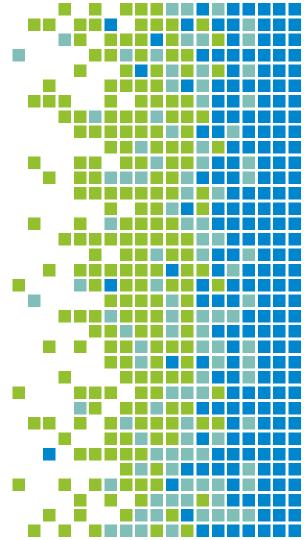


# PRACE Course: Intermediate MPI

9-11 November 2022

**MPI** Point-to-Point Communication





## Messages

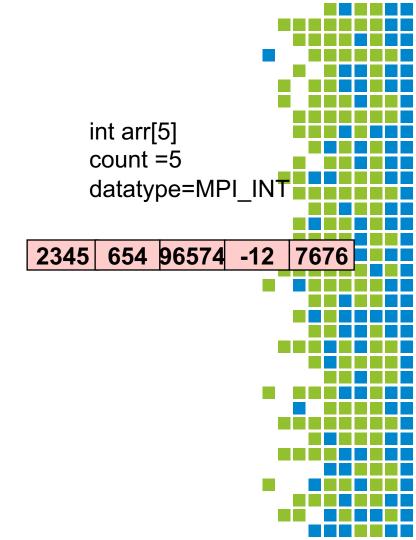
- A message contains a number of elements of some particular datatype.
- MPI datatypes:
  - Basic datatype.
  - Derived datatypes
- C types are different from Fortran types.
- Datatype handles are used to describe the type of the data in the memory.

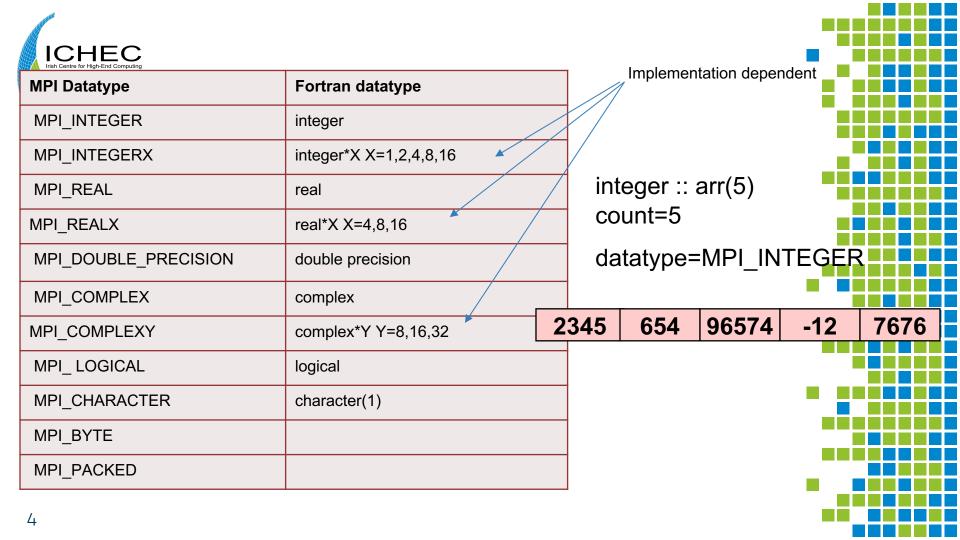
Example: message with 5 integers

2345	654	96574	-12	7676



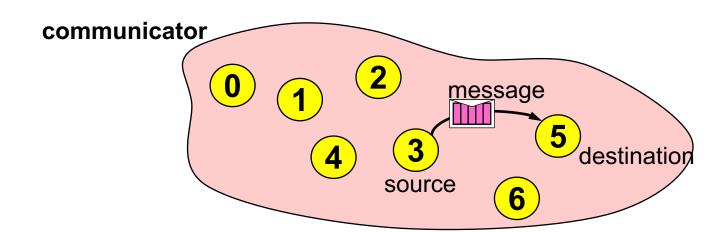
MPI Datatype	C datatype
MPI_CHAR	signed char
MPI_SHORT	signed short int
MPI_INT	signed int
MPI_LONG	signed long int
MPI_UNSIGNED_CHAR	unsigned char
MPI_UNSIGNED_SHORT	unsigned short int
MPI_UNSIGNED	unsigned int
MPI_UNSIGNED_LONG	unsigned long int
MPI_FLOAT	float
MPI_DOUBLE	double
MPI_LONG_DOUBLE	long double
MPI_BYTE	
MPI_PACKED	





## ICHEC Point-to-Point Communication

- Communication between two processes.
- Source process sends message to destination process.
- Communication takes place within a communicator, e.g., MPI\_COMM\_WORLD.
- Processes are identified by their ranks in the communicator





## Sending a Message

int MPI\_Send(void \*buf, int count, MPI\_Datatype datatype, int dest, int tag, MPI\_Comm comm)

Fortran:

```
MPI_SEND(buf, count, datatype, dest, tag, comm, ierror)

TYPE(*), DIMENSION(..) :: buf

TYPE(MPI_Datatype) :: datatype; TYPE(MPI_Comm) :: comm
integer :: count, dest, tag; integer,optional :: ierror
```

- buf is the starting point of the message with count elements, each described with datatypeHandle.
- dest is the rank of the destination process within the communicator comm.
- tag is an additional nonnegative integer piggyback information, additionally transferred with the message.



## Receiving a Message

int MPI\_Recv(void \*buf, int count, MPI\_Datatype datatype, int source, int tag, MPI\_Comm comm, MPI\_Status \*status)

Fortran:

```
MPI_RECV(buf, count, datatype, source, tag, comm, status, ierroff)

TYPE(*), DIMENSION(..) :: buf; TYPE(MPI_Status) :: status

TYPE(MPI_Datatype) :: datatype; TYPE(MPI_Comm) :: committee integer :: count, source, tag; integer,optional :: ierror
```

- buf/count/datatype describe the receive buffer.
- Receiving the message sent by process with rank source in comm.
- Envelope information is returned in status if not MPI\_STATUS\_IGNORE

# Requirements for Point-to-Point Communications

For a communication to succeed:

- Sender must specify a valid destination rank.
- Receiver must specify a valid source rank.
- The communicator must be the same.
- Tags must match.
- Message datatypes must match.
- Message matching rule: receives only if <u>comm</u>, <u>tag</u>, and <u>type</u> match.
- Receiver's buffer must be large enough



## **Example - One Ping**

```
#include <stdio.h>
#include <mpi.h>
                                                           Ping: Process 0 sends a message to
                                                           process 1
int main(int argc, char **argv){
                                                           Run with two processes
  int myRank, ierror, arr[5];
  MPI Status stat;
  ierror=MPI_Init(&argc,&argv);
  ierror=MPI Comm_rank(MPI_COMM_WORLD,&myRank);
  if (myRank == 0) {
    arr[0]=2345; arr[1]=654; arr[2]=96574; arr[3]=-12; arr[4]=7676;
    MPI Send(arr, 5, MPI_INT, 1, 100, MPI_COMM_WORLD);
  } else if (myRank == 1) {
    MPI Recv(arr, 5, MPI INT, 0, 100, MPI COMM WORLD, &stat);
  ierror=MPI Finalize();
  return 0;
```





### **Wildcards**

- Receiver can wildcard.
- To receive from any source
  - source = MPI\_ANY\_SOURCE
- To receive from any tag:
  - tag = MPI\_ANY\_TAG
- Actual source and tag are returned in the receiver's status parameter.
- Use only when necessary and beneficial. It is much safer to specify the source and tag when you know them
- An MPI application can tell the MPI library that it will never use MPI\_ANY\_SOURCE and/or MPI\_ANY\_TAG on this communicator by setting an assertion.



## **Communication Envelope**

- Envelope information is returned from MPI\_RECV in status.
- C:

```
MPI_Status status;
status.MPI_SOURCE
status.MPI_TAG
count via MPI Get count()
```

Fortran:

```
TYPE(MPI_Status) :: status
status%MPI_SOURCE
status%MPI_TAG
count via MPI_GET_COUNT()
```

From: **source** rank tag To: destination rank item-1 item-2 ≿count" item-3 item-4 elements item-n



## **Receive Message Count**

- C: int MPI\_Get\_count(MPI\_Status \*status, MPI\_Datatype datatype, ir \*count)
- Fortran:

```
MPI_GET_COUNT(status, datatype, count, ierror)
    TYPE(MPI_Status) :: status
    TYPE(MPI_Datatype) :: datatype
    integer :: count; integer,optional :: ierror
```

returns integer, 0 or MPI\_UNDEFINED.



### **Communication Modes**

- Send communication modes:
  - synchronous send
  - buffered [asynchronous] send
  - standard send
  - ready send

- MPI\_Ssend
- MPI Bsend
- MPI\_Send
- MPI\_Rsend
- Receiving all modes MPI\_Recv
- All in blocking and nonblocking forms.



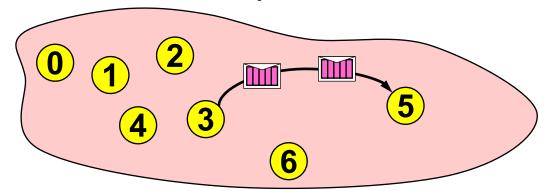
## ICHEOMMunication Modes - Definition

Sender modes	Definition	Notes
Synchronous send MPI_Ssend	Only completes when the receive has started	risk of deadlock risk of serialization risk of waiting —> idle time
Buffered send MPI_Bsend	Always completes (unless an error occurs), irrespective of receiver	needs application-defined buffer to be declared with MPI_BUFFER_ATTACH/DET ACH
Standard send MPI_Send	Standard send. Either synchronous or buffered	Uses an internal buffer
Ready send MPI_Rsend	May be started only if the matching receive is already posted!	highly dangerous!
Receive MPI_Recv	Completes when the message (data) has arrived	Same routine for all communication modes



## ICHEC Message Order Preservation

- Rule for messages on the same connection, i.e., same communicator, source, and destination rank:
- Messages do not overtake each other.
- This is true even for non-synchronous sends.



MPI messages are non-overtaking: if one process send two messages to another, then they will be received in the order they were sent



## Timing in MPI

- C: double MPI\_Wtime( void );
- Fortran: DOUBLE PRECISION MPI WTIME()
- The elapsed (wall-clock) time between two points:

```
double t1, t2;
t1 = MPI_Wtime();
t1 = MPI_Wtime();
t1 = MPI_Wtime ()
... work to be timed ...
t2 = MPI_Wtime();
t2 = MPI_Wtime();
printf("Elapsed time is %f\n", t2 - t1 );
Print*, 'Elapsed time is', t2-t1
```

MPI\_Wtick(): Returns the number of seconds between processor clock ticks.





### **Practical**

• Practical 1: pingpong benchmark





### **Deadlock**

 Code in each MPI process: MPI\_Ssend(..., right\_rank, ...)

MPI\_Recv( ..., left\_rank, ...)

Will block and never return, because MPI\_Recv cannot be called in the right-hand MPI process

- Reorganise the communications, i.e. First even processes send odd processes receive, Then odd processes send even processes receive.(Inefficient)
- Use MPI\_SendRecv(), which combines send and recv in a single deadlock-free call



## **Example - Deadlock**

 Problem (Unless MPI\_Send/MPI\_Recv is buffered)

#### Solutions:

 Reverse the order of one of the send/receive pairs

```
if (myRank == 0) {
         MPI_Send(a, 5, MPI_INT, 1, 100, MPI_COMM_WORLD);
         MPI_Recv(b, 5, MPI_INT, 1, 101, MPI_COMM_WORLD, &status);
} else if (myRank == 1) {
         MPI_Send(b, 5, MPI_INT, 0, 101, MPI_COMM_WORLD);
         MPI_Recv(a, 5, MPI_INT, 0, 100, MPI_COMM_WORLD, &status);
}
```

```
if (myRank == 0) {
         MPI_Send(a, 5, MPI_INT, 1, 100, MPI_COMM_WORLD);
         MPI_Recv(b, 5, MPI_INT, 1, 101, MPI_COMM_WORLD, &status);
} else if (myRank == 1) {
         MPI_Recv(a, 5, MPI_INT, 0, 100, MPI_COMM_WORLD, &status);
         MPI_Send(b, 5, MPI_INT, 0, 101, MPI_COMM_WORLD);}
```

Use non-blocking communication

## ICHEC Non-Blocking Communication

Separate communication into three phases:

- Initiate non-blocking communication
  - returns Immediately
  - routine name starting with MPI\_I...
- Do some work
  - "latency hiding"
- Wait for non-blocking communication to complete, i.e.
  - The send buffer is read out or
  - The receive buffer is filled in.
- Overlap communication with computation
- Better performance



## **Non-Blocking Examples**

Non-blocking send

MPI\_Isend(...)

doing some other work

MPI\_Wait(...)

Non-blocking receive

MPI\_Irecv(...)

doing some other work

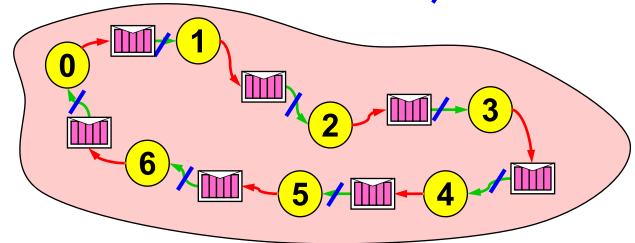
MPI\_Wait(...)

= waiting until operation locally completed



## **Non-Blocking Send**

- Initiate non-blocking send
- in the ring example: Initiate non-blocking send to the right neighbor
- Do some work:
- in the ring example: Receiving the message from left neighbor
- Now, the message transfer can be completed
- Wait for non-blocking send to complete /





```
Non-Blocking Send
```

int MPI Isend(void \*buf, int count, MPI Datatype datatype, int dest, in

```
tag, MPI_Comm comm, MPI_Request * request)

• Fortran:

MPI_ISEND(buf, count, datatype, dest, tag, comm, request, ierror)

TYPE(*), DIMENSION(..) :: buf

TYPE(MPI_Datatype) :: datatype

TYPE(MPI_Comm) :: comm

TYPE(MPI_Request) :: request

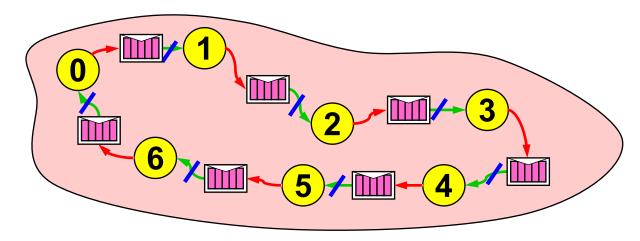
integer :: count, dest, tag; integer,optional :: ierror
```

- Request: used to check the status of send or to wait for its completion
- MPI Isend + MPI Wait immediately after is equivalent to MPI Send



## **Non-Blocking Receive**

- Initiate non-blocking receive
- → in the ring example: Initiate non-blocking receive from left neighbor
- Do some work:
- in the ring example: Sending the message to the right neighbor
- Now, the message transfer can be completed
- Wait for non-blocking receive to complete





## **Non-Blocking Receive**

int MPI\_Irecv(void \*buf, int count, MPI\_Datatype datatype, int source, int tag, MPI Comm comm, MPI Request \* request)

Fortran:

```
MPI_IRECV(buf, count, datatype, source, tag, comm, request, ier TYPE(*), DIMENSION(..) :: buf
TYPE(MPI_Datatype) :: datatype
TYPE(MPI_Comm) :: comm
TYPE(MPI_Request) :: request
integer :: count, dest, tag; integer,optional :: ierror
```

A blocking send can be used with a non-blocking receive, and vice-versa



## Completion

- C
  - MPI\_Wait(MPI Request \*request, MPI\_Status \*status);
    MPI\_Test (MPI Request \*request, int \*flag, MPI\_Status \*status);
- Fortran:

- buf must not be used between Isend/Irecv and Wait
- one must
  - wait or
  - loop with TEST until request is completed, i.e., flag == 1 or.true.



## **Example - Deadlock**

 Problem (Unless MPI\_Send/MPI\_Recv is buffered)

#### Solutions:

- Reverse the order of or e of the send/receive pairs
- Use non-blocking communication

```
if (myRank == 0) {
         MPI_Send(a, 5, MPI_INT, 1, 100, MPI_COMM_WORLD);
         MPI_Recv(b, 5, MPI_INT, 1, 101, MPI_COMM_WORLD, &status);
} else if (myRank == 1) {
         MPI_Send(b, 5, MPI_INT, 0, 101, MPI_COMM_WORLD);
         MPI_Recv(a, 5, MPI_INT, 0, 100, MPI_COMM_WORLD, &status);
}
```

```
if (myRank == 0) {
         MPI_Isend(a, 5, MPI_INT, 1, 100, MPI_COMM_WORLD,&request);
         MPI_Recv(b, 5, MPI_INT, 1, 101, MPI_COMM_WORLD, &status);
         MPI_Wait(&request, &status);
} else if (myRank == 1) {
         MPI_Recv(a, 5, MPI_INT, 0, 100, MPI_COMM_WORLD, &status);
         MPI_Send(b, 5, MPI_INT, 0, 101, MPI_COMM_WORLD);}
```



## ICHEC Blocking and Non-Blocking

Send and receive can be blocking or non-blocking.

Send Mode	Blocking Function	Nonblocking Function
Standard	MPI_Send	MPI_Isend
Synchronous	MPI_Ssend	MPI_Issend
Ready	MPI_Rsend	MPI_Irsend
Buffered	MPI_Bsend	MPI_lbsend

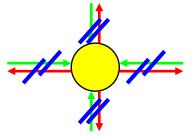
- Issend + Wait is equivalent to blocking call: Ssend
- MPI Sendrecv: Irecv + Send + Wait



## Multiple Non-Blocking Communications

You have several request handles:

- Wait or test for completion of one message
  - MPI\_Waitany / MPI\_Testany
- Wait or test for completion of all messages
  - MPI\_Waitall / MPI\_Testall
- Wait or test for completion of at least one message
  - MPI\_Waitsome / MPI\_Testsome





## **Summary**

- Deadlock-free code with Nonblocking communication
- Code harder to debug and maintain.

