

On a new Quantum Rényi Divergence

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[J. Math. Phys. 54, 122203 (2013)]

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[arXiv:1306.1586]

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[arXiv:1309.3228]

Outline

- ⦿ Motivation and Contributions
- ⦿ Definition
- ⦿ Elementary Properties
- ⦿ Operational Interpretation (hypothesis testing)
- ⦿ Application to Strong Converse (classical data over entanglement-breaking channels)
- ⦿ Conditional Quantum Rényi Entropy

Motivation: QIP 2014 Talks and Posters using this Quantum Rényi Divergence

12:20h The second laws of quantum thermodynamics

Fernando Brandao, Michal Horodecki, Jonathan Oppenheim , Nelly Ng, Stephanie Wehner

09:30h Strong converses for quantum channel capacities

Andreas Winter

11:50h Entanglement sampling and applications

Frédéric Dupuis, Omar Fawzi , Stephanie Wehner

11:20h Robust protocols for securely expanding randomness and distributing keys using untrusted quantum devices

Carl Miller, Yaoyun Shi

Salman Beigi. Sandwiched Renyi Divergence Satisfies Data Processing Inequality

Koenraad Audenaert, [Nilanjana Datta](#) and Felix Leditzky. A limit of the quantum Renyi divergence

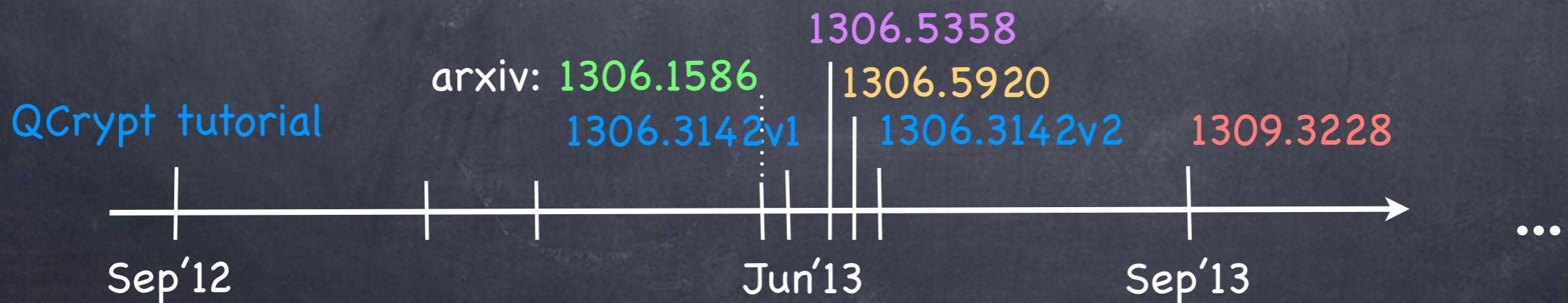
Milan Mosonyi. Renyi divergences and the classical capacity of finite compound channels

[Manish Gupta](#) and [Mark Wilde](#). Multiplicativity of completely bounded p-norms implies a strong converse for entanglement-assisted capacity

Patrick Coles and Fabian Furrer. Entropic Error-Disturbance Relations

Contributions

- ⦿ New Rényi conditional entropy proposed, specializes to known quantities (Team Müller-Lennert)
- ⦿ Independent discovery of the divergence as a proof tool for strong converse (Team Wilde)
- ⦿ Properties and conjectures released (Team Blue)
- ⦿ All conjectures resolved (Frank&Lieb, Beigi, Team Blue)
- ⦿ Operational interpretation in hypothesis testing discovered (Team Mosonyi)



Classical Rényi Divergence

- Given two probability distributions P and Q :

$$D_\alpha(P||Q) := \frac{1}{\alpha - 1} \log \left(\sum_x P(x)^\alpha Q(x)^{1-\alpha} \right)$$

- Operational significance in information theory, for example in the study of error exponents and cutoff rates.
- Versatile tool in proofs, for example to derive one-shot bounds.

Quantum Rényi Relative Entropy

- Petz investigated quantum generalizations of Csiszár's f -divergences.
- For two quantum states ρ and σ :

$$D_\alpha(\rho\|\sigma) := \frac{1}{\alpha - 1} \log \text{tr} (\rho^\alpha \sigma^{1-\alpha})$$

- Desirable properties in the range $\alpha \in [0, 2]$ due to the operator concavity/convexity of $f : t \mapsto t^\alpha$
- Operational significance in the direct part of quantum hypothesis testing

Quantum Rényi Divergence / Sandwiched Rényi Divergence

- For $\alpha \in (0, 1) \cup (1, \infty)$ and $\rho, \sigma \geq 0, \rho \neq 0$,

$$\tilde{D}_\alpha(\rho\|\sigma) := \frac{1}{\alpha - 1} \log \frac{\text{tr} \left(\left(\sigma^{\frac{1-\alpha}{2\alpha}} \rho \sigma^{\frac{1-\alpha}{2\alpha}} \right)^\alpha \right)}{\text{tr}(\rho)}$$

- For $\alpha > 1$ and if σ is not invertible, set $\sigma' = \sigma + \xi \mathbb{I}$ and take the limit $\xi \rightarrow 0$.

- generalized inverse if $\sigma \gg \rho$
- $+\infty$ otherwise



Quantum Rényi Divergence / Sandwiched Rényi Divergence

- For the following: $\text{tr}(\rho) = 1$
- Divergence is related to the Schatten norm

$$\tilde{D}_\alpha(\rho\|\sigma) = \frac{\alpha}{\alpha - 1} \log \left\| \sigma^{\frac{1-\alpha}{2\alpha}} \rho \sigma^{\frac{1-\alpha}{2\alpha}} \right\|_\alpha$$

$$\frac{\alpha}{1 - \alpha} \log \left\| \text{submarine sandwich} \right\|_\alpha$$



Two-Parameter Family

- Both definitions can be seen as special cases of a two-parameter family of divergences:
(Jaksic et al. / Audenaert&Datta)

$$\bar{D}_{\alpha,z}(\rho\|\sigma) := \frac{1}{\alpha - 1} \log \text{tr} \left(\left(\sigma^{\frac{1-\alpha}{2z}} \rho^{\frac{\alpha}{z}} \sigma^{\frac{1-\alpha}{2z}} \right)^z \right)$$

- Clearly, $\bar{D}_{\alpha,1} \equiv D_\alpha$ and $\bar{D}_{\alpha,\alpha} \equiv \tilde{D}_\alpha$.

Limits and Special Cases of

$$\tilde{D}_\alpha(\rho\|\sigma) := \frac{1}{\alpha - 1} \log \text{tr} \left(\left(\sigma^{\frac{1-\alpha}{2\alpha}} \rho \sigma^{\frac{1-\alpha}{2\alpha}} \right)^\alpha \right)$$

⦿ **Quantum Relative Entropy, $\alpha \rightarrow 1$:**

$$D(\rho\|\sigma) = \tilde{D}_1(\rho\|\sigma) = \text{tr} (\rho(\log \rho - \log \sigma))$$

⦿ **Datta's Max Relative Entropy, $\alpha \rightarrow \infty$:**

$$D_{\max}(\rho\|\sigma) = \tilde{D}_\infty(\rho\|\sigma) = \inf \left\{ \lambda \in \mathbb{R} \mid \rho \leq \exp(\lambda)\sigma \right\}$$

⦿ **Fidelity:** $\tilde{D}_{\frac{1}{2}}(\rho\|\sigma) = 2 \log F(\rho, \sigma)$

⦿ **Collision:** $\tilde{D}_2(\rho\|\sigma) = \log \text{tr} \left(\rho \sigma^{-\frac{1}{2}} \rho \sigma^{-\frac{1}{2}} \right)$

Properties of

$$\tilde{D}_\alpha(\rho\|\sigma) := \frac{1}{\alpha - 1} \log \text{tr} \left(\left(\sigma^{\frac{1-\alpha}{2\alpha}} \rho \sigma^{\frac{1-\alpha}{2\alpha}} \right)^\alpha \right)$$

- Continuous in ρ and σ
- Additivity: $\tilde{D}_\alpha(\rho \otimes \tau \| \sigma \otimes \omega) = \tilde{D}_\alpha(\rho\|\sigma) + \tilde{D}_\alpha(\tau\|\omega)$
- Positive Definite: For two quantum states,
 $\tilde{D}_\alpha(\rho\|\sigma) \geq 0$ with equality iff* $\rho = \sigma$.
- Monotonically non-decreasing in α (also: Beigi)
- Scaling: $\tilde{D}_\alpha(a\rho\|b\sigma) = \tilde{D}_\alpha(\rho\|\sigma) + \log \frac{a}{b}$
- *Domination: $\sigma' \geq \sigma$ implies $\tilde{D}_\alpha(\rho\|\sigma') \leq \tilde{D}_\alpha(\rho\|\sigma)$

Data-Processing Inequality

- For any completely positive trace-preserving map \mathcal{E} and $\alpha \geq 1/2$, we have

$$\tilde{D}_\alpha(\rho\|\sigma) \geq \tilde{D}_\alpha(\mathcal{E}(\rho)\|\mathcal{E}(\sigma))$$



- Müller-Lennert et al. / Wilde et al. : $\alpha \in [1, 2]$
- Frank+Lieb: $\alpha \geq 1/2$
- Beigi: $\alpha \geq 1$
- (Mosonyi+Ogawa: $\alpha \geq 1$)

Asymptotic Achievability ($\alpha > 1$)

- For any sequence of measurement maps \mathcal{M}_n (quantum-to-classical channels), we have

$$\tilde{D}_\alpha(\rho\|\sigma) \geq \frac{1}{n} \tilde{D}_\alpha(\mathcal{M}_n(\rho^{\otimes n})\|\mathcal{M}_n(\sigma^{\otimes n}))$$

- There exists a sequence \mathcal{M}_n^* such that

$$\tilde{D}_\alpha(\rho\|\sigma) = \lim_{n\rightarrow\infty} \frac{1}{n} \tilde{D}_\alpha(\mathcal{M}_n^*(\rho^{\otimes n})\|\mathcal{M}_n^*(\sigma^{\otimes n}))$$

- This allows to lift many properties from the classical domain.

Hypothesis Testing

- State discrimination using POVM $\{T, \mathbb{I} - T\}$

$$\alpha_n(T) := \text{tr} (\rho^{\otimes n} (\mathbb{I} - T))$$
$$\beta_n(T) := \text{tr} (\sigma^{\otimes n} T)$$

- Critical rate (quantum Stein's Lemma):

$$\lim_{n \rightarrow \infty} \max \left\{ -\frac{1}{n} \log \beta_n(T) \mid T : \alpha_n(T) \leq \varepsilon \right\} = D(\rho || \sigma)$$

- Error exponents: what happens to $\alpha_n(T)$ if

$$-\frac{1}{n} \log \beta_n(T) \neq D(\rho || \sigma) ?$$

Quantum Hoeffding Bound

- We are interested in the quantity

$$\alpha_{n,r} := \min \left\{ \alpha_n(T) \mid T : \beta_n(T) \leq \exp(-nr) \right\}$$

- Rate below critical rate (Hayashi'07/Nagaoka'06):

$$\lim_{n \rightarrow \infty} -\frac{1}{n} \log \alpha_{n,r} = \sup_{0 < \alpha < 1} \frac{\alpha - 1}{\alpha} \left(r - D_\alpha(\rho \|\sigma) \right)$$

if $r < D(\rho \|\sigma)$

- Yields operational interpretation of “old” Rényi relative entropy

Strong Converse Regime

- If $r > D(\rho\|\sigma)$, we expect $1 - \alpha_{n,r}$ to drop exponentially in n .

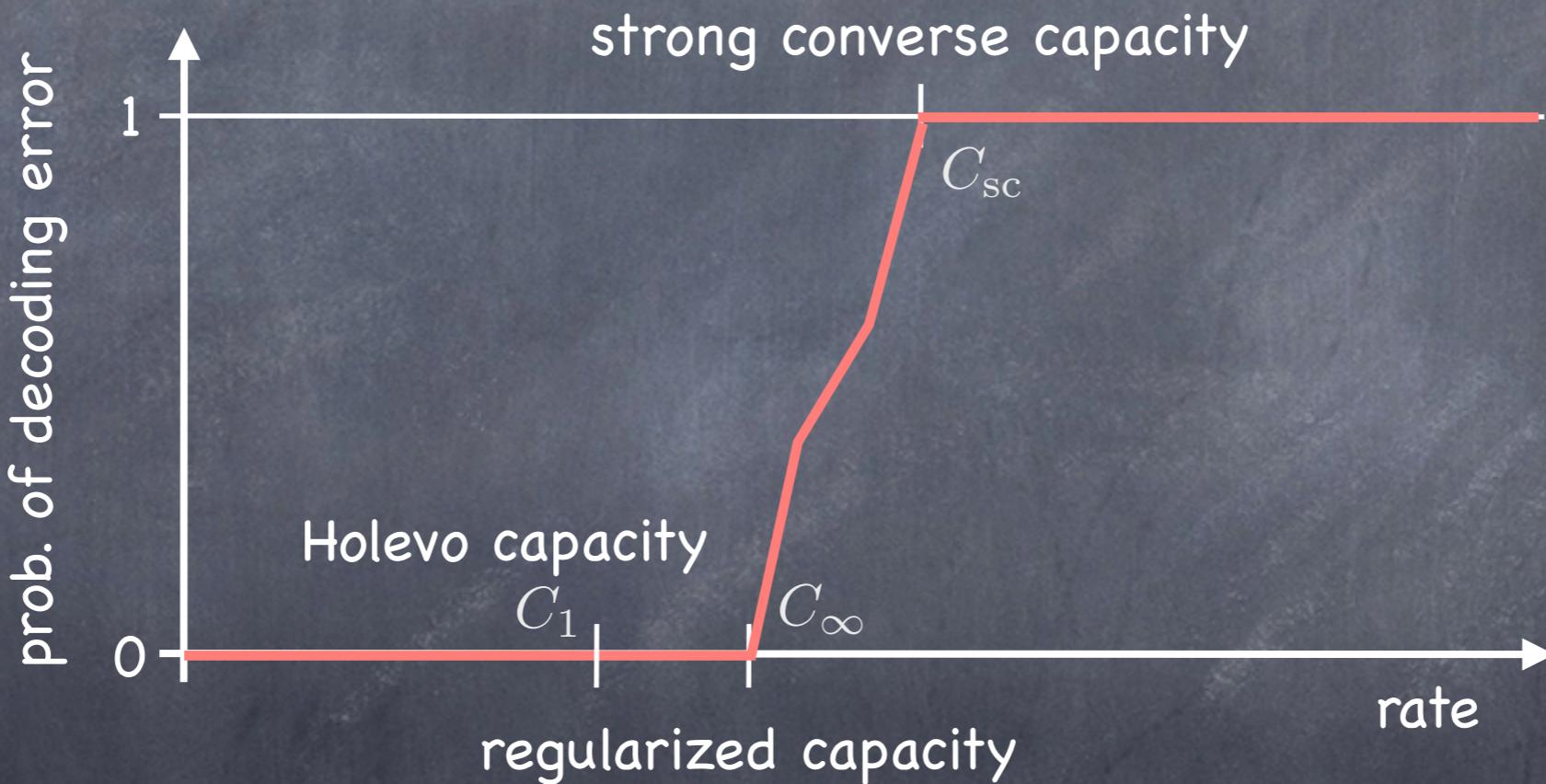
- We show that

$$\lim_{n \rightarrow \infty} -\frac{1}{n} \log(1 - \alpha_{n,r}) = \sup_{\alpha > 1} \frac{\alpha - 1}{\alpha} \left(r - \tilde{D}_\alpha(\rho\|\sigma) \right)$$

- Yields an operational interpretation of the “new” Rényi divergence
- More direct interpretation (without optimization) using cut-off rates

Strong Converse Capacity

- Minimum of all R , such that every code with rate exceeding R leads to an asymptotically vanishing probability of successful decoding.



Entanglement-breaking (EB) and Hadamard channels: Holevo and regularized capacity agree

α -Information Radius

- For a single use of the channel, the success probability is bounded by, for any $\alpha > 1$,
(proof due to Polyanskiy&Verdú'10 / Sharma&Warsi'13)

$$p_{\text{succ}} \leq M^{\frac{1-\alpha}{\alpha}} \exp\left(\frac{\alpha-1}{\alpha} \tilde{R}_\alpha(\mathcal{W}_{A \rightarrow B})\right)$$

where M is the number of different messages,
and the α -information radius is

$$\tilde{\chi}_\alpha(\mathcal{W}_{A \rightarrow B}) := \inf_{\sigma_B} \sup_{\rho_A} \tilde{D}_\alpha(\mathcal{W}_{A \rightarrow B}(\rho_A) \parallel \sigma_B).$$

- Holevo capacity: $C_1 = \tilde{\chi}_1(\mathcal{W}_{A \rightarrow B})$

Sub-Additivity \Rightarrow Strong Converse

- For n uses of the channel and $M = \exp(nR)$:

$$p_{\text{succ}} \leq \exp\left(\frac{\alpha - 1}{\alpha} \left(\tilde{\chi}_\alpha(\mathcal{W}_{A \rightarrow B}^{\otimes n}) - nR \right)\right)$$

- Assume sub-additivity of information radius:

$$\tilde{\chi}_\alpha(\mathcal{W}_{A \rightarrow B}^{\otimes n}) \leq n\tilde{\chi}_\alpha(\mathcal{W}_{A \rightarrow B})$$

- For any $R > \tilde{\chi}_1(\mathcal{W}_{A \rightarrow B})$, we find $\alpha > 1$ such that $R > \tilde{\chi}_\alpha(\mathcal{W}_{A \rightarrow B})$ (by continuity)

- Thus, $p_{\text{succ}} \leq \exp\left(n\frac{\alpha - 1}{\alpha} (\tilde{\chi}_\alpha(\mathcal{W}_{A \rightarrow B}) - R)\right)$

Sub-additivity \Rightarrow all three capacities agree

α -Information Radius for EB Channels is Sub-Additive

- Using the EB CPM $\widehat{\mathcal{W}}(\cdot) = \sigma^{\frac{1-\alpha}{2\alpha}} \mathcal{W}(\cdot) \sigma^{\frac{1-\alpha}{2\alpha}}$,

$$\tilde{D}_\alpha(\mathcal{W}(\rho) \|\sigma) = \frac{\alpha}{\alpha-1} \log \|\widehat{\mathcal{W}}(\rho)\|_\alpha$$

- Maximal output norm of EB CPM (with any CPM $\widehat{\mathcal{E}}$) is multiplicative (King'03):

$$\max_{\rho_{AA'}} \|(\widehat{\mathcal{E}} \otimes \widehat{\mathcal{W}})(\rho_{AA'})\|_\alpha = \max_{\rho_A, \rho_{A'}} \|\widehat{\mathcal{E}}(\rho_A)\|_\alpha \cdot \|\widehat{\mathcal{W}}(\rho_{A'})\|_\alpha$$

α -Information Radius for EB Channels is Sub-Additive (II)

$$\begin{aligned}\tilde{\chi}_\alpha(\mathcal{E} \otimes \mathcal{W}) &= \min_{\sigma_{BB'}} \max_{\rho_{AA'}} \tilde{D}_\alpha((\mathcal{E} \otimes \mathcal{W})(\rho_{AB}) \middle\| \sigma_{BB'}) \\ &\leq \min_{\sigma_B, \sigma_{B'}} \max_{\rho_{AA'}} \tilde{D}_\alpha((\mathcal{E} \otimes \mathcal{W})(\rho_{AB}) \middle\| \sigma_B \otimes \sigma_{B'}) \\ &= \min_{\sigma_B, \sigma_{B'}} \frac{\alpha}{\alpha - 1} \log \max_{\rho_{AA'}} \|(\widehat{\mathcal{E}} \otimes \widehat{\mathcal{W}})(\rho_{AA'})\|_\alpha \\ &= \min_{\sigma_B, \sigma_{B'}} \frac{\alpha}{\alpha - 1} \log \max_{\rho_A, \rho_{A'}} \|\widehat{\mathcal{E}}(\rho_A)\|_\alpha \|\widehat{\mathcal{W}}(\rho_{A'})\|_\alpha \\ &= \tilde{\chi}_\alpha(\mathcal{E}) + \tilde{\chi}_\alpha(\mathcal{W})\end{aligned}$$

Conditional Rényi Entropy

- For a bipartite state ρ_{AB} , we define

$$\tilde{H}_\alpha(A|B)_\rho := \sup_{\sigma_B} -\tilde{D}_\alpha(\rho_{AB} \parallel \mathbb{I}_A \otimes \sigma_B).$$



- von Neumann Entropy: $H \equiv \tilde{H}_1$



- Min-Entropy (Renner): $H_{\min} \equiv \tilde{H}_\infty$

- Max-Entropy (König et al.): $H_{\max} \equiv \tilde{H}_{\frac{1}{2}}$

- Collision Entropy (Renner): \tilde{H}_2



- Classical case corresponds to Arimoto's conditional Rényi entropy



Duality and Uncertainty

- ⦿ For a tripartite pure state ρ_{ABC} , we find

$$\tilde{H}_\alpha(A|B) + \tilde{H}_\beta(A|C) = 0, \quad \frac{1}{\alpha} + \frac{1}{\beta} = 2.$$

- ⦿ Includes known relations for Min-/Max-Entropy and von Neumann entropy.
- ⦿ Implies full side information generalization of Massen-Uffink uncertainty relations:

$$\tilde{H}_\alpha(X|B) + \tilde{H}_\beta(Y|C) \geq \log \frac{1}{c}$$

The proof is implied by properties of the conditional entropy together with a proof framework due to Coles et al.'12 (also Berta++'10, Tomamichel&Renner'11)