Índice 1. Setup			1. Setup Template corto	
3	Range queries 3.1. Range queries comunes	3 3 4 4 4	return 0; } correr_interactivo.sh: Compilar y ejecutar \$1 clear && make -s \$1 && ./\$1	
4 4 4	Grafos 1.1. Toposort de un DAG 1.2. DAG condensado 1.3. Bipartite check 1.4. Encontrar puentes y articulaciones 1.5. Menores caminos	5 5 6 6 6	Makefile CC = g++ CPPFLAGS = -Wall -g \ -fsanitize=undefined -fsanitize=bounds \ -std=c++17 -00 correr_archivo.sh: Compilar y ejecutar \$1 con el input \$2	
	Programacion Dinamica 5.1. LIS (Longest Increasing Subsequence)	7 7	clear && make -s \$1 && ./\$1 \$2 compilar.sh: Compilar \$1 y mostrar primeras \$2 lineas de error	
6 6	Matemática 5.1. Aritmética	8 8 8 9	clear && make -s \$1 2>&1 head -\$2 Template completo #include <bits stdc++.h=""></bits>	
7	Estructuras locas 7.1. Disjoint set union	9 9 10	<pre>using namespace std; #define forall(it,v) for (auto it = begin(v); it != end(v); it++) #define forr(i,a,b) for(int i = int(a); i < int(b); i++) #define forn(i,n) forr(i,0,n) #define all(v) begin(v), end(v)</pre>	
	Sin categorizar Brainstorming	10 11	<pre>#define all(v)</pre>	

```
templAB ostream& operator << (ostream& o, pair<A,B>& p) { return o <<
   → p.fst << " " << p.snd; }</pre>
templT ostream& operator << (ostream& o, vector <T>& v) { forall(it,v
   → ) { o << *it << " ": } return o: }</pre>
using ll = long long;
int main (int argc, char** argv) { ios::sync_with_stdio(0); cin.tie

→ (0); cout.tie(0); if (argc == 2) freopen("input", "r", stdin);

   return 0;
    STL
2.1. Vector
2.1.1. Busqueda binaria (lower_bound)
Primer igual
templT int primer_igual (vector < T > & arr, T x) {
    auto it = lower_bound(all(arr), x);
    if (it == arr.end() || *it != x) return -1;
    return it - arr.begin();
Último igual
templT int ultimo_igual (vector<T>& arr, T x) {
    if (arr.begin() == arr.end()) return -1;
    auto it = prev(upper_bound(all(arr), x));
    if (*it != x) return -1;
    return it - arr.begin();
}
Primer mayor
templT int primer_mayor (vector < T > & arr, T x) {
    auto it = upper_bound(all(arr), x);
    if (it == arr.end()) return -1;
    return it - arr.begin();
}
Último menor
templT int ultimo_menor (vector<T>& arr, T x) {
    if (arr.begin() == arr.end()) return -1;
    auto it = prev(lower_bound(all(arr), x));
    if (*it >=) return -1;
```

return it - arr.begin();

2.1.2. Operaciones de conjuntos y modificación

Funciones que modifican rangos

Función	Params	Ejemplo
copy	first last result	B.resize(A.size()); copy(all(A), B)
fill	first last val	memo.resize(MAXN); fill(all(memo), -1)
rotate	first middle last	<pre>rotate(begin(A), begin(A) + 3, end(A));</pre>

Operaciones de conjuntos con vectors ordenados (lineal)

```
// Siempre hacer resize al final asi:

vector<int> A = { 5, 10, 15, 20, 25};
vector<int> B = {10, 20, 30, 40, 50};

vector<int> U(A.size() + B.size());

auto it = set_union(all(A), all(B), begin(U));

U.resize(it - U.begin());
```

Función	Descripción
set_union	Unión
set_intersection	Intersección
set_difference	Elementos que están en el primero y no en el segundo
set_symmetric_difference	Elementos que están en uno pero no los dos (como el xor)

2.2. Set indexado (policy based)

Set indexado

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;

templT struct SetIndexado {
    tree<
        T, null_type, less<T>,
        rb_tree_tag, tree_order_statistics_node_update
    > s;
    void add (T x) { s.insert(x); }
    int idx (T x) { return s.order_of_key(x); }
    bool has (T x) { return s.find(x) != ms.end(); }
    T ith (int i) { return *s.find_by_order(i); }
};
```

Multiset indexado

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
```

```
using namespace __gnu_pbds;
templT struct MultisetIndexado {
    int t = 0: tree<</pre>
        pair<T, int>, null_type, less<pair<T, int>>,
        rb_tree_tag, tree_order_statistics_node_update
    > ms;
    void add (T x) { ms.insert(mp(x, t++)); }
                 x) { return ms.order_of_key(mp(x, -1)); }
         nle (T
         nleq (T x) { return ms.order_of_key(mp(x, INT_MAX)); }
         cnt (T x) { return nleq(x) - nle(x); }
         ith (int i) { return (*ms.find_by_order(i)).fst; }
};
      Truquitos con la STL
2.3.1. Compresión de coordenadas
Para números enteros (con lower_bound)
// Obtener valor original con D[A[i]]
vector<ll> CompCoordenadas (vector<ll>& A) {
   int N = A.size();
   vector<ll> D = A;
   sort(all(D)):
   D.resize(unique(all(D)) - D.begin());
   forn(i, N) A[i] = lower_bound(all(D), A[i]) - D.begin();
   return D;
}
Versión genérica (con map)
templT map <T, int > CompCoordenadas (vector <T > & A) {
   map<T, int> ord;
   int n = 0;
   for (auto v : A) ord[v];
   for (auto& e : ord) e.snd = n++;
   return ord;
}
2.3.2. Intervalos consecutivos (simular cortar un palito son set)
struct IntervalosConsecutivos {
    set <int > I;
    map<int, int> L;
    IntervalosConsecutivos (int i, int j) {
        I.insert(i);
```

I.insert(j);

L[i - i] ++;

}

```
void cortar (int k) {
        int i = *prev(I.lower_bound(k));
        int j = *(I.lower_bound(k));
        L[i - i]--:
        if (L[j-i] == 0) L.erase(j-i);
        L[k - i]++;
        L[j - k] ++;
        I.insert(k);
    int max_intervalo () {
        return (*L.rbegin()).fst;
};
     Range queries
3.1. Range queries comunes
3.1.1. Suma estático (prefix + diff arrays)
Range sum (prefix array)
templT vector<T> prefix_array (vector<T>& A) {
    vector <T> P(A.size());
    P[0] = A[0];
    forn(i, P.size() - 1) P[i+1] = P[i] + A[i+1];
    return P;
}
// Retorna A[i] + ... + A[j]
templT T query_prefix_array (vector<T>& P, int i, int j) {
    T res = P[i];
    if (i > 0) res -= P[i-1];
    return res;
}
Range update (diff array)
templT vector<T> diff_array (vector<T>& A) {
    vector <T> D(A.size());
    D[O] = A[O];
    forn(i, D.size() - 1) D[i+1] = A[i+1] - A[i];
    return D;
}
// Aplica +x en A[i] ... A[j]
templT void update_diff_array (vector < T > & D, int i, unsigned j, T x)
   \hookrightarrow {
   D[i] += x;
    if (j + 1 < D.size()) D[j+1] -= x;</pre>
}
```

3.1.2. Suma dinámico (fenwick tree) Range sum point set using FT = 11; using Fenwick = unordered_map <int , FT>; FT FT_prefix (Fenwick& A, int i) { FT res = 0: for (int j = i; $j \ge 0$; j = j & (j + 1), j--) res += A[j]; } void FT_add (Fenwick& A, int N, int i, FT x) { for (; i < N; i = i | (i + 1)) A[i] += x;FT FT_sum (Fenwick& A, int i, int j) { return FT_prefix(A, j) - FT_prefix(A, i - 1); void FT_set (Fenwick& A, int N, int i, FT x) { FT_add(A, N, i, - FT_sum(A, i, i)); $FT_add(A, N, i, + x);$ } Range add point get using FT = 11; using Fenwick = unordered_map<int, FT>; FT FT_prefix (Fenwick& A, int i) { FT res = 0;for (int j = i; $j \ge 0$; j = j & (j + 1), j--) res += A[j]; return res; void FT_update (Fenwick& A, int N, int i, FT x) { for (; i < N; i = i | (i + 1)) A[i] += x;FT FT_get (Fenwick& A, int i) { return FT_prefix(A, i); void FT_add (Fenwick& A, int N, int i, int j, FT x) { FT_update(A, N, i, x); $FT_update(A, N, j+1, -x);$ } 3.1.3. Range minimum query (RMQ) (sparse table + segment tree) RMQ estático 1D (sparse table) using ST = int; using SparseT = vector<vector<ST>>; SparseT ST_build (vector < ST > & A, int N) { SparseT res(20, vector < ST > (N));

res[0] = A;

forr(w, 1, 20) forn(i, N - (1 << w) + 1)

```
res[w][i] = min(res[w - 1][i], res[w - 1][i + (1 << (w - 1))
    return res;
}
ST ST_rmq (SparseT& S, int i, int j) {
    int w = 63 - \_builtin\_clzll(j - i + 1);
    return min(S[w][i], S[w][j - (1 << w) + 1]);</pre>
}
RMQ + point set (segment tree)
3.2. Segment tree point set
Template
struct STNode {
   // Completar
};
STNode operator * (STNode a, STNode b) {
   // Completar
const STNode ST_ID = {
   // Completar
}
using STree = vector < STNode >;
STree segtree_build (STree& hojas) {
   int N = hojas.size();
   STree S(N \ll 1);
   forn(i, N) S[i + N] = hojas[i];
   for (int i = N - 1; i; i--) S[i] = S[i << 1] * <math>S[i << 1 | 1];
   return S:
}
STNode segtree_query (STree& S, int i, int j) {
   int N = S.size() >> 1;
   STNode res = ST_ID;
   int 1 = i + N;
   int r = j + N + 1;
   for (; 1 < r; 1 >>= 1, r >>= 1) {
      if (1 \& 1) res = res * S[1++];
      if (r \& 1) res = res * S[--r]:
   }
   return res;
}
void segtree_update (STree& S, int i, STNode x) {
   int N = S.size() >> 1;
```

```
S[i += N] = x:
   for (; i > 1; i >>= 1) S[i >> 1] = S[i] * S[i ^ 1];
}
RMQ
struct STNode { int val; };
STNode operator * (STNode a, STNode b) { return (a.val < b.val) ? a :
   \hookrightarrow b; }
const STNode ST_ID = { INT_MAX };
vector < int > A = { ... };
vector < STNode > hojas(A.size());
transform(all(A), begin(hojas), [](int x) { return (STNode){x}; });
STree T = segtree_build(hojas);
XOR.
struct STNode { int val; };
STNode operator * (STNode a, STNode b) { return { a.val ^ b.val }; }
const STNode ST_ID = { 0 };
vector < int > A = { ... };
vector < STNode > hojas(A.size());
transform(all(A), begin(hojas), [](int x) { return (STNode){x}; });
STree T = segtree_build(hojas);
     Grafos
```

4.1. Toposort de un DAG

```
using AdjList = vector<vector<int>>;
vector<int> Toposort (AdjList& G) {
    int N = G.size();
    vector<int> indegree(N), res;
    forn(u, N) for (int v : G[u]) indegree[v]++;
    // Elegir crierio de priorizacion cambiando el orden en el que se
       \hookrightarrow sacan
    // (por defecto el menor)
    using Bag = priority_queue<int, vector<int>, greater<int>>;
    Bag bag;
    forn(u, N) if(indegree[u] == 0) bag.push(u);
    while (bag.size()) {
        int u = bag.top();
        bag.pop();
        res.push_back(u);
        for (int v : G[u]) {
```

```
indegree[v]--;
            if (indegree[v] == 0) bag.push(v);
        }
    }
    return res;
}
4.2. DAG condensado
using AdjList = vector < vector < int >>;
AdjList DAGCondensado (AdjList& G) {
   int N = G.size();
   vector < bool > visitado(N);
   vector < int > orden;
   function < void(int) > get_orden = [&](int u) -> void {
      visitado[u] = true;
      for (int v : G[u]) if (!visitado[v]) get_orden(v);
      orden.pb(u);
   };
   forn(u, N) if (!visitado[u]) get_orden(u);
   reverse(all(orden));
   AdjList T(N);
   forn(u, N) for (int v : G[u]) T[v].pb(u);
   vector<int> comp, raiz(N), raices;
   function < void(int) > extraer_comp = [&](int u) -> void {
      visitado[u] = true;
      comp.pb(u);
      for (int v : T[u]) if (!visitado[v]) extraer_comp(v);
   };
   visitado.assign(N, false);
   for (int u : orden) if (!visitado[u]) {
      extraer_comp(u);
      int r = comp.front();
      for (int v : comp) raiz[v] = r;
      raices.pb(r);
      comp.clear();
   }
   // Opcion 1: hacer compresion de coordenadas: O(nlogn) lento
   int c = 0;
   map < int , int > coords;
   for (int r : raices) coords[r];
   for (auto& e : coords) e.snd = c++;
   AdjList SCC(raices.size());
   forn(u, N) for (int v : G[u]) {
      int ru = coords[raiz[u]], rv = coords[raiz[v]];
```

Floyd-Warshall

templT using Matriz = vector<vector<T>>;

const 11 INF = LLONG_MAX / 4;

```
if (ru != rv) SCC[ru].pb(rv);
  return SCC;
  // Opcion 2: no hacer compresion y devolver raices (rapido)
  // AdjList SCC(N);
  // forn(u, N) for (auto [v, w] : G[u]) {
     // int ru = raiz[u], rv = raiz[v];
     // if (ru != rv) SCC[ru].pb({rv, w});
     // else (RC[ru]) += R(w);
  // }
}
    Bipartite check
```

```
using AdjList = vector<vector<int>>;
bool EsBipartito (AdjList& G) {
    vector < int > color(G.size(), -1);
    color[0] = 0;
    queue < int > bag;
    for (bag.push(0); bag.size();) {
        int u = bag.front();
        bag.pop();
        for (int v : G[u]) {
            if (color[u] == color[v]) return false;
            if (color[v] == -1) {
                color[v] = 1 - color[u];
                bag.push(v);
            }
        }
    }
    return true;
```

Encontrar puentes y articulaciones

```
using AdjList = vector<vector<int>>;
using Edge = pair<int, int>;
pair < vector < Edge >, vector < int >> GetPuentesArticulaciones (AdjList&G)
   \hookrightarrow {
   int N = G.size(), time = 0;
   vector < bool > visitado(N);
   vector <int> tin(N, -1), tlow(N, -1), articulaciones;
   vector < Edge > puentes;
   function < void(int, int) > dfs = [&](int u, int p) -> void {
      visitado[u] = true;
      tin[u] = tlow[u] = time++;
      int hijos = 0;
      for (int v : G[u]) {
```

```
if (v == p) continue;
         if (visitado[v]) tlow[u] = min(tlow[u], tin[v]);
         else {
             dfs(v, u);
             hijos++;
             tlow[u] = min(tlow[u], tlow[v]);
             if (tlow[v] > tin[u]) puentes.pb({u,v});
             if (tlow[v] >= tin[u] && p != -1) articulaciones.pb(u);
         }
      if (p == -1 && hijos > 1) articulaciones.pb(u);
   };
   forn(r, N) if (!visitado[r]) dfs(r, -1);
   return mp(puentes, articulaciones);
}
4.5. Menores caminos
Dijkstra
struct Hedge { ll weight; int node; };
bool operator < (const Hedge& a, const Hedge& b) { return a.weight >
    \hookrightarrow b.weight; }
using AdjList = vector < vector < Hedge >>;
void Dijkstra (AdjList& G, int s, vector<11>& dist, vector<int>&
    → parent) {
   int N = G.size();
   dist.assign(N, LLONG_MAX);
   dist[s] = 0;
   parent.assign(N, -1);
   parent[s] = s;
   priority_queue < Hedge > bag;
   for (bag.push({0, s}); bag.size();) {
      auto [d, u] = bag.top();
      bag.pop();
      if (d > dist[u]) continue;
      for (auto [w, v] : G[u]) {
         11 \text{ relax} = d + w;
         if (relax < dist[v]) {</pre>
             dist[v] = relax;
             parent[v] = u;
             bag.push({relax, v});
         }
```

```
void FloydWarshall (Matriz<11>& D) {
   int N = D.size();
   forn(u, N) D[u][u] = 0;
   forn(k, N) forn(u, N) forn(v, N) if (D[u][k] < INF) if (D[k][v] <
      \hookrightarrow INF)
      D[u][v] = min(D[u][v], D[u][k] + D[k][v]);
   // Opcional: chequear ciclos negativos
   forn(u, N) forn(v, N) forn(k, N) if (D[u][k] < INF && D[k][k] < 0
      \hookrightarrow && D[k][v] < INF)
      D[u][v] = -INF;
}
    Programacion Dinamica
5.1. LIS (Longest Increasing Subsequence)
Obtener largo del LIS
// Usa compresion de coordenadas y segtree point set RMQ (tomar el
   \hookrightarrow maximo)
int LIS (vector<int>& A) {
   int N = A.size();
   auto C = Compress(A);
   vector < STNode > hojas(N, {0});
   STree dp = segtree_build(hojas);
   segtree_update(dp, C[A[0]], {1});
   forr(i, 1, N) {
      int x = C[A[i]];
      int subres = 0;
      if (x-1 \ge 0) subres = segtree_query(dp, 0, x-1).val;
      segtree_update(dp, x, {1 + subres});
   }
   return segtree_query(dp, 0, N - 1).val;
Construir LIS lexicograficamente menor
struct STNode { int len, idx, val, parent; };
bool operator < (STNode a, STNode b) {</pre>
   if (a.len != b.len) return a.len < b.len:</pre>
   return a.val > b.val;
STNode operator * (STNode a, STNode b) { return max(a,b); }
const STNode ST_ID = { -INT_MAX, -1, INT_MAX, -1 };
```

vector<int> LIS (vector<int>& A) {

```
int N = A.size();
   auto C = Compress(A);
   STNode def = \{0, -1, INT_MAX, -1\};
   vector < STNode > hojas(N, def);
   STree dp = segtree_build(hojas);
   vector < STNode > res(N);
   res[0] = \{1, 0, A[0], -1\};
   segtree_update(dp, C[A[0]], {1, 0, A[0], -1});
   forr(i, 1, N) {
      int x = C[A[i]];
      STNode subres = def;
      if (x-1 \ge 0) subres = segtree_query(dp, 0, x-1);
      STNode r = {1 + subres.len, i, A[i], subres.idx};
      segtree_update(dp, x, r);
      res[i] = r:
   }
   vector < int > path;
   STNode best = *max_element(all(res));
   STNode x:
   for (x = best; x.parent != -1; x = res[x.parent]) path.pb(x.idx);
   path.pb(x.idx);
   reverse(all(path));
   return path;
}
LIS en arbol (largo del LIS de raiz a cada nodo)
// Usa compresion de coordenadas y segtree point set RMQ (tomar el
   \hookrightarrow maximo)
using AdjList = vector < vector < int >>;
vector<int> LIS (AdjList& G, vector<int>& valor_nodo, int root = 0) {
   int N = valor_nodo.size();
   auto C = Compress(valor_nodo);
   STNode def = { 0 };
   vector < STNode > hojas(N, def);
   STree dp = segtree_build(hojas);
   vector<int> res(N);
   segtree_update(dp, C[valor_nodo[root]], {1});
   function < void(int) > dfs = [&](int u) {
      int x = C[valor_nodo[u]];
      int old = segtree_query(dp, x, x).val;
      int subres = \{0\};
      if (x-1 \ge 0) subres = segtree_query(dp, 0, x-1).val;
      segtree_update(dp, x, {1 + subres});
```

```
Notebook
      res[u] = segtree_query(dp, 0, N-1).val;
      for (int v : G[u]) dfs(v);
      segtree_update(dp, x, {old});
   };
   dfs(root);
   return res;
    Matemática
6.1. Aritmética
Techo de la división
       #define ceildiv(a,b) ((a + b - 1) / b)
Piso de la raiz cuadrada
using ll = long long;
ll isqrt (ll x) {
    11 s = 0;
    for (11 k = 1 \ll 30; k; k \gg 1)
        if ((s+k) * (s+k) <= x) s += k;
    return s;
}
Piso del log2
       #define log2fl(x) (x ? 63 - __builtin_clzll(x) : -1)
Aritmética en Zp
const 11 mod = 1e9 + 7;
11 resta_mod (11 a, 11 b) { return (a - b + mod) % mod; }
ll pow_mod (ll x, ll n) {
    11 \text{ res} = 0;
    while (n) {
        if (n % 2) res = res * x % mod;
        n /= 2;
```

11 div_mod (11 a, 11 b) { return a * pow_mod(b, mod - 2) % mod; }

x = x * x % mod;

} return res;

}

6.2. Teoria de numeros

```
Criba
```

```
struct Criba {
    bool c[1000001]; vector<int> p;
    Criba () {
        p.reserve(1<<16);
        for (int i = 2; i <= 1000000; i++) if (!c[i]) {</pre>
             p.pb(i);
            for (int j = 2; i*j \le 1000000; j++) c[i*j] = 1;
        }
    }
    bool isprime (int x) {
        for (int i = 0, d = p[i]; d*d \le x; d = p[++i])
            if (!(x % d)) return false;
        return x \ge 2;
    }
};
Phollards Rho
11 gcd(ll a, ll b){return a?gcd(b %a, a):b;}
//COMPILAR CON G++20
ll mulmod(ll a, ll b, ll m) {
 return ll(__int128(a) * b % m);
}
ll expmod (ll b, ll e, ll m){\frac{1}{0}} \log b
        if(!e) return 1;
        11 q= expmod(b,e/2,m); q=mulmod(q,q,m);
        return e %2? mulmod(b,q,m) : q;
}
bool es_primo_prob (ll n, int a)
        if (n == a) return true;
        11 s = 0, d = n-1;
        while (d \% 2 == 0) s++, d/=2;
        11 x = expmod(a,d,n);
        if ((x == 1) || (x+1 == n)) return true;
        forn (i, s-1){
                 x = mulmod(x, x, n);
                 if (x == 1) return false;
                 if (x+1 == n) return true;
        return false:
```

```
bool rabin (ll n){ //devuelve true si n es primo
                       return false;
        if (n == 1)
        const int ar[] = \{2,3,5,7,11,13,17,19,23\};
        forn (j,9)
                if (!es_primo_prob(n,ar[j]))
                        return false;
        return true;
}
11 \text{ rho}(11 \text{ n})
    if( (n & 1) == 0 ) return 2;
    11 x = 2 , y = 2 , d = 1;
    11 c = rand() % n + 1;
    while( d == 1 ){
        x = (mulmod(x, x, n) + c) %n;
        y = (mulmod(y, y, n) + c) %n;
        y = (mulmod(y, y, n) + c) %n;
        if(x - y >= 0) d = gcd(x - y, n);
        else d = gcd(y - x, n);
    return d==n? rho(n):d;
map<ll,ll> prim;
void factRho (ll n){ //O (lg n)^3. un solo numero
        if (n == 1) return;
        if (rabin(n)){
                prim[n]++;
                return;
        11 factor = rho(n);
        factRho(factor);
        factRho(n/factor);
}
```

6.3. Geometria

Template geometria

```
Sca operator * (Vec a, Vec b) { return a.x * b.x + a.y * b.y; }
Sca operator ^ (Vec a, Vec b) { return a.x * b.y + a.y * b.x; }
bool operator < (Vec a, Vec b) { return (a.x != b.x) ? (a.x < b.x) :
   \hookrightarrow (a.y < b.y); }
ostream& operator << (ostream &o, Vec& p) { auto x = mp(p.x, p.y);
   → return o << x; }</pre>
Sca norma2 (Vec p) { return p.x * p.x + p.y * p.y; }
struct pto{
        11 x, y; int t;
        pto(ll x=0, ll y=0, int t = -1):x(x),y(y), t(t){}
        pto operator-(pto a){return pto(x-a.x, y-a.y);}
        11 operator*(pto a){return x*a.x+y*a.y;}
        11 operator^(pto a){return x*a.y-y*a.x;}
        bool left(pto q, pto r){return ((q-*this)^(r-*this))>0;}
        bool operator < (const pto &a) const{return x < a.x | | (x == a.x &&
            \hookrightarrow v < a.v);}
  bool operator == (pto a) {return x == a.x && y == a.y;}
};
//stores convex hull of P in S, CCW order
//left must return >=0 to delete collinear points!
void CH(vector<pto>& P, vector<pto> &S){
        S.clear();
        sort(P.begin(), P.end());//first x, then y
        forn(i, sz(P)){//lower hull
                 while (sz(S) \ge 2 \&\& S[sz(S) - 1] . left(S[sz(S) - 2], P[i]))
                    ⇔ S.pop_back();
                 S.pb(P[i]);
        S.pop_back();
        int k=sz(S);
        dforn(i, sz(P)){//upper hull
                 while(sz(S) >= k+2 \&\& S[sz(S)-1].left(S[sz(S)-2], P[i])
                    → ])) S.pop_back();
                 S.pb(P[i]);
        S.pop_back();
}
```

7. Estructuras locas

7.1. Disjoint set union

```
struct DSU {
    vector < int > p, w; int nc;
    DSU (int n) {
        nc = n, p.resize(n), w.resize(n);
        forn(i,n) p[i] = i, w[i] = 1;
```

```
int get (int x) { return p[x] == x ? x : p[x] = get(p[x]); }
    void join (int x, int y) {
        x = get(x), y = get(y);
        if (x == y) return;
        if (w[x] > w[y]) swap(x,y);
        p[x] = y, w[y] += w[x];
    bool existe_camino (int x, int y) { return get(x) == get(y); }
};
7.2. Binary trie
struct BinaryTrieVertex { vector<int> next = {-1, -1}; };
using BinaryTrie = vector < BinaryTrieVertex >;
void binary_trie_add (BinaryTrie& trie, int x) {
    int v = 0;
    for (int i = 31; i >= 0; i--) {
        bool b = (x & (1 << i)) > 0;
        if (trie[v].next[b] == -1) {
            trie[v].next[b] = trie.size();
            trie.emplace_back();
        v = trie[v].next[b];
    }
}
int binary_trie_max_xor (BinaryTrie& trie, int x) {
    int v = 0, res = 0;
    for (int i = 31; i >= 0; i--) {
        bool b = (x & (1 << i)) > 0;
        if (trie[v].next[!b] != -1) {
            v = trie[v].next[!b];
            if (!b) res |= (1 << i);
       }
        else {
            v = trie[v].next[ b];
            if (b) res |= (1 << i);
    } return res;
// Inicializar asi:
BinaryTrie trie(1);
```

8. Sin categorizar

Búsqueda binaria sobre un predicado

```
int primer_true (int i, int j, function < bool(int) > P, int def) {
   while (j - i > 1) {
      int m = i + ((j - i) >> 1);
      P(m) ? j = m : i = m;
   if (P(i)) return i;
   if (P(j)) return j;
   return def;
}
int ultimo_false (int i, int j, function < bool(int) > P, int def) {
   while (i - i > 1) {
      int m = i + ((j - i) >> 1);
      P(m) ? j = m : i = m;
   if (!P(j)) return j;
   if (!P(i)) return i;
   return def;
}
Enumerar subconjuntos de un conjuto con bitmask
// Imprimir representaciones en binario de todos los numeros "[0,
    \hookrightarrow .... 2^N-1]"
forn(mask, (1 << N)) {
    forn(i, N) cout << "01"[(mask & (1 << i)) > 0] << "\0\n"[i == N]
}
// Iterar por los bits de cada subconjunto
forn(mask, (1 << N)) {
    forn(i, N) {
        bool on = (mask & (1 << i)) > 0;
        if (on) { ... }
        else { ... }
}
Hashing Rabin Karp
using ll = long long;
const ll primo = 27, MAX_PRIME_POW = 1e6;
11 prime_pow[MAX_PRIME_POW];
void get_prime_pow () {
    prime_pow[0] = 1;
    forn(i, MAX_PRIME_POW) prime_pow[i+1] = prime_pow[i] * primo %
        \hookrightarrow mod:
}
vector<ll> get_rolling_hash (string& s) {
```

9. Brainstorming

- Graficar como puntos/grafos
- Pensarlo al revez
- ¿Que propiedades debe cumplir una solución?
- Si existe una solución, ¿existe otra más simple?
- ¿Hay electiones independientes?
- ¿El proceso es parecido a un algoritmo conocido?
- \blacksquare Si se busca calcular f(x) para todo x, calcular cuánto contribuye x a f(y) para los otros y
- Definiciones e identidades: ¿que significa que un array sea palindromo? (ejemplo)