

Recap: LCS longest Common Subsequence

s_1 : _____
 s_2 : _____ } pick
chars
that are
in same
order in
both the
strings.

DP if $m == -1$ || $n == -1$:

O/p: length
of LCS.

* if $memo[m][n] != -1$: ~~return~~ $memo[m][n]$ int $memo[m+1][n+1] = \{-1\}$
if $s_1[m] == s_2[n]$:

~~$memo[m][n] = return$~~ $1 + \underline{lcs}(s_1, s_2, m-1, n-1)$

T.C. $O(m \cdot n)$

else
 ~~$memo[m][n] = return$~~ $\max \{ \underline{lcs}(s_1, s_2, m+1, n), \underline{lcs}(s_1, s_2, m, n+1) \}$
~~return~~ $memo[m][n]$

5 variations : (1) len longest common substring \rightarrow if \rightarrow keep track max len
 \rightarrow jump \rightarrow return 0.

(2) SCS, shortest common supersequence s_1, s_2
 \downarrow
len s.t. s_1 and s_2 are subsequences of SCS.

EASY

$$\text{len} = |m + n - \text{lcs}|$$

ops: (3) min no. of operations to convert $s_1 \rightarrow s_2$.
1 insertion }
1 deletion }
 $\# \text{ min ops} = |m + n - 2 * \text{lcs}|$

(4) longest palindromic subsequence.

$$\hookrightarrow \text{lcs}(s, s.\text{reverse}())$$

(5) min ops required to make s palindrome.

$$\# \text{ min ops} = |m - \text{lcs}| \rightarrow \text{either insertions or deletions.}$$

Tough. the actual LCS. | content with the length.

LCS itself !!

⑥ Give me the LCS. $s_1 = \underline{a}\underline{b}\underline{a}\underline{c}\underline{d}$ ⑥

$s_2 = \underline{a}\underline{a}\underline{b}\underline{d}\underline{c}\underline{c}\underline{d}$ ⑧

} LCS $\rightarrow \underline{4}$
 $\{a, b, c, d\}$
 $\{a, a, c, d\}$
 one of these !!

Take help of memo table.

final memo

	0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0	0
1	0	1	1	1	1	1	1	1
2	0	1	1	2	2	2	2	2
3	0	1	2	2	2	2	2	2
4	0	1	2	2	2	3	3	3
5	0	1	2	2	3	3	3	4

for (int i=1, i<=n, i++)

for (int j=1, j<=n, j++)

a a c d

*

Knapsack \rightarrow tabulation | No recursion

	0	1	2	3	4	5	6	7	8	9
0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	0	0	0	0	0	0
2	0	0	0	0	0	0	0	0	0	0
3	0	0	0	0	0	0	0	0	0	0
4	0	0	0	0	0	0	0	0	0	0
5	0	0	0	0	0	0	0	0	0	0
6	0	0	0	0	0	0	0	0	0	0
7	0	0	0	0	0	0	0	0	0	0
8	0	0	0	0	0	0	0	0	0	0
9	0	0	0	0	0	0	0	0	0	0

for (i=1, i<=n, i++)

for (j=1, j<=C, j++)

if $w[i] > j$:

$memo[i][j] = memo[i-1][j]$

else

$memo[i][j] = \max(memo[i-1][j], p[i] +$

$memo[i-1][j - w[i]])$

Manual fill. $m+1$
 $n+1$

1. Base cond.

if str1 finishes || str2 finishes:

return 0

2. if $s_1[i] == s_2[j]$:

$memo[i][j] =$

$1 + memo[i-1][j-1]$

else:

$memo[i][j] =$

$\max(memo[i-1][j], memo[i][j-1])$

Tabulation X

Unless Recursion.

LCS | use memo.

$i=m, j=n \rightarrow$ address of final cell \rightarrow end cell

while ($i > 0 \parallel j > 0$):

if $s_1[i] == s_2[j]$:

ans += $s_1[i]$ \rightarrow char

$i-- \quad j--$

else:

if $j > 0$: $j--$

else: $i--$

\rightarrow reach 1st cell

\rightarrow char. move diag.

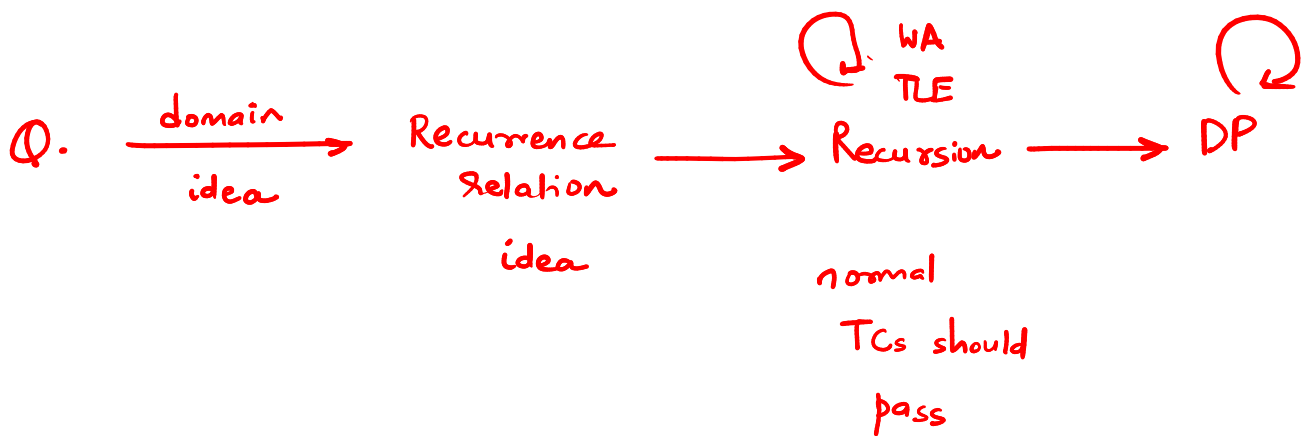
\rightarrow otherwise i or j

leetcode

LCS

\downarrow

+



⑦ longest repeating sequence :

LRS $\rightarrow 3 \quad \{a, b, d\}$.

+ 1 condition in if.

$s_1 = aabbbcd$

$LCS(s, s, m, n)$:

if $m == 0 \parallel n == 0$:
return 0

⑧ SCS \rightarrow string itself!

⑨ This was an independent question.

Today

it is a variation!

if $s[m] == s[n]$ and $m \neq n$:

$\rightarrow 1 + LCS(s, s, m-1, n-1)$

else

$\rightarrow \max(LCS(s, s, m, n-1), LCS(s, s, m-1, n))$

arr = [10, 9, 2, 5, 3, 7, 101, 18]

subsequence = [2, 3, 7, 18]

Longest Increasing Subsequence.

len = 4.

increasing



sorted.

$$LIS = LCS(arr, \text{sort}(arr))$$