

Quantum de Finetti theorems under local measurements

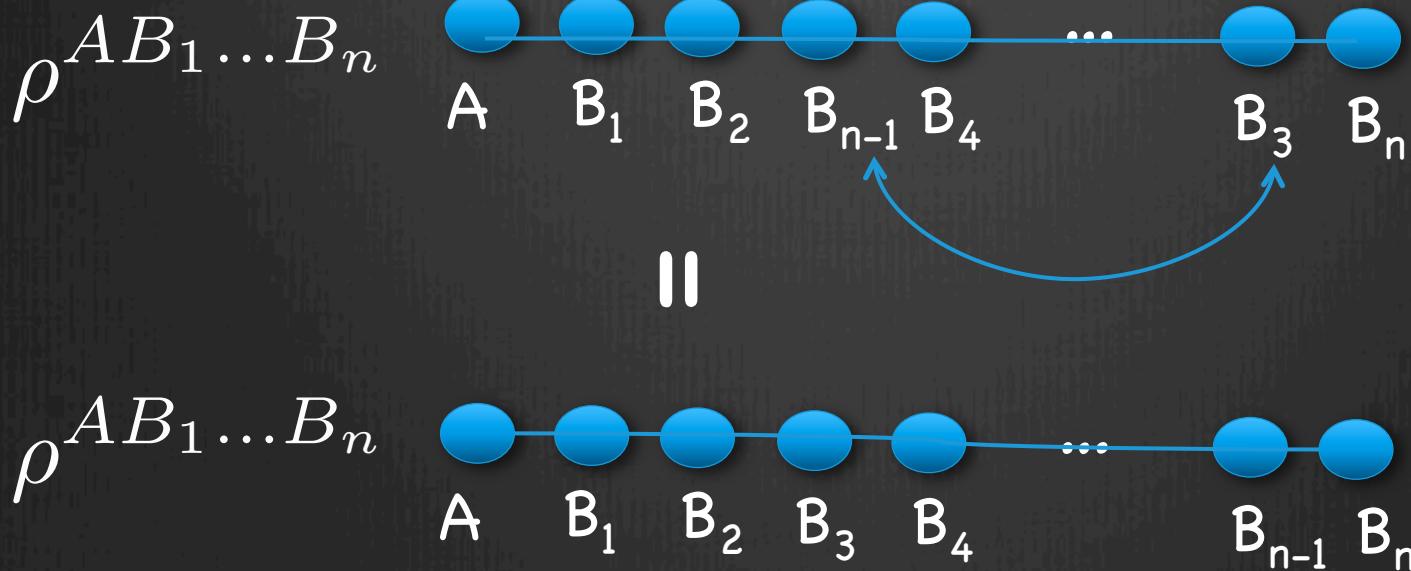
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based on [arXiv:1210.6367](https://arxiv.org/abs/1210.6367)
joint work with Fernando Brandão (ETH)

Symmetric States

$\rho^{AB_1\dots B_n}$ is permutation symmetric in the B subsystems if for every permutation π ,

$$\rho^{AB_1\dots B_n} = \rho^{AB_{\pi(1)}\dots B_{\pi(n)}}$$



Quantum de Finetti Theorem

Theorem [Christandl, Koenig, Mitchison, Renner '06]

Given a state $\rho^{AB_1\dots B_n}$ symmetric under exchange of $B_1\dots B_n$,
there exists μ such that

$$\left\| \rho^{AB_1\dots B_k} - \int \mu(d\sigma) \rho_\sigma \otimes \sigma^{\otimes k} \right\|_1 \leq \frac{d^2 k}{n}$$

builds on work by [Størmer '69], [Hudson, Moody '76], [Raggio, Werner '89]
[Caves, Fuchs, Sachs '01], [Koenig, Renner '05]

Proof idea:

Perform an informationally complete measurement of $n-k$ B systems.

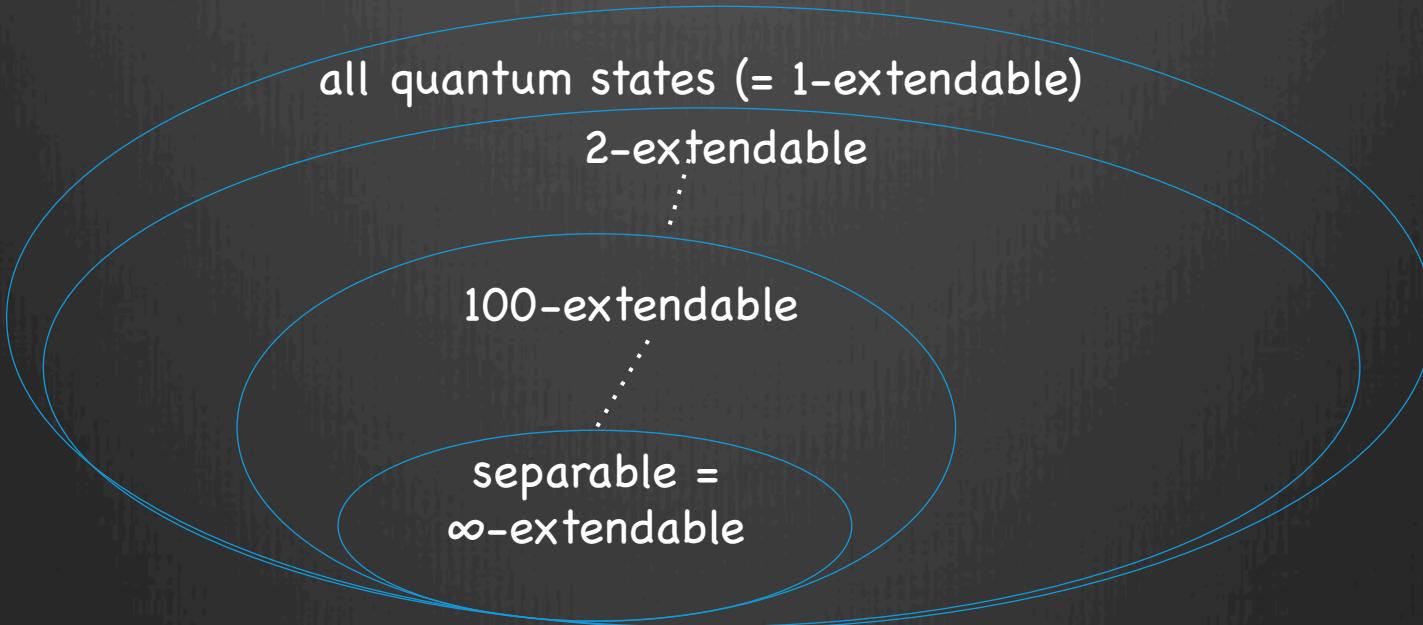
Applications:

information theory: tomography, QKD, hypothesis testing

algorithms: approximating separable states, mean-field theory

Quantum de Finetti Theorem as Monogamy of Entanglement

Definition: ρ^{AB} is **n-extendable** if there exists an extension $\rho^{AB_1 \dots B_n}$ with $\rho^{AB} = \rho^{AB_i}$ for each i.



Algorithms: Can search/optimize over n-extendable states in time $d^{O(n)}$.

Question: How close are n-extendable states to separable states?

Quantum de Finetti theorem

Theorem [Christandl, Koenig, Mitchison, Renner '06]

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there exists μ such that

$$\left\| \rho^{AB_1\dots B_k} - \int \mu(d\sigma) \rho_\sigma \otimes \sigma^{\otimes k} \right\|_1 \leq \frac{d^2 k}{n}$$

Difficulty:

1. Parameters are, in many cases, **too weak**.
2. They are also essentially **tight**.

Way forward:

1. Change definitions (of error or i.i.d.)
2. Obtain better scaling

relaxed/improved versions

Two examples known:

1. Exponential de Finetti Theorem: [Renner '07]
error term $\exp(-\Omega(n-k))$.

Target state convex combination of “almost i.i.d.” states.

2. measure error in 1-LOCC norm [Brandão, Christandl, Yard '10]
For error ε and $k=1$, requires $n \sim \varepsilon^{-2} \log|A|$.

This talk
improved de Finetti theorems for local
measurements

main idea use information theory

$$\log |A| \geq$$

$$I(A:B_1 \dots B_n) = I(A:B_1) + I(A:B_2|B_1) + \dots + I(A:B_n|B_1 \dots B_{n-1})$$

repeatedly uses chain rule: $I(A:BC) = I(A:B) + I(A:C|B)$

$$\rightarrow I(A:B_t|B_1 \dots B_{t-1}) \leq \log(|A|)/n \text{ for some } t \leq n.$$

If $B_1 \dots B_n$ were classical, then we would have

$$\rho^{AB} = \rho^{AB_t} = \sum_i \pi_i \rho_i^{AB}$$

Question:
How to make $B_1 \dots B_n$ classical?

distribution
on $B_1 \dots B_{t-1}$

≈ separable
≈ product state
(cf. Pinsker ineq.)

Answer: measure!

Fix a measurement $M: B \rightarrow Y$.

$I(A:B_t | B_1 \dots B_{t-1}) \leq \varepsilon$ for the measured state $(\text{id} \otimes M^{\otimes n})(\rho)$.

Then

- ρ^{AB} is hard to distinguish from $\sigma \in \text{Sep}$ if we first apply $(\text{id} \otimes M)$
- $\|(\text{id} \otimes M)(\rho - \sigma)\| \leq \text{small}$ for some $\sigma \in \text{Sep}$.

Theorem

Given a state $\rho^{AB_1 \dots B_n}$ symmetric under exchange of $B_1 \dots B_n$,
and $\{\Lambda_i\}$ a collection of operations from $A \rightarrow X$,

$$\min_{\sigma \in \text{Sep}} \max_M \mathbb{E} \left\| (\Lambda_i^A \otimes M^B)(\rho^{AB} - \sigma^{AB}) \right\|_1 \leq \sqrt{\frac{2 \ln |X|}{n}}$$

Cor: setting $\Lambda = \text{id}$ recovers [Brandão, Christandl, Yard '10] 1-LOCC result.

advantages/extensions

Theorem

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1. Simpler proof and better constants
2. Bound depends on $|X|$ instead of $|A|$ (can be ∞ dim)
3. Applies to general non-signalling distributions
4. There is a multipartite version (multiply error by k)
5. Efficient “rounding” (i.e. σ is explicit)
6. Symmetry isn’t required (see Fernando’s talk on Thursday)

applications

- nonlocal games
Adding symmetric provers “immunizes” against entanglement / non-signalling boxes. (Caveat: needs uncorrelated questions.)
Conjectured improvement would yield NP-hardness for 4 players.
- BellQMA(poly) = QMA
Proves Chen-Drucker $\text{SAT} \in \text{BellQMA}_{\log(n)}(\sqrt{n})$ protocol is optimal.
- pretty good tomography [Aaronson '06]
on permutation-symmetric states (instead of product states)
- convergence of Lasserre hierarchy for polynomial optimization
see also 1205.4484 for connections to small-set expansion

open questions

- Is $\text{QMA}(2) = \text{QMA}$? Is $SAT \in \text{QMA}_{\sqrt{n}}(2)_{1,1/2}$ optimal?
(Would follow from replacing 1-LOCC with SEP-YES.)
- Can we reorder our quantifiers to obtain

$$\min_{\sigma \in \text{Sep}} \mathbb{E} \max_i \max_M \left\| (\Lambda_i^A \otimes M^B)(\rho^{AB} - \sigma^{AB}) \right\|_1 \leq \sqrt{\frac{2 \ln |X|}{n}} ?$$

(no-signalling analogue is FALSE assuming $P \neq NP$)

- The usual de Finetti questions:
 - better counter-examples
 - how much does it help to add PPT constraints?