

# The finiteness of security (Tutorial)

Renato Renner

$\varepsilon > 0$   
(Tutorial)

Renato Renner

# Rationale

$$\varepsilon > 0$$

- Security is always finite.
- It is therefore crucial to understand how to quantify it.

# Epsilon-security



## Certificate

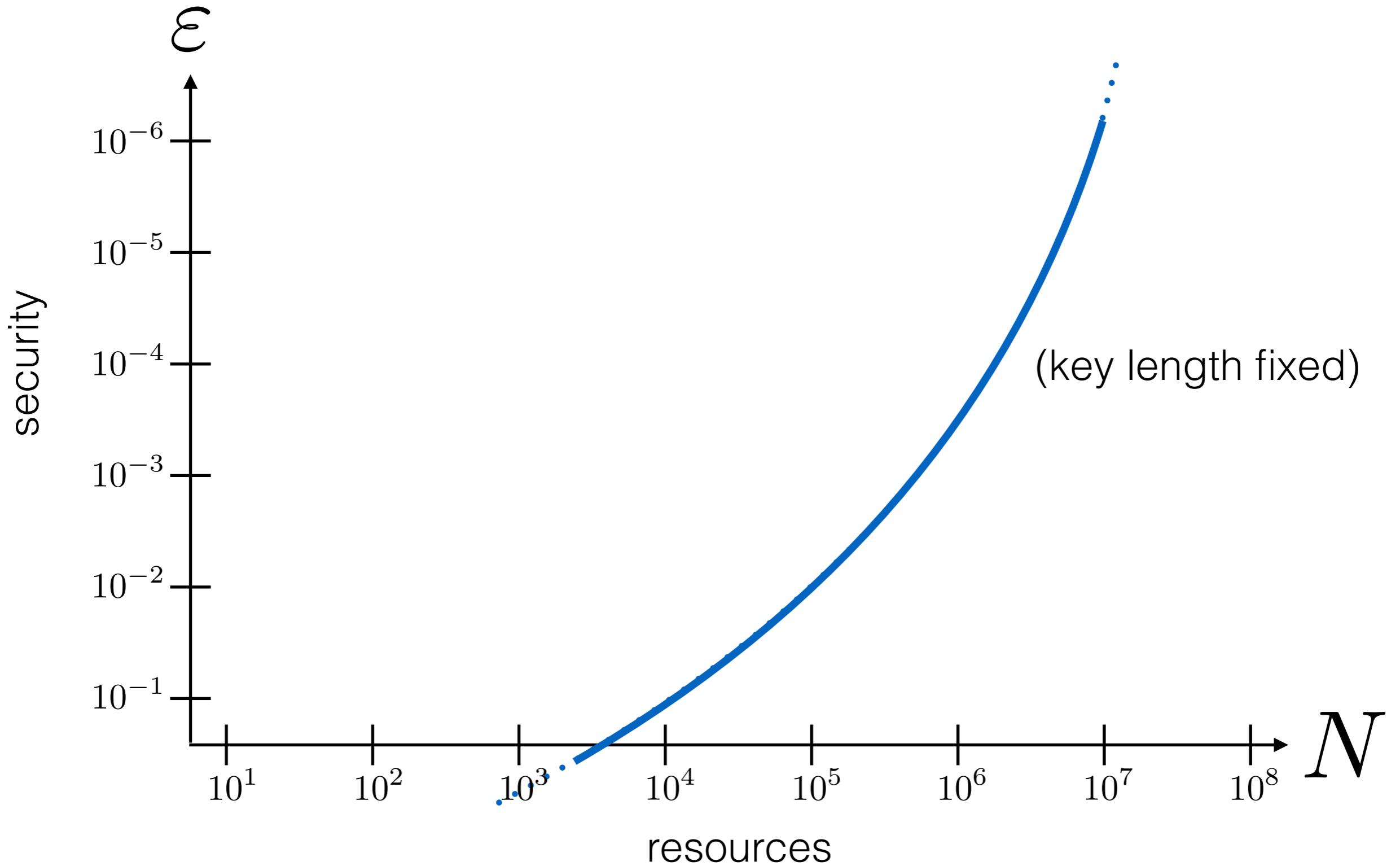
The keys generated  
by this device have  
security

$$\varepsilon = 10^{-8}$$



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# “Finite-size effects” sound rather boring ...



... but the epsilon is hotly debated



Debate at the “HotPI” conference, Hunan University, Changsha

# The debate is still ongoing ...

## Misconception in Theory of Quantum Key Distribution -Reply to Renner-

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(Dated: August 16, 2018)

It has been pointed out by Yuen that the security theory of quantum key distribution(QKD) guided by Shor-Preskill theory has serious defects, in particular their key rate theory is not correct. Theory groups of QKD tried to improve several defects. Especially, Renner employed trace distance

# ... and is basically about the epsilon

## On the Foundations of Quantum Key Distribution — Reply to Renner and Beyond\*

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August 4, 2018

### **Abstract**

*In a recent note ([arXiv:1209.2423](https://arxiv.org/abs/1209.2423)) Renner claims that the criticisms of Hirota and Yuen on the security foundation of quantum key distribution arose from a logical mistake. In this paper it is shown that Renner*

# What is the problem?

has been repeatedly given in [2-5]. Rather, Renner made a fundamental error in [7-8] which has become the standard interpretation of the trace distance criterion  $d$  widely employed in QKD. This incorrect interpretation leads to the current prevalent QKD security claim that the generated key  $K$  has a probability  $p \geq 1 - d$  of being ideal [9-11]. In actuality,  $K$  is not

# Security claim

## Certificate

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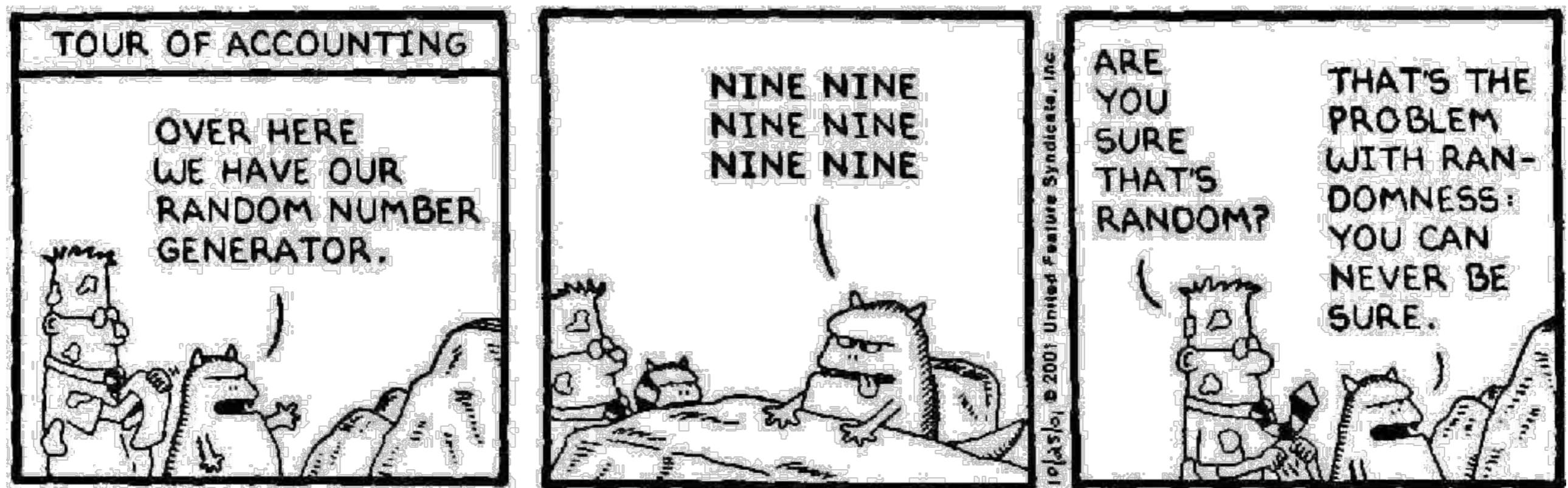
Operational meaning:  
“An adversary cannot  
gain any information  
about the secret,  
except with probability  
epsilon.”



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# Where does the epsilon come from?

From statistical fluctuations in the random choices .



© Dilbert by Scott Adams

# Risk that adversary makes correct guesses

Recall that QKD protocols involve various random choices.

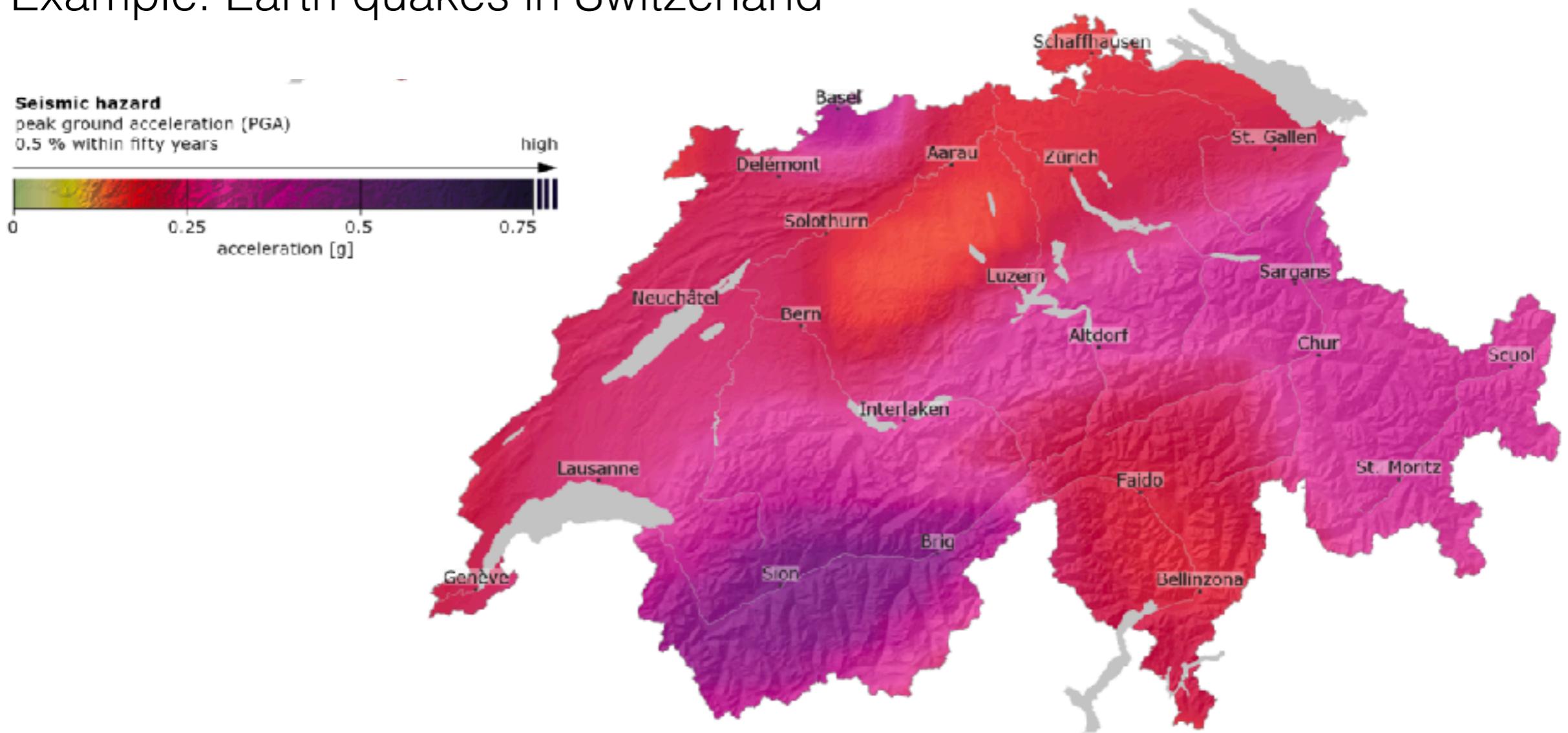
Alice's bit string.....	1	0	1	0	0	1	1	1	0	1	0	1	1	0	0
Alice's random basis.....															
Photons Alice sends.....	↓	↔	↑	↔	↔	↑	↓	↑	↔	↓	↔	↑	↓	↔	↔
Bob's random bases.....	R	D	D	D	R	R	D	R	R	D	R	R	D	D	R
Bob's rectilinear table.....	1					1				0					0
Bob's diagonal table.....		0		1						1		0			
Bob's guess.....												'Rectilinear'			
Alice's reply.....												'You win'			
Alice sends her original bit string to certify....	1	0	1	0	0	1	1	1	0	1	0	1	1	0	0
Bob's rectilinear table.....	1					1				0					0
Bob's diagonal table.....		0		1						1		0			

From Bennett and Brassard, Quantum cryptography: Public key distribution and coin tossing (1984)

Note: This risk cannot be reduced to zero.

# Risks can have different levels of severeness

Example: Earth quakes in Switzerland



Operational meaning:

The probability of experiencing this is 0.5 % within fifty years.

# Epsilon-security is “all or nothing”



## Certificate

Any key generated by this device is secure, except with probability

$$\varepsilon = 10^{-8}$$



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# Epsilon-security is common in engineering



## Certificate

A DBA does not occur,  
except with probability

$$\varepsilon = 10^{-6}$$

per year.



Note: epsilon cannot be reduced to zero.

# Epsilon-security is common in engineering



SECURITY™

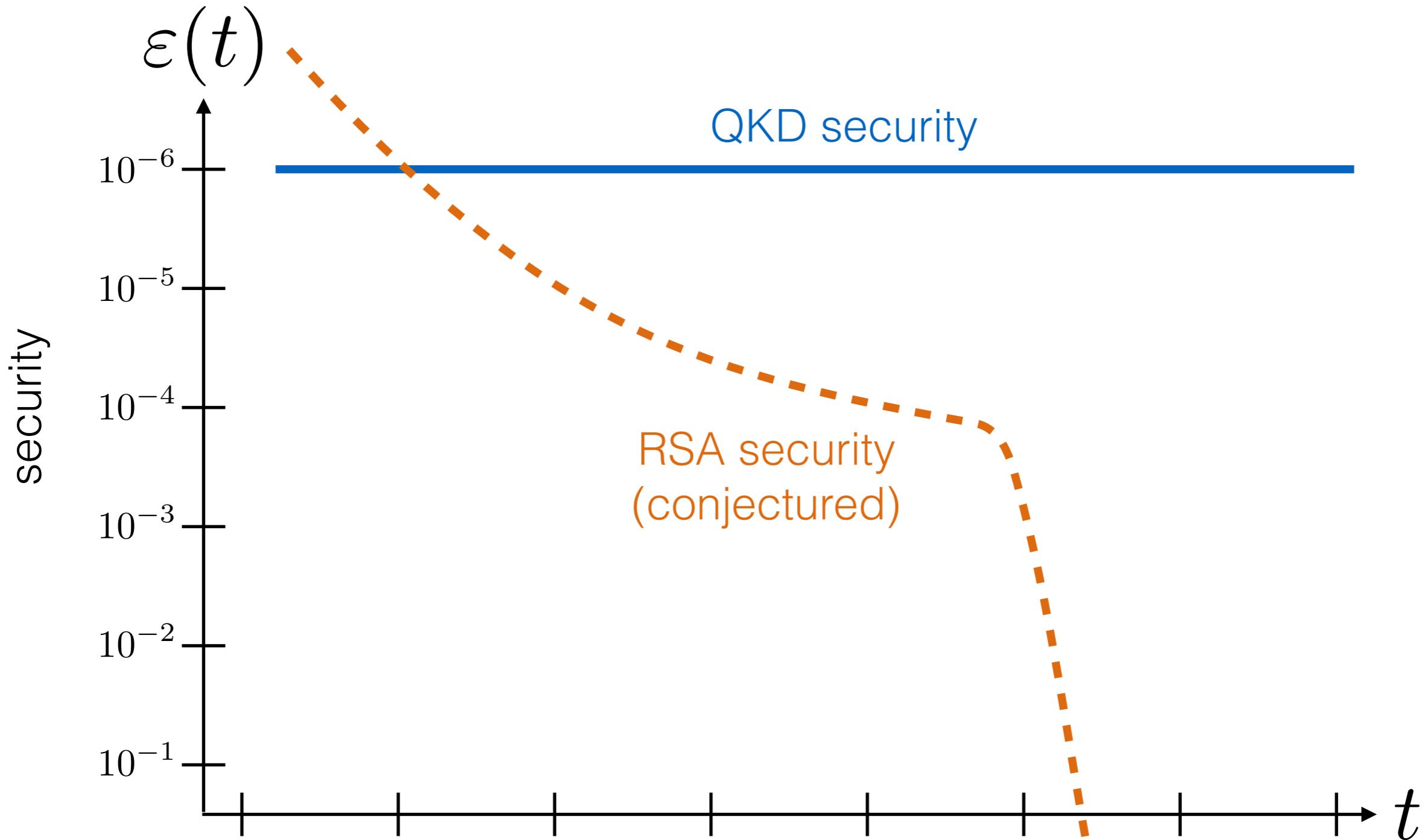
## Certificate

The keys generated by RSA remain secure for time  $t$ , except with probability

$$\varepsilon(t)$$

which is related to the probability that large numbers can be factored in time  $\text{poly}(t)$ .

# Quantum versus computational cryptography



# How is epsilon defined?

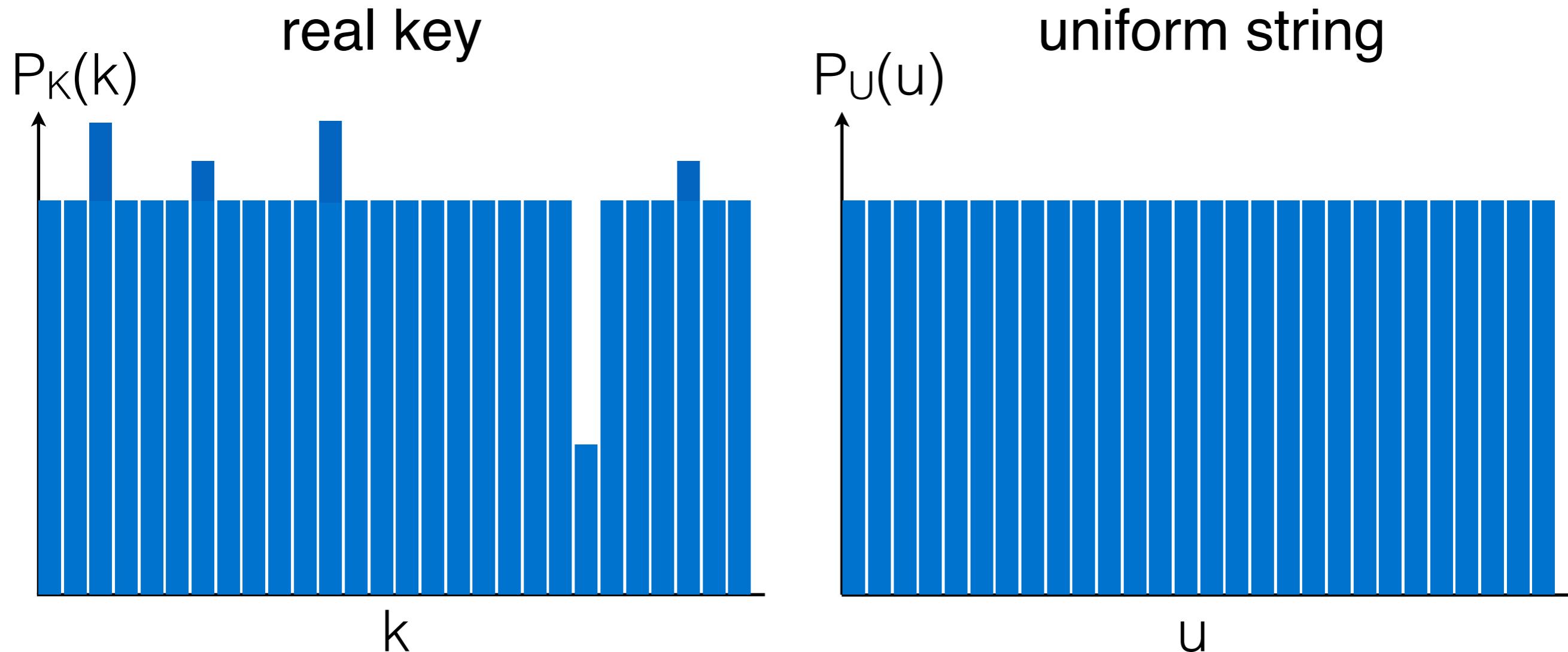
## Certificate

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 $\varepsilon = 10^{-8}$



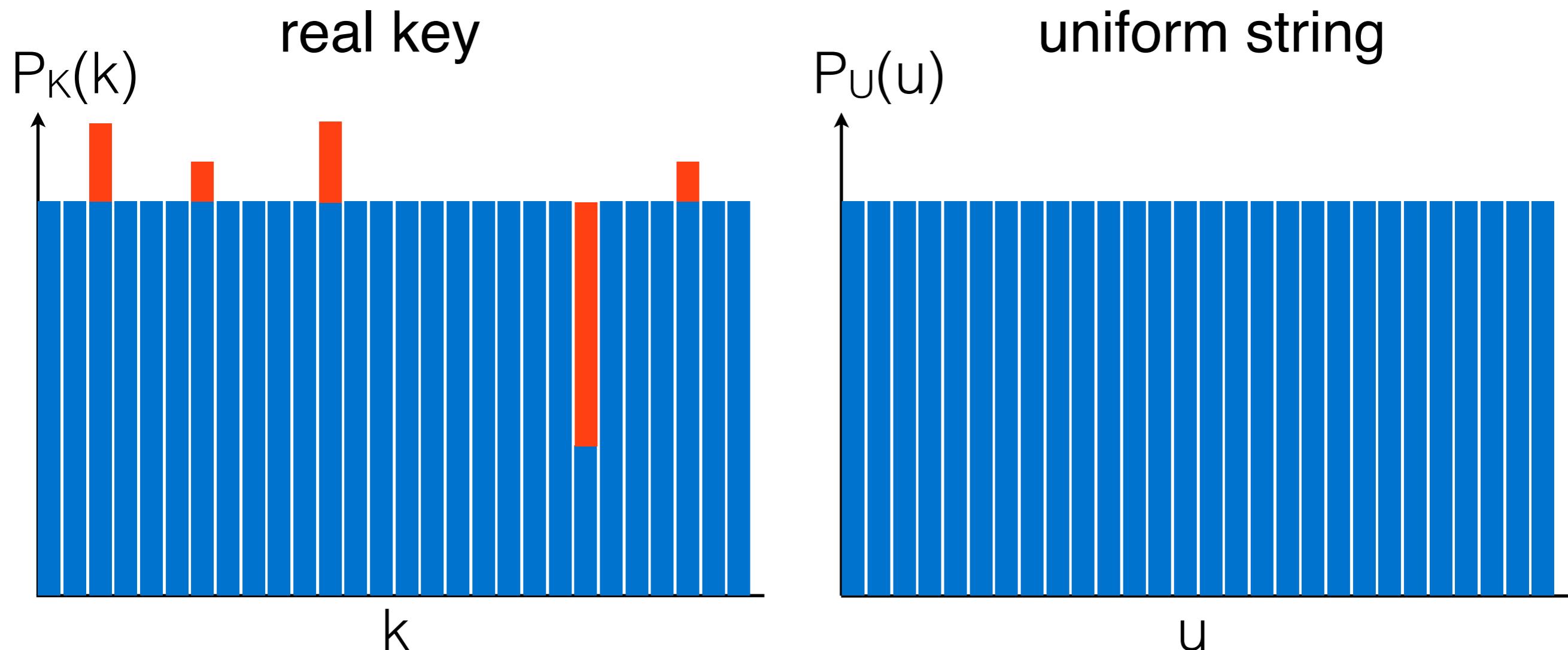
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# Technical definition (without Eve)



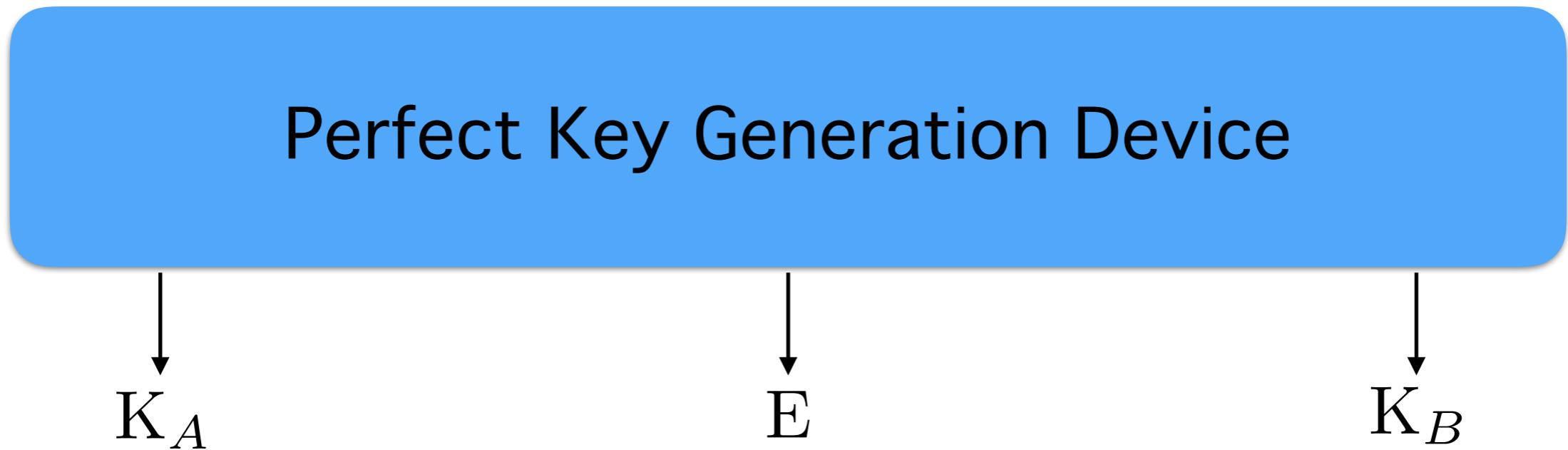
$\epsilon$  Trace distance between probability distribution of real key  $K$  and uniformly distributed string  $U$ .

# Technical definition (without Eve)



$\epsilon$  corresponds to weight of red area.

# Real world / ideal world paradigm

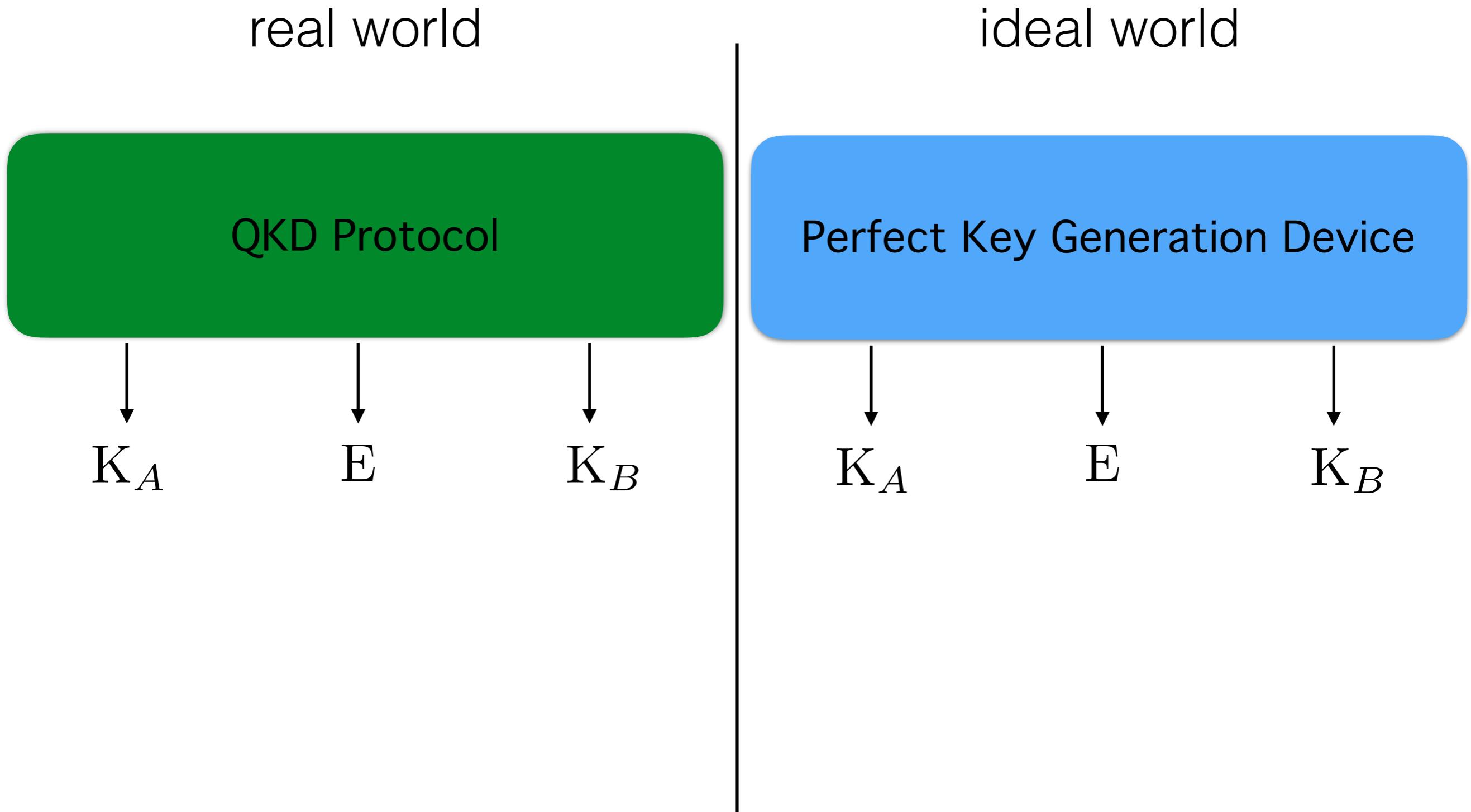


## Requirements

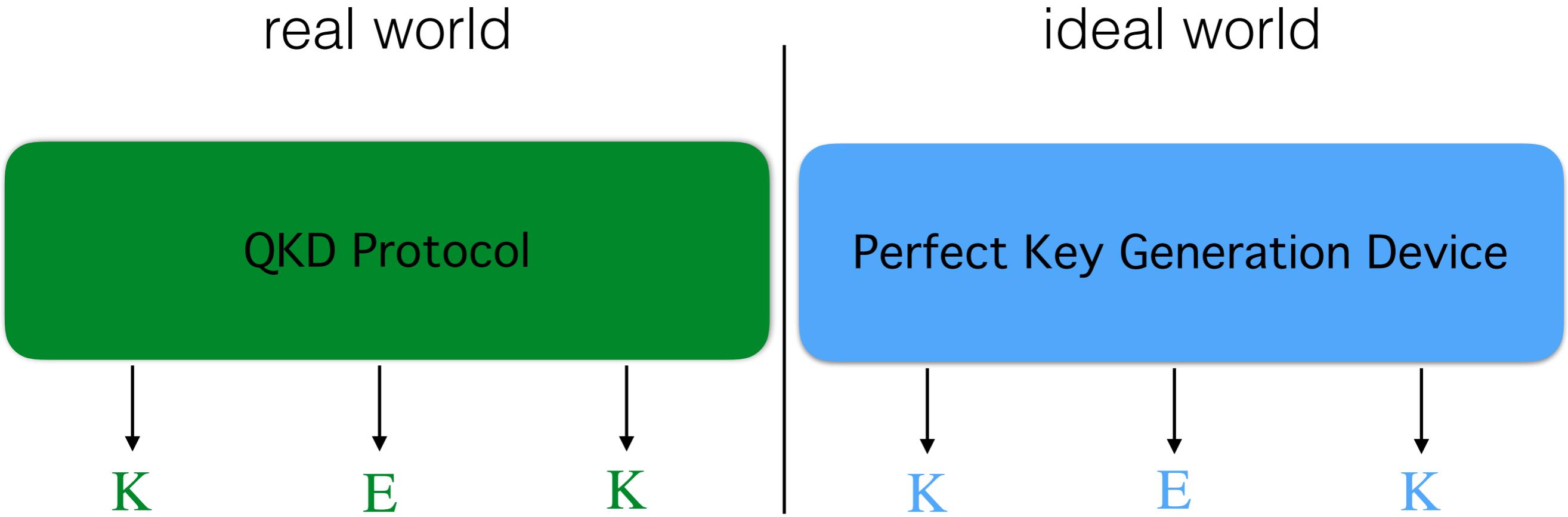
*Correctness:*  $K_A = K_B = K$

*Secrecy:*  $K$  uniformly distributed and independent of  $E$

# Real world / ideal world paradigm

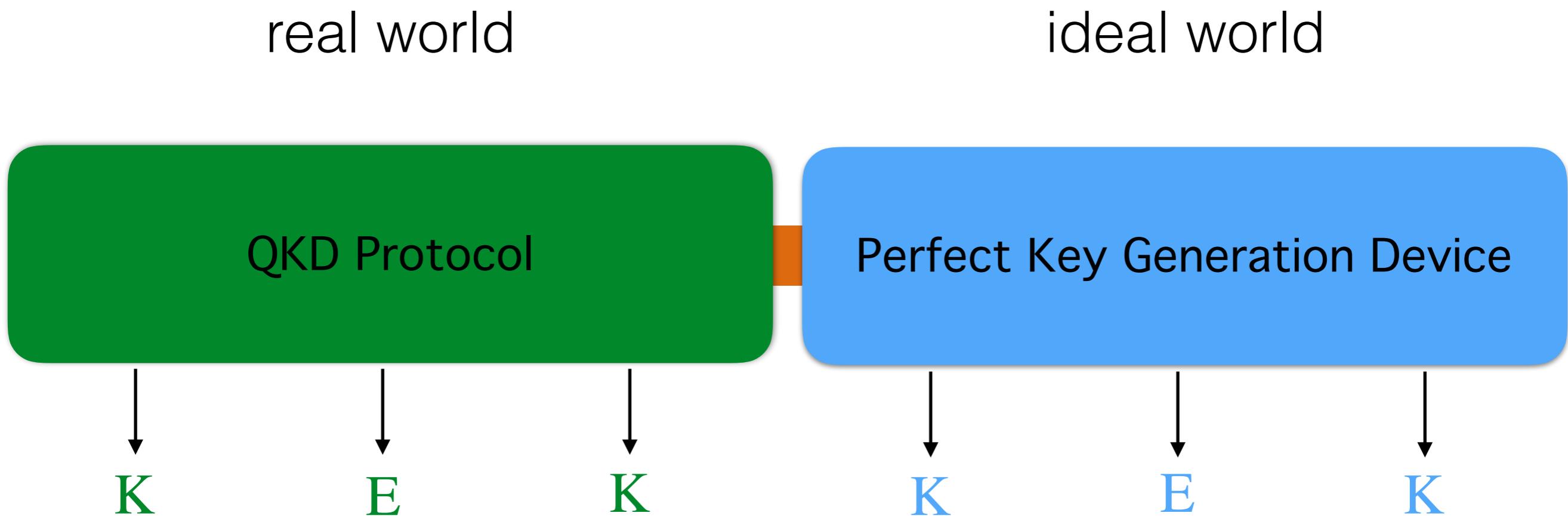


# Real world / ideal world paradigm



Definition: The Protocol is  $\varepsilon$ -secure if  $P_{KE}$  and  $P'_{KE}$  have trace distance  $\varepsilon$  from each other.

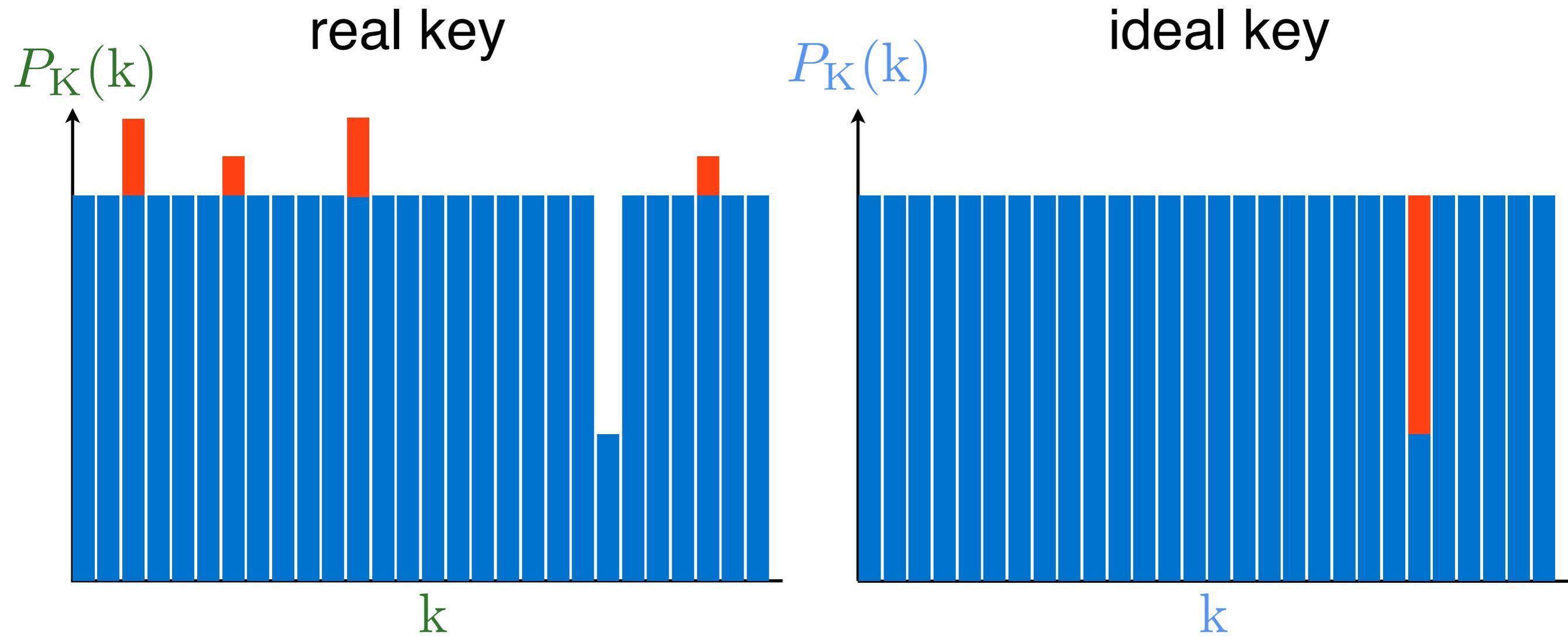
# What does epsilon mean operationally?



**Theorem:** If the Protocol is  $\varepsilon$ -secure then there exists a joint distribution such that, with probability at least  $1 - \varepsilon$ ,

$$K = K \text{ and } E = E$$

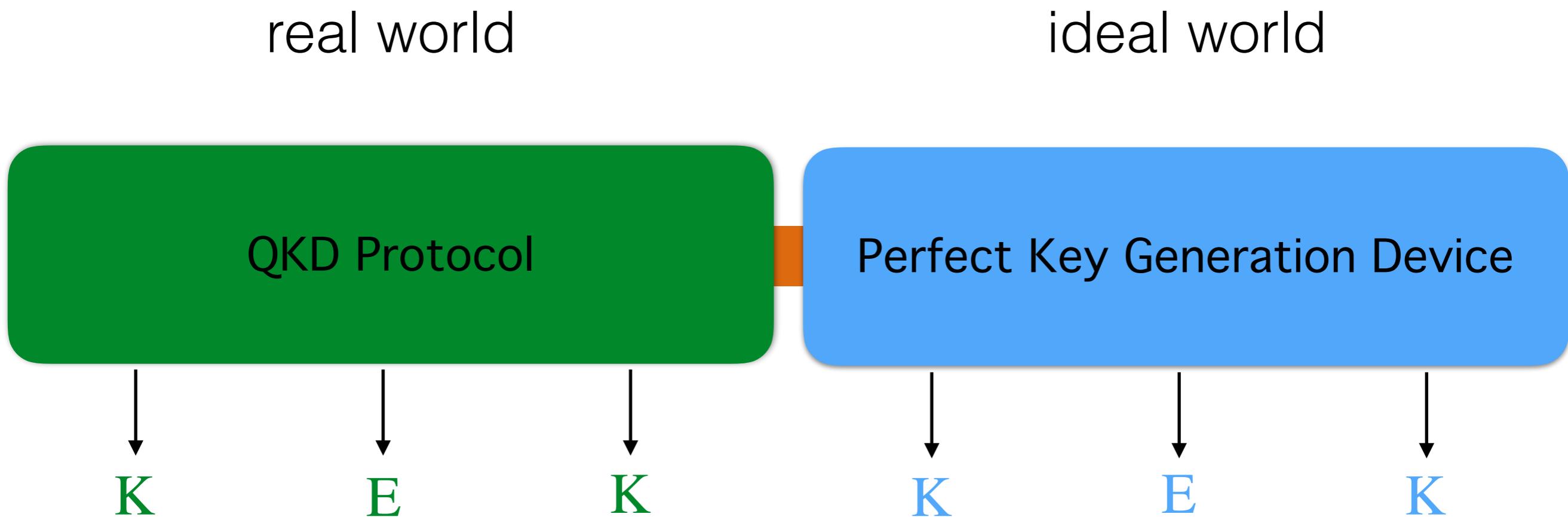
# Proof idea



Recall: Red area corresponds to trace distance  $\varepsilon$ .

Idea: Define  $K = K$ , except when red.

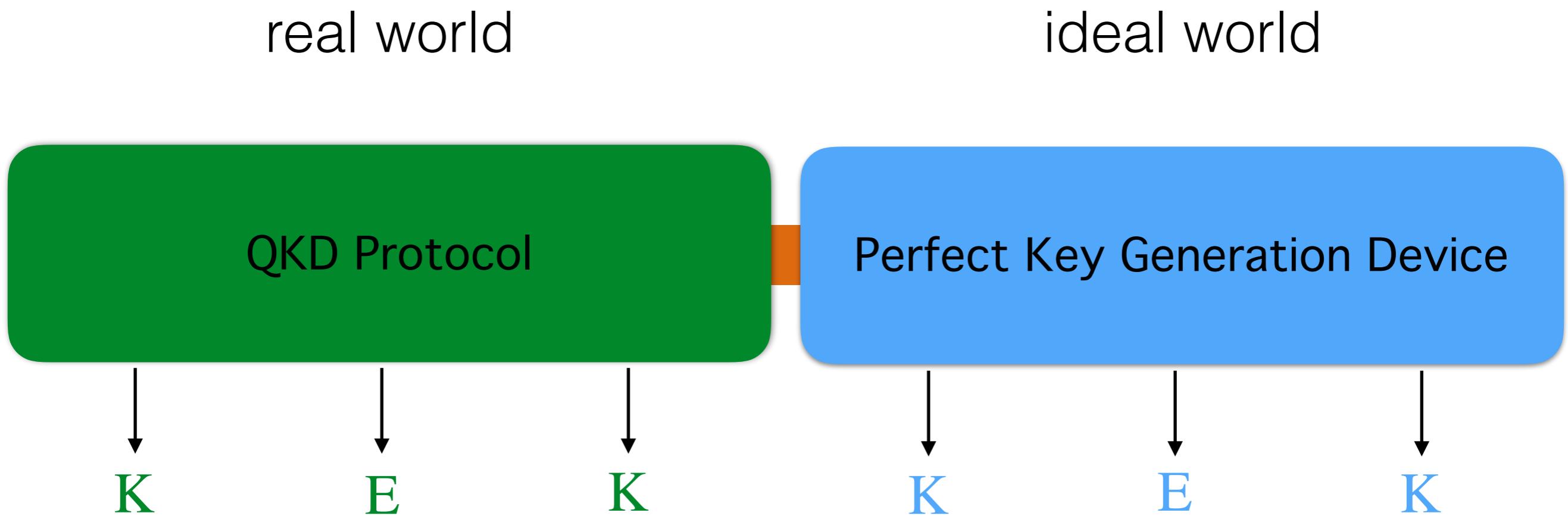
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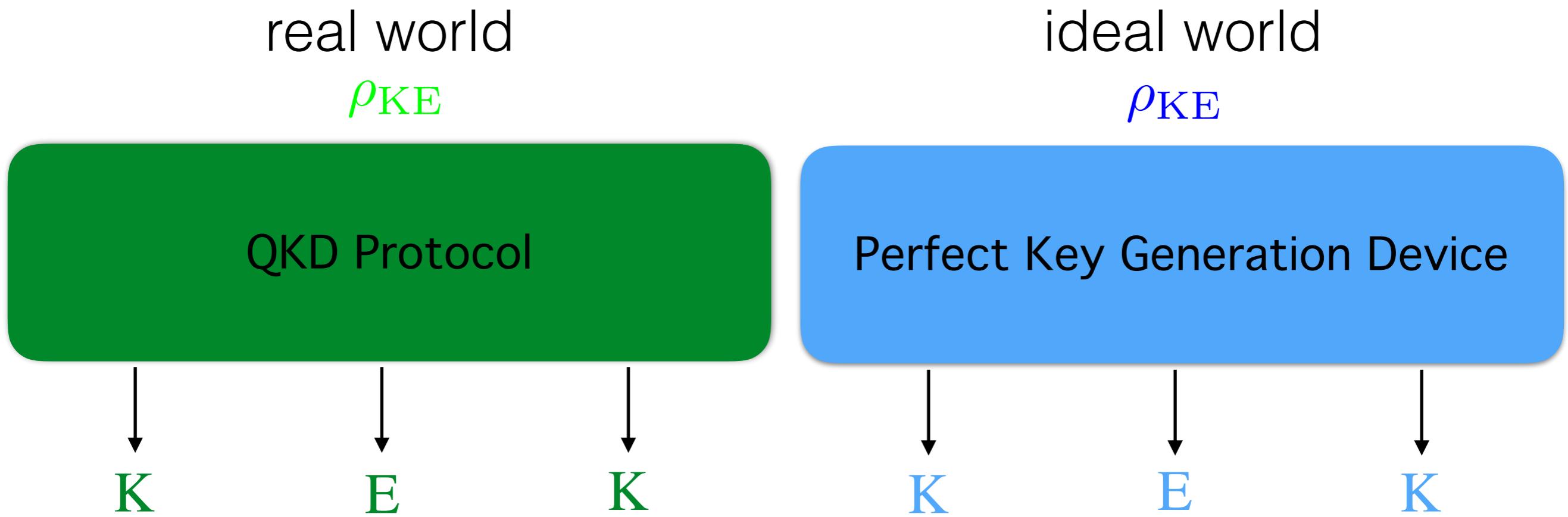
$$K = K \text{ and } E = E$$

# What does epsilon mean operationally?



Interpretation: If the Protocol is  $\epsilon$ -secure then the probability that it behaves differently from a perfect device is at most  $\epsilon$ .

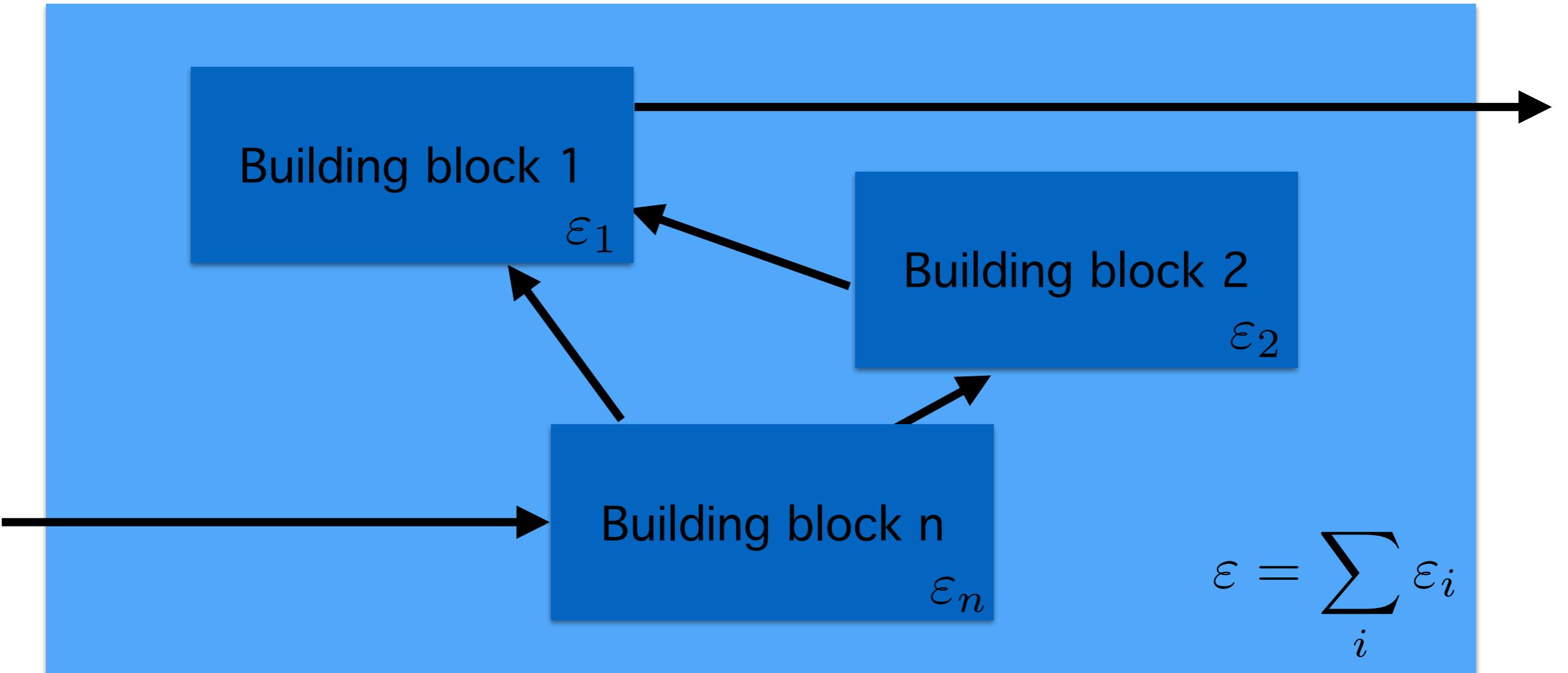
# Quantum version



**Theorem:** If the Protocol is  $\varepsilon$ -secure then there exists a state  $\rho_{KE}$  and events  $\Omega$  and  $\Omega'$  with probability  $1 - \varepsilon$  s.t.

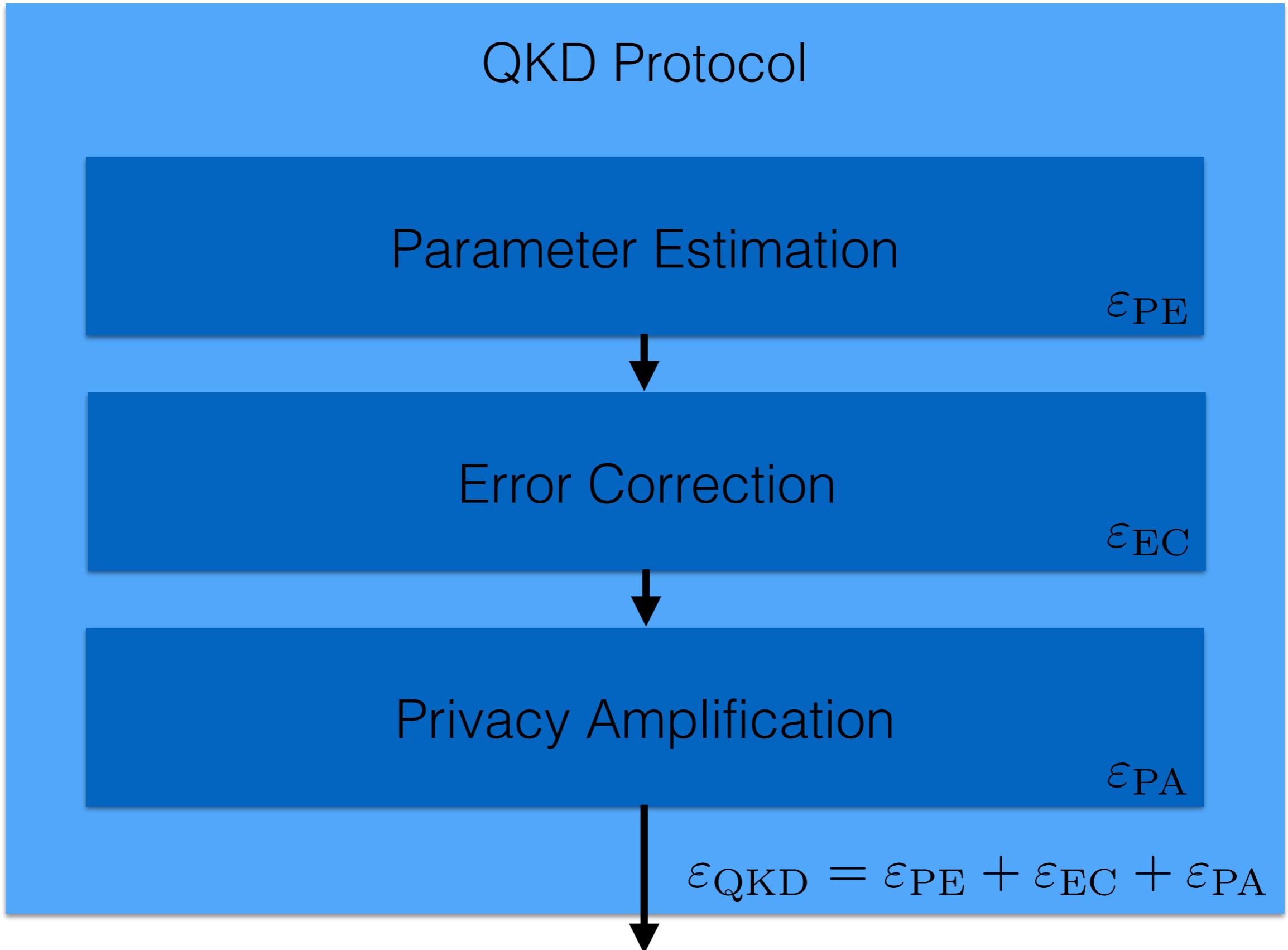
$$\rho_{KE}|\Omega = \rho_{KE} \quad \text{and} \quad \rho_{KE}|\Omega' = \rho_{KE}$$

# Composability: Summation rule



Epsilons add up (because failure probabilities add up).

# Contributions to epsilon



# Example: Privacy Amplification

raw key with min-entropy at least  $h$



Privacy Amplification



final  $n$ -bit key

Theorem: Given a raw key with min-entropy at least equal to  $h$ , the  $n$ -bit key is uniform, except with probability

$$\varepsilon_{\text{PA}} = 2^{-\frac{1}{2}(h-n)}$$

# How to choose epsilon?



failure probability  
per key generated

$$\varepsilon = 10^{-12}$$

(recommended value)

number of keys  
generated

$$N = 10^9$$

upper bound on  
failure

$$p_{\text{fail}} = 1/1000$$

$$\varepsilon = 2 \cdot 10^{-5}$$

$$N = 50$$

$$p_{\text{fail}} = 1/1000$$

failure probability  
per year

number of years in  
operation

upper bound on  
failure

# Summary

$$\varepsilon > 0$$

- Security is always finite.
- It is therefore crucial to understand how to quantify it.

# Questions?

$$\epsilon > 0$$