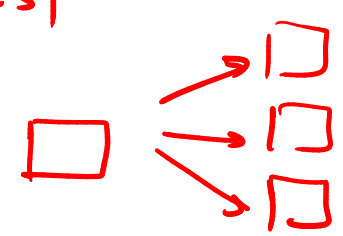
AWS / Cloud

master-slave



X cyclic dependency

n/w of computers



DAG

① Create a graph  $\rightarrow$  ensure no cycle. | ensure DAG.

② install the libraries (eg pip, npm)

What should be the order to install lib?

1 3 2 4 5 6 8 7 9      npm i pkg

sorted order in DAG  $\rightarrow$  topological sort. | no unique  
to po sort.

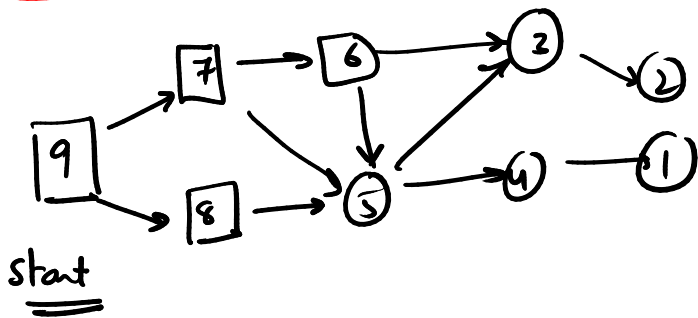
\* Cond<sup>n</sup> to install package X, dependencies should be there  
before it.

DFS.

hint a DS.

Stack.

How?



pseudo :      DFS  $\rightarrow$  Backtrack.

Void topoSort (g) :

stack s;    visited[n] = {false}

for i=0 i<n i++

if !visited[i]

topoSortVisit (g, i, visited, s)

// print the stack

while !s.empty():    print (s.top())  
                             s.pop()

topoSortVisit(g, u, visited, s):

visited[u] = true

for each v in g[u]:

if !visited[v]:

topoSortVisit

(g, v, ...)

3 2 5 4,  
          

$\rightarrow$  s.push(u)

