COURSE 9

4.3. Iterative methods for solving linear systems

Because of round-off errors, direct methods become less efficient than iterative methods for large systems (>100 000 variables).

An iterative scheme for linear systems consists of converting the system

$$Ax = b \tag{1}$$

to the form

$$x = b' - Bx$$
.

After an initial guess for $x^{(0)}$, the sequence of approximations of the solution $x^{(0)}, x^{(1)}, ..., x^{(k)}, ...$ is generated by computing

$$x^{(k)} = b' - Bx^{(k-1)}$$
, for $k = 1, 2, 3, ...$

4.3.1. Jacobi iterative method

Consider the $n \times n$ linear system,

$$\begin{cases} a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n = b_1 \\ a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n = b_2 \\ \dots \\ a_{n1}x_1 + a_{n2}x_2 + \dots + a_{nn}x_n = b_n, \end{cases}$$

where we assume that the diagonal terms $a_{11}, a_{22}, ..., a_{nn}$ are all nonzero.

We begin our iterative scheme by solving each equation for one of the variables:

$$\begin{cases} x_1 = u_{12}x_2 + \dots + u_{1n}x_n + c_1 \\ x_2 = u_{21}x_1 + \dots + u_{2n}x_n + c_2 \\ \dots \\ x_n = u_{n1}x_1 + \dots + u_{nn-1}x_{n-1} + c_n, \end{cases}$$

where $u_{ij}=-rac{a_{ij}}{a_{ii}},\ c_i=rac{b_i}{a_{ii}},\ i=1,...,n.$

Let $x^{(0)} = (x_1^{(0)}, x_2^{(0)}, ..., x_n^{(0)})$ be an initial approximation of the solution. The k+1-th approximation is:

$$\begin{cases} x_1^{(k+1)} = u_{12}x_2^{(k)} + \dots + u_{1n}x_n^{(k)} + c_1 \\ x_2^{(k+1)} = u_{21}x_1^{(k)} + u_{23}x_3^{(k)} + \dots + u_{2n}x_n^{(k)} + c_2 \\ \dots \\ x_n^{(k+1)} = u_{n1}x_1^{(k)} + \dots + u_{nn-1}x_{n-1}^{(k)} + c_n, \end{cases}$$

for k = 0, 1, 2, ...

An algorithmic form:

$$x_i^{(k)} = \frac{b_i - \sum\limits_{j=1, j \neq i}^n a_{ij} x_j^{(k-1)}}{a_{ii}}, \ i = 1, 2, ..., n, \ \text{for } k \ge 1.$$

The iterative process is terminated when a convergence criterion is satisfied.

Stopping criterions: $\left|x^{(k)}-x^{(k-1)}\right|<\varepsilon$ or $\frac{\left|x^{(k)}-x^{(k-1)}\right|}{\left|x^{(k)}\right|}<\varepsilon$, with $\varepsilon>0$ - a prescribed tolerance.

Example 1 Solve the following system using the Jacobi iterative method. Use $\varepsilon = 10^{-3}$ and $x^{(0)} = (0\ 0\ 0\ 0)$ as the starting vector.

$$\begin{cases} 7x_1 - 2x_2 + x_3 & = 17 \\ x_1 - 9x_2 + 3x_3 - x_4 & = 13 \\ 2x_1 + 10x_3 + x_4 & = 15 \\ x_1 - x_2 + x_3 + 6x_4 & = 10. \end{cases}$$

These equations can be rearranged to give

$$x_1 = (17 + 2x_2 - x_3)/7$$

$$x_2 = (-13 + x_1 + 3x_3 - x_4)/9$$

$$x_3 = (15 - 2x_1 - x_4)/10$$

$$x_4 = (10 - x_1 + x_2 - x_3)/6$$

and, for example,

$$x_1^{(1)} = (17 + 2x_2^{(0)} - x_3^{(0)})/7$$

$$x_2^{(1)} = (-13 + x_1^{(0)} + 3x_3^{(0)} - x_4^{(0)})/9$$

$$x_3^{(1)} = (15 - 2x_1^{(0)} - x_4^{(0)})/10$$

$$x_4^{(1)} = (10 - x_1^{(0)} + x_2^{(0)} - x_3^{(0)})/6.$$

Substitute $x^{(0)} = (0,0,0,0)$ into the right-hand side of each of these equations to get

$$x_1^{(1)} = (17 + 2 \cdot 0 - 0)/7 = 2.428 571 429$$

 $x_2^{(1)} = (-13 + 0 + 3 \cdot 0 - 0)/9 = -1.444 444 444$
 $x_3^{(1)} = (15 - 2 \cdot 0 - 0)/10 = 1.5$
 $x_4^{(1)} = (10 - 0 + 0 - 0)/6 = 1.666 666 667$

and so $x^{(1)} = (2.428\ 571\ 429, -1.444\ 444\ 444, 1.5, 1.666\ 666\ 667)$. The Jacobi iterative process:

$$x_1^{(k+1)} = \left(17 + 2x_2^{(k)} - x_3^{(k)}\right) / 7$$

$$x_2^{(k+1)} = \left(-13 + x_1^{(k)} + 3x_3^{(k)} - x_4^{(k)}\right) / 9$$

$$x_3^{(k+1)} = \left(15 - 2x_1^{(k)} - x_4^{(k)}\right) / 10$$

$$x_4^{(k+1)} = \left(10 - x_1^{(k)} + x_2^{(k)} - x_3^{(k)}\right) / 6, \qquad k \ge 1.$$

We obtain a sequence that converges to

$$\mathbf{x}^{(9)} = (2.000127203, -1.000100162, 1.000118096, 1.000162172).$$

4.3.2. Gauss-Seidel iterative method

Almost the same as Jacobi method, except that each x-value is improved using the most recent approx. of the other variables.

For a $n \times n$ system, the k + 1-th approximation is:

$$\begin{cases} x_1^{(k+1)} = u_{12}x_2^{(k)} + \dots + u_{1n}x_n^{(k)} + c_1 \\ x_2^{(k+1)} = u_{21}x_1^{(k+1)} + u_{23}x_3^{(k)} + \dots + u_{2n}x_n^{(k)} + c_2 \\ \dots \\ x_n^{(k+1)} = u_{n1}x_1^{(k+1)} + \dots + u_{nn-1}x_{n-1}^{(k+1)} + c_n, \end{cases}$$

with $k=0,1,2,...;\;u_{ij}=-\frac{a_{ij}}{a_{ii}},\;c_i=\frac{b_i}{a_{ii}},\;i=1,...,n$ (as in Jacobi method).

Algorithmic form:

$$x_i^{(k)} = \frac{b_i - \sum_{j=1}^{i-1} a_{ij} x_j^{(k)} - \sum_{j=i+1}^{n} a_{ij} x_j^{(k-1)}}{a_{ii}}$$

for each i = 1, 2, ...n, and for $k \ge 1$.

Stopping criterions: $\left|x^{(k)}-x^{(k-1)}\right|<\varepsilon$, or $\frac{\left|\mathbf{x}^{(k)}-\mathbf{x}^{(k-1)}\right|}{\left|\mathbf{x}^{(k)}\right|}<\varepsilon$, with ε - a prescribed tolerance, $\varepsilon>0$.

Remark 2 Because the new values can be immediately stored in the location that held the old values, the storage requirements for x with the Gauss-Seidel method is half than that for Jacobi method and the rate convergence is faster.

Example 3 Solve the following system using the Gauss-Seidel iterative method. Use $\varepsilon = 10^{-3}$ and $\mathbf{x}^{(0)} = (0\ 0\ 0\ 0)$ as the starting vector.

$$\begin{cases} 7x_1 - 2x_2 + x_3 & = 17 \\ x_1 - 9x_2 + 3x_3 - x_4 & = 13 \\ 2x_1 + 10x_3 + x_4 & = 15 \\ x_1 - x_2 + x_3 + 6x_4 & = 10 \end{cases}$$

We have

$$x_1 = (17 + 2x_2 - x_3)/7$$

$$x_2 = (-13 + x_1 + 3x_3 - x_4)/9$$

$$x_3 = (15 - 2x_1 - x_4)/10$$

$$x_4 = (10 - x_1 + x_2 - x_3)/6,$$

and, for example,

$$x_1^{(1)} = (17 + 2x_2^{(0)} - x_3^{(0)})/7$$

$$x_2^{(1)} = (-13 + x_1^{(1)} + 3x_3^{(0)} - x_4^{(0)})/9$$

$$x_3^{(1)} = (15 - 2x_1^{(1)} - x_4^{(0)})/10$$

$$x_4^{(1)} = (10 - x_1^{(1)} + x_2^{(1)} - x_3^{(1)})/6,$$

which provide the following Gauss-Seidel iterative process:

$$\begin{split} x_1^{(k+1)} &= \left(17 + 2x_2^{(k)} - x_3^{(k)}\right)/7 \\ x_2^{(k+1)} &= \left(-13 + x_1^{(k+1)} + 3x_3^{(k)} - x_4^{(k)}\right)/9 \\ x_3^{(k+1)} &= \left(15 - 2x_1^{(k+1)} - x_4^{(k)}\right)/10 \\ x_4^{(k+1)} &= \left(10 - x_1^{(k+1)} + x_2^{(k+1)} - x_3^{(k+1)}\right)/6, \quad \text{for } k \ge 1. \end{split}$$

Substitute $\mathbf{x}^{(0)} = (0,0,0,0)$ into the right-hand side of each of these equations to get

$$x_1^{(1)} = (17 + 2 \cdot 0 - 0)/7 = 2.428 571 429$$

 $x_2^{(1)} = (-13 + 2.428 571 429 + 3 \cdot 0 - 0)/9 = -1.1746031746$
 $x_3^{(1)} = (15 - 2 \cdot 2.428 571 429 - 0)/10 = 1.0142857143$
 $x_4^{(1)} = (10 - 2.428 571 429 - 1.1746031746 - 1.0142857143)/6$
 $= 0.8970899472$

and so

$$\mathbf{x}^{(1)} = (2.428571429 - 1.1746031746, 1.0142857143, 0.8970899472).$$

Similar procedure generates a sequence that converges to

$$\mathbf{x}^{(5)} = (2.000025, -1.000130, 1.000020.0.999971).$$

4.3.3. Relaxation method

In case of convergence, the Gauss-Seidel method is twice faster than Jacobi method. The convergence can be more improved using **relaxation method (SOR method) (**SOR=Succesive Over Relaxation)

Algorithmic form of the method:

$$x_i^{(k)} = \frac{\omega}{a_{ii}} \left(b_i - \sum_{j=1}^{i-1} a_{ij} x_j^{(k)} - \sum_{j=i+1}^n a_{ij} x_j^{(k-1)} \right) + (1 - \omega) x_i^{(k-1)}$$

for each i = 1, 2, ...n, and for k > 1.

For $0<\omega<1$ the procedure is called **under relaxation method,** that can be used to obtain convergence for systems which are not convergent by Gauss-Siedel method.

For $\omega>1$ the procedure is called **over relaxation method**, that can be used to accelerate the convergence for systems which are convergent by Gauss-Siedel method.

By Kahan's Theorem follows that the method converges for $0 < \omega < 2$.

Remark 4 For $\omega = 1$, relaxation method is Gauss-Seidel method.

Example 5 Solve the following system, using relaxation iterative method Use $\varepsilon = 10^{-3}$, $\mathbf{x}^{(0)} = (1\ 1\ 1)$ and $\omega = 1.25$,

$$4x_1 + 3x_2 = 24$$

 $3x_1 + 4x_2 - x_3 = 30$
 $- x_2 + 4x_3 = -24$

We have

$$x_1^{(k)} = 7.5 - 0.937x_2^{(k-1)} - 0.25x_1^{(k-1)}$$

$$x_2^{(k)} = 9.375 - 9.375x_1^{(k)} + 0.3125x_3^{(k-1)} - 0.25x_2^{(k-1)}$$

$$x_3^{(k)} = -7.5 + 0.3125x_2^{(k)} - 0.25x_3^{(k-1)}, \text{ for } k \ge 1.$$

The solution is (3, 4, -5).

4.3.4 The matriceal formulations of the iterative methods

Split the matrix A into the sum

$$A = D + L + U,$$

where D is the diagonal of A, L the lower triangular part of A, and U the upper triangular part of A. That is,

$$D = \begin{bmatrix} a_{11} & 0 & \cdots & 0 \\ 0 & \cdots & \ddots & \vdots \\ \vdots & \cdots & \cdots & 0 \\ 0 & \cdots & 0 & a_{nn} \end{bmatrix}, \quad L = \begin{bmatrix} 0 & \cdots & 0 \\ a_{21} \vdots & & \ddots \\ \vdots & \cdots & \ddots & \vdots \\ a_{n1} & \cdots & a_{n,n-1} & 0 \end{bmatrix},$$

$$U = \begin{bmatrix} 0 & a_{12} & \cdots & a_{1n} \\ \vdots & \cdots & \ddots & \vdots \\ & & \cdots & a_{n-1,n} \\ 0 & \cdots & & 0 \end{bmatrix}$$

The system Ax = b can be written as

$$(D+L+U)\mathbf{x} = \mathbf{b}.$$

The Jacobi method in matriceal form is given by:

$$D\mathbf{x}^{(k)} = -(L+U)\mathbf{x}^{(k-1)} + \mathbf{b}$$

the Gauss-Seidel method in matriceal form is given by:

$$(D+L)\mathbf{x}^{(k)} = -U\mathbf{x}^{(k-1)} + \mathbf{b}$$

and the relaxation method in matriceal form is given by:

$$(D + \omega L)x^{(k)} = ((1 - \omega)D - \omega U)x^{(k-1)} + \omega b$$

Convergence of the iterative methods

Remark 6 The convergence (or divergence) of the iterative process in the Jacobi and Gauss-Seidel methods does not depend on the initial guess, but depends only on the character of the matrices themselves. However, a good first guess in case of convergence will make for a relatively small number of iterations.

A sufficient condition for convergence:

Theorem 7 (Convergence Theorem) If A is strictly diagonally dominant, then the Jacobi, Gauss-Seidel and relaxation methods converge for any choice of the starting vector $\mathbf{x}^{(0)}$.

Example 8 Consider the system of equations

$$\begin{bmatrix} 3 & 1 & 1 \\ -2 & 4 & 0 \\ -1 & 2 & -6 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 4 \\ 1 \\ 2 \end{bmatrix}.$$

The coefficient matrix of the system is strictly diagonally dominant since

$$|a_{11}| = |3| = 3 > |1| + |1| = 2$$

 $|a_{22}| = |4| = 4 > |-2| + |0| = 2$
 $|a_{33}| = |-6| = 6 > |-1| + |2| = 3$.

Hence, if the Jacobi or Gauss-Seidel method are used to solve the system of equations, they will converge for any choice of the starting vector $\mathbf{x}^{(0)}$

Example 9 Consider the linear system

$$\left[\begin{array}{cc} 4 & 1 \\ 2 & 5 \end{array}\right] \left[\begin{array}{c} x \\ y \end{array}\right] = \left[\begin{array}{c} 3 \\ 1 \end{array}\right]$$

Perform two iterations of the Jacobi and the Gauss-Seidel methods' to this system beginning with the vector x = [3, 11]. (Solutions of the system are 7/9 and -1/9).

5. Numerical methods for solving nonlinear equations in $\ensuremath{\mathbb{R}}$

Let $f: \Omega \to \mathbb{R}, \ \Omega \subset \mathbb{R}$. Consider the equation

$$f(x) = 0, \quad x \in \Omega. \tag{2}$$

We attach a mapping $F: D \to D, D \subset \Omega^n$ to this equation.

Let $(x_0,...,x_{n-1}) \in D$. Using F and the numbers $x_0,x_1,...,x_{n-1}$ we construct iteratively the sequence

$$x_0, x_1, ..., x_{n-1}, x_n, ...$$
 (3)

with

$$x_i = F(x_{i-n}, ..., x_{i-1}), \quad i = n,$$
 (4)

The problem consists in choosing F and $x_0, ..., x_{n-1} \in D$ such that the sequence (3) to be convergent to the solution of the equation (2).

Definition 10 The procedure of approximation the solution of equation (2) by the elements of the sequence (3), computed as in (4), is called F-method.

The numbers $x_0, x_1, ..., x_{n-1}$ are called **the starting points** and the k-th element of the sequence (3) is called an approximation of k-th order of the solution.

If the set of starting points has only one element then the F-method is **an one-step method**; if it has more than one element then the F-method is **a multistep method**.

Definition 11 If the sequence (3) converges to the solution of the equation (2) then the F-method is convergent, otherwise it is divergent.

Definition 12 Let $\alpha \in \Omega$ be a solution of the equation (2) and let $x_0, x_1, ..., x_{n-1}, x_n, ...$ be the sequence generated by a given F-method. The number p having the property

$$\lim_{x_i \to \alpha} \frac{\alpha - F(x_{i-n+1}, \dots, x_i)}{(\alpha - x_i)^p} = C \neq 0, \quad C = constant,$$

is called the order of the F-method.

We construct some classes of F-methods based on the interpolation procedures.

Let $\alpha \in \Omega$ be a solution of the equation (2) and $V(\alpha)$ a neighborhood of α . Assume that f has inverse on $V(\alpha)$ and denote $g := f^{-1}$. Since

$$f\left(\alpha\right)=0$$

it follows that

$$\alpha = g(0)$$
.

This way, the approximation of the solution α is reduced to the approximation of g(0).

Definition 13 The approximation of g by means of an interpolating method, and α by the value of g at point zero is called **inverse interpolation procedure**.

5.1. One-step methods

Let F be a one-step method, i.e., for a given x_i we have $x_{i+1} = F(x_i)$.

Remark 14 If p = 1 the convergence condition is |F'(x)| < 1.

If p>1 there always exists a neighborhood of α where the F-method converges.

All information on f are given at a single point, the starting value \Rightarrow we are lead to Taylor interpolation.

Theorem 15 Let α be a solution of equation (2), $V(\alpha)$ a neighborhood of α , $x, x_i \in V(\alpha)$, f fulfills the necessary continuity conditions.

Then we have the following method, denoted by F_m^T , for approximating α :

$$F_m^T(x_i) = x_i + \sum_{k=1}^{m-1} \frac{(-1)^k}{k!} [f(x_i)]^k g^{(k)}(f(x_i)),$$
 (5)

where $g = f^{-1}$.

Proof. There exists $g = f^{-1} \in C^m[V(0)]$. Let $y_i = f(x_i)$ and consider Taylor interpolation formula

$$g(y) = (T_{m-1}g)(y) + (R_{m-1}g)(y)$$

with

$$(T_{m-1}g)(y) = \sum_{k=0}^{m-1} \frac{1}{k!} (y - y_i)^k g^{(k)}(y_i),$$

and $R_{m-1}g$ is the corresponding remainder.

Since $\alpha = g(0)$ and $g \approx T_{m-1}g$, it follows

$$\alpha \approx (T_{m-1}g)(0) = x_i + \sum_{k=1}^{m-1} \frac{(-1)^k}{k!} y_i^k g^{(k)}(y_i).$$

Hence,

$$x_{i+1} := F_m^T(x_i) = x_i + \sum_{k=1}^{m-1} \frac{(-1)^k}{k!} [f(x_i)]^k g^{(k)}(f(x_i))$$

is an approximation of $\alpha,$ and F_m^T is an approximation method for the solution $\alpha.$ \blacksquare

Concerning the order of the method F_m^T we state:

Theorem 16 If $g = f^{-1}$ satisfies condition $g^{(m)}(0) \neq 0$, then $\operatorname{ord}(F_m^T) = m$.

Proof. Bibliography ■

Remark 17 We have an upper bound for the absolute error in approximating α by x_{i+1} :

$$\left|\alpha - F_m^T(x_i)\right| \le \frac{1}{m!} [f(x_i)]^m M_m g, \text{ with } M_m g = \sup_{y \in V(0)} \left|g^{(m)}(y)\right|.$$

Particular cases.

1) Case m = 2.

$$F_2^T(x_i) = x_i - \frac{f(x_i)}{f'(x_i)}$$
.

This method is called **Newton's method** (the tangent method). Its order is 2.