

# Data compression

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# Data Compression

- Data in memory have used fixed length for representation
- For data transfer (in particular), this method is inefficient.
- For speed and storage efficiencies, data symbols should use the minimum number of bits possible for representation.
- Methods Used For Compression:
  - Encode high probability symbols with fewer bits
    - Shannon-Fano, Huffman, UNIX compact
  - Encode sequences of symbols with location of sequence in a dictionary
    - PKZIP, ARC, GIF, UNIX compress, V.42bis
  - Lossy compression
    - JPEG and MPEG

# Variable Length Bit Codings

- Suppose 'A' appears 50 times in text, but 'B' appears only 10 times
- ASCII coding assigns 8 bits per character, so total bits for 'A' and 'B' is  $60 * 8 = 480$
- If 'A' gets a 4-bit code and 'B' gets a 12-bit code, total is  $50 * 4 + 10 * 12 = 320$

*Compression rules:*

- Use minimum number of bits
- No code is the prefix of another code
- Enables left-to-right, unambiguous decoding

# Variable Length Bit Codings

- No code is a prefix of another
  - For example, can't have 'A' map to 10 and 'B' map to 100, because 10 is a prefix (the start of) 100.
- Enables left-to-right, unambiguous decoding
  - That is, if you see 10, you know it's 'A', not the start of another character.

# Huffman code

- Constructed by using a code tree, but starting at the leaves
- A compact code constructed using the binary Huffman code construction method

## Huffman code Algorithm

- ① **Make a leaf node for each code symbol**
  - ◆ Add the generation probability of each symbol to the leaf node
- ② **Take the two leaf nodes with the smallest probability and connect them into a new node**
  - ◆ Add 1 or 0 to each of the two branches
  - ◆ The probability of the new node is the sum of the probabilities of the two connecting nodes
- ③ **If there is only one node left, the code construction is completed. If not, go back to (2)**

# Demo

- 65demo-huffman.ppt

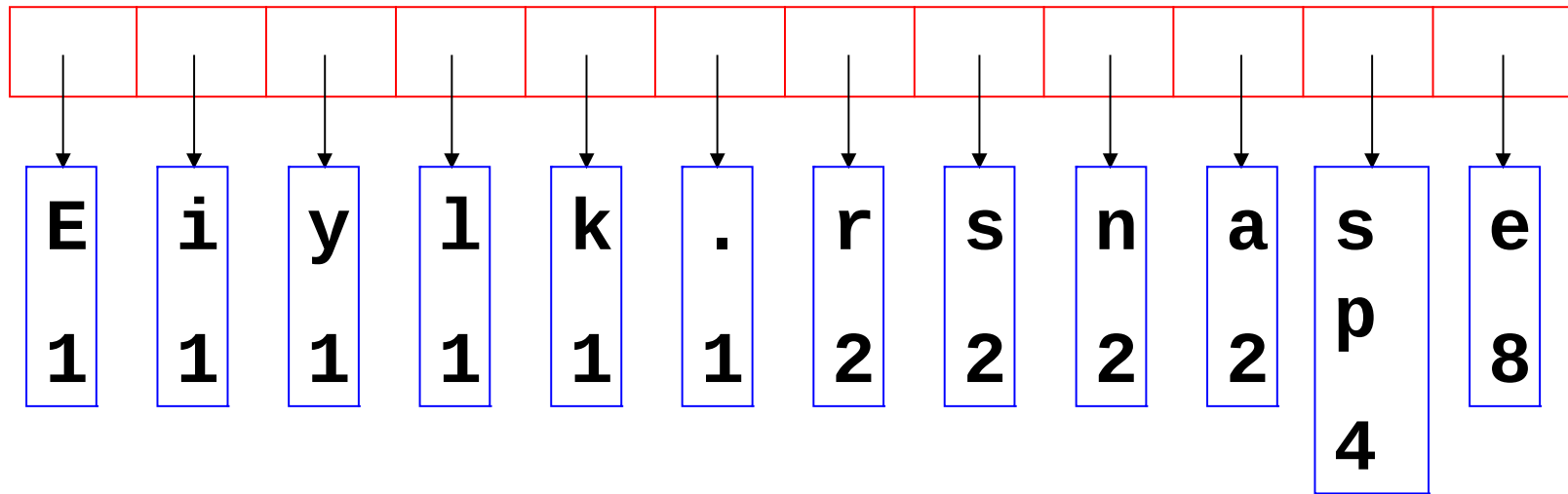
# Compress a text

- Consider the following short text:  
*Eerie eyes seen near lake.*
- Count up the occurrences of all characters in the text

Char	Freq.	Char	Freq.	Char	Freq.
E	1	y	1	k	1
e	8	s	2	.	1
r	2	n	2		
i	1	a	2		
space	4	l	1		

# Building a Tree

- The queue after inserting all nodes



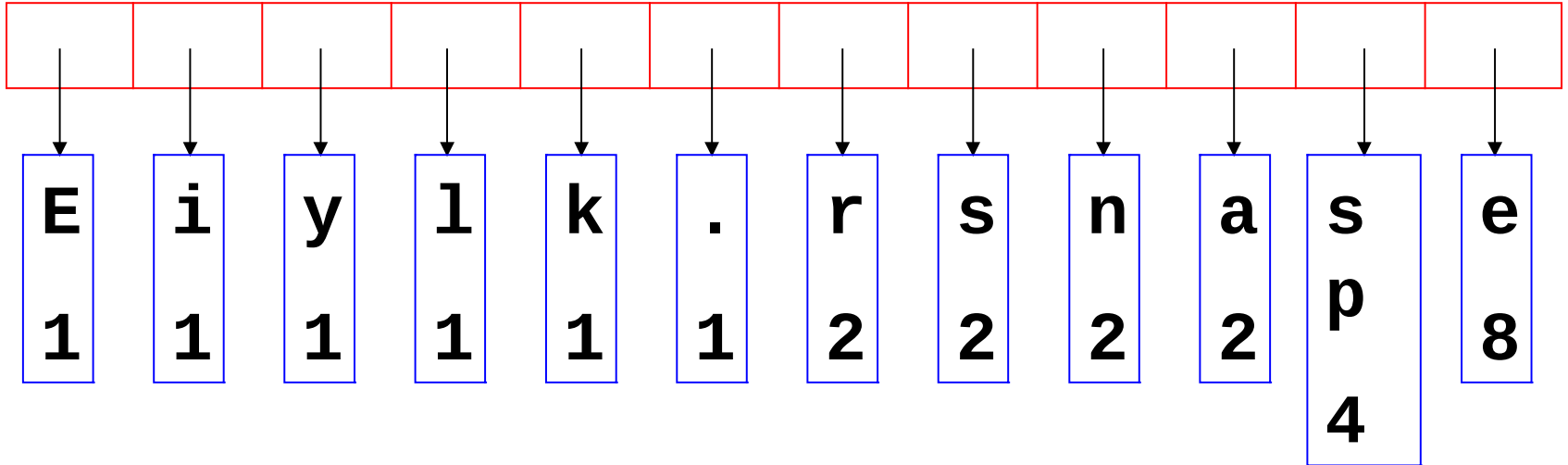
- Null Pointers are not shown



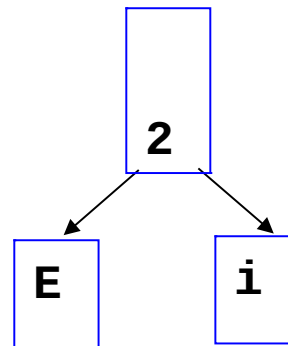
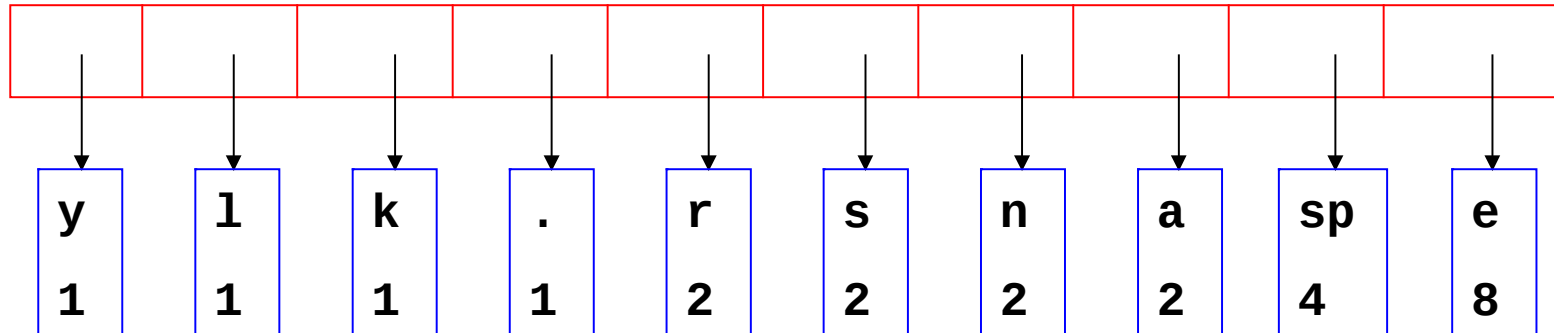
# Building a Tree

- While priority queue contains two or more nodes
  - Create new node
  - Dequeue node and make it left subtree
  - Dequeue next node and make it right subtree
  - Frequency of new node equals sum of frequency of left and right children
  - Enqueue new node back into queue

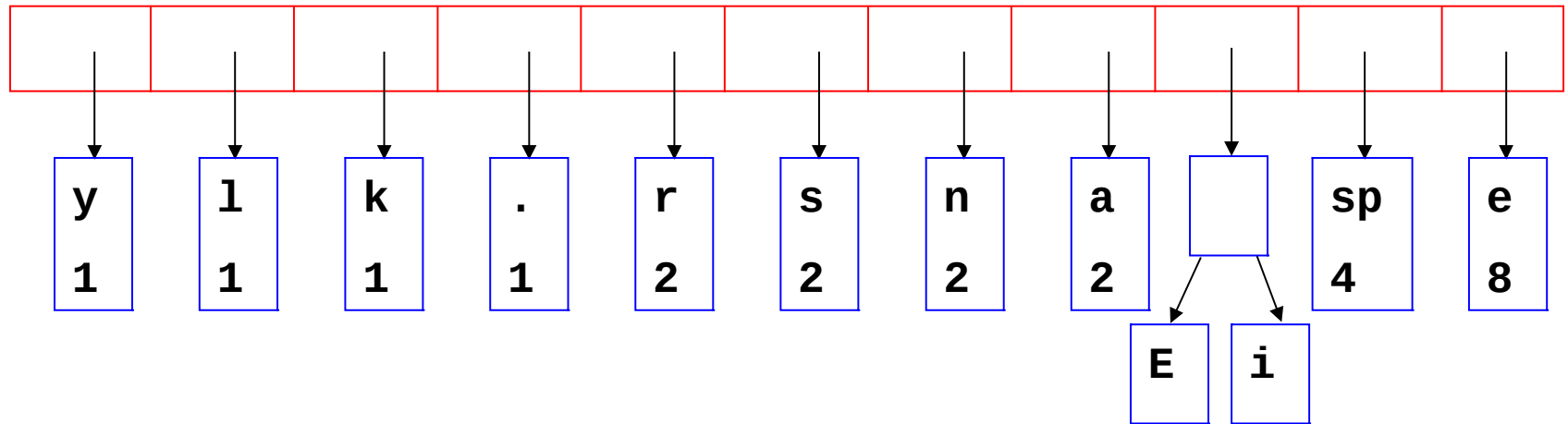
# Building a Tree



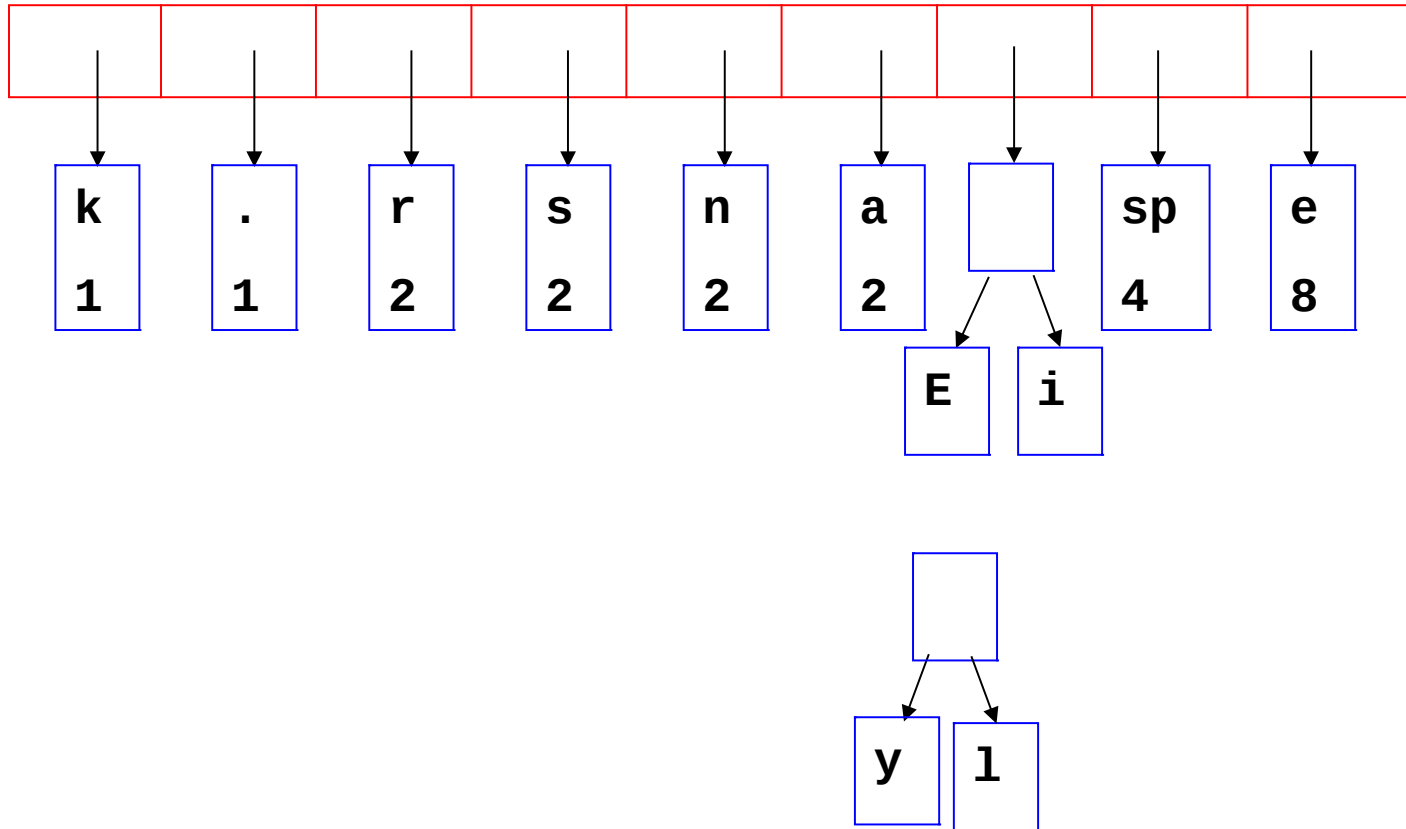
# Building a Tree



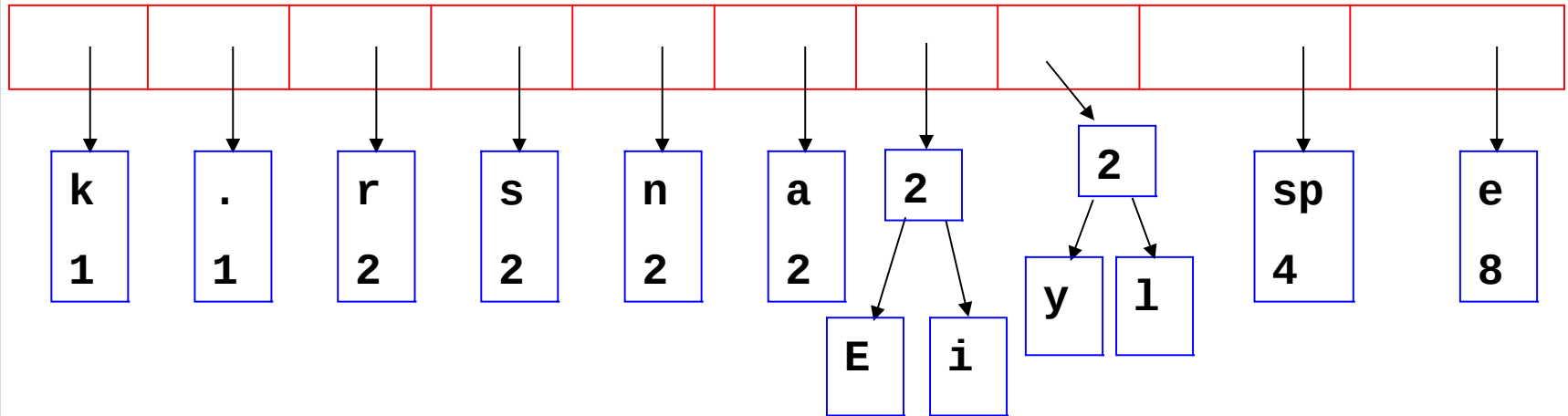
# Building a Tree



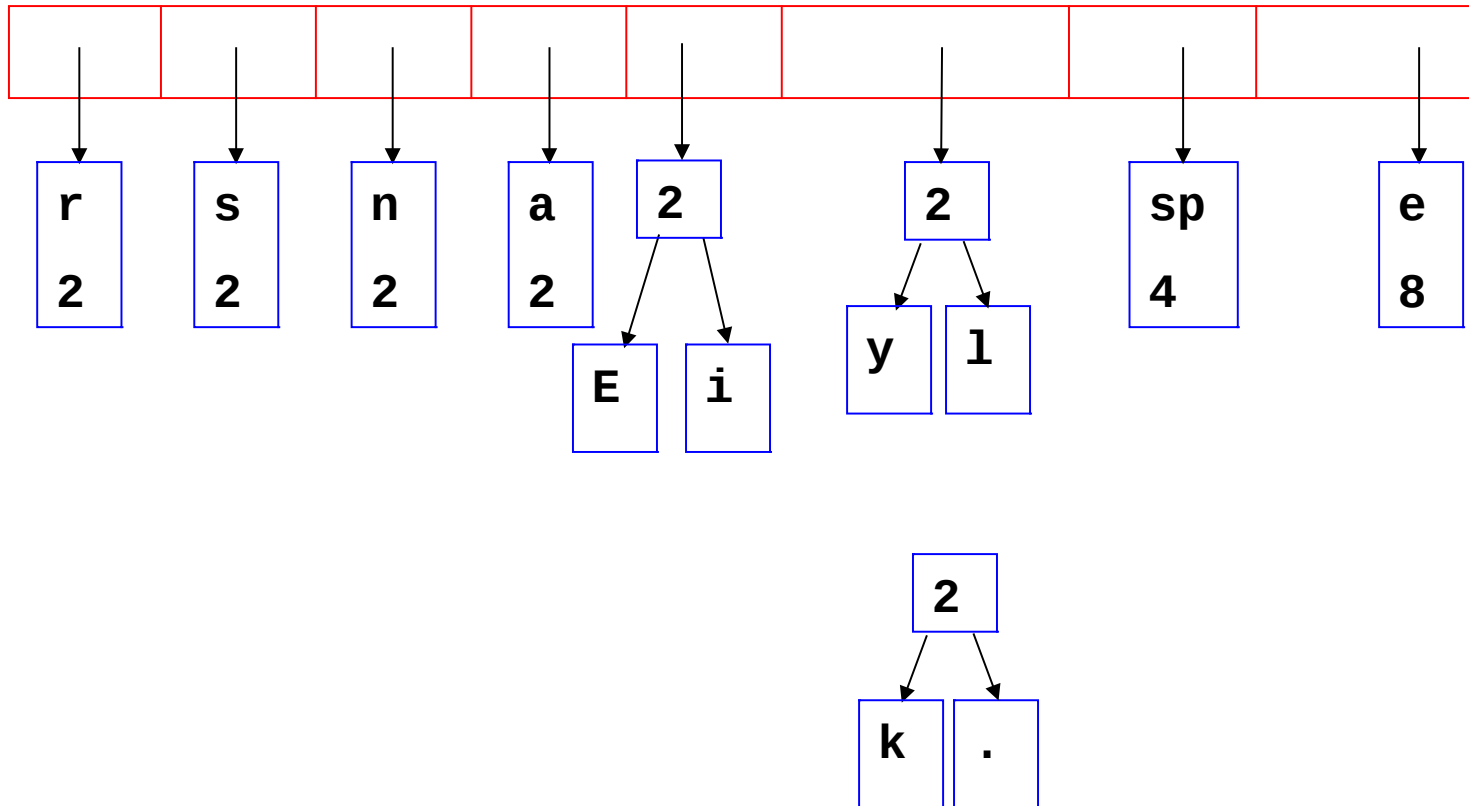
# Building a Tree



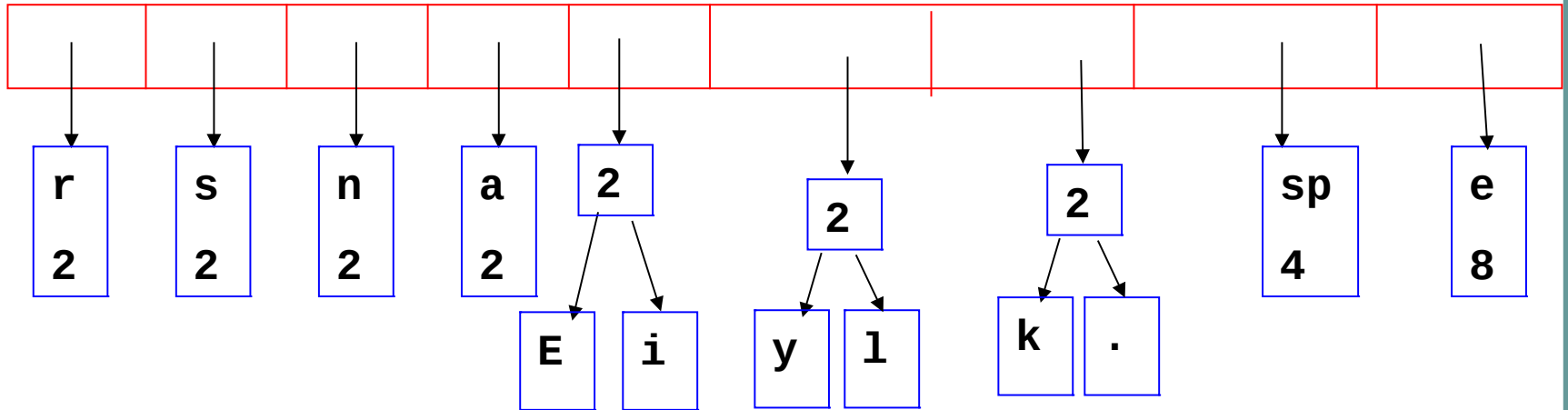
# Building a Tree



# Building a Tree



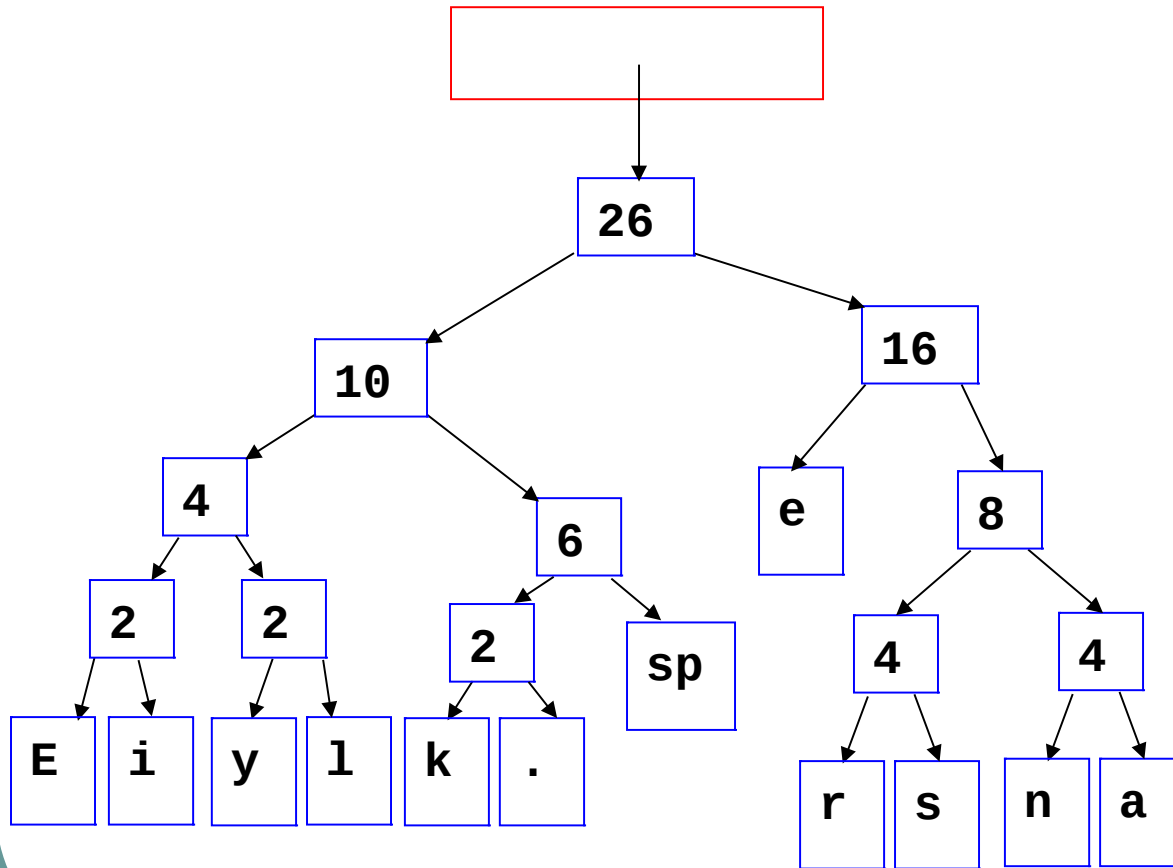
# Building a Tree



- To continue ...



# At the end



After enqueueing this node there is only one node left in priority queue.

# How to implement ?

- Reuse JRB to represent the tree
  - Each new node is created as a JRB node
  - The edges are directional from the parents to the children.
  - Two edges are created and marked using label 0 or 1 when a parent node is created.
- Reuse Dlist or JRB to represent the priority queue
  - A queue node contains a key as the frequency of the related node in the tree
  - The queue node's value is a pointer referencing to the node in the tree

# Quiz 1

- Reuse the graph API defined in previous class to write a function that builds a Huffman tree from a string as the following

```
typedef struct {
```

```
    Graph graph;
```

```
    JRB root;
```

```
} HuffmanTree;
```

```
HuffmanTree makeHuffman (char * buffer, int size);
```

# Huffman code table

- In order to compress the data string, we have to build a code table from the Huffman tree. The following data structure is used to represent the code table

```
typedef struct {  
    int size;  
    char bits[2];  
} Coding;  
Coding huffmanTable[256];
```

- `huffmanTable['A']` give the coding of 'A'. If the coding's size = 0, the character 'A' is not present in the text. `bits` contains the huffman code (sequence of bits) of the given character.

# Quiz 2

- Write a function to create the Huffman code table from a Huffman tree
  - `void createHuffmanTable(HuffmanTree htree, Coding* htable);`
- Write a function to compress a text buffer to a Huffman sequence.
  - `void compress(char * buffer, int size, char* huffman, int* nbit);`
- The buffer contains *size* characters. After compressing, the huffman buffer contains *nbit* bits for output.
- In order to write this function, you should create a function to add a new character into the huffman buffer as the following
  - `void addHuffmanChar(char * ch, Coding* htable, char* huffman, int* nbit);`