

Exercise n°1:

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① The standard formulation of (P) is :
$$\begin{cases} \min_x c^T x \\ \text{s.t. } Ax - b = 0 \\ -x \leq 0 \end{cases}$$

The Lagrangian is given by :

$$\begin{aligned} \mathcal{L}(x, \lambda, \nu) &= c^T x - \lambda^T (Ax - b) + \nu^T (Ax - b) \\ &= (c^T + \nu^T A - \lambda^T) x - \nu^T b \\ &= (A^T \nu + c - \lambda)^T x - b^T \nu \end{aligned}$$

$$\text{Then, } g(\lambda, \nu) = \inf_{x \in \mathbb{R}^d} \mathcal{L}(x, \lambda, \nu) = \begin{cases} -b^T \nu & \text{if } A^T \nu + c - \lambda = 0 \\ -\infty & \text{otherwise} \end{cases}$$

$$\begin{aligned} \text{Thus, the dual of (P) is : } & \begin{cases} \max_{(\lambda, \nu) \in \mathbb{R}^q \times \mathbb{R}^m} g(\lambda, \nu) \\ \text{s.t. } \lambda \geq 0 \end{cases} \Leftrightarrow \begin{cases} \max_{\nu \in \mathbb{R}^m} -b^T \nu \\ \text{s.t. } \lambda \geq 0 \text{ and } A^T \nu + c = \lambda \end{cases} \\ & \Leftrightarrow \boxed{\begin{cases} \max_{\nu} b^T \nu \\ \text{s.t. } A^T \nu \leq c \end{cases}} \end{aligned}$$

I replaced ν with $(-\nu)$

We can notice that the dual of (P) is (D).

② The standard form of (D) is :
$$\begin{cases} \min_y -b^T y \\ \text{s.t. } A^T y - c \leq 0 \end{cases}$$

The Lagrangian is given by :

$$\begin{aligned} \mathcal{L}(y, \lambda) &= -b^T y + \lambda^T (A^T y - c) \\ &= (A\lambda - b)^T y - \lambda^T c = (A\lambda - b)^T y - c^T \lambda \end{aligned}$$

$$\text{Then, } g(\lambda) = \inf_{y \in \mathbb{R}^n} \mathcal{L}(y, \lambda) = \begin{cases} -c^T \lambda & \text{if } A\lambda = b \\ -\infty & \text{otherwise} \end{cases}$$

$$\text{Thus, the dual of (D) is : } \begin{cases} \max_{\lambda} g(\lambda) \\ \text{s.t. } \lambda \geq 0 \end{cases} \Leftrightarrow \boxed{\begin{cases} \min_{\lambda} c^T \lambda \\ \text{s.t. } A\lambda = b \\ \lambda \geq 0 \end{cases}}$$

Therefore, the dual of (D) is (P).

③ We can notice that the "Self-Dual" problem is $(P) \cap (D)$, (P) and (D) are independent. So, $\text{Dual}((P) \cap (D)) = \text{Dual}((P)) \cap \text{Dual}((D)) = (P) \cap (D)$.

More rigorously, let us compute the Lagrangian:

$$\begin{aligned} \mathcal{L}(x, y, \lambda_1, \lambda_2, \nu) &= c^T x - b^T y - \lambda_1^T x + \lambda_2^T (A^T y - c) + \nu^T (Ax - b) \\ &= (c^T - \lambda_1^T + \nu^T A) x + (\lambda_2^T A^T - b^T) y - c^T \lambda_2 - b^T \nu \\ &= (A^T \nu + c - \lambda_1)^T x + (A\lambda_2 - b)^T y - c^T \lambda_2 - b^T \nu \end{aligned}$$

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$$\begin{aligned} \text{Then } g(\lambda_1, \lambda_2, v) &= \inf_{x, y} \mathcal{L}(x, y, \lambda_1, \lambda_2, v) \\ &= \left(\inf_x (A^T v + c - \lambda_1)^T x \right) + \left(\inf_y (A \lambda_2 - b)^T y \right) - c^T \lambda_2 - b^T v. \\ &= \begin{cases} -c^T \lambda_2 - b^T v & \text{if } A^T v + c = \lambda_1 \text{ and } A \lambda_2 = b \\ -\infty & \text{otherwise.} \end{cases} \end{aligned}$$

Then, the dual of "Self-Dual" is given by:

$$\begin{cases} \max_{(\lambda, v)} & -c^T \lambda_2 - b^T v \\ \text{s.t.} & A^T v + c \geq 0 \\ & A \lambda_2 = b \\ & \lambda_2 \geq 0 \end{cases}$$

I just substituted v with $(-v)$

which is equivalent to

$$\begin{cases} \max_{(\lambda, v)} & -c^T \lambda + b^T v \\ \text{s.t.} & -A^T v + c \geq 0 \\ & A \lambda = b \\ & \lambda \geq 0 \end{cases}$$

which is equivalent to

$$\begin{cases} \min_{(\lambda, v)} & c^T \lambda - b^T v \\ \text{s.t.} & A^T v \leq c \\ & A \lambda = b \\ & \lambda \geq 0 \end{cases}$$

Thus, the dual of "Self-Dual" is "Self-Dual". i.e. the problem is self-dual

④ a) Since (x^*, y^*) is optimal for "Self-Dual", then, $\begin{cases} Ax^* = b \\ x^* \geq 0 \end{cases}$ and $A^T y^* \leq c$.

So, (P) and (D) are feasible.

• Let \bar{x} a feasible solution for (P). Suppose with absurd that $c^T \bar{x} < c^T x^*$.
Then, $c^T \bar{x} - b^T y^* < c^T x^* - b^T y^*$: impossible because (x^*, y^*) is optimal for "Self-Dual".

Then, x^* is optimal for (P).

• Similarly, let \bar{y} a feasible solution for (D). Suppose with absurd that $b^T \bar{y} > b^T y^*$.
Then, $c^T x^* - b^T \bar{y} < c^T x^* - b^T y^*$: impossible because (x^*, y^*) is optimal for "Self-Dual".

Then, y^* is optimal for (D).

b). If we assume that strong duality is applicable for LPs: since (D) is the dual of (P) and both are feasible LPs, so, $d^* = p^*$ i.e. $c^T x^* = b^T y^*$ i.e. $c^T x^* - b^T y^* = 0$.

• Actually, we don't need strong duality. In fact, weak duality is sufficient. In fact, since (D) is the dual of (P), with weak duality, we get $b^T y^* \leq c^T x^*$. Besides, (P) is the dual of (D), so, with weak duality, we have $c^T x^* \leq b^T y^*$. Hence, $c^T x^* = b^T y^*$.

Therefore, the optimal value of "Self-Dual" is exactly 0.

Exercise n°2:

(3)

① Let $F: x \in \mathbb{R}^d \mapsto y^T x - \|x\|_1$ for $y \in \mathbb{R}^d$. Let $(\vec{e}_k)_{1 \leq k \leq d}$ the canonical basis of \mathbb{R}^d .

• Suppose that $\exists k \in [1, d]$ such that $y_k > 1$:

Then, $\forall t > 0$, $F(t\vec{e}_k) = t(y_k - 1) \xrightarrow{t \rightarrow +\infty} +\infty \Rightarrow F$ is unbounded.

• Suppose that $\exists k \in [1, d]$ such that $y_k < -1$:

Then, $\forall t > 0$, $F(-t\vec{e}_k) = -t(y_k + 1) \xrightarrow{t \rightarrow +\infty} +\infty \Rightarrow F$ is unbounded.

Thus, F is bounded $\Rightarrow \forall k \in [1, d]$, $|y_k| \leq 1$ i.e. $\|y\|_\infty \leq 1$.

• Now, let $y \in \mathbb{R}^d$ such that $\|y\|_\infty \leq 1$:

$$\forall i \in [1, d]; x_i y_i \leq |x_i y_i| \leq \|y\|_\infty |x_i|$$

$$\text{Then, } \sum_{i=1}^d x_i y_i \leq \sum_{i=1}^d \|y\|_\infty |x_i| = \|y\|_\infty \|x\|_1 \leq \|x\|_1$$

$$\text{i.e. } F(x) \leq 0.$$

$$\text{Moreover, } F(0) = 0.$$

$$\text{Then, } \sup_{x \in \mathbb{R}^d} F(x) = 0$$

Thus, the conjugate of $\|x\|_1$ is given by:

$$f^*: \mathbb{R}^d \rightarrow \mathbb{R}$$

$$y \mapsto \begin{cases} 0 & \text{if } \|y\|_\infty \leq 1 \\ +\infty & \text{otherwise.} \end{cases}$$

$$\textcircled{2} \text{ (RLS): } \min_{x \in \mathbb{R}^d} \|Ax - b\|_2^2 + \|x\|_1 \Leftrightarrow \begin{cases} \min_{(x,y) \in \mathbb{R}^d \times \mathbb{R}^m} \|y\|_2^2 + \|x\|_1 \\ \text{subject to } y = Ax - b. \end{cases}$$

* The Lagrangian is given by:

$$\mathcal{L}(x, y, v) := \|y\|_2^2 + \|x\|_1 + v^T (y - (Ax - b))$$

$$= \|y\|_2^2 + v^T y + \|x\|_1 - (A^T v)^T x + v^T b$$

$$g(v) := \inf_{x,y} \mathcal{L}(x, y, v)$$

$$= \inf_y (\|y\|_2^2 + v^T y) + \inf_x (\|x\|_1 - (A^T v)^T x) + v^T b$$

$$\text{• We notice that: } \inf_x (\|x\|_1 - (A^T v)^T x) = - \sup_x ((A^T v)^T x - \|x\|_1) = -f^*(A^T v)$$

$$= \begin{cases} 0 & \text{if } \|A^T v\|_\infty \leq 1 \\ -\infty & \text{otherwise.} \end{cases}$$

$$\text{• Let } \varphi: y \in \mathbb{R}^m \mapsto \|y\|_2^2 + v^T y = y^T y + v^T y.$$

$$\varphi \text{ is differentiable; } \nabla \varphi(y) = 2y + v. \text{ Then, } \nabla \varphi(y) = 0 \Leftrightarrow y = -\frac{1}{2}v.$$

$$\varphi \text{ is a convex function, then, } \varphi \text{ reaches its global minimum at } y = -\frac{1}{2}v$$

$$\text{and } \inf_y \|y\|_2^2 + v^T y = -\frac{1}{4}v^T v.$$

Thus, $g(v) = \begin{cases} -\frac{1}{4} v^T v + b^T v & \text{if } \|A^T v\|_\infty \leq 1 \\ -\infty & \text{otherwise} \end{cases}$

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Then, the dual of (RLS) is: $\begin{cases} \max_v & -\frac{1}{4} v^T v + b^T v \\ \text{s.t.} & \|A^T v\|_\infty \leq 1 \end{cases}$

This is equivalent to: $\begin{cases} \min_v & \frac{1}{4} v^T v + b^T v \\ \text{s.t.} & \|A^T v\|_\infty \leq 1 \end{cases}$

Exercise m°3:

① We notice that the problem (Sep2) is equivalent to $\begin{cases} \min_{(w, z)} & \frac{1}{n} \mathbb{1}^T z + \frac{\tau}{2} \|w\|_2^2 \\ \text{s.t.} & z_i \geq \mathcal{L}(w, x_i, y_i) \quad \forall i \in \llbracket 1, m \rrbracket. \end{cases}$

Let (w^*, z^*) is an optimal solution for (Sep2).

Suppose, that $\exists \bar{w} \in \mathbb{R}^d$ such that: $\frac{1}{m} \sum_{i=1}^m \mathcal{L}(\bar{w}, x_i, y_i) + \frac{\tau}{2} \|\bar{w}\|_2^2 < \frac{1}{m} \sum_{i=1}^m \mathcal{L}(w^*, x_i, y_i) + \frac{\tau}{2} \|w^*\|_2^2$ (by absurd)

Let $\bar{z} \in \mathbb{R}^m$ such that $\bar{z}_i := \mathcal{L}(\bar{w}, x_i, y_i) \quad \forall i \in \llbracket 1, m \rrbracket$

Then, (\bar{w}, \bar{z}) is a feasible solution for (Sep2).

Moreover, we have: $\frac{1}{n} \mathbb{1}^T \bar{z} + \frac{\tau}{2} \|\bar{w}\|_2^2 = \frac{1}{m} \sum_{i=1}^m \mathcal{L}(\bar{w}, x_i, y_i) + \frac{\tau}{2} \|\bar{w}\|_2^2 < \frac{1}{m} \sum_{i=1}^m \mathcal{L}(w^*, x_i, y_i) + \frac{\tau}{2} \|w^*\|_2^2$
 $< \frac{1}{m} \sum_{i=1}^m \underbrace{\mathcal{L}(w^*, x_i, y_i)}_{\leq z_i^*} + \frac{\tau}{2} \|w^*\|_2^2$
 $< \frac{1}{m} \mathbb{1}^T z^* + \frac{\tau}{2} \|w^*\|_2^2$
 \Rightarrow Absurd because (w^*, z^*) is optimal for (Sep2).

Therefore, $\forall w \in \mathbb{R}^d$, $\frac{1}{m} \sum_{i=1}^m \mathcal{L}(w, x_i, y_i) + \frac{\tau}{2} \|w\|_2^2 \geq \frac{1}{m} \sum_{i=1}^m \mathcal{L}(w^*, x_i, y_i) + \frac{\tau}{2} \|w^*\|_2^2$
 ie w^* is an optimal solution for (Sep1).

Thus, (Sep2) solves (Sep1).

② Let $L := \text{diag}(y) X$ where $\text{diag}(y) = \begin{pmatrix} y_1 & & 0 \\ & \ddots & \\ 0 & & y_m \end{pmatrix} \in \mathbb{R}^{m \times m}$ and $X = (x_1, \dots, x_m)^T \in \mathbb{R}^{n \times d}$

which means that $[Lw]_i = y_i w^T x_i \quad \forall i \in \llbracket 1, m \rrbracket$.

The Lagrangian is given by:

$\mathcal{L}(w, z, \lambda, \pi) = \frac{1}{m\tau} \mathbb{1}^T z + \frac{1}{2} \|w\|_2^2 - \sum_{i=1}^m \lambda_i (z_i - 1 + y_i w^T x_i) - \pi^T z$
 $= \frac{1}{m\tau} \mathbb{1}^T z + \frac{1}{2} \|w\|_2^2 - \lambda^T z + \mathbb{1}^T \lambda - \lambda^T Lw - \pi^T z$

$$\mathcal{L}(w, z, \lambda, \pi) = \left(\frac{1}{mC} \mathbf{1}^T - \lambda^T - \pi^T \right) z + \frac{1}{2} \|w\|_2^2 - \lambda^T L w + \mathbf{1}^T \lambda$$

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$$= \left(\frac{1}{mC} \mathbf{1} - \lambda - \pi \right)^T z + \frac{1}{2} \|w\|_2^2 - \lambda^T L w + \mathbf{1}^T \lambda$$

$$g(\lambda, \pi) = \inf_{w, z} \mathcal{L}(w, z, \lambda, \pi)$$

$$= \left(\inf_z \left(\frac{1}{mC} \mathbf{1} - \lambda - \pi \right)^T z \right) + \inf_w \left(\frac{1}{2} \|w\|_2^2 - \lambda^T L w \right) + \mathbf{1}^T \lambda$$

$$\bullet \inf_z \left(\frac{1}{mC} \mathbf{1} - \lambda - \pi \right)^T z = \begin{cases} 0 & \text{if } \frac{1}{mC} \mathbf{1} - \pi = \lambda \\ -\infty & \text{otherwise} \end{cases}$$

$$\bullet \text{ Let } f: w \mapsto \frac{1}{2} \|w\|_2^2 - \lambda^T L w$$

$$\nabla f(w) = w - L^T \lambda \quad ; \quad f \text{ is convex} \quad \text{Then, } \inf_w \frac{1}{2} \|w\|_2^2 - \lambda^T L w = -\frac{1}{2} \lambda^T L L^T \lambda$$

$$\text{Thus, } g(\lambda, \pi) = \begin{cases} -\frac{1}{2} \lambda^T L L^T \lambda + \mathbf{1}^T \lambda & \text{if } \lambda = \frac{1}{mC} \mathbf{1} - \pi \\ -\infty & \text{otherwise} \end{cases}$$

$$\text{Therefore the dual is given by } \begin{cases} \max_{(\lambda, \pi)} g(\lambda, \pi) \\ \text{st. } \lambda \geq 0 \\ \pi \geq 0 \end{cases}$$

$$\Leftrightarrow \begin{cases} \max_{\lambda} -\frac{1}{2} \lambda^T L L^T \lambda + \mathbf{1}^T \lambda \\ \text{st. } \lambda \geq 0 \\ \lambda \leq \frac{1}{mC} \mathbf{1} \end{cases}$$

$$\Leftrightarrow \boxed{\begin{cases} \min_{\lambda} \frac{1}{2} \lambda^T L L^T \lambda - \mathbf{1}^T \lambda \\ \text{st. } \lambda \geq 0 \\ \lambda \leq \frac{1}{mC} \mathbf{1} \end{cases}}$$

Exercise n°4:

NB: C is a matrix, $C \in \mathbb{R}^{p \times n}$, $d \in \mathbb{R}^p$, $\mathcal{P} = \{a \in \mathbb{R}^n \mid Ca \leq d\}$

(6)

We notice that $\sup_{a \in \mathcal{P}} a^T x$ is the optimal value of the following LP: $(*) \begin{cases} \max_a x^T a \\ \text{s.t. } Ca \leq d \end{cases}$
(where the variable is a).

The associated Lagrangian is: $\mathcal{L}(a, \lambda) = -x^T a + \lambda^T (Ca - d)$
 $= (C^T \lambda - x)^T a - d^T \lambda$

$$\Rightarrow g(\lambda) = \inf_a \mathcal{L}(a, \lambda) = \begin{cases} -d^T \lambda & \text{if } C^T \lambda = x \\ -\infty & \text{otherwise} \end{cases}$$

Then, the dual of $(*)$ is $\begin{cases} \max_{\lambda \in \mathbb{R}^p} -d^T \lambda \\ \text{s.t. } C^T \lambda = x \\ \lambda \geq 0 \end{cases} \iff (**) \begin{cases} \min_{\lambda \in \mathbb{R}^p} d^T \lambda \\ \text{s.t. } C^T \lambda = x \\ \lambda \geq 0 \end{cases}$

Since we are dealing with LPs here, we can apply the strong duality theorem.

Thus, $\sup_{a \in \mathcal{P}} a^T x$ is the optimal value of $(**)$.

Therefore, $\sup_{a \in \mathcal{P}} a^T x \leq b \iff \exists \lambda^* \in \mathbb{R}^p$ such that $\begin{cases} d^T \lambda^* \leq b \\ C^T \lambda^* = x \\ \lambda^* \geq 0 \end{cases}$

Thus, $\begin{cases} \min_{x \in \mathbb{R}^n} c^T x \\ \text{s.t. } \sup_{a \in \mathcal{P}} a^T x \leq b \end{cases} \iff \begin{cases} \min_{x \in \mathbb{R}^n} c^T x \\ \text{s.t. } d^T z \leq b \\ C^T z = x \\ z \geq 0 \end{cases}$

Exercise n°5:

$$\begin{aligned} \textcircled{1} \quad \mathcal{L}(x, \lambda, \nu) &= c^T x + \lambda^T (Ax - b) - \sum_{i=1}^m \nu_i x_i (1 - x_i) \\ &= c^T x + \lambda^T (Ax - b) - \nu^T x + x^T \text{diag}(\nu) x \\ &= x^T \text{diag}(\nu) x + (A^T \lambda + c - \nu)^T x - b^T \lambda \end{aligned}$$

• If $\exists i \in \llbracket 1, m \rrbracket$ such that $\nu_i < 0$:
 $\forall t > 0, \mathcal{L}(t \vec{e}_i, \lambda, \nu) = -\nu_i t^2 + ((A^T \lambda)_i + \nu_i + c_i) t - b^T \lambda \xrightarrow{t \rightarrow +\infty} -\infty \therefore$ unbounded

Then, to search $\inf_x \mathcal{L}(x, \lambda, \nu)$, we must have $\boxed{\nu \geq 0}$.

• Let $J \subseteq \llbracket 1, m \rrbracket$ such that $\forall i \notin J, \nu_i = 0$ and $\forall i \in J, \nu_i > 0$. (J can be \emptyset).

We denote $J = \{i_1 < \dots < i_k\}$ and $\nu_J = (\nu_{i_1}, \dots, \nu_{i_k})^T$.

Similarly, we denote $\mu_J = (\mu_{i_1}, \dots, \mu_{i_k})^T$ for a given vector $\mu \in \mathbb{R}^d$.

and $\mu_{\bar{J}} = (\mu_j, j \notin J)^T$.

$$\text{So, } \mathcal{L}(x, \lambda, \nu) = x_J^T \text{diag}(\nu_J) x_J + ((A^T \lambda)_J + c_J - \nu_J)^T x_J + ((A^T \lambda)_{\bar{J}} + c_{\bar{J}}) x_{\bar{J}} - b^T \lambda$$

Suppose that $\exists j \notin J$ such that $(A^T \lambda)_j + c_j \neq 0$:

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$$\forall t \in \mathbb{R}, \mathcal{L}(t \vec{e}_j, \lambda, \nu) = ((A^T \lambda)_j + c_j) t$$

$$\square \text{ If } (A^T \lambda)_j + c_j > 0 : \text{ then } \lim_{t \rightarrow -\infty} \mathcal{L}(t \vec{e}_j, \lambda, \nu) = -\infty$$

$$\square \text{ If } (A^T \lambda)_j + c_j < 0 : \text{ then } \lim_{t \rightarrow +\infty} \mathcal{L}(t \vec{e}_j, \lambda, \nu) = -\infty$$

Thus, to define $\inf_x \mathcal{L}(x, \lambda, \nu)$, we must have $(A^T \lambda)_j + c_j = 0 \ \forall j \notin J$ (ie $\forall j$ such that $\nu_j = 0$)

• Now, we assume that $(A^T \lambda)_{\bar{J}} + c_{\bar{J}} = 0$ (ie $\forall j \notin J, (A^T \lambda)_j + c_j = 0$):

$$\mathcal{L}(x, \lambda, \nu) = x_{\bar{J}}^T \text{diag}(\nu_{\bar{J}}) x_{\bar{J}} + \left((A^T \lambda)_{\bar{J}} + c_{\bar{J}} - \nu_{\bar{J}} \right)^T x_{\bar{J}} - b^T \lambda$$

By definition, $\nu_{\bar{J}} > 0 \Rightarrow \text{diag}(\nu_{\bar{J}}) \in S_{|\bar{J}|}^{++}$ (in particular, it is invertible).

$$g(\lambda, \nu) = \inf_x \mathcal{L}(x, \lambda, \nu) = \inf_{x_{\bar{J}}} \underbrace{x_{\bar{J}}^T \text{diag}(\nu_{\bar{J}}) x_{\bar{J}} + \left((A^T \lambda)_{\bar{J}} + c_{\bar{J}} - \nu_{\bar{J}} \right)^T x_{\bar{J}}}_{\text{we note it } f(x_{\bar{J}})} - b^T \lambda$$

We want to minimize a quadratic function. Thanks to the calculus done in the former exercises, we get: $\arg\min_{x_{\bar{J}}} f(x_{\bar{J}}) = x_{\bar{J}}^*$ where $[x_{\bar{J}}^*]_j = -\frac{(A^T \lambda)_j + c_j - \nu_j}{\nu_j}, \forall j \in \bar{J}$.

$$\text{Thus, } \inf_{x_{\bar{J}}} f(x_{\bar{J}}) = f(x_{\bar{J}}^*) = -\frac{1}{4} \sum_{j \in \bar{J}} \frac{((A^T \lambda)_j + c_j - \nu_j)^2}{\nu_j} - b^T \lambda$$

* For $\nu \geq 0$: to have simple notations, we will define, for $i \in \llbracket 1, m \rrbracket$, $\frac{((A^T \lambda)_i + c_i - \nu_i)^2}{\nu_i}$ even if $\nu_i = 0$. In fact, if $\nu_i = 0$ and $(A^T \lambda)_i + c_i \neq 0$, then $\frac{((A^T \lambda)_i + c_i - \nu_i)^2}{\nu_i} = +\infty$ which is compatible to the fact that if $(A^T \lambda)_{\bar{J}} + c_{\bar{J}} \neq 0$, $\inf_x \mathcal{L}(x, \lambda, \nu) = -\infty$.

The other "pathologic" case for the formula $\frac{((A^T \lambda)_i + c_i - \nu_i)^2}{\nu_i}$ is when $\nu_i = 0$ and $(A^T \lambda)_i + c_i = 0$, but, in this case, $\frac{((A^T \lambda)_i + c_i - \nu_i)^2}{\nu_i} = \frac{0^2}{0}$, which can be interpreted as 0.

$$\Rightarrow \text{Thus, when } \nu \geq 0, \sum_{j \in \bar{J}} \frac{((A^T \lambda)_j + c_j - \nu_j)^2}{\nu_j} = \sum_{i=1}^m \frac{((A^T \lambda)_i + c_i - \nu_i)^2}{\nu_i}$$

$$\text{Therefore, } g(\lambda, \nu) = \begin{cases} -\frac{1}{4} \sum_{i=1}^m \frac{((A^T \lambda)_i + c_i - \nu_i)^2}{\nu_i} - b^T \lambda & \text{if } \nu \geq 0 \\ -\infty & \text{otherwise} \end{cases}$$

$$\text{Then, the dual is given by: } \begin{cases} \max_{\lambda, \nu} g(\lambda, \nu) \\ \text{s.t. } \lambda \geq 0 \end{cases} \Leftrightarrow \begin{cases} \max_{\lambda, \nu} -\frac{1}{4} \sum_{i=1}^m \frac{((A^T \lambda)_i + c_i - \nu_i)^2}{\nu_i} - b^T \lambda \\ \text{s.t. } \lambda \geq 0 \\ \nu \geq 0 \end{cases}$$

Let A_i^T denote the i th row of A . So, $(A^T \lambda)_i = A_i^T \lambda_i$ (8)

Our purpose is to simplify $\max_{\lambda, v \geq 0} -\frac{1}{4} \sum_{i=1}^n \frac{(A_i^T \lambda + c_i - v_i)^2}{v_i} - b^T \lambda$ (for $v \geq 0$).

We notice that:

$$\max_{\lambda, v \geq 0} -\frac{1}{4} \sum_{i=1}^n \frac{(A_i^T \lambda + c_i - v_i)^2}{v_i} - b^T \lambda = \max_{\lambda} \left(\max_{v \geq 0} -\frac{1}{4} \sum_{i=1}^n \frac{(A_i^T \lambda + c_i - v_i)^2}{v_i} - b^T \lambda \right)$$

$$= \max_{\lambda} \left(\frac{1}{4} \sum_{i=1}^n \left(\max_{v_i \geq 0} -\frac{(A_i^T \lambda + c_i - v_i)^2}{v_i} \right) - b^T \lambda \right).$$

Using the hint, $\max_{v_i \geq 0} -\frac{(A_i^T \lambda + c_i - v_i)^2}{v_i} = 4 \min \{0, c_i + A_i^T \lambda\}$

Therefore, the dual is given by:

$$\begin{cases} \max_{\lambda} & \sum_{i=1}^n \min \{0, c_i + A_i^T \lambda\} - b^T \lambda \\ \text{s.t.} & \lambda \geq 0 \end{cases} \quad (A)$$

(2). We start by computing the dual of (2): $\begin{cases} \min c^T x \\ \text{s.t.} & Ax \leq b \\ & 0 \leq x_i \leq 1 \quad \forall i \in [1, n] \end{cases}$

• The Lagrangian is:

$$\begin{aligned} \mathcal{L}(x, \lambda, v, w) &= c^T x + \lambda^T (Ax - b) - v^T x + w^T (x - \mathbb{1}) \\ &= (A^T \lambda + c - v + w)^T x - b^T \lambda - \mathbb{1}^T w. \end{aligned}$$

So, $g(\lambda, v, w) = \inf_x \mathcal{L}(x, \lambda, v, w) = \begin{cases} -b^T \lambda - \mathbb{1}^T w & \text{if } A^T \lambda + c + w = v \\ -\infty & \text{otherwise.} \end{cases}$

Then, the dual of (2) is:

$$\begin{cases} \max_{\lambda, v, w} & g(\lambda, v, w) \\ \text{s.t.} & \lambda \geq 0 \\ & v \geq 0 \\ & w \geq 0 \end{cases} \Leftrightarrow \begin{cases} \max_{\lambda, w} & -\mathbb{1}^T w - b^T \lambda \\ \text{s.t.} & \lambda \geq 0 \\ & w \geq 0 \\ & A^T \lambda + c + w \geq 0 \end{cases}$$

$$\Leftrightarrow \begin{cases} \max_{\lambda, w} & \mathbb{1}^T w - b^T \lambda \\ \text{s.t.} & \lambda \geq 0 \\ & w \leq \min \{0, A^T \lambda + c\} \end{cases} \quad (B)$$

We notice that (A) and (B) are equivalent. Thus, they lead to the same optimal value.

Recall that (A) is the dual of $\begin{cases} \min c^T x \\ \text{s.t.} & Ax \leq b \\ & x_i(1-x_i) = 0 \quad \forall i \in [1, n] \end{cases}$ and (B) is the dual of $\begin{cases} \min c^T x \\ \text{s.t.} & Ax \leq b \\ & 0 \leq x_i \leq 1 \end{cases}$.

Therefore, the lower bound obtained via Lagrangian relaxation (3) and via the LP relaxation (2) are the same.