

# PH502: Scientific Programming Concepts

Irish Centre for High End Computing (ICHEC)

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- In this lecture we will continue to examine algorithms.
- We have already mentioned that speed and memory usage can be correlated inversely.
- However also some designs are simpler to code but maybe inefficient.

- Insertion sort uses an *Incremental approach*: Sort array  $A[1, \dots, j - 1]$ , insert  $A[j]$ , get the sorted array  $A[1, \dots, j]$ .
- Design Techniques:
  - ▶ Brute-Force
  - ▶ Divide and Conquer
  - ▶ Greedy Method
  - ▶ Backtracking
  - ▶ Dynamic Programming

- Trivial naive method that tries every possible solution and check which is the best one
- Straightforward and easiest approach to apply
- Useful for solving small size problems, should always reach the solution.
- Not always efficient
- Very useful when writing a test routine to check the correctness of more efficient algorithms
- Examples:
  - ▶ Iterative Algorithm for Factorial
  - ▶ Linear search
  - ▶ Selection sort
  - ▶ Bubble sort

- Break up problem into smaller parts of the same problem until each is small enough to be easily solved.
- Solve each part recursively.
- Combine solutions to sub-problems into overall solution.
- It is important that the subproblems are independent.
- Examples:
  - ▶ Recursive Algorithm for Factorial
  - ▶ Binary Search
  - ▶ Merge Sort

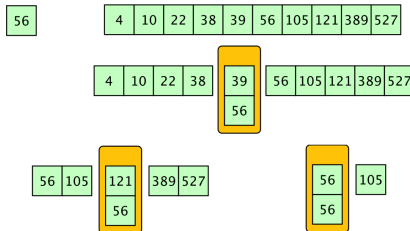
## Example: Binary Search

**Input:** A list of integers sorted in ascending order  $A = \{a_1, a_2, \dots, a_n\}$ ,  $k$

**Output:** True if the integer is found; False otherwise.

**Procedure:** Select the middle element on the list and compare; if not found, discard half of the list where the element is definitely not placed.

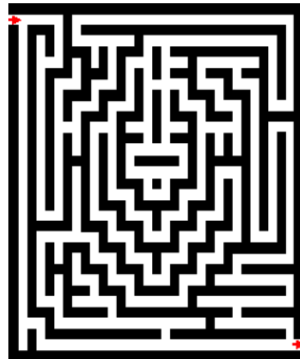
```
begin=1, end=n
while begin ≤ end do
    mid=⌊((begin+end)/2)⌋
    if k=A[mid] then
        return TRUE
    else if k<A[mid] then
        end=mid-1
    else
        end=mid+1
    end if
end while
return FALSE
```



- The solution is constructed through a sequence of steps. At each step we choose the locally optimal solution.
- Always makes the choice that looks best at the moment and adds it to the current subsolution.
- Simple and straightforward, Easy to invent, easy to implement and most of the time quite efficient
- Many problems cannot be solved correctly by greedy approach
- Mainly used to solve optimization problems
- Examples:
  - ▶ Job scheduling problems
  - ▶ The Knapsack problem: A thief robbing a store and can carry a maximal weight of  $w$  into their knapsack. There are  $n$  items and  $i^{\text{th}}$  item weigh  $w_i$  and is worth  $v_i$  dollars. What items should thief take?
  - ▶ The coin exchange problem: Pay money back to customer using fewest number of coins

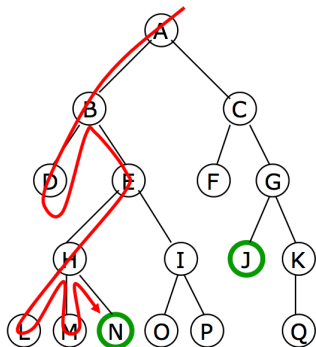
- A choice taken may be wrong, in which case the prospective solution is “backtracked” to undo the mistake
- A methodical way of trying out various sequences of decisions, until you find one that “works”
- Example: Maze Path Finder

```
while not at destination do  
  Choose a path X  
  if You have been in X then  
    Choose another path Y  
     $X \leftarrow Y$   
  end if  
  Move to X  
end while
```





# Example: Depth-First Searching



- It is a combination of brute-force and backtracking.
- It starts at the root and explores nodes from there, looking for a goal node
- It explores a path all the way to a leaf before backtracking and exploring another path
- Node are explored in the order A B D E H L M **N** I O P C F G **J**

- An inverse divide-and-conquer approach
- The smallest sub-problems are firstly solved and their solutions cached;
- The final solution arises from the solutions to each sub-problem
- A bottom-up approach
- Difference from the classical Divide and Conquer is subproblems are solved only once and solutions are stored for reuse.
- Applications: Bioinformatics, Control theory, Information theory, Operations research, ...
- Examples:
  - ▶ Viterbi Algorithm for Markov Models
  - ▶ Bellman-Ford for finding shortest path in networks
  - ▶ The Manhattan Tourist Problem

## Example: Fibonacci Sequence

- **Definition:** The first two elements in the sequence are respectively 0 and 1. Every single element is the sum of the previous two elements in the sequence.

$$F_n = F_{n-1} + F_{n-2}, F_0 = 0, F_1 = 1$$

Dynamic Programming:

DynFib(n)

Table[0]=0

Table[1]=1

**for** i=2 to n **do**

    Table[i]=Table[i-1]+Table[i-2]

**end for**

return Table[n]

Divide and Conquer:

DivConFib(n)

**if** n >= 2 **then**

    return DivConFib(n-1)+DivConFib(n-2)

**end if**

return